

Practical - 1

Aim:- Implement linear search to find an item in the list.

Theory:-

Linear search.

Linear search is one of the simplest searching algorithm in which targeted item is sequentially searched with each item in the list. It is worst searching algorithm with worse case time complexity. It is a force approach on the other hand in case of an ordered list, instead of searching the list in sequence. A binary search is used which will start by examining the middle term.

Linear search is a technique to compare each and every element with the key element to be found, if both matches, the algorithm returns that element found and its position is also found.

Unsorted :

Algorithm:

Create an empty list and assign to a variable.

Accept the total no. of elements to be inserted into the list from the user say 'n'.

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Use for loop for adding the elements into the list.

Point the new list.

Accept an element from the user that to be searched in the list.

Use for loop in a range from '0' to the total no. of elements to search the element from the list.

Use if loop that the elements in the list equal to the element accepted from user.

Show the output of the given algorithm

```

input("linear search")
a = []
n = int(input("Enter the range"))
for s in range(0, n):
    s = int(input("Enter a number"))
    a.append(s)
print(a)

c = int(input("Enter a number to search"))
for i in range(0, n):
    if a[i] == c:
        print("Number found at this position")
        break
    else:
        print("Number not found")

```

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OUTPUT:

Enter the range: 3

Enter a number: 1

[1]

Enter a number: 2

[1, 2]

Enter a number: 3

[1, 2, 3]

Enter a number to search: 2

Number found at this position: 1.

2) Seated linear search
 Sorting means to arrange the elements in increasing and decreasing order.

algo:-

create empty list and assign it is a variable
 Accept total no. of elements to be inserted into the list from user, say 'n'.

use for loop for using append() method to add the element in the list.

use sort() method to sort the accepted elements and assign it is in increasing order the list then print the list.

use if statement to give the range in which element is found in given range the display "Element not found".

or use else statement , if element is not found in range then satisfy the given condition.

for loop in range from 0 to the total no. before doing this element to be searched from user using input and search number statement.

tops of trees suitable for
use as fuel ("fuel tree crop")
for use as a hedge (2, 3)
as fuel, pallieter or lumber,
or apparel
fuel (3)
as a fuel (trees of various sizes
for use as a hedge (2, 3)
fuel (2, 3)
fuel ("woodland land at
some distance from
the sea")
fuel ("wood and lumber"
2, 3)
fuel and lumber (3)
fuel (2, 3)

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for the input and output of above
go

Aim:- Implement binary search to find a search no. in the list.

Theory :- BINARY SEARCH

Binary search is also known as half interval search, logarithmic search or binary chop. It is a search algorithm that finds the position of a target value within a sorted array. If you are looking for the number which is at either end of the list then you need to search entire list in linear search, which is time consuming. This can be avoided by using binary fashion search.

Algorithm:-

Create empty list and assign it to a variable. Using input method, accept the range of given list.

Use for loop, add elements in list using append method.

Use sort() method to sort the accepted elements in increasing order. Assign it to a list after sorting.

the range "])

are wrong "])

the numbers to search")

```
a = [int(input("Enter elements of array: ")),  
n = int(input("Enter size of array: ")),  
for b = int(input("Enter element to search: ")),  
    b = input("Enter element to search: ")),  
    a.append(b),  
    a.sort(),  
    print(a),  
    print("Elements are sorted"),  
    s = int(input("Enter element to search: ")),  
    if (s == a[0] or s == a[n-1]):  
        print("Element not found"),  
  
else:  
    f = 0  
    l = n-1  
    for i in range(0, n):  
        m = int((f+l)/2),  
        print(m),  
        if (s == a[m]):  
            print("Element found at", m),  
            break,  
        else if (s < a[m]):  
            l = m-1  
        else:  
            f = m+1
```

- 5) Use if loop to give the range in which element is found in given range, then display a message "element not found".
- 6) Then use else statement, if statement is found in range then satisfy the below condition
- 7) Accept an argument and log of the elements the list as array is starting from 0 hence it is initialize unless than total count.
- 10) If statement in list and still the element to be searched is not found then find the middle element (m).
- 11) Else if the item to be searched is still less than the middle term then Initialize last(l) = mid(m) - 1.
Else Initialize - first (l) = mid (m) - 1.
- 2) Repeat till you found the element Sixth step of above algorithm

Output

>>> enter the range: 3
>>> enter the numbers: 2.
[2]
>>> enter the number: 1
[1, 2].
>>> enter the number: 4
[1, 2, 4].
>>> enter the number to search: 3
>>> element not found

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Bubble sort.

