

HIGH PERFORMANCE STENCIL CODE GENERATION WITH LIFT

Bastian Hagedorn | Larisa Stoltzfus | Michel Steuwer | Sergei Gorlatch | Christophe Dubach



WWU
MÜNSTER

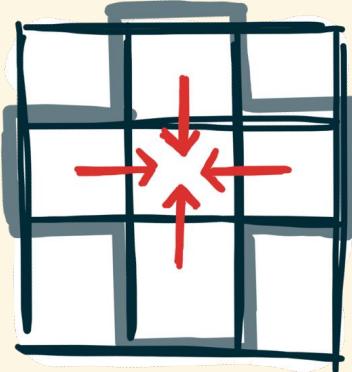


THE UNIVERSITY
of EDINBURGH



University
of Glasgow

WHY STENCIL COMPUTATIONS?



Stencil computations are a class of kernels which update *neighboring* array elements according to a fixed pattern, called *stencil*.

Frequently occur in:



Medical Imaging



Machine Learning



Physics Simulations



PDE Solvers

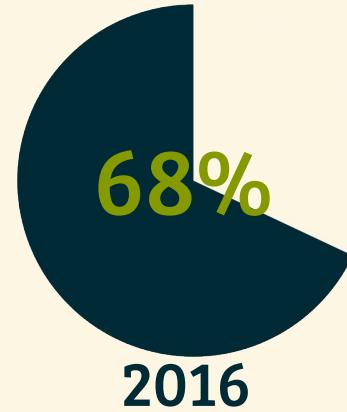
WHY STENCIL COMPUTATIONS?

Stencil compute time:

HPC Center
München



HPC Center
Stuttgart



Frequently occur in:



Medical Imaging



Machine Learning



Physics Simulations



PDE Solvers

YET ANOTHER STENCIL PAPER?



2005

2007

2009

2011

2013

2015

2018

...

ICS'05

ICS'09

CGO'12

CGO'15

CLUSTER'17

PLDI'07

SC'10

CLUSTER'13

WOLFHPC'16

CGO'18

DOMAIN SPECIFIC LANGUAGES

PATUS

Pochoir

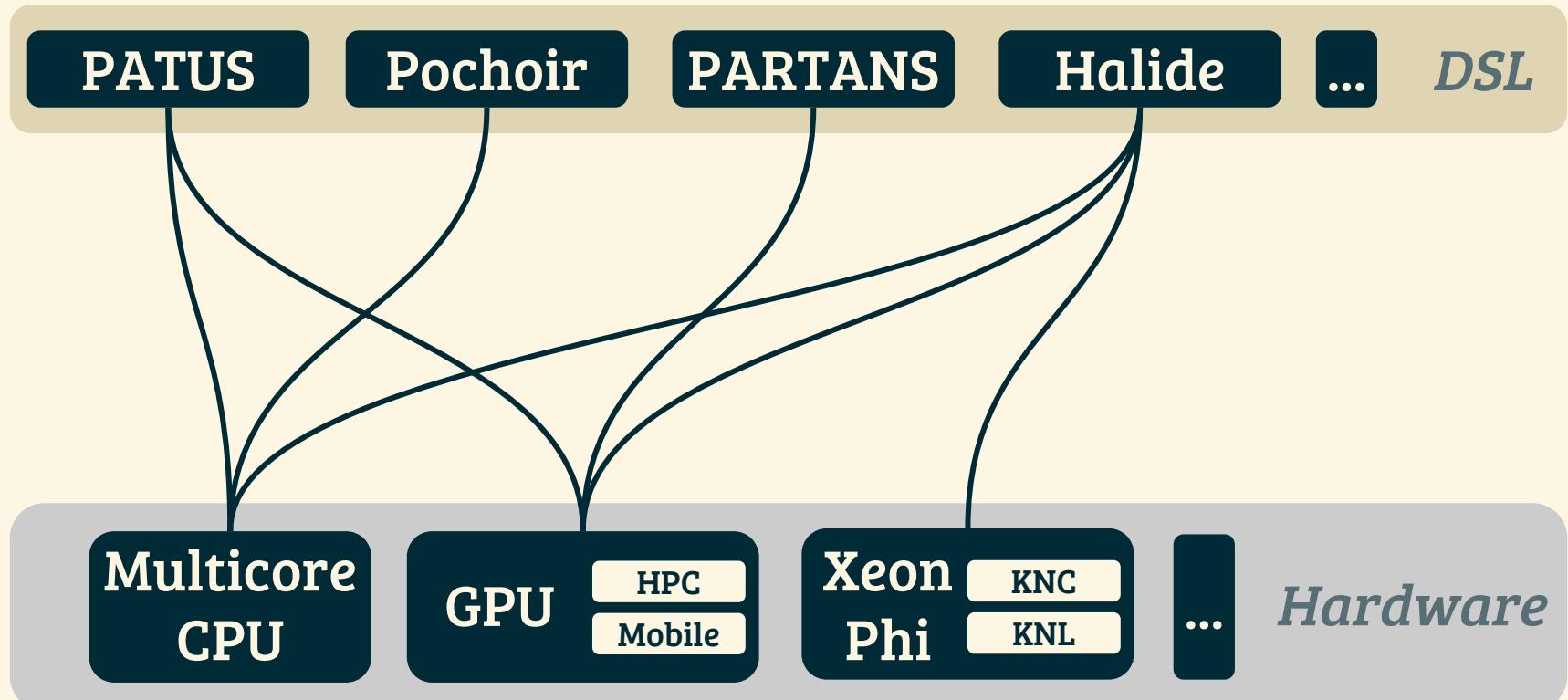
PARTANS

Halide

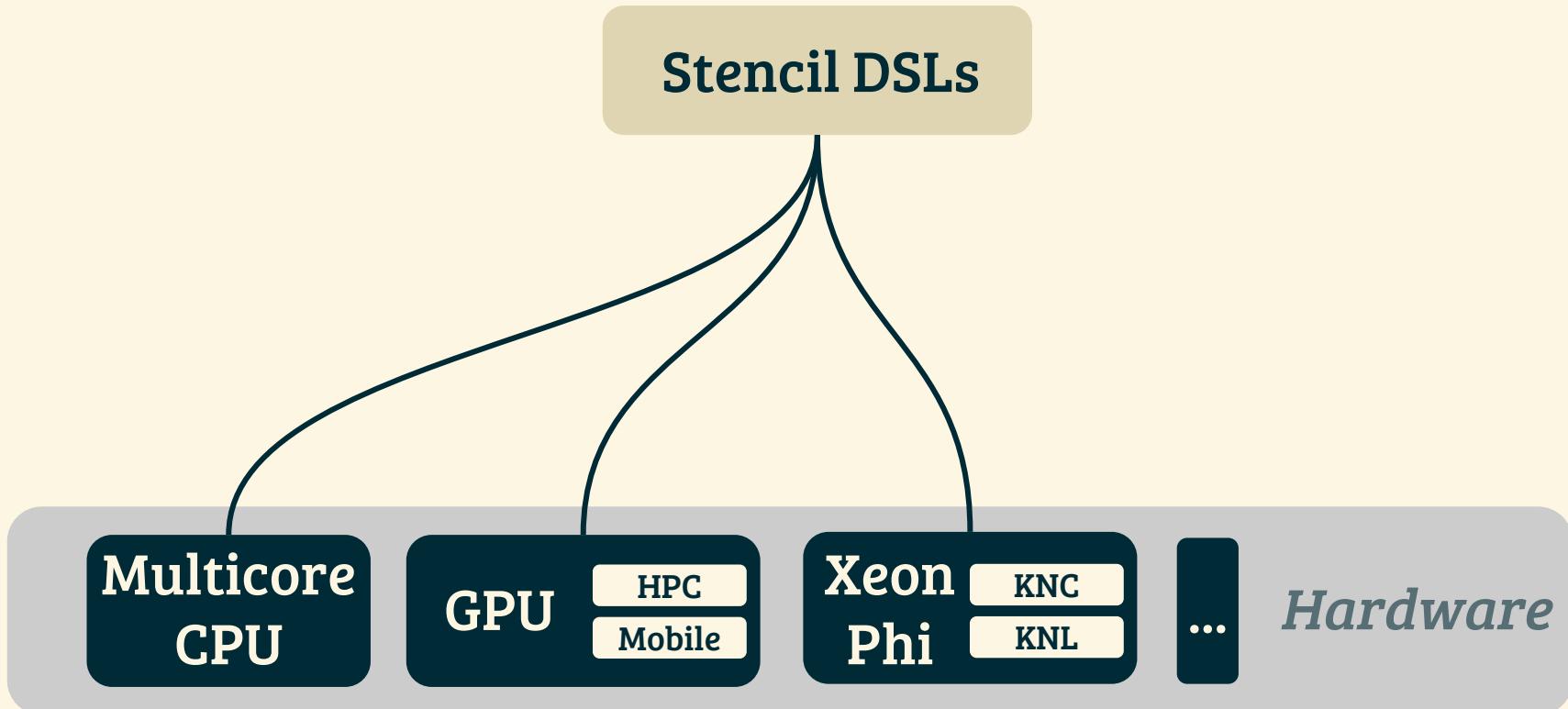
...

