ALGORITHMS FOR THE MANUSCRIPT "HAMMING AND GILBERT BOUNDS FOR POLY-ALPHABETIC CODES WITH THE L1 DISTANCE"

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ABSTRACT. This document presents several SageMath implementations useful to apply the main results of the manuscript entitled "Hamming and Gilbert bounds for poly-alphabetic codes with the Manhattan distance".

Algorithms to compute $\eta_r(\mathcal{X})$ and $\gamma_r(\mathcal{X})$ and \overline{B}_r

Algorithm ?? provides the code in SageMath to calculate $\eta_r(\mathcal{X})$.

To use the *Python* functions given in algorithms in this note, one can import them into a Sage-Math worksheet by writing load("GitHub_L1Codes.py")¹ at the beginning of the worksheet.

```
Data: \mathbf{r} \in \mathbb{Z}_{\geq 0} and the Python list \mathbf{m} = [m_1, ..., m_n] where \mathcal{X} = \prod_{i=1}^n [0, m_i - 1]. Result: \operatorname{eta}(\mathbf{r}, \mathbf{m}) computes \eta_r(\mathcal{X}).

def \operatorname{eta}(\mathbf{r}, \mathbf{m}):
    if \mathbf{r} = 0:
        return 1

\mathbf{m} = \operatorname{len}(\mathbf{m}); indices = \operatorname{range}(1, \mathbf{n} + 1); \operatorname{eta} = 0

for delta in \operatorname{range}(\mathbf{r} + 1):
    for J in Subsets(indices):
        sum_m_J = \operatorname{sum}(\mathbf{m}[i-1]) for i in J)
        if delta - \operatorname{sum}_m = \operatorname{binomial}(\mathbf{n} + \operatorname{delta} - \operatorname{sum}_m = \operatorname{delta} - \operatorname{delta
```

Algorithm 1: A Python function that applies Theorem 3.8 to compute $\eta_r(\mathcal{X})$ in SageMath.

Let $x \in Inm(\mathcal{X})$, $n \in \mathbb{Z}_{\geq 1}$, $\delta \in \mathbb{Z}_{>0}$, p_{δ} be the coefficient of degree δ of p(t,x). The function in Algorithm ?? uses the functions given in Algorithms ??, ?? to provide the SageMath code to

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¹The document "GitHub_L1Codes.py" can be donwloaded from https://github.com/egarcia-claro/L1Codes

```
calculate p_{\delta} for \mathcal{X} = \prod_{i=1}^{n} [0, m_i - 1] when E = \{i \in [n] : 2 \mid m_i\} is arbitrary.
  Data: delta \in \mathbb{Z}_{\geq 0} and the Python list \mathbf{m} = [m_1, ..., m_n] where \mathcal{X} = \prod_{i=1}^n [0, m_i - 1].
  Result: p_delta_o(delta, m) computes the coefficient p_{\delta} of degree \delta of
            p(t, \text{ any innermost point}) \text{ when } E = \{i \in [n] : 2 \mid m_i\} = \emptyset.
  def p_delta_o(delta, m):
        if delta==0:
             return 1
        n=len(m); l = [ceil((mi - 1) / 2) for mi in m]
        indices=range(1, n+1)
       E = \{i \text{ for } i \text{ in indices if } m[i-1] \% 2 == 0\}
        partial=sum(li for li in l)
       p_delta = 0
        if len(E) != 0:
             raise ValueError("The m_i's must be odd")
        elif delta == partial:
             return 2**n
       P_n= Subsets(indices)
        for J in P_n:
             if len(J) > 0:
                  for A in Subsets(J):
                       sum_1_A = sum(1[i-1] + 1 for i in A)
                       if delta - sum_l_A >= 0:
                            binom_term = binomial(len(J) + delta - sum_l_A
                                          - 1, delta - sum_l_A)
                            p_{delta} += (-1)**(n - len(J) + len(A))
                                       * 2**len(J) * binom_term
```

Algorithm 2: An algorithm that applies Theorem 3.10 (part 1) to compute p_{δ} in SageMath.

return p_delta

```
Data: delta \in \mathbb{Z}_{>0} and the Python list \mathbf{m} = [m_1, ..., m_n] where \mathcal{X} = \prod_{i=1}^n [0, m_i - 1].
Result: p_delta_o(delta, m) computes the coefficient p_{\delta} of degree \delta of
        p(t, \text{ any innermost point}) \text{ when } E = \{i \in [n] : 2 \mid m_i\} = [n].
def p_delta_e(delta, m):
    if delta==0:
         return 1
    n=len(m); l = [ceil((mi - 1) / 2) for mi in m]
    indices=range(1, n+1)
    E = \{i \text{ for } i \text{ in indices if } m[i-1] \% 2 == 0\}
    partial=sum(li for li in l)
    p_delta = 0
    if len(E) != n:
         raise ValueError("The m_i's must be even")
    if delta == partial:
         return 1
    P_n = Subsets(indices)
    for J in P_n:
         if len(J) > 0: # Ensure J is non-empty
             J_c = set(indices) - set(J)
             u_J_{delta} = (-1)**(len(J)) if delta ==
                           sum(l[i-1] for i in J_c) else 0
             term_sum = u_J_delta
             for A in Subsets(J):
                  if len(A) > 0:
                       for B in Subsets(A):
                           sum_1_B_J_c = sum(1[i-1] \text{ for i in B})
                                         + sum(l[i-1] for i in J_c)
                           if delta - sum_l_B_J_c >= 0:
                                binom_term = binomial(len(A) + delta
                                             - sum_l_B_J_c - 1, delta
                                             - sum_1_B_J_c)
                                term_sum += (-1)**(len(J) - len(A)
                                            + len(B)) * 2**len(A)
                                            * binom_term
             p_delta += term_sum
    return p_delta
```

Algorithm 3: An algorithm that applies Theorem 3.10 (part 2) to compute p_{δ} in SageMath.