Aim: Implementation of bubble sort program on given list.

Theory: Bubblesort is based on the idea of repeatedly comparing pairs of adjacent elements and then swapping their position if they exist in the wrong order. This is the simplest form of sorting available. In this we sort the given elements in ascending or descending order by comparing two adjacent elements at a time.

Algorithm:

1. Bubble sort algorithm starts by comparing the first two elements of an array and swapping if necessary.

If we want to sort the elements of array in ascending order then first element is greater than second then we need to swap the element.

If the element is smaller than second then we again second third elements are compared and swapped, if it is necessary and this process go on until last and second last element is compared and swapped.

```

print("Bubble sort algorithm"))
a = []
b = int(input("Enter no. of elements"))
for s in range(0, b):
    s = int(input("Enter the element"))
    a.append(s)
print(a)

n = len(a)
for i in range(0, n):
    for j in range(n-1):
        if a[i] < a[j]:
            temp = a[j]
            a[j] = a[i]
            a[i] = temp
print("elements after sorting are:", a)

```

Output:

```

>>> Bubble sort algorithm
Enter number of elements: 3
Enter the element: 8

```

[8]

Enter the element: 9

[8, 9]

Enter the element: 1

[8, 9, 1]

Elements after sorting are: [1, 8, 9]

There are n elements to be sorted then
the process mentioned above should be
repeated $n-1$ times to get the required
result.

State the output and input of above
algorithm of bubble sort stepwise

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Quick Sort

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Aim: Implement Quick Sort to the given list.

Theory: The quick sort is a recursive algorithm based on the divide and conquer technique.

Algorithm:

Step1: Quick sort first selects a value, which is called pivot value, first element serve as our first pivot value. Since we know that first will eventually end up as last in that list.

Step2: The partition process will happen next. It will find the split point and at the same time move our items to the appropriate side of the list, either less than or greater than pivot value.

~~def~~ quick (alist):
 help(alist, 0, len(alist)-1)

def help (alist, first, last):
 if first < last:

split = part (alist, first, last)
 help(alist, first, split-1)
 help(alist, split+1, last)

def part (alist, first, last):

pivot = alist [first]

l = first + 2

r = last

done = False

while not done:

while l <= r and list [l] == pivot:

l = l + 1

while alist [r] >= pivot and r >= l:

r = r - 1

if r < l:

done = True,

else:

t = alist [r]

alist [r] = alist [l]

alist [l] < t.

t = alist [first]

alist [first] = alist [r]

alist [r] = t.

r = ~~return~~ r

a list []

Enter range for the list

Step 5: At the point where rightmark becomes leftmark we stop. The position of rightmark is now the split point.

Step 6: The pivot value can be exchanged with contents of split point and piv is now

Step 8: The quicksort function invokes a recursive function, quicksort(piv)

Step 9: Quick sort helper begins with some base case for range sort.

Step 10: The partition function implements the described partitioning scheme.

```
for b in range(0, x)
    b = int(input("enter the elements"))
    list = append
    n = len(list)
quick(list)
print(list)
```

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Output:

Enter range for the list: 4

Enter the element: 6

Enter the element: 8

Enter the element: 4

Enter the element: 9

[4, 4, 6, 9].

Aim: Implementation of stacks using python list

Theory: A stack is a linear data structure that can be represented in the real world in the stack or cup. The elements in the stack are added or removed only from one position, i.e., the topmost position. Thus, the stack works on the LIFO principle as the element that was inserted last will be removed first. A stack can be implemented using array as well as linked list.