DSL

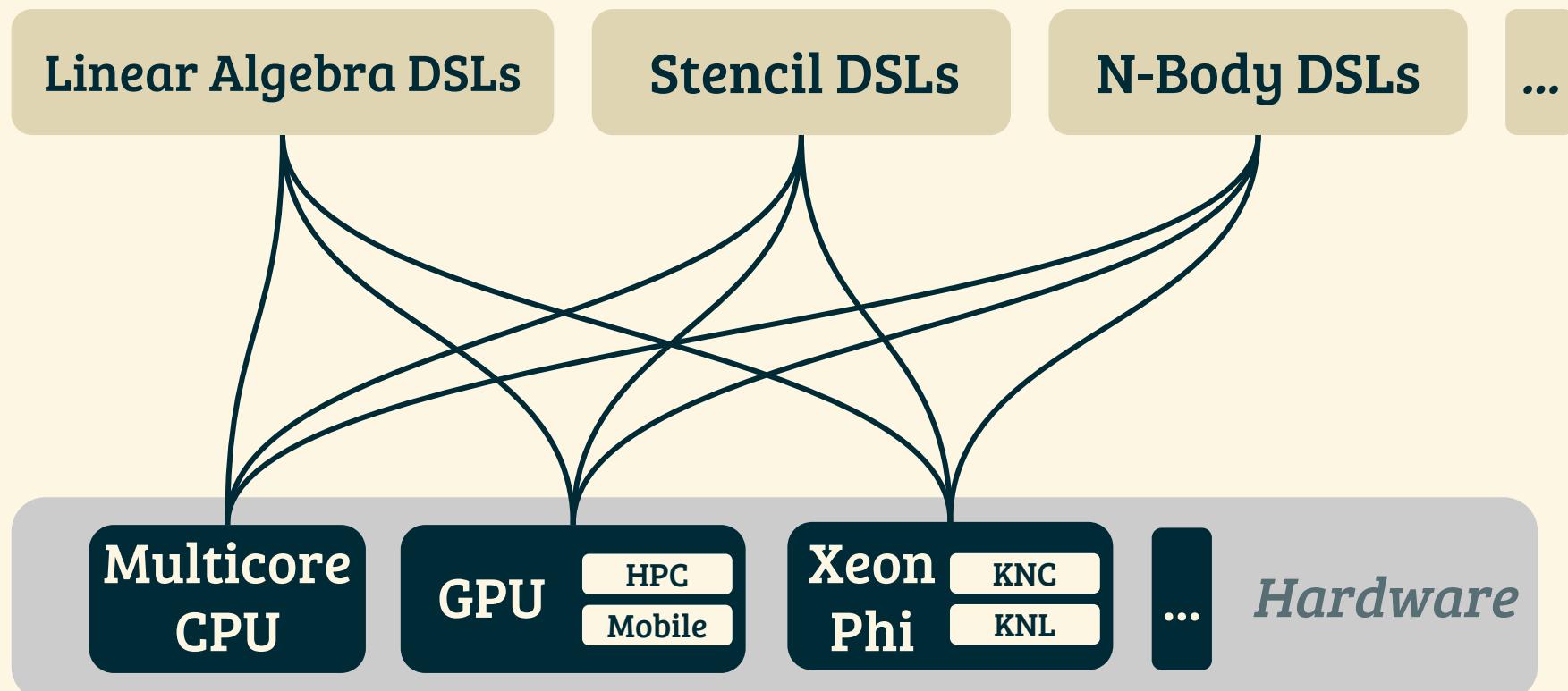
EXPLOITING DOMAIN KNOWLEDGE



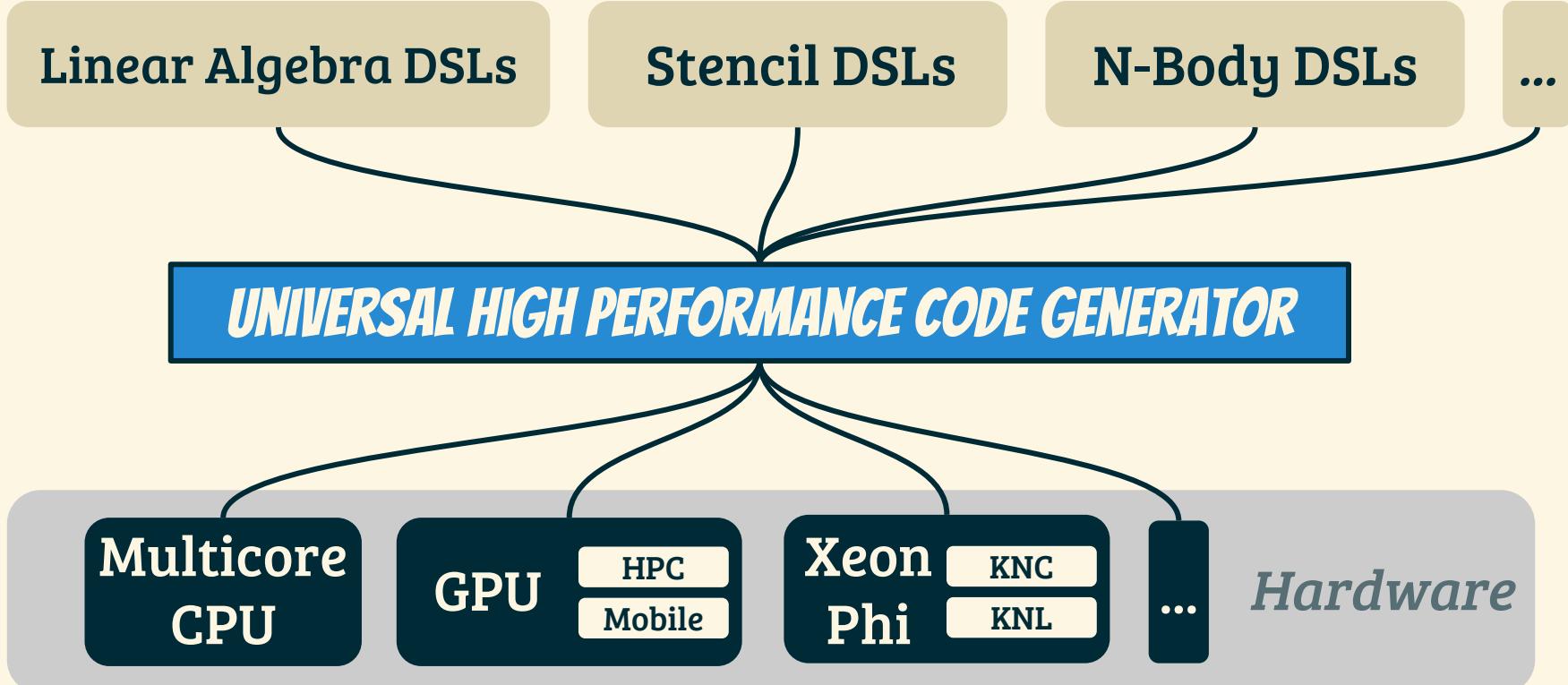
EXPLOITING DOMAIN KNOWLEDGE



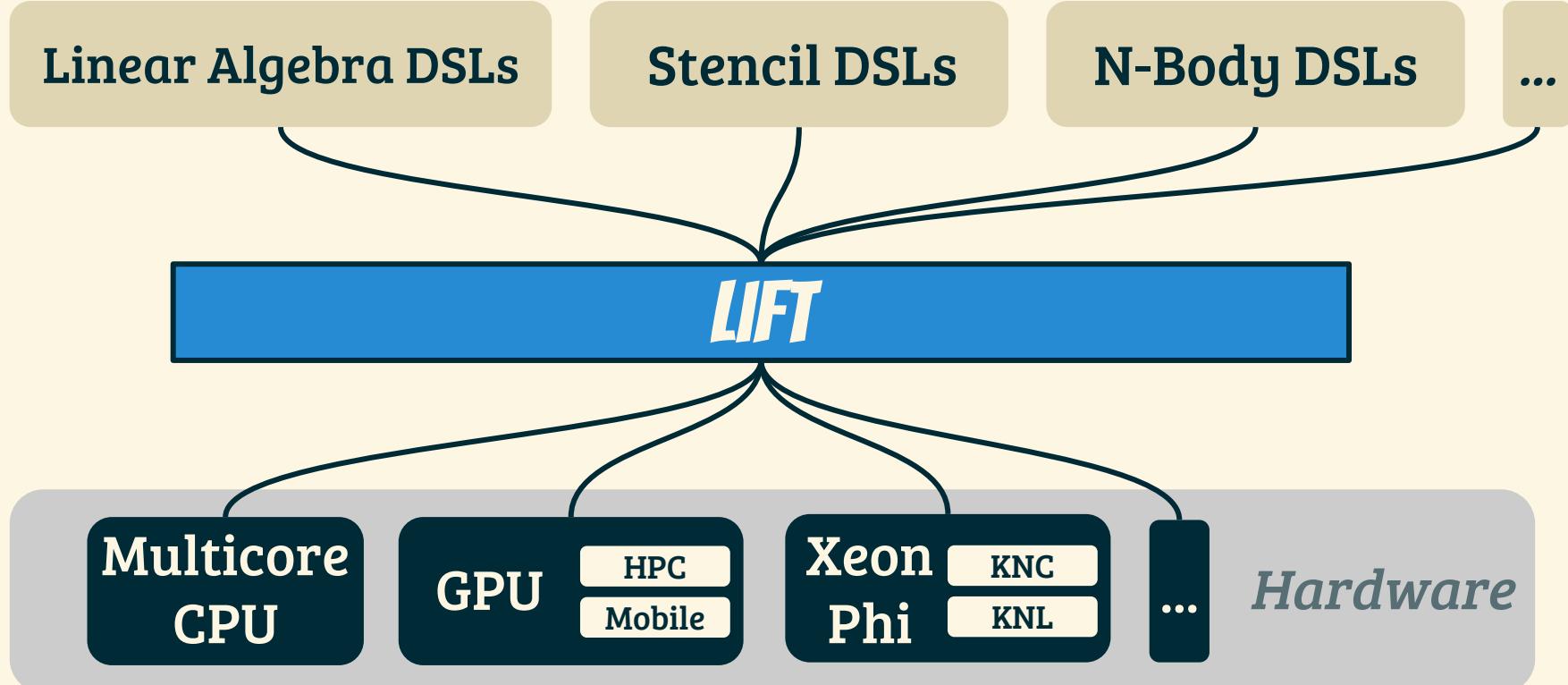
EXPLOITING DOMAIN KNOWLEDGE



APPROACHING PERFORMANCE PORTABILITY



APPROACHING PERFORMANCE PORTABILITY



LIFT

DSL

DSL

DSL

**Multicore
CPU**

GPU

HPC
Mobile

**Xeon
Phi**

KNC
KNL

...

Hardware

LIFT

DSL

DSL

DSL

High-Level IR

Multicore
CPU

GPU

HPC
Mobile

Xeon
Phi

KNC
KNL

...

Hardware

LIFT

DSL

DSL

DSL

High-Level IR

Explore Optimizations
by rewriting



Multicore
CPU

GPU

HPC
Mobile

Xeon
Phi

KNC
KNL

...

Hardware

LIFT

DSL

DSL

DSL

High-Level IR

Explore Optimizations
by rewriting



Low-Level Program

Multicore
CPU

GPU

HPC
Mobile

Xeon
Phi

KNC
KNL

...

Hardware

LIFT

DSL

DSL

DSL

High-Level IR

Explore Optimizations
by rewriting

Low-Level Program

Multicore
CPU

GPU

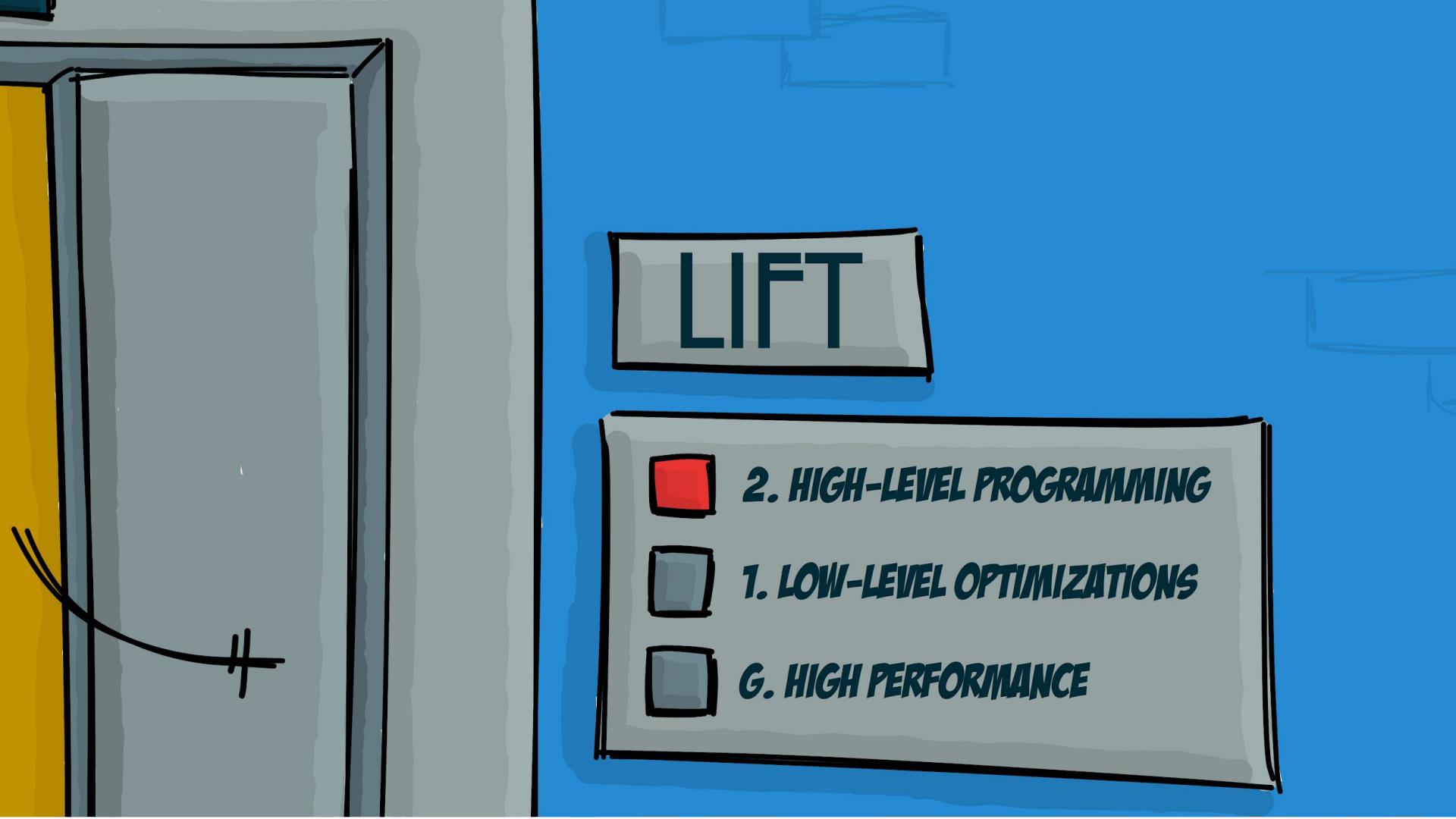
HPC
Mobile

Xeon
Phi

KNC
KNL

...

Hardware



LIFT

2. HIGH-LEVEL PROGRAMMING

1. LOW-LEVEL OPTIMIZATIONS

G. HIGH PERFORMANCE

LIFT'S HIGH-LEVEL PRIMITIVES

map($\square \rightarrow \square$) 

reduce(\oplus) 

split(n) 

join 

zip 

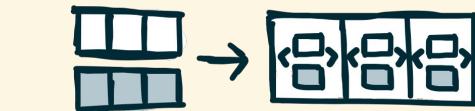
LIFT'S HIGH-LEVEL PRIMITIVES

map($\square \rightarrow \square$) 

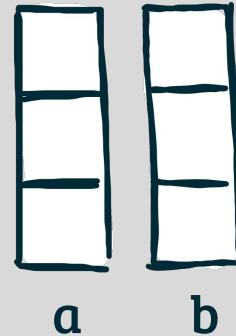
reduce(\oplus) 

split(n) 

join 

zip 

dotproduct.lift



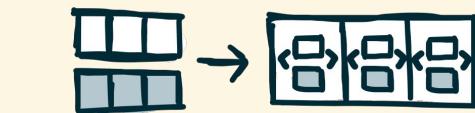
LIFT'S HIGH-LEVEL PRIMITIVES

map($\square \rightarrow \square$) 

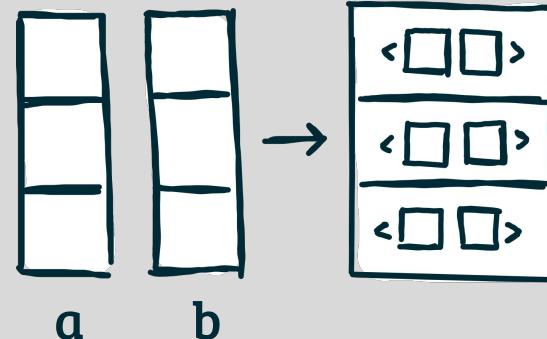
reduce(\oplus) 

split(n) 

join 

zip 

dotproduct.lift



zip(a, b)

LIFT'S HIGH-LEVEL PRIMITIVES

map($\square \rightarrow \square$) 

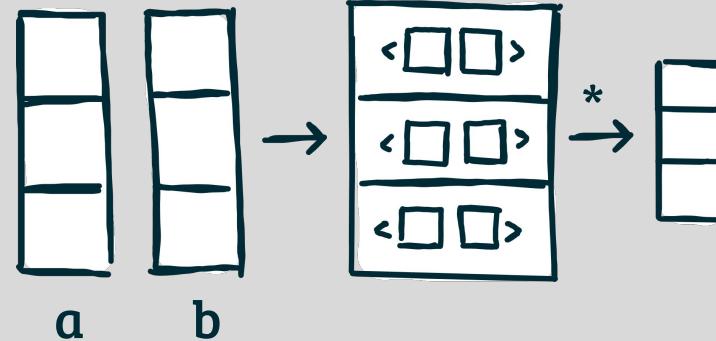
reduce(\oplus) 

split(n) 

join 

zip 

dotproduct.lift



map(, zip(a,b))*

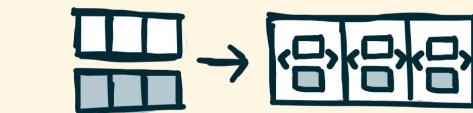
LIFT'S HIGH-LEVEL PRIMITIVES

map($\square \rightarrow \square$) 

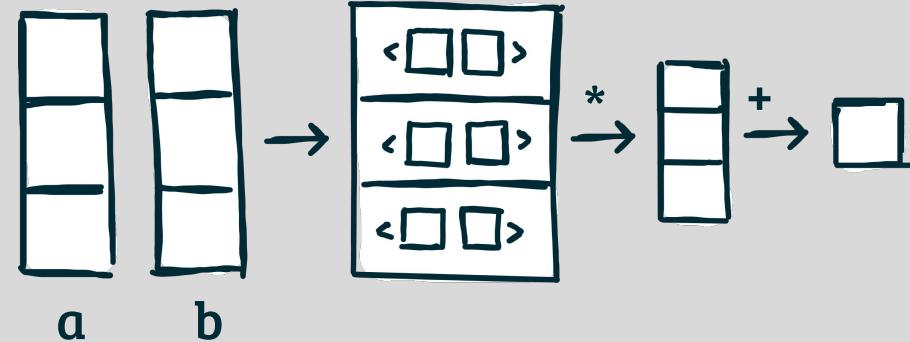
reduce(\oplus) 

split(n) 

join 

zip 

dotproduct.lift



reduce($+$, 0 , *map*(* , *zip*(a , b)))

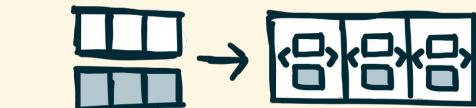
LIFT'S HIGH-LEVEL PRIMITIVES

map($\square \rightarrow \square$) 

reduce(\oplus) 

split(n) 

join 

zip 

stencil.lift?