Algorithm 4: An algorithm that applies Theorem 3.10 (part 3) to compute p_{δ} in SageMath. It works even when $E = \emptyset$ or E = [n], because p_delta_e(0, [])=p_delta_o(0, [])=1.

```
Data: \mathbf{r} \in \mathbb{Z}_{\geq 0} and the Python list \mathbf{m} = [m_1,...,m_n] where \mathcal{X} = \prod_{i=1}^n [0,m_i-1]. Result: gamma(r, m) computes \gamma_r(\mathcal{X}). def gamma(r, m): if r==0: return 1
```

Algorithm 5: Since $\gamma_r(\mathcal{X}) = p(1,x)_{\leq r} = 1 + \sum_{\delta=1}^r p_{\delta}$ (by Lemma 3.2, part 2, and Theorem 3.6) and $p_{\delta} = p_{\text{delta}(\text{delta, m})}$ (where p_delta(delta, m) is given as in Algorithm ??), this function computes $\gamma_r(\mathcal{X})$.

Algorithms ?? and ?? compute the polynomials P(t, a) and $\overline{p}(t)$ described in Lemma 3.2, parts 1 and 2, respectively.

```
Data: j,m\in\mathbb{Z}_{\geq 0} and x\in\mathcal{X}.

Result: sphere\_size(j,x,m) computes |S_j(x)| computes in [0,m-1].

def sphere\_size(j,x,m):
1\_min = max(x,m-(x+1)); 1\_max = min(x,m-(x+1))

if j == 0:
return 1
elif 0 < j <= 1\_max:
return 2
elif 1\_max < j <= 1\_min:
return 1
else:
return 0
```

Algorithm 6: An algorithm that applies Lemma 3.5 to compute the size of an *j*-sphere in [0, m-1] in SageMath.

```
Data: the Python list \mathbf{m} = [m_1, ..., m_n] where \mathcal{X} = \prod_{i=1}^n [0, m_i - 1], and a \in \mathcal{X}.

Result: compute_p_t(t,a,m) computes the local distance enumerator polynomial p(t,a).

def compute_p_t(t,a, m):
    product = 1

for i in range(len(a)):
    l_i = max(a[i], m[i] - (a[i] + 1))
    sum_expr = sum(sphere_size(j, a[i], m[i]) * t**j for j in range(l_i + 1))
    product *= sum_expr
    return product
```

Algorithm 7: An algorithm that applies Lemma 3.2 (part 1) to compute the local distance enumerator polynomial p(t, a) in SageMath.

Algorithm 8: An algorithm that applies Lemma 3.1 (part 2) to compute a set of representative of orbit in *SageMath*.

```
Data: m \in \mathbb{Z}_{>0} and a \in \mathfrak{T}, where \mathfrak{T} is a set of representatives of orbits under the action of
       G = C_2^n \rtimes S_n on \mathcal{X} = [0, m-1]^n.
Result: compute_stabilizer_size(a, m) computes |Stab_G(a)|.
def compute_stabilizer_size(a, m):
    from collections import Counter
    from math import factorial
    n = len(a); I_a = set()
    for i in range(n):
         reflected_value = m - (a[i] + 1)
         if reflected_value in a:
             I_a.add(i)
    multiplicities = Counter(a)
    product_factorial = 1
    for count in multiplicities.values():
         product_factorial *= factorial(count)
    stabilizer_size = (2 ** len(I_a)) * product_factorial
```

Algorithm 9: An algorithm that applies Lemma 3.1 (part 3) to compute the size of the stabilizer of an element (in \mathfrak{T}) in SageMath.

return stabilizer_size

```
Data: t is the variable in R = \text{PolynomialRing}(\mathbb{QQ}, t), m \in \mathbb{Z}_{>0}, n \in \mathbb{Z}_{>0}
Result: compute_p_bar(t, m, n) computes the polynomial \overline{p}(t) \in \mathbb{Q}[t] of \mathcal{X} = [0, m-1]^n.

def compute_p_bar(t, m, n):
    from math import factorial

T = \text{generate}_r\text{epresentative}_o\text{rbits}(m,n); \text{ total}_s\text{um} = 0
for a in T:
    stab_size = compute_stabilizer_size(a, m)
    orbit_size = \mathbb{QQ}(\text{factorial}(n) * (2**n)) / \text{stab}_s\text{size}
    p_t_a = compute_p_t(t, a, [m]*n)
    total_sum += orbit_size * p_t_a

return total_sum / m**n
```

Algorithm 10: An algorithm that applies Lemma 3.2 (part 2) to compute $\overline{p}(t)$, this is a generating function for the average size of an sphere (in $\mathcal{X} = [0, m-1]^n$) in SageMath.

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