Algorithm:

1. Create a stack with instance variable items.
2. Define init method with self argument and initialize the initial value and the initialize to an empty list.
3. Define methods push and pop Under the class stack.
4. Use if statement to give the function that if length of given list is greater than the range of list then print stack is full.
5. Or else print statement as insert the element into the stack and initialize the values.
6. Push method used to insert the element but pop is used to delete the element from the stack.

abhishek

```
>>> s.push(20)
>>> s.pop()
[20, 0, 0, 0, 0]
>>> s.pop()
Data = 20
>>> s.push(10)
[0, 0, 0, 0, 0]
>>> s.push(20)
>>> s.push(30)
>>> s.push(40)
>>> s.push(50)
>>> s.pop()
[0, 20, 30, 40, 50]
```

- 1. Push method used to insert the element.
- 2. Pop method used to delete the element by stack.
- 3. If it is pop method, value is less than zero and stack is empty or else delete the element from stack at topmost position.
- 4. First and checks whether the no. of elements
- 5. Attach the input and above algorithm.

```

print "abhishek"
class Stack:
    global tos
    def __init__(self):
        self.l = [0, 0, 0, 0, 0, 0]
        self.tos = 1
    def push(self, data):
        n = len(self.l)
        if self.tos == n - 1:
            print "stack is full."
        else:
            self.tos = self.tos + 1
            self.l[self.tos] = data
    def pop(self):
        if self.tos < 0:
            print "stack is empty."
        else:
            k = self.l[self.tos]
            print "Data = ", k
            self.tos = self.tos - 1
    def peek(self):
        if self.tos < 0:
            print ("stack empty")
        else:
            p = self.l[self.tos]
            print ("top element = ", p)

```

S = stack()

Implementing a queue using Python list

Ques is a linear data structure which has 2 differences from and rears. Implementing a queue using list is the simplest as the python list provides built functions to perform the specified operations of queue. It specifies operations of the queue. It based on the principle that a new element is inserted rear and element at front is deleted while is front. In simple terms, a queue can be described as a data structure based on first in out principle.

len(): Creating a new empty queue.

enq(): Insert an element at the rear of queue, and similar to that of insert, ored using tail.

deq(): Returns the element which was at s. The front is moved to successive element. Operation cannot remove element if it is empty.

```
class queue:  
    global r  
    global f  
    def __init__(self):  
        self.r = 0  
        self.f = 0  
        self.l = (0, 0, 0)  
    def enqueue(self, data):  
        n = len(self.l)  
        if self.r < n:  
            self.l[self.r] = data  
            self.r += 1  
            print("element inserted - ", data)  
        else:  
            print("queue is full")  
    def dequeue(self):  
        n = len(self.l)  
        if self.f < n:  
            print(self.l[self.f])  
            self.l[self.f] = self.l[n]  
            self.f += 1  
            print("Element deleted - ...")  
        else:  
            print("queue is empty")
```

Algorithm:

- Step1: Define a class queue and assign global variable to the deque's init() method with self argument. In init(), assign or initialize the front and rear with the help of the self arguments.
- Step2: Define an empty list and define enqueue() function with x arguments assign the length of empty list or arguments assign the length of empty list.
- Step3: Use if Statement front length is equal to rear queue is full or else insert the element from x or display that queue element added successfully increment by 1.
- Step5: Now call the dequeue() function and give the element that has to be added in the empty list by using enqueue() and point the list after adding and same for deleting.

Output:

- ⇒ Q-add (10)
element insert ... 10
- ⇒ Q-add (20)
element insert ... 20
- ⇒ Q-add (3)
element insert ... 3
- ⇒ Q-add (4)
queue is full
- ⇒ Q-remove
10 element deleted

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Aim: Program on postfix expression.

using stack in environment.

The postfix expression is free of any parentheses. We take one possibility of the operation in the program. A given postfix expression can easily be evaluated using stacks.

Algorithm:

Step 1: Define evaluate as function then create an empty stack in python.

Step 2: Convert the string to a list by using the string method split.

```

def evaluate(S):
    k = S.split(' ')
    n = len(k)
    stack = []
    for i in range(n):
        if k[i] == digit():
            stack.append(int(k[i]))
        elif k[i] == '+' or digit():
            stack.append(int(k[i])))
        elif k[i] == '-':
            stack.append(int(k[i])))
        elif k[i] == '*':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) * int(a))
        elif k[i] == '/':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) / int(a))
    else:
        a = stack.pop()
        b = stack.pop()
        stack.append(int(b) / int(a))
    return stack.pop()

S = "3 * 5 + 2"
n = evaluate(S)
print("The evaluated value is:", n)
print("abhishek")

```

1. Perform the arithmetic operation. Push the result back on the stack.
2. After the input expression has been completely processed, the result is on the stack.
3. Print the result of storing after the eval. of positive.
4. Attach the output and input of the above alg.

output: The evaluated value is π . 50
Abstract:

1. π is a constant number.
2. π is an irrational number.