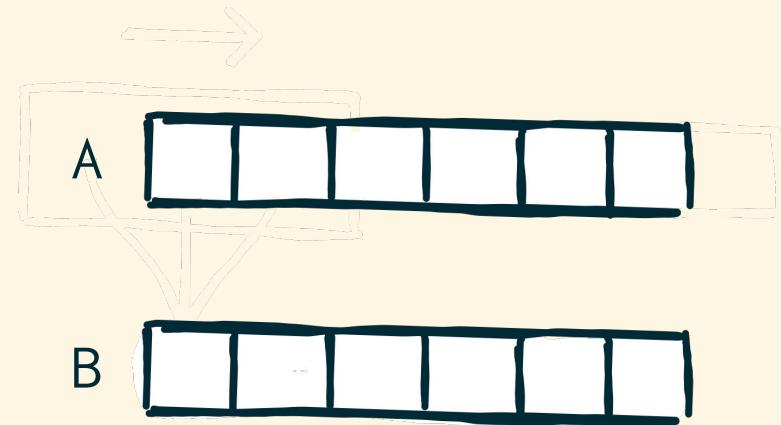
Can we express stencil computations in Lift?

Let's look at a simple stencil example...

WHAT ARE STENCIL COMPUTATIONS?

3-point-stencil.c

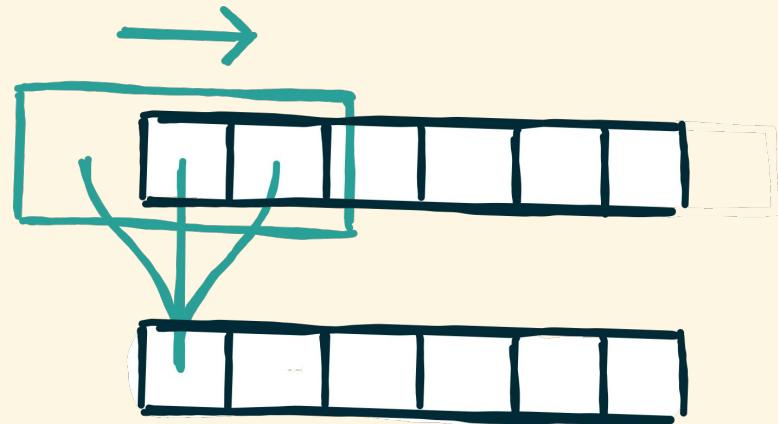
```
for (int i = 0; i < N ; i++) {  
    int sum = 0;  
    for ( int j = -1; j <= 1; j++) {  
        int pos = i + j;  
        pos = pos < 0 ? 0 : pos;  
        pos = pos > N - 1 ? N - 1 : pos;  
        sum += A[ pos ]; }  
    B[ i ] = sum ; }
```



WHAT ARE STENCIL COMPUTATIONS?

3-point-stencil.c

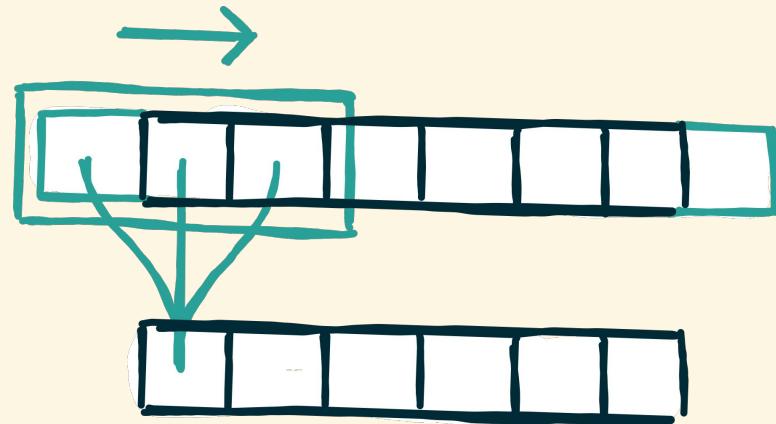
```
for (int i = 0; i < N ; i++) {  
    int sum = 0;  
    for ( int j = -1; j <= 1; j++) {  
        int pos = i + j;  
        pos = pos < 0 ? 0 : pos;  
        pos = pos > N - 1 ? N - 1 : pos;  
        sum += A[ pos ]; }  
    B[ i ] = sum ; }
```



WHAT ARE STENCIL COMPUTATIONS?

3-point-stencil.c

```
for (int i = 0; i < N ; i++) {  
    int sum = 0;  
    for ( int j = -1; j <= 1; j++) {  
        int pos = i + j;  
        pos = pos < 0 ? 0 : pos;  
        pos = pos > N - 1 ? N - 1 : pos;  
        sum += A[ pos ]; }  
    B[ i ] = sum ; }
```



STENCIL COMPUTATIONS IN LIFT

map($\square \rightarrow \square$) 

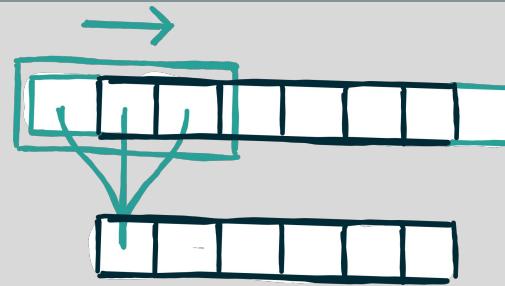
reduce(\oplus) 

split(n) 

join 

zip 

3-point-stencil.lift



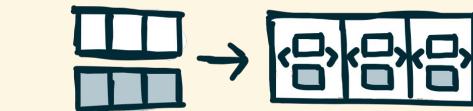
STENCIL COMPUTATIONS IN LIFT

map($\square \rightarrow \square$) 

reduce(\oplus) 

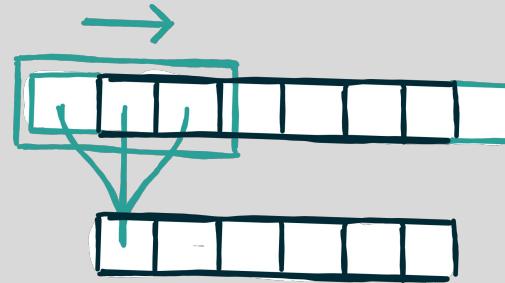
split(n) 

join 

zip 

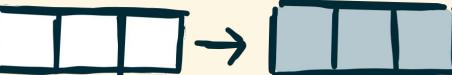
stencil 

3-point-stencil.lift



Add specialized primitive: Job done?

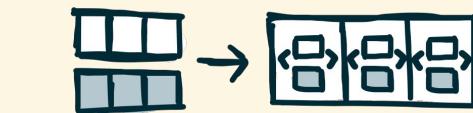
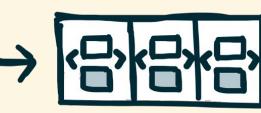
STENCIL COMPUTATIONS IN LIFT

map($\square \rightarrow \square$)  \rightarrow 

reduce(\oplus)  \rightarrow 

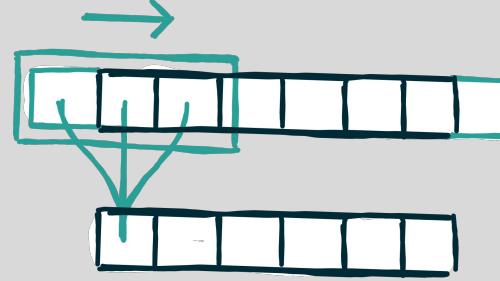
split(n)  \rightarrow 

join  \rightarrow 

zip  \rightarrow 

stencil  \rightarrow 

3-point-stencil.lift



Add specialized primitive: Job done?

🚫 **No Reuse**

of existing primitives and optimizations

🚫 **Domain-specific**

rather than generic

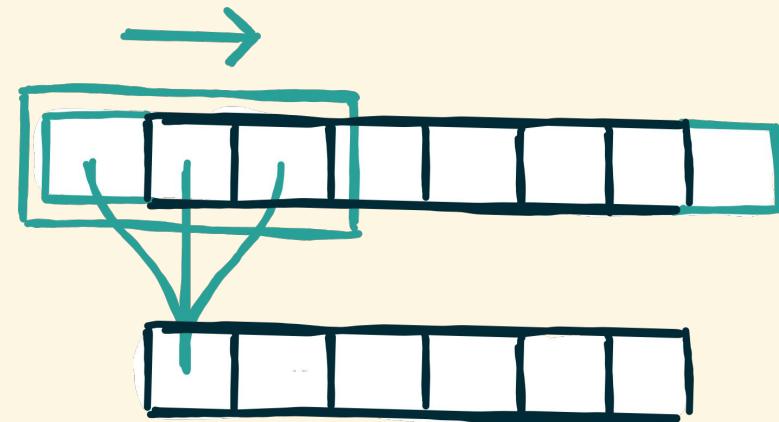
🚫 **Multidimensional?**

is it composable?

DECOMPOSING STENCIL COMPUTATIONS

3-point-stencil.c

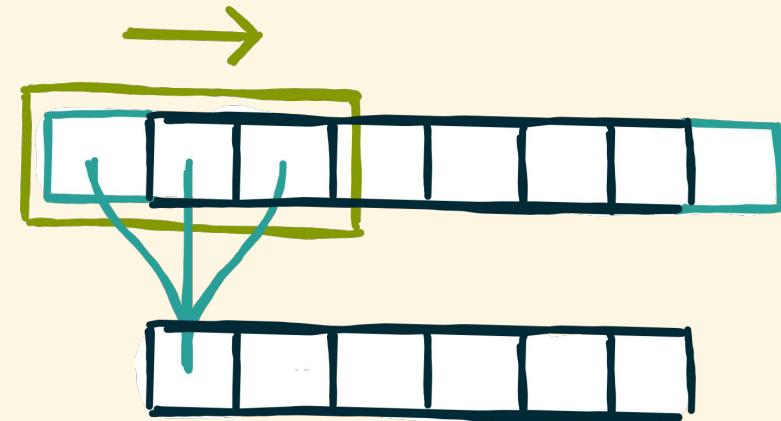
```
for (int i = 0; i < N ; i++) {  
    int sum = 0;  
    for ( int j = -1; j <= 1; j++) {  
        int pos = i + j;  
        pos = pos < 0 ? 0 : pos;  
        pos = pos > N - 1 ? N - 1 : pos;  
        sum += A[ pos ]; }  
    B[ i ] = sum ; }
```



DECOMPOSING STENCIL COMPUTATIONS

3-point-stencil.c

```
for (int i = 0; i < N ; i++) {  
    int sum = 0;  
    for ( int j = -1; j <= 1; j++) { // ( a )  
        int pos = i + j;  
        pos = pos < 0 ? 0 : pos;  
        pos = pos > N - 1 ? N - 1 : pos;  
        sum += A[ pos ]; }  
    B[ i ] = sum ; }
```

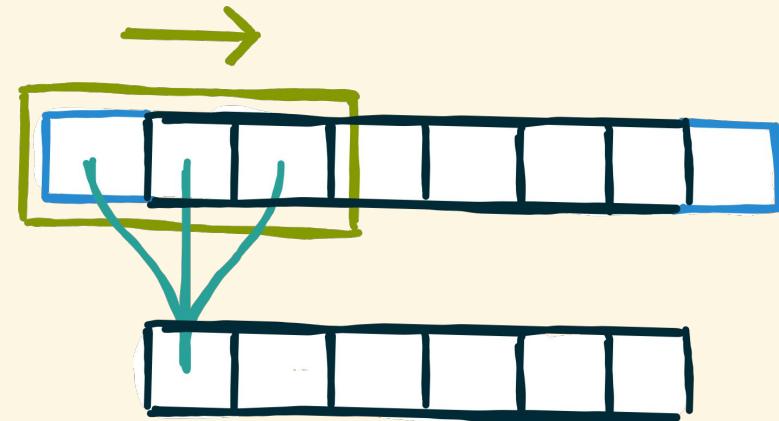


(a) access neighborhoods for every element

DECOMPOSING STENCIL COMPUTATIONS

3-point-stencil.c

```
for (int i = 0; i < N ; i++) {  
    int sum = 0;  
    for ( int j = -1; j <= 1; j++) { // ( a )  
        int pos = i + j;  
        pos = pos < 0 ? 0 : pos;           // ( b )  
        pos = pos > N - 1 ? N - 1 : pos;  
        sum += A[ pos ]; }  
    B[ i ] = sum ; }
```

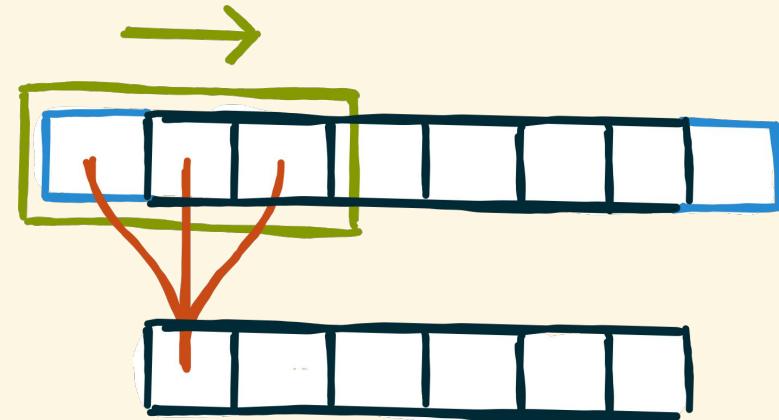


- (a) access neighborhoods for every element
- (b) specify boundary handling

DECOMPOSING STENCIL COMPUTATIONS

3-point-stencil.c

```
for (int i = 0; i < N ; i++) {  
    int sum = 0;  
    for ( int j = -1; j <= 1; j++) { // ( a )  
        int pos = i + j;  
        pos = pos < 0 ? 0 : pos;           // ( b )  
        pos = pos > N - 1 ? N - 1 : pos;  
        sum += A[ pos ]; }               // ( c )  
    B[ i ] = sum ; }
```

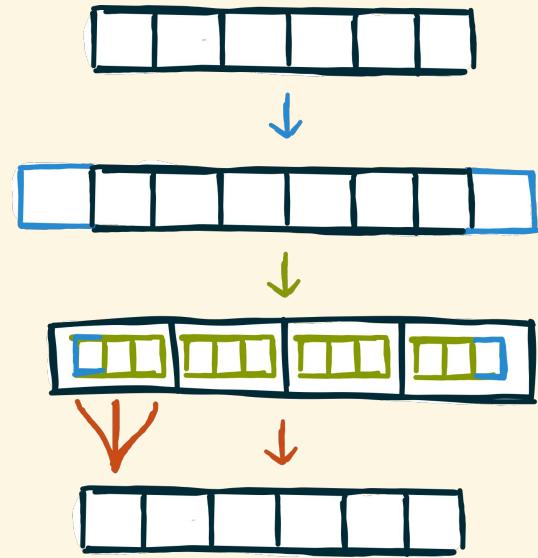


- (a) access neighborhoods for every element
- (b) specify boundary handling
- (c) apply stencil function to neighborhoods

DECOMPOSING STENCIL COMPUTATIONS

3-point-stencil.c

```
for (int i = 0; i < N ; i++) {  
    int sum = 0;  
    for ( int j = -1; j <= 1; j++) { // ( a )  
        int pos = i + j;  
        pos = pos < 0 ? 0 : pos;           // ( b )  
        pos = pos > N - 1 ? N - 1 : pos;  
        sum += A[ pos ]; }               // ( c )  
  
    B[ i ] = sum ; }
```

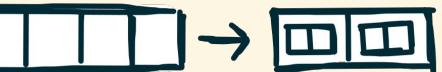


- (a) access neighborhoods for every element
- (b) specify boundary handling
- (c) apply stencil function to neighborhoods

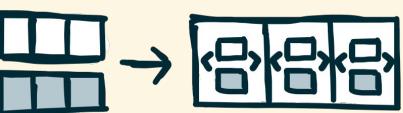
STENCIL COMPUTATIONS IN LIFT

map($\square \rightarrow \square$) 

reduce(\oplus) 

split(n) 

join 

zip 

3-point-stencil.lift



$\downarrow ???$



$\downarrow ???$

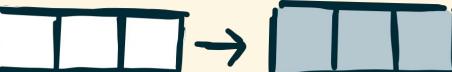


\swarrow

$\downarrow ???$



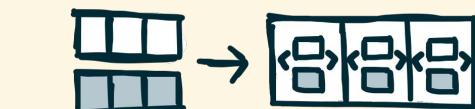
STENCIL COMPUTATIONS IN LIFT

map($\square \rightarrow \square$) 

reduce(\oplus) 

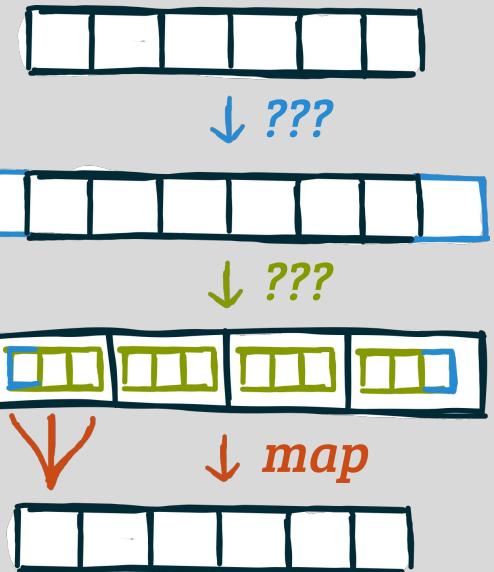
split(n) 

join 

zip 

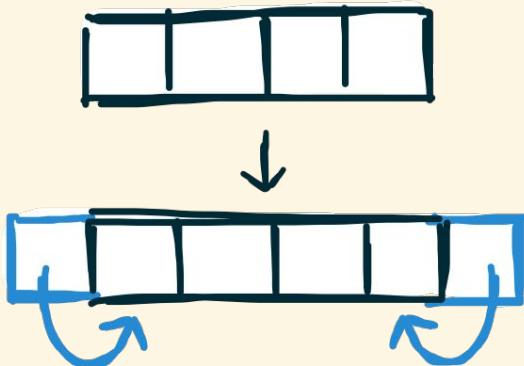
3-point-stencil.lift

- ✓ **Reuse map**
allows to reuse existing rewrite rules
- ✓ **Simplicity**
one primitive per task
- ✓ **Multidimensional**
easily composable

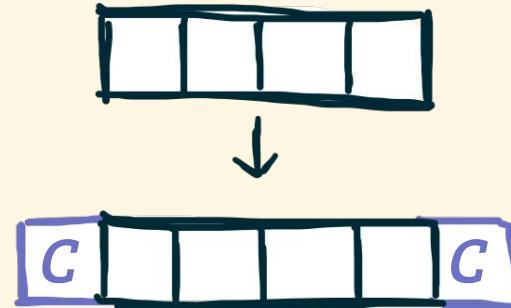


BOUNDARY HANDLING USING PAD

pad (reindexing)



pad (constant)



pad-reindexing.lift

```
clamp(i, n) = (i < 0) ? 0 :  
((i >= n) ? n-1:i)
```

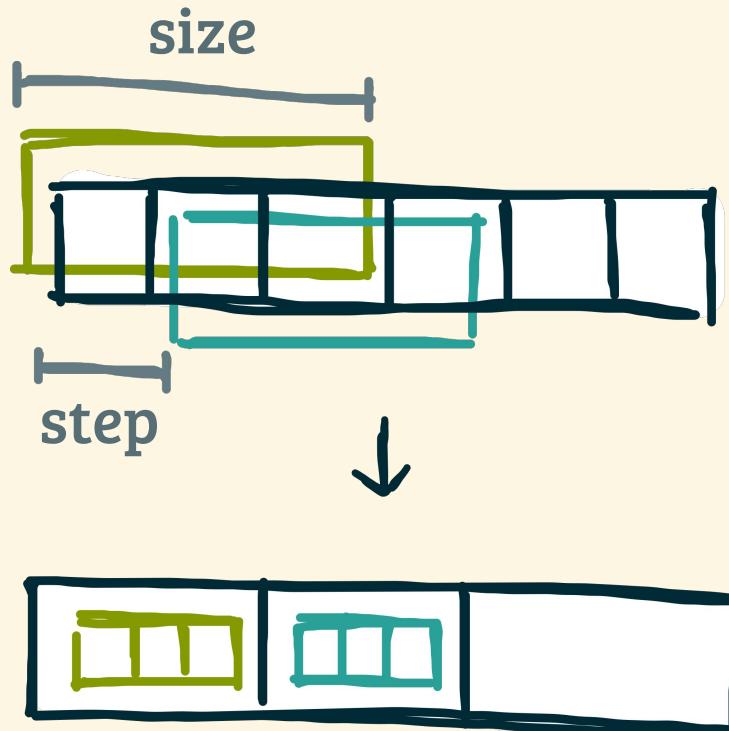
```
pad(1,1,clamp, [a,b,c,d]) =  
[a,a,b,c,d,d]
```

pad-constant.lift

```
constant(i, n) = C
```

```
pad(1,1,constant, [a,b,c,d]) =  
[C,a,b,c,d,C]
```

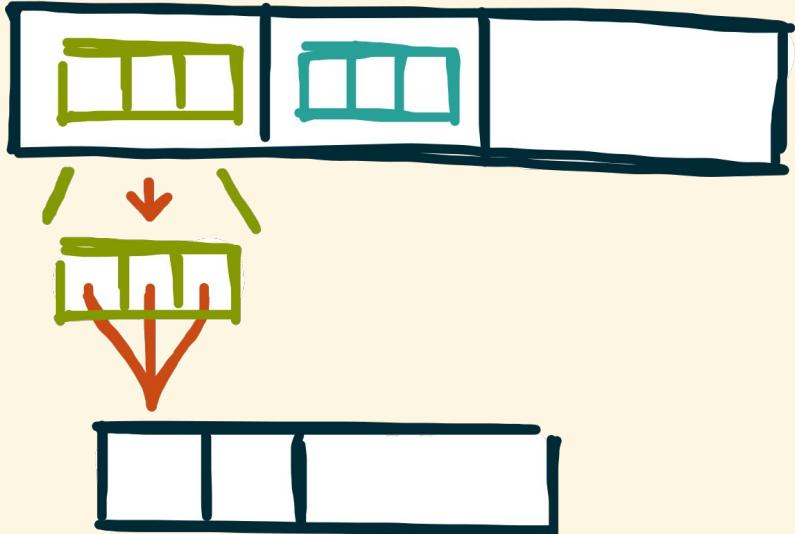
NEIGHBORHOOD CREATION USING SLIDE



slide-example.lift

```
slide(3,1,[a,b,c,d,e]) =  
[[a,b,c],[b,c,d],[c,d,e]]
```

APPLYING STENCIL FUNCTION USING MAP



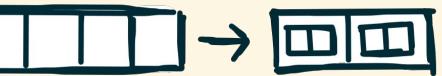
sum-neighborhoods.lift

```
map(nbh =>  
    reduce(add, 0.0f, nbh))
```

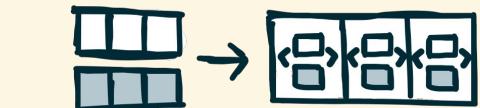
PUTTING IT TOGETHER

map($\square \rightarrow \square$) 

reduce(\oplus) 

split(n) 

join 

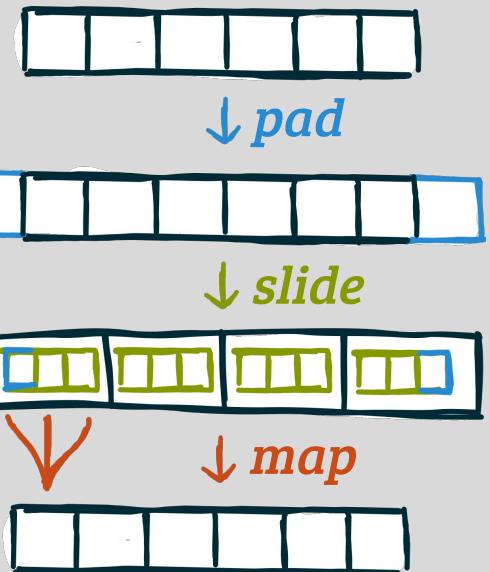
zip 

pad(l,r,b) 

slide(n,s) 

stencil1D.lift

```
def stencil1D =  
  fun(A =>  
    map(reduce(add, 0.0f),  
        slide(3,1,  
              pad(1,1,clamp,A))))
```



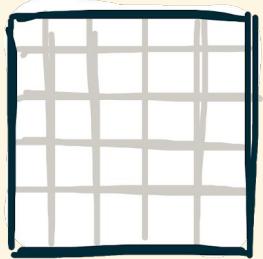
MULTIDIMENSIONAL STENCIL COMPUTATIONS

are expressed as compositions of intuitive, generic 1D primitives

Decompose to Re-Compose

MULTIDIMENSIONAL STENCIL COMPUTATIONS

are expressed as compositions of intuitive, generic 1D primitives

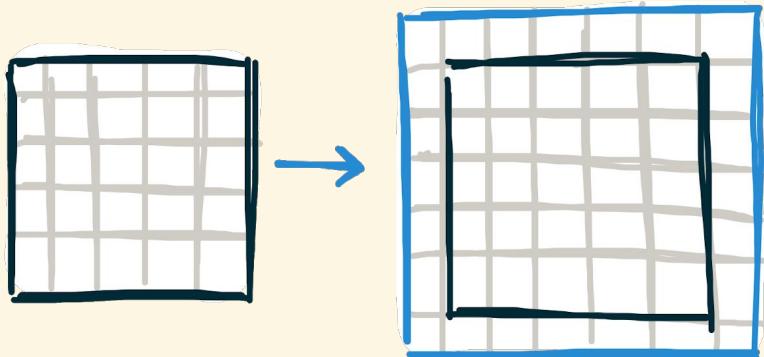


Decompose to Re-Compose

MULTIDIMENSIONAL STENCIL COMPUTATIONS

are expressed as compositions of intuitive, generic 1D primitives

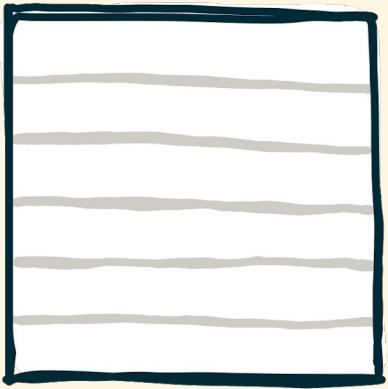
Decompose to Re-Compose



$\text{pad}_2(1, 1, \text{clamp}, \text{input})$

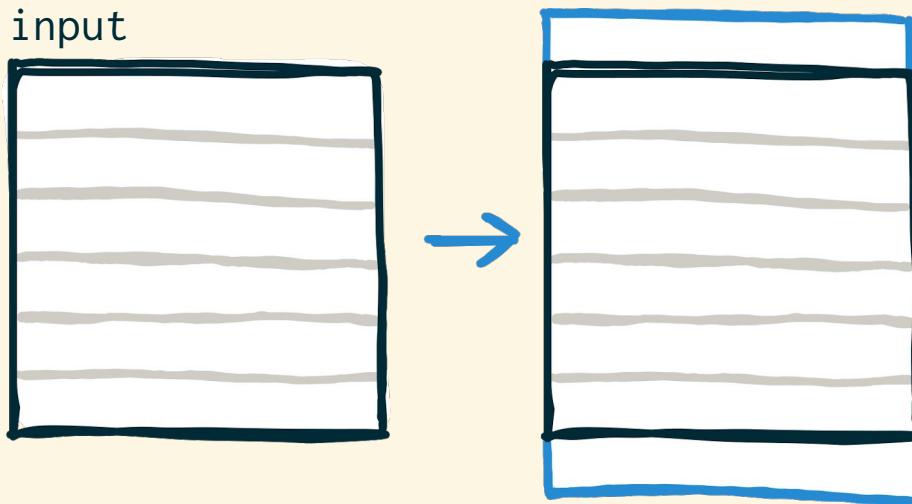
MULTIDIMENSIONAL BOUNDARY HANDLING USING pad_2

input



$pad_2 =$

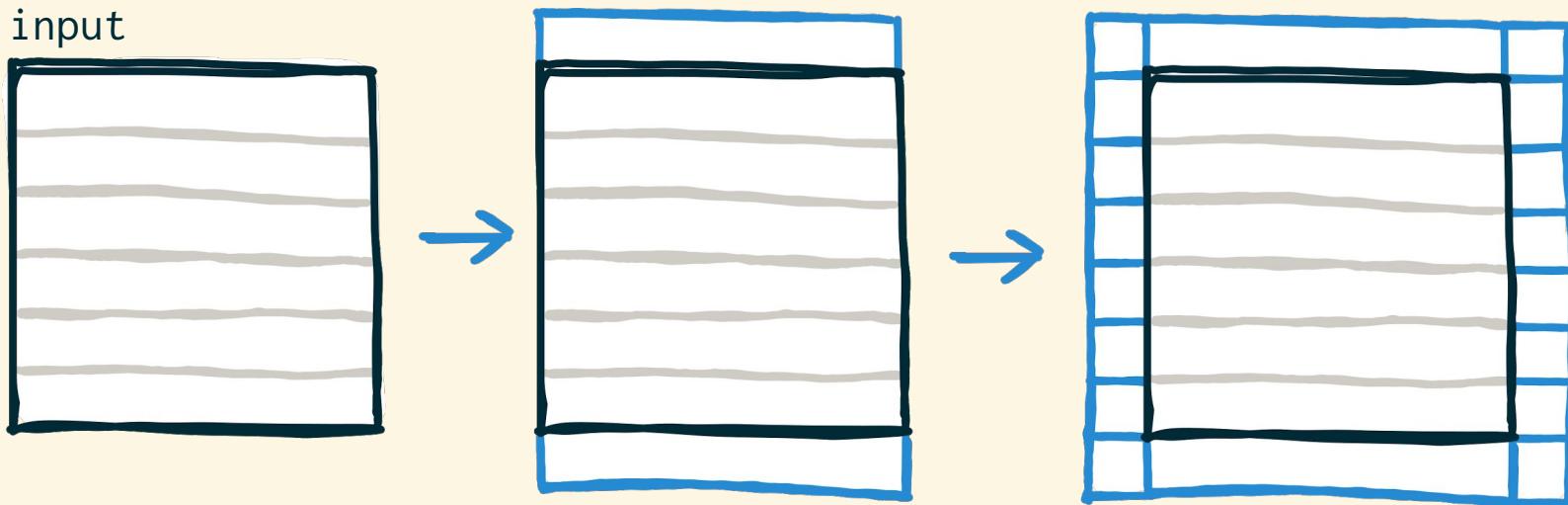
MULTIDIMENSIONAL BOUNDARY HANDLING USING pad_2