3. π is a transcendental number.
4. π is a real number.

5. π is a rational number.
6. π is a complex number.

7. π is a rational number.
8. π is a complex number.

9. π is a rational number.
10. π is a complex number.

11. π is a rational number.
12. π is a complex number.

13. π is a rational number.
14. π is a complex number.

15. π is a rational number.
16. π is a complex number.

17. π is a rational number.
18. π is a complex number.

19. π is a rational number.
20. π is a complex number.

21. π is a rational number.
22. π is a complex number.

23. π is a rational number.
24. π is a complex number.

25. π is a rational number.
26. π is a complex number.

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Aim: Implementation of Single Linked List by adding the nodes from user's deposition.

Algorithm:

Traversing of linked list means all the nodes in the linked list is in order to perform some operation on them.

The entire linked list means can be accessed using first nodes of the linked list i.e. first node of the linked list in term of referred by the pointer of the linked list.

Thus the entire linked list can be traversed using the nodes which is referred by head pointer.

We should not use the head pointer to traverse

all code:

```
class node:
    global data
    global next
    def __init__(self, item):
        self.data = item
        self.next = None

class linked list:
    global S
    def __init__(self):
        self.S = None
        self.S = newnode
    else:
        head = self.S
        while head.next != None:
            head = head.next
            head.next = newnode
    def addB(self, item):
        newnode = node(item)
        if self.S == None:
            self.S = newnode
        else:
            newnode.next = self.S
            self.S = newnode

    def display(self):
        head = self.S
        while head.next != None:
            print(head.data)
            head = head.next

    def delete(self):
        head = self.S
        if self.S == None:
            print("List is empty")
        else:
            head = self.S
            while head.next != None:
                if head.next.next == None:
                    head.next = None
                    break
                head = head.next
```

- 2
- i) Similarly traverse the rest of W3G of nodes and
traversed using while loop.
 - ii) Attach the coding of I/P or O/P of the
attained algorithm

```
s = linkedlist()
s.add(50)
s.add(L(80))
s.add(L(70))
s.add(L(85))
s.add(L(40))
s.add(B(30))
s.add(20)
s.display()
```

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output:

20

30

40

50

60

70

80

Aim:- Implementation of merge sort

Theory:- Like quicksort, merge sort is an divide and conquer algorithm. It divides I/P array into halves call itself for the two halves and then merges the 2 sorted halves. The merge is done. is used for merging two halves. The merge (arr, l, m, r) is long process that assures that $arr[l:m]$ and $arr[m+1:r]$ are sorted and merges the two sorted sub-arrays into one. Once the size becomes 1, the merge process comes into action and starts merging array books till the complete array is merged.

Applications:

- 1) Merge cost is useful for sorting linked lists.
- 2) Merge sort access data sequentially and therefore random access is low.

Code:

```
def mergesort(arr):
    if len(arr) > 1:
        mid = len(arr) // 2
        left_half = arr[:mid]
        right_half = arr[mid:]
        merge_sort(left_half)
        merge_sort(right_half)
        i = j = k = 0
        while i < len(left_half) and j < len(right_half):
            if left_half[i] < right_half[j]:
                arr[k] = left_half[i]
                i += 1
            else:
                arr[k] = right_half[j]
                j += 1
            k += 1
        while i < len(left_half):
            arr[k] = left_half[i]
            i += 1
            k += 1
        while j < len(right_half):
            arr[k] = right_half[j]
            j += 1
            k += 1
```

Inversion count problem

- ③ Used in external sorting.

Merge sort is more efficient than quick for some types of lists if the data to be sorted can only be efficiently accessed sequentially as is the case for sequential access data structure like arrays.

```

arr = [27, 89, 70, 55, 62, 99, 43, 14, 10]
print ("random list : ", arr)
momesort (arr)
print ("In merge sorted list : ", arr)

```

Output:

random list : [77, 81, 70, 55, 62, 99, 43, 14, 10]
~~MERGESORTED~~ [10, 14, 27, 43, 55, 62, 81, 99]

Practicals

Aim: Implementation of sets using python.