$pad_2 =$

$pad(1, r, b, \text{input})$

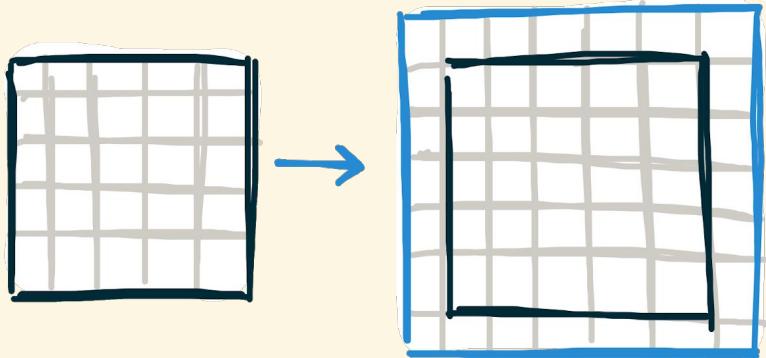
MULTIDIMENSIONAL BOUNDARY HANDLING USING PAD_2


$$\text{pad}_2 = \text{map}(\text{pad}(l, r, b, \text{pad}(l, r, b, \text{input})))$$

MULTIDIMENSIONAL STENCIL COMPUTATIONS

are expressed as compositions of intuitive, generic 1D primitives

Decompose to Re-Compose

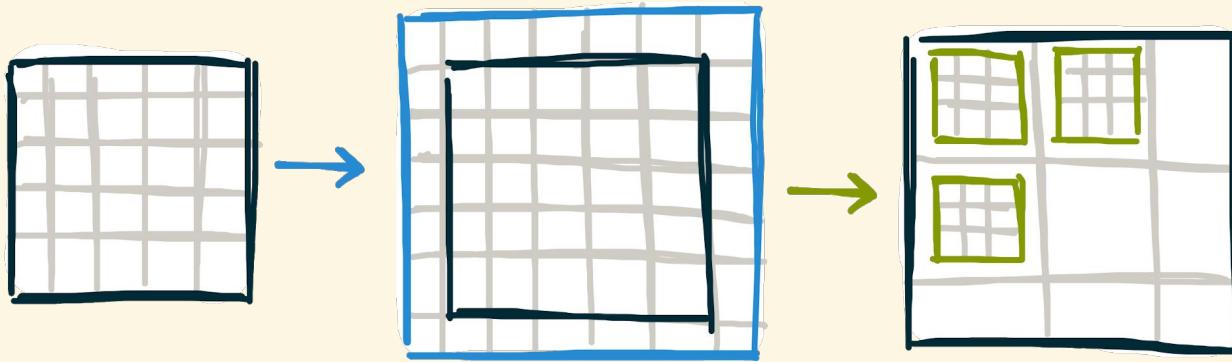


$pad_2(1, 1, clamp, input)$

MULTIDIMENSIONAL STENCIL COMPUTATIONS

are expressed as compositions of intuitive, generic 1D primitives

Decompose to Re-Compose

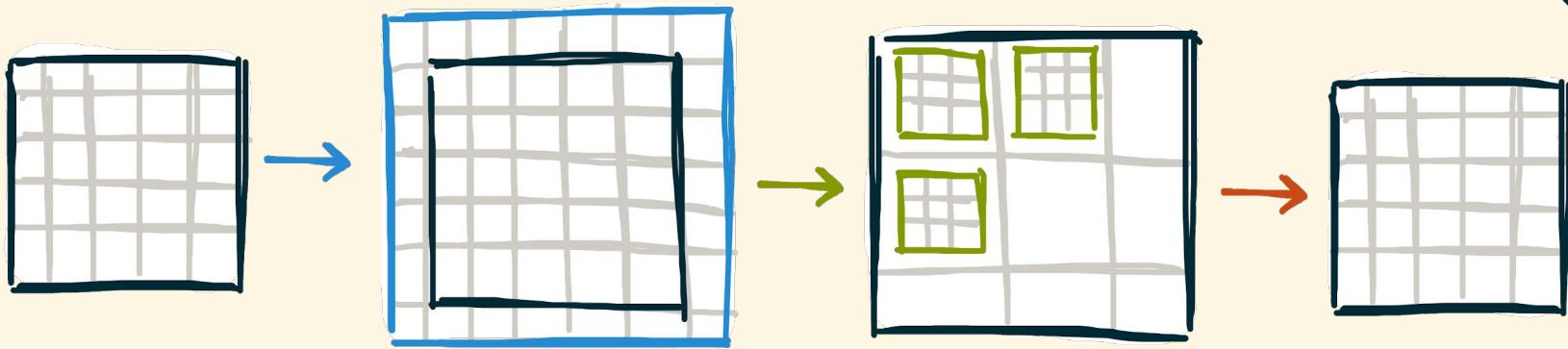


slide₂(3,1, pad₂(1,1,clamp,input))

MULTIDIMENSIONAL STENCIL COMPUTATIONS

are expressed as compositions of intuitive, generic 1D primitives

Decompose to Re-Compose

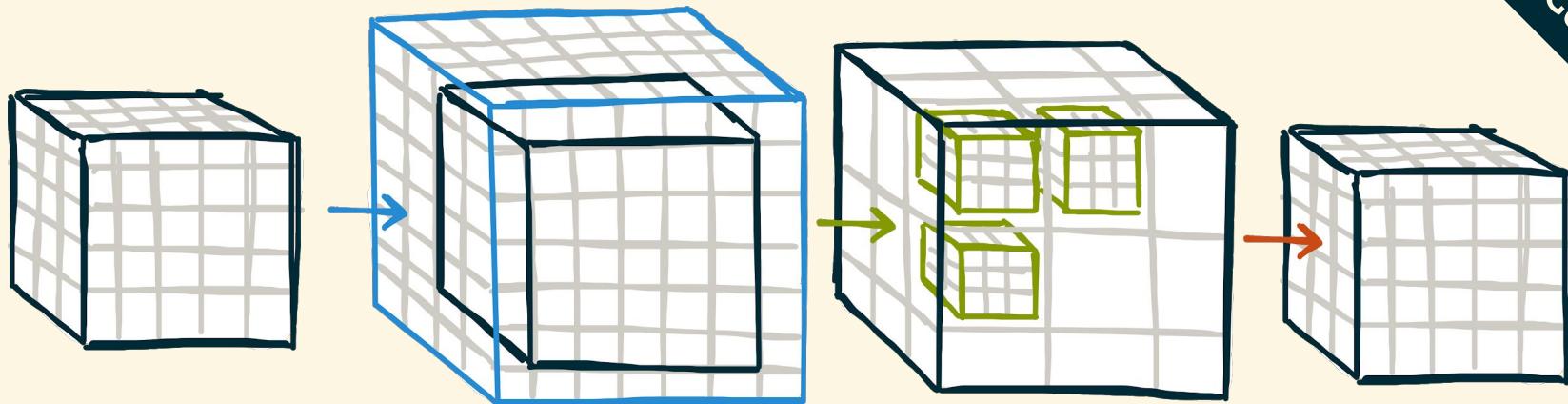


$\text{map}_2(\text{sum}, \text{slide}_2(3,1, \text{pad}_2(1,1, \text{clamp}, \text{input})))$

MULTIDIMENSIONAL STENCIL COMPUTATIONS

are expressed as compositions of intuitive, generic 1D primitives

Decompose to Re-Compose

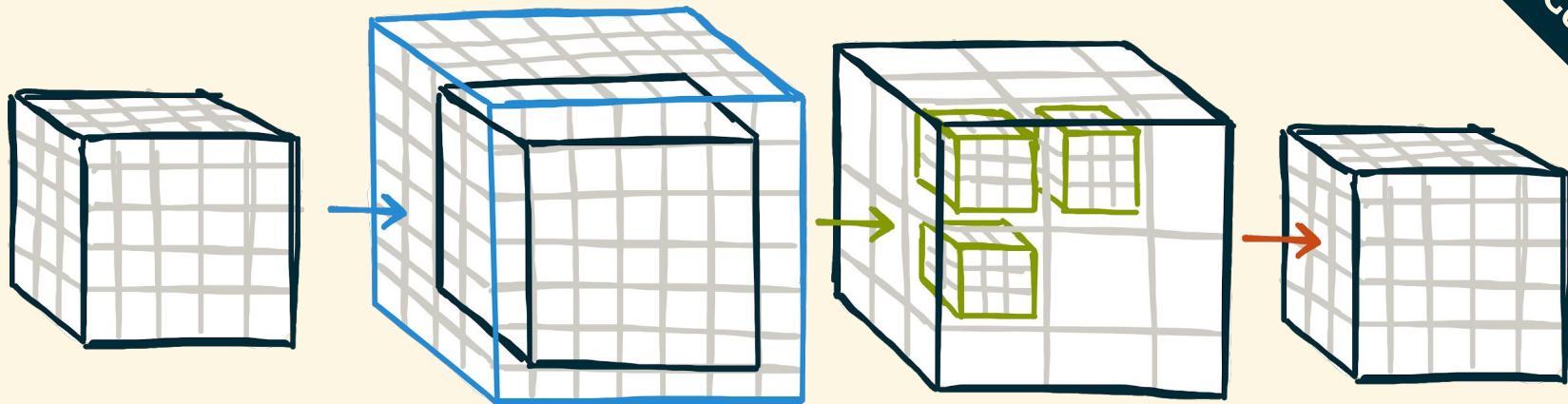


$\text{map}_3(\text{sum}, \text{slide}_3(3,1, \text{pad}_3(1,1,\text{clamp},\text{input})))$

MULTIDIMENSIONAL STENCIL COMPUTATIONS

are expressed as compositions of intuitive, generic 1D primitives

Decompose to Re-Compose



$\text{map}_3(\text{sum}, \text{slide}_3(3, 1, \text{pad}_3(1, 1, \text{clamp}, \text{input})))$

Advantages:



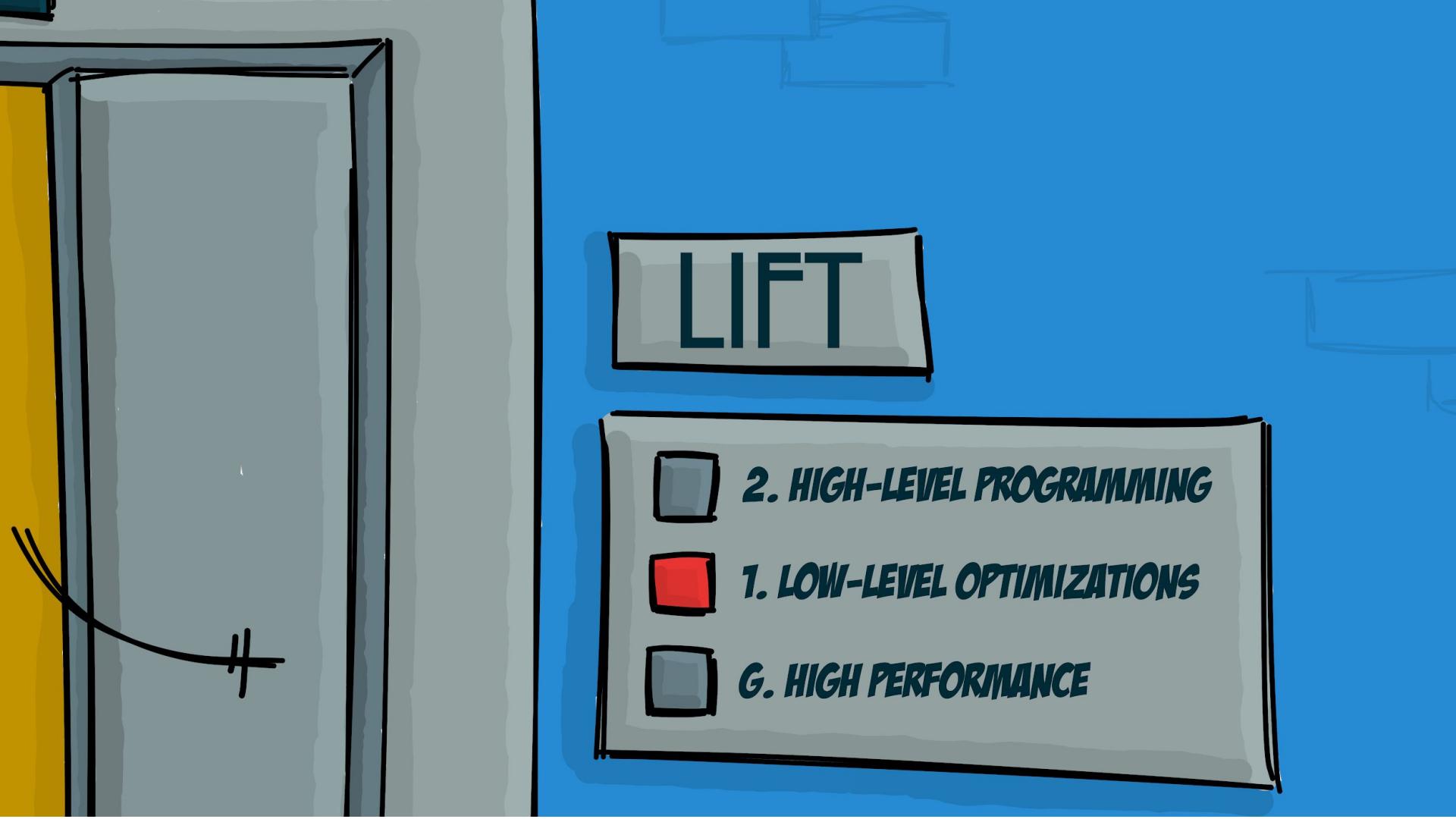
Compact Language



Reuse Rewrites



Simple Compilation



LIFT

2. HIGH-LEVEL PROGRAMMING

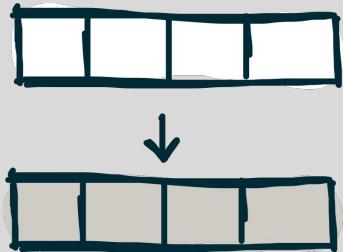
1. LOW-LEVEL OPTIMIZATIONS

G. HIGH PERFORMANCE

REUSING EXISTING REWRITE RULES

Divide & Conquer

map(f, A)



REUSING EXISTING REWRITE RULES

Divide & Conquer

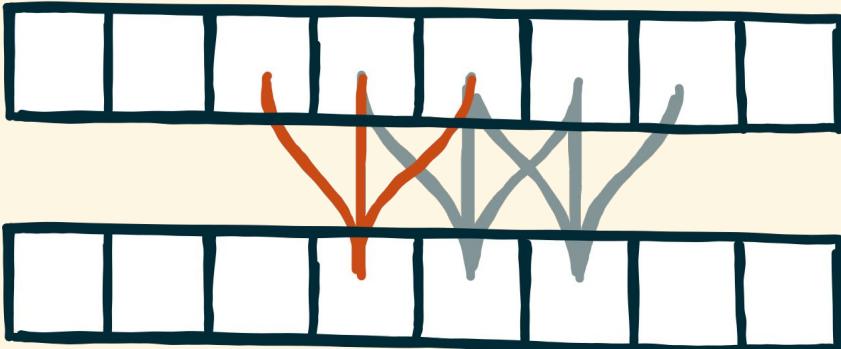
$map(f, A)$



$join(map(map(f),$
 $split(n, A)))$



OPTIMIZATION: OVERLAPPED TILING



Exploit Locality

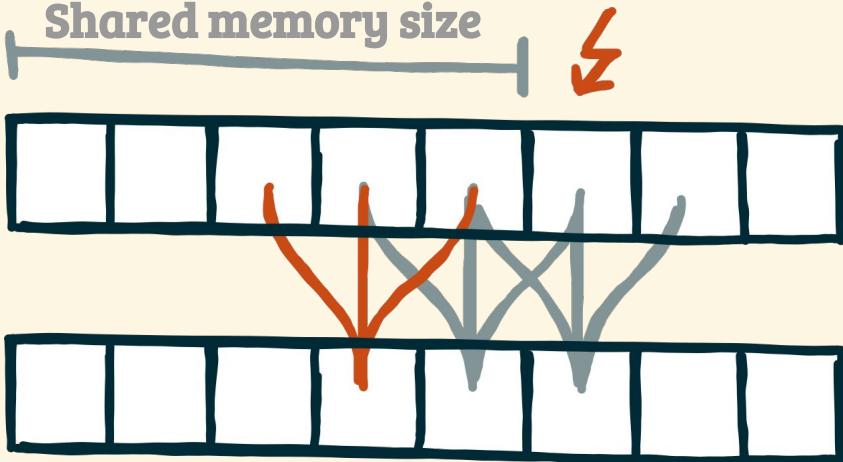
Close neighborhoods share elements that can be grouped in tiles



Shared Memory

Fast memory can be used to cache tiles

OPTIMIZATION: OVERLAPPED TILING



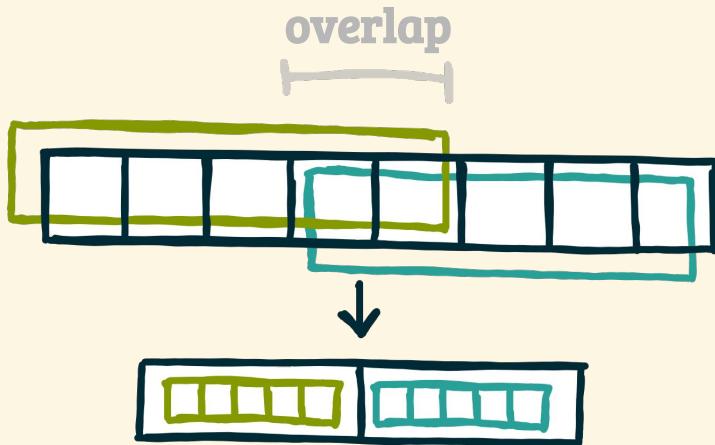
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OPTIMIZATION: OVERLAPPED TILING



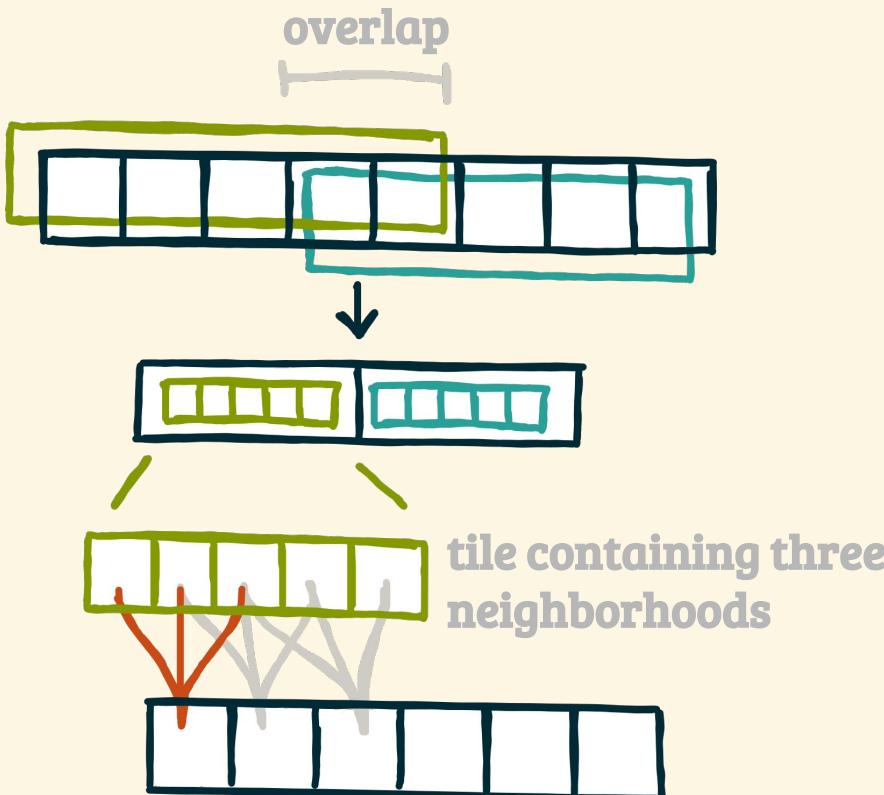
✓ Exploit Locality

Close neighborhoods share elements that can be grouped in tiles

✓ Shared Memory

Fast memory can be used to cache tiles

OPTIMIZATION: OVERLAPPED TILING



✓ Exploit Locality

Close neighborhoods share elements that can be grouped in tiles

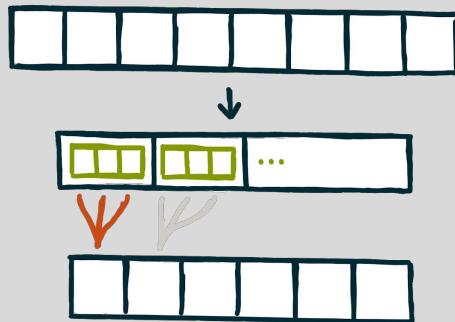
✓ Shared Memory

Fast memory can be used to cache tiles

OVERLAPPED TILING AS A REWRITE RULE

overlapped tiling rule

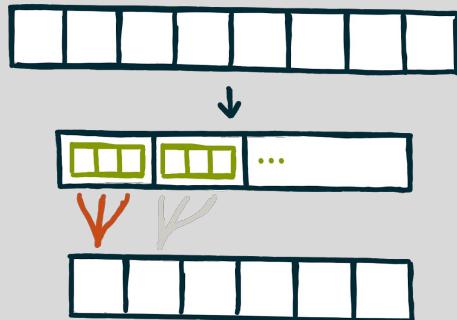
`map(f, slide(3,1,input))`



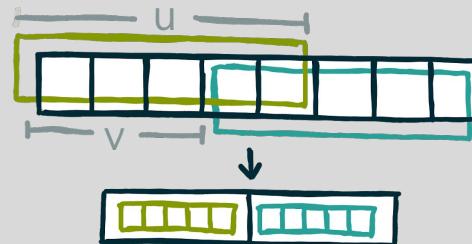
OVERLAPPED TILING AS A REWRITE RULE

overlapped tiling rule

map(f, slide(3,1,input)) ↪



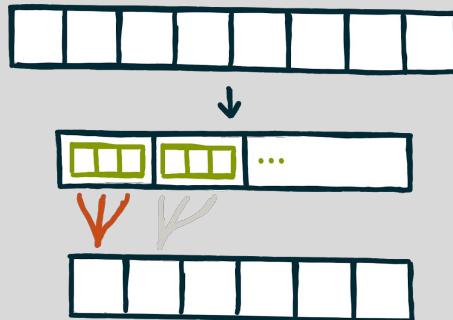
slide(u,v,input)



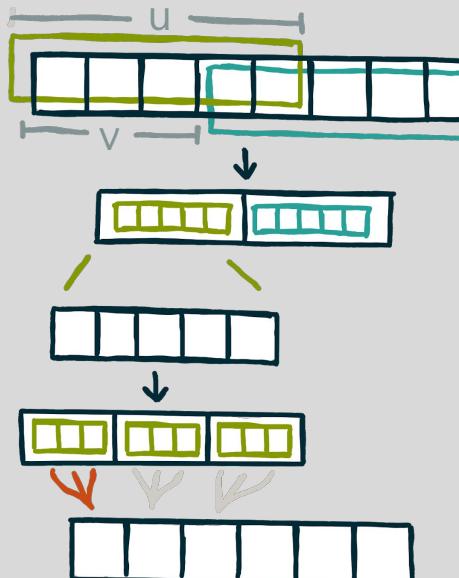
OVERLAPPED TILING AS A REWRITE RULE

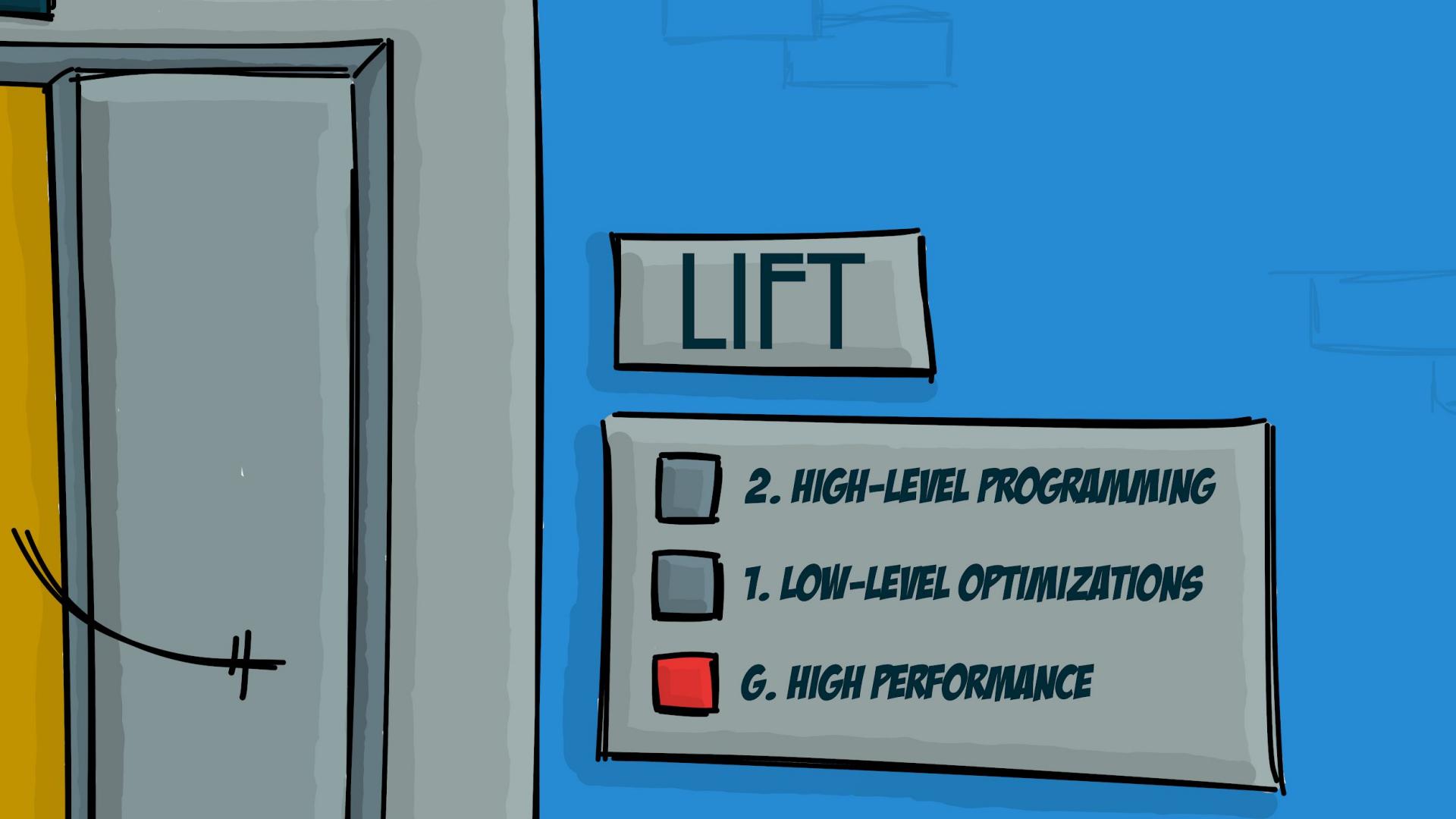
overlapped tiling rule

$\text{map}(f, \text{slide}(3,1,\text{input}))$



$\text{join}(\text{map}(\text{tile} \Rightarrow$
 $\text{map}(f, \text{slide}(3,1,\text{tile})),$
 $\text{slide}(u,v,\text{input})))$