Algorithm:

- i) Define 2 empty sets 0 and 1 and set 2 now providing the range of above 2 sets.
- ii) Now, add() method used for addition the elements according to given range then print the after addition.
- iii) Find the Union and Intersection of above 2 sets by using h(and), i(or) methods. Print the sets of union and intersection of set 3.
- iv) Use if statements to find out the subset and superset of set 3 and set 4. Display the above sets.
- v) Display that element in set 3 is not in set 4 using mathematical operation
- vi) Use disjoint() to check that anything is common in element is present or not. If not then display that it's mutually exclusive

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code:

```
print ("Abhishek 1843")
set1 = set()
set2 = set()
for i in range (8, 15)
    set1.add(i)
for i in range (1, 12)
    set2.add(i)
print ("set1:", set1)
print ("set2:", set2)
print ("\n")
set3 = set1 | set2
print ("Union of set1 and set2: Set3", set3)
set4 = set1 & set2
print ("Intersection of set1 and set2: Set4", set4)
print ("\n")
if set3 > set4:
    print ("Set3 is superset of set4")
elif set3 < set4:
    print ("Set3 is same as set4")
else:
    print ("Set3 is same as set4")
if set4 < set3:
    print ("Set4 is subset of set3")
print ("\n")
set5 = set3 - set4
print ("Elements of set3 and not in set4: set5")
print ("\n")
```

if set u is disjoint (set & s).
print ("set U and set S are naturally
exclusive in")

Set S. clear ()
print ("after applying clear set S is empty")
point ("set S", setS)

Output:

>>> Abhijit 1843
set 1: { 8, 9, 10, 11, 12, 13, 14 }
set 2: { 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11 }
Union of set 1 and set 2: sets { 1, 2, 3, 4, 5, 6,
8, 9, 10, 11, 12, 13, 14 }

Intersection of set 1 and set 2: set 4 { 8, 9, 10,
11 }

set 3 is superset of set 4
set 4 is subset of set 3

Aim: Program based on binary search by implementing in order, pre-order and post-order traversal.

Theory:

Binary tree is a tree which supports maximum of 2 children for any node within the tree, thus any particular node can have 0 or 1 or 2 children such that it is ordered such that one child is identified as left child and other as right child.

Order: i) Transverse the left subtree, the left subtree intern might have left and right subtrees.

Visit the root node.

Transverse the left ~~subtree~~. The left subtree doesn't have left and right subtree.

Transverse the right subtree repeat it.

Post order: i) Transverse the left subtree. The left subtree then right have left and right subtree.

Transverse the right subtree
and the root node.

Class node:

global π

global α

def __init__(self, IP, E):

SP1F.l = None

SP1F.delta = l

SP1F.r = None

class tree:

global root

def __init__(self):

if self.root == None

self.root = node(va)

else:

newnode = node(val)

h = self.root

while True:

if newnode.data < h.data:

if h.l == None:

h.l = newnode

else:

h.l = newnode

print(newnode.data, "data on",

left of", h.data)

break

else:

if h.r == None:

h.r = newnode

else:

h.r = newnode

print(newnode.data, "added",

to right of",

h.data)

break

Algorithm

Define class node and define .init() method with 2 arguments. Initialize the value is the method.

Again define a class BST that is Binary tree with init() method with split argument and assign the root is none.

define add() method for adding the node. define a variable p that $p = \text{node}(\text{value})$

Use if statement for checking the condition that root is none then use else statement. For if node is less than the main node then put or assign that in left side.

Use while loop for checking the node is less or greater than the main node and break loop if it is not satisfying.

if statement within that else statement check that node is greater than main root so put it into right side.

```

det preOrder (self, start)
if start != None:
    print (start.data)
    self.Preorder (start.l)
    self.Preorder (start.r)

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det inorder (self, start):
if start != None:
    self.inorder (start.l)
    print (start.data)
    self.inorder (start.r)

det postorder (self, start):
if start != None:
    self.inorder (start.l)
    self.inorder (start.r)
    print (start.data)

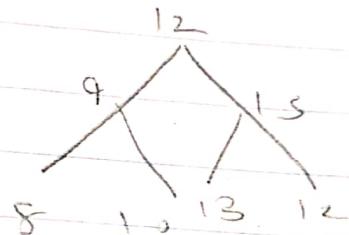
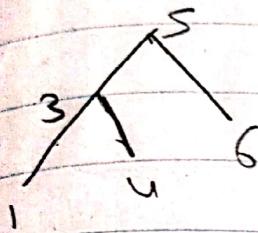
```