LIFT

2. HIGH-LEVEL PROGRAMMING

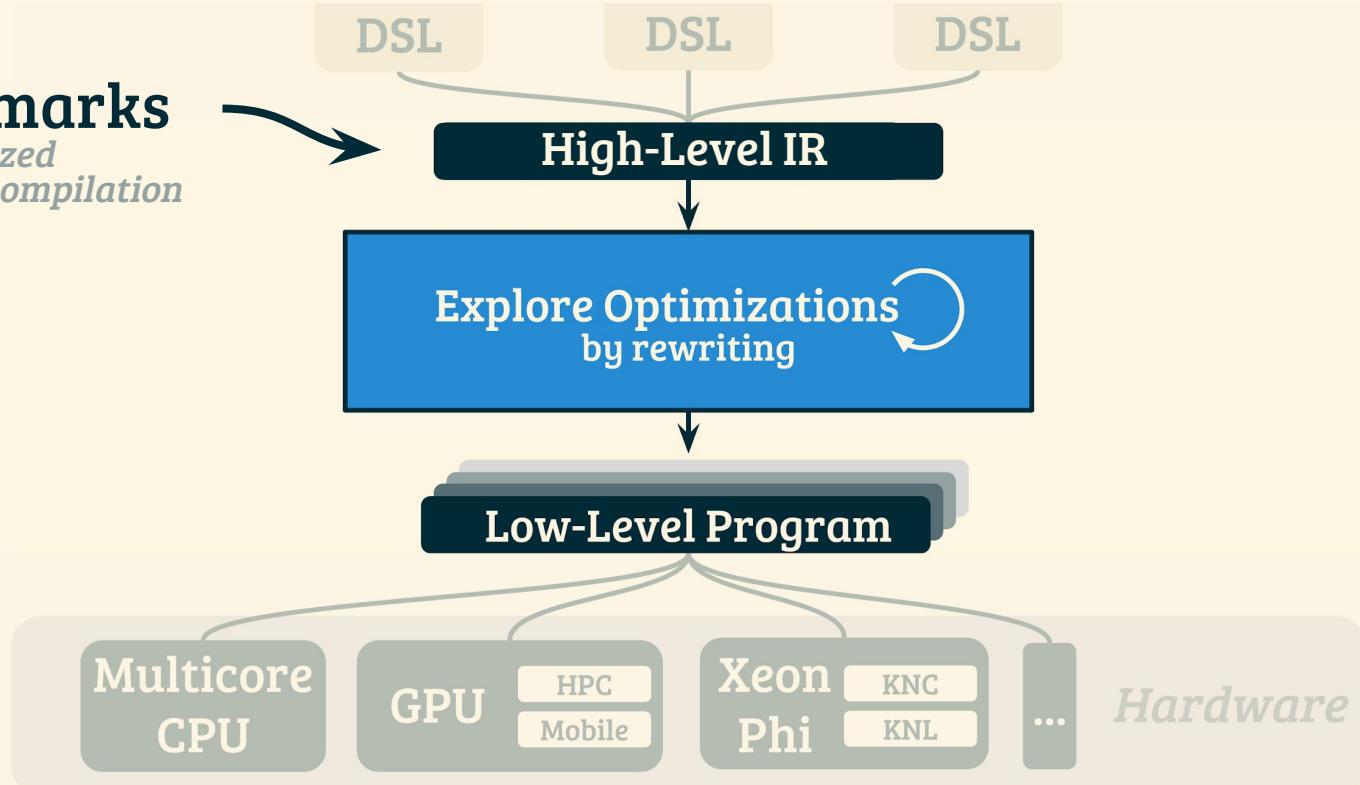
1. LOW-LEVEL OPTIMIZATIONS

G. HIGH PERFORMANCE

EXPERIMENTAL EVALUATION



14 Benchmarks
*6 hand-optimized
8 polyhedral compilation*



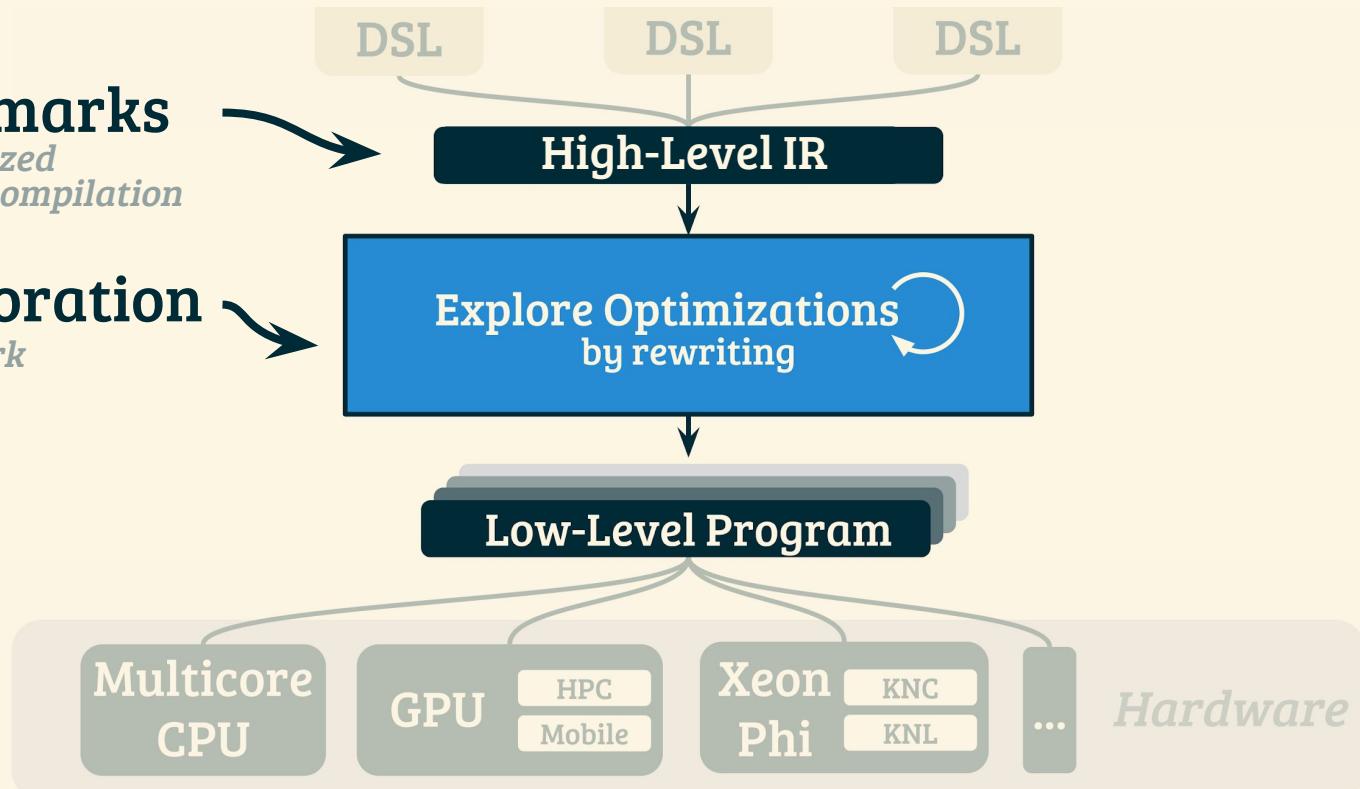
EXPERIMENTAL EVALUATION



14 Benchmarks
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< 3h Exploration
per benchmark



EXPERIMENTAL EVALUATION



14 Benchmarks

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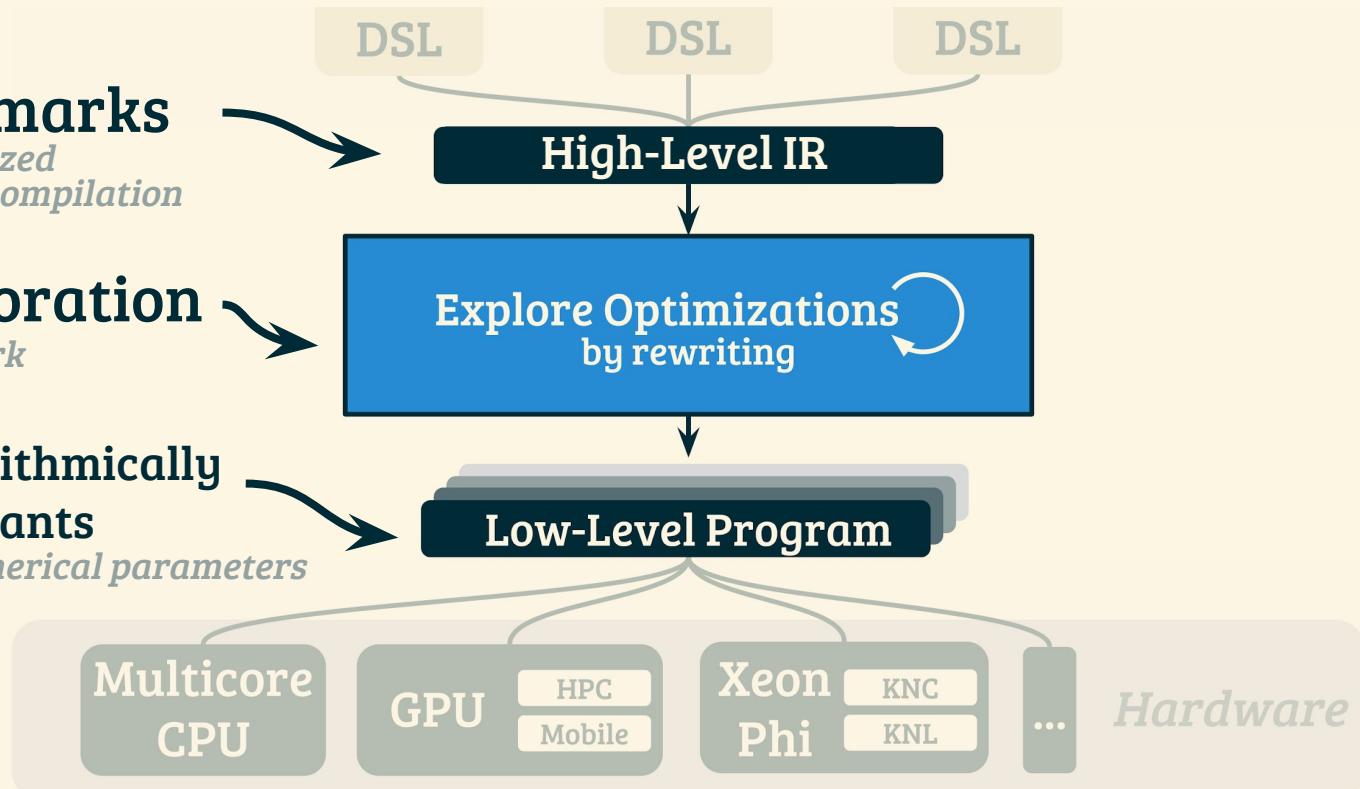
< 3h Exploration

per benchmark



**up to 20 algorithmically
different variants**

+ auto-tuning of numerical parameters



EXPERIMENTAL EVALUATION



14 Benchmarks

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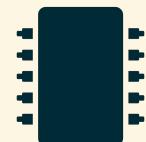
< 3h Exploration

per benchmark



up to 20 algorithmically different variants

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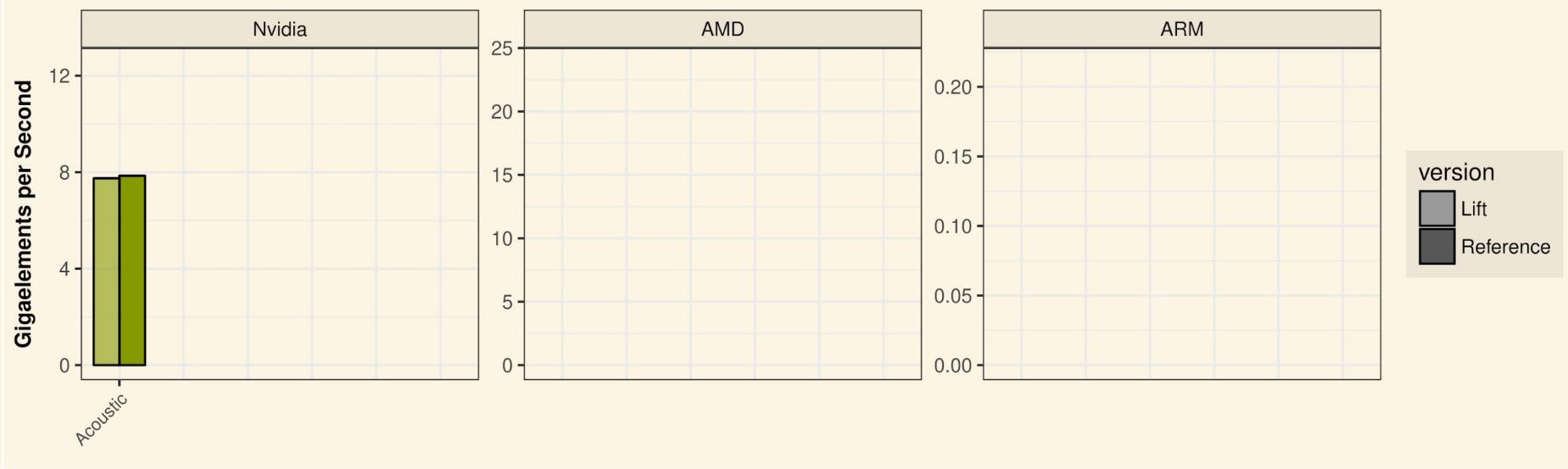
3 GPU Architectures