$T = \text{Tree}()$
 $T = \text{add}(10)$
 $T = \text{add}(80)$
 $T = \text{add}(10)$
 $T = \text{add}(85)$
 $T = \text{add}(10)$
 $T = \text{add}(48)$
 $T = \text{add}(60)$
 $T = \text{add}(80)$
 $= \text{add}(80)$
 $= \text{add}(15)$
 $\cdot \text{add}(12)$
print ("preorder")
+ preorder (T.root)
** ("inorder")
inorder (T.root)
+ ("~~postorder~~ Abfrage")

In inorder, the statement first left, root and then right node used for giving that condition.

For postorder. In right and then go else for root, assign left then

Display the output ad input of above algorithm
In order: (LVR).



~~3 4 5 6 7 8 9 10 11 12 13 15~~

3 4 5 6 7 8 9 10 11 12 13 15

Output:

80 added on left of 100

70 added on left of 80

65 added on left of 70

10 added on left of 10

preorder

100

80

70

10

15

17

85

inorder

10

15

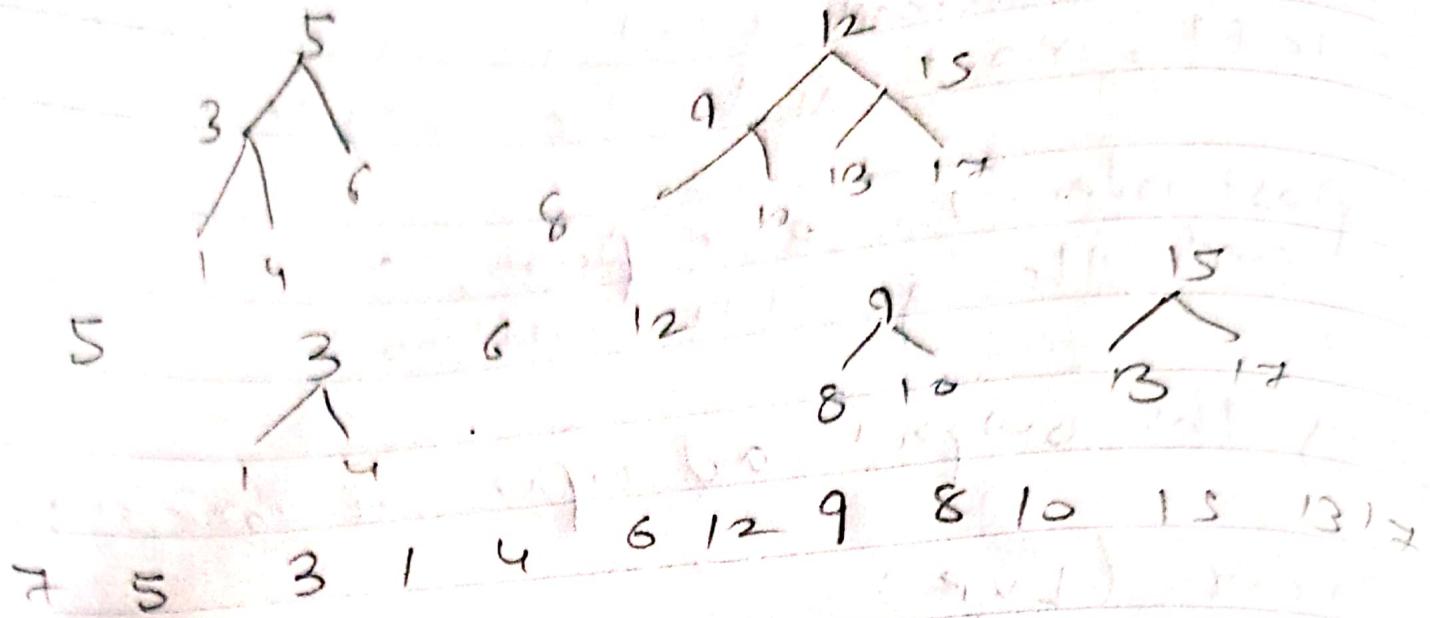
17

70

85

100

Preorder (VLR)



post order (LRV)