*2 Desktop GPUs
1 Mobile GPU*



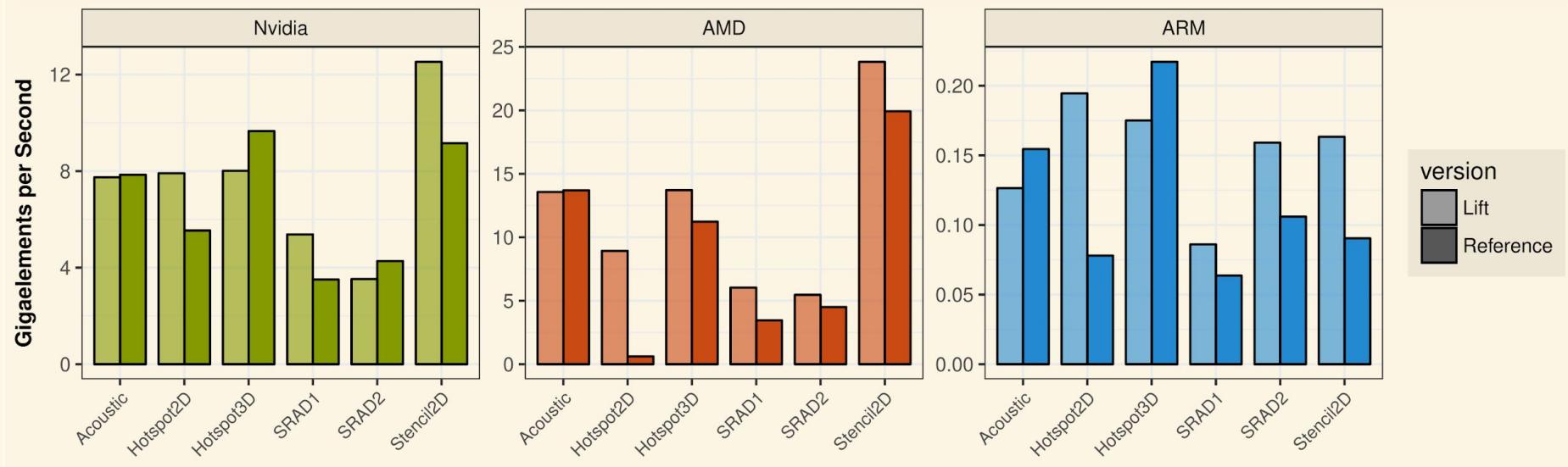
COMPARISON WITH HAND-OPTIMIZED CODES

higher is better



COMPARISON WITH HAND-OPTIMIZED CODES

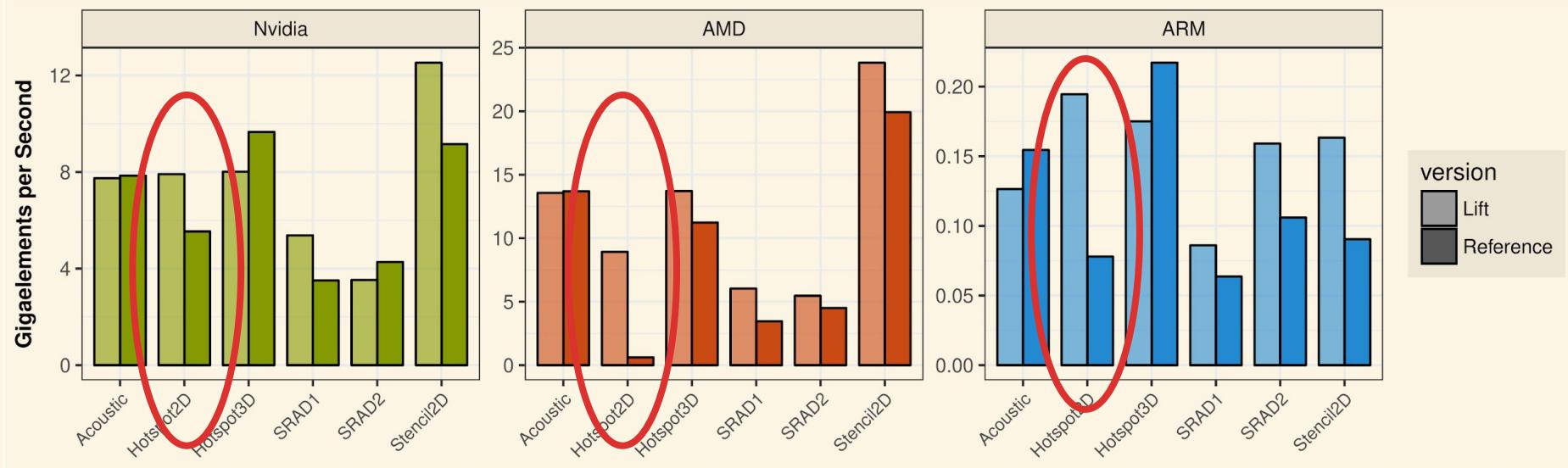
higher is better



Lift achieves the same performance
as hand optimized code

COMPARISON WITH HAND-OPTIMIZED CODES

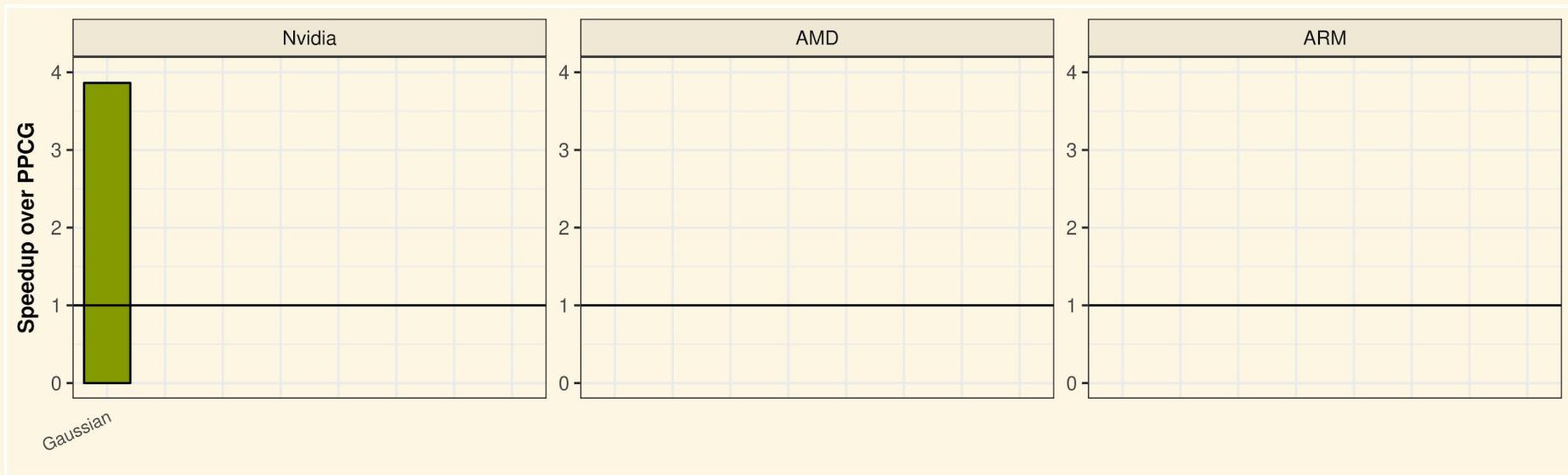
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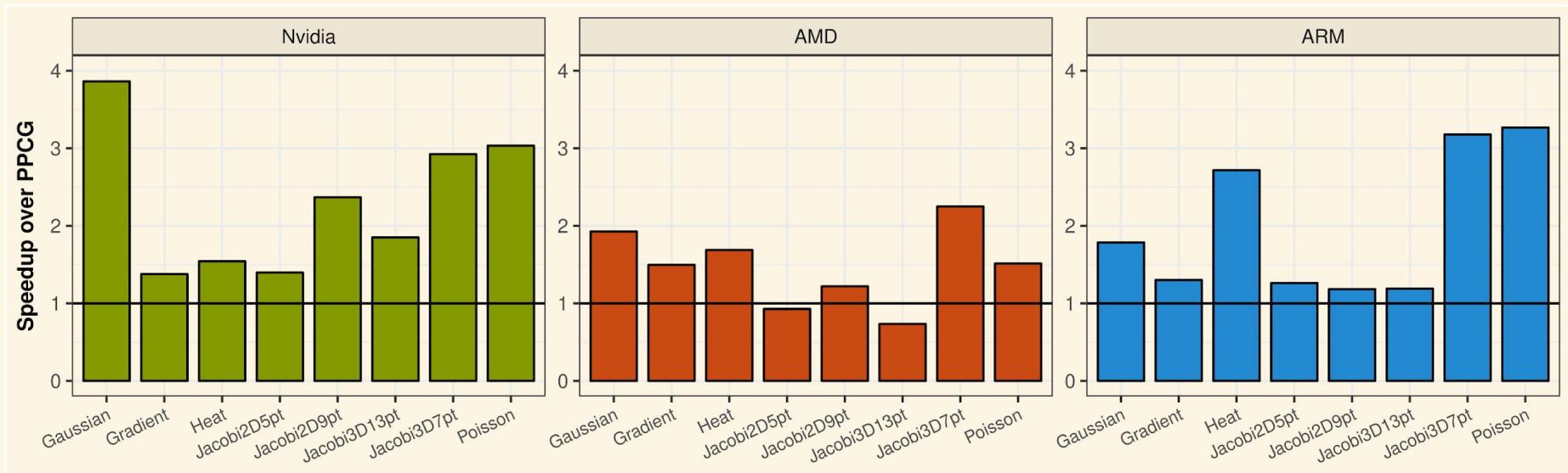
COMPARISON WITH POLYHEDRAL COMPIRATION

higher is better



COMPARISON WITH POLYHEDRAL COMPIRATION

higher is better



**Lift outperforms state-of-the-art
optimizing compilers**

STENCIL COMPUTATIONS IN LIFT

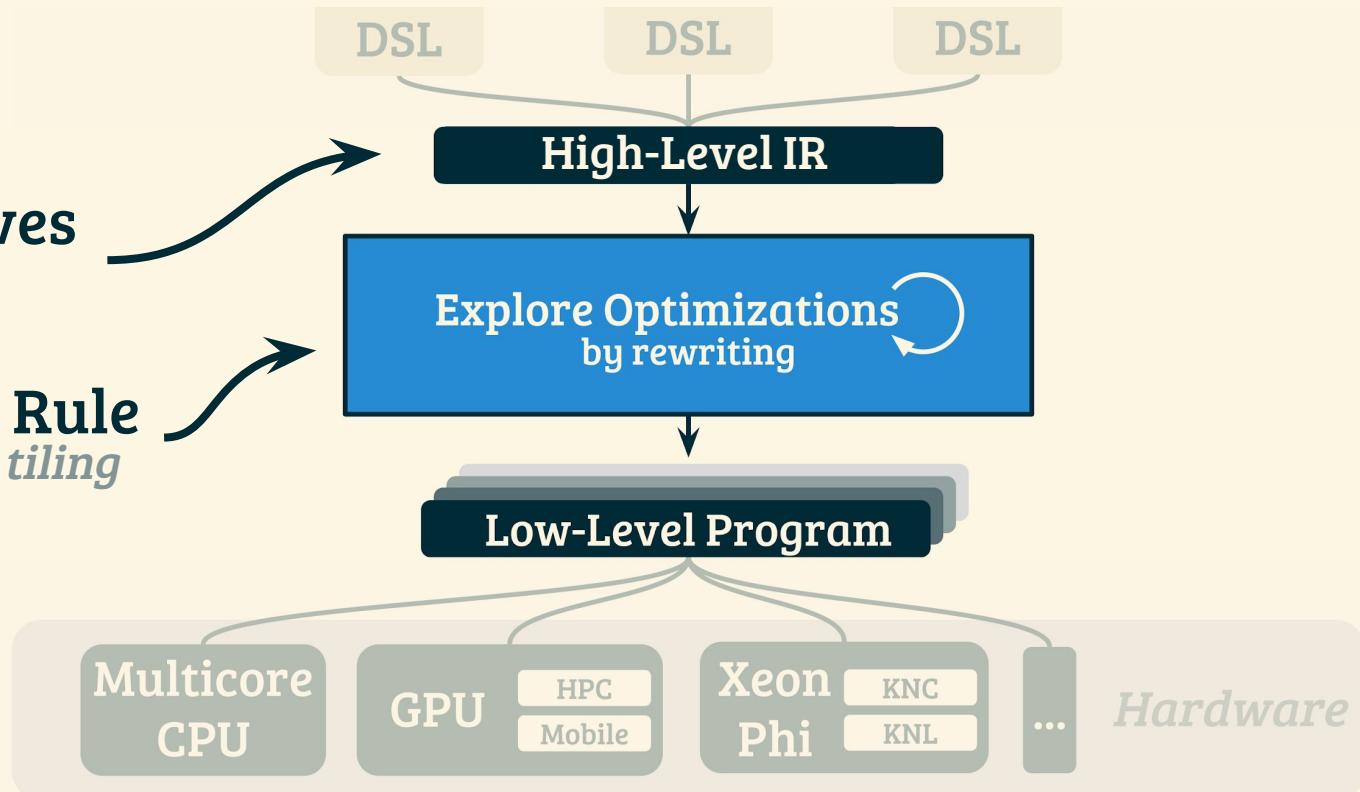
We added:



2 Primitives
pad, slide



1 Rewrite Rule
overlapped tiling



A STEP FURTHER

towards multi-kernel stencil computations



conv

```
def conv =  
  fun(A =>  
    map2D(multWeights,  
           slide2D(3,1,  
                    pad2D(1,1,clamp,A))))
```



blur

grad (x)

grad (y)

mm

mm

mm

conv

conv

conv

det

trace

harris



A STEP FURTHER

towards multi-kernel stencil computations

conv

```
def conv =  
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stage 1

stage 2



blur

grad (x)

grad (y)

mm

mm

mm

conv

conv

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harris

A STEP FURTHER

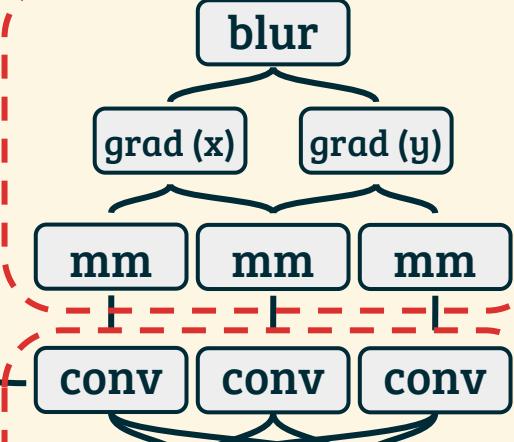
towards multi-kernel stencil computations



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```

stage 1



stage 2



Kernel Fusion via Rewriting



Guided Search Space Exploration



LIFT IS OPEN SOURCE!



more info at:

lift-project.org

” Paper



CGO Artifact



Source Code



Best Paper Award (CGO'18)

Bastian Hagedorn: b.hagedorn@wwu.de

TO BE CONTINUED...