

The function `mystery()` computes a value based on the structure of a binary tree. To understand the behavior of the `mystery()` function, let's follow these steps:

```
public int mystery(Node x){
    if (x == null) return 0;
    return mystery(x.left) + mystery(x.right);
}
```

Function Breakdown:

- The base case is when the input node `x` is `null`, in which case the function returns `0`.
- For any other node, the function recursively calls itself on the left and right child nodes and adds the results of these two recursive calls.

This implies that the function is traversing the entire tree but never directly adds any value corresponding to the node itself, meaning it doesn't accumulate values based on node content, only on the structure (i.e., the nodes it encounters).

Formulating a hypothesis

`mystery()` returns the **number of non-leaf nodes** in the binary tree. This hypothesis is formed because:

- Leaf nodes have both left and right children as `null`, and when `x == null`, the function returns `0`.
- For every non-leaf node, the function will continue traversing through its children, accumulating returns until it reaches leaf nodes.

Let's consider a few sample binary trees and compute the value returned by `mystery()`:

Tree 1: Single node

1

Result: `mystery(1) = 0 + 0 = 0`

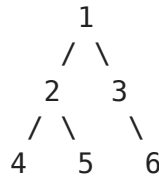
Tree 2: Three nodes

```

  1
 / \
2   3
```

Result: `mystery(1) = mystery(2) + mystery(3) = (0 + 0) + (0 + 0) = 0`

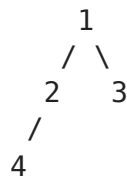
Tree 3: More complex tree



Result: $\text{mystery}(1) = \text{mystery}(2) + \text{mystery}(3) = (\text{mystery}(4) + \text{mystery}(5)) + (\text{mystery}(6) + 0) = ((0 + 0) + (0 + 0)) + ((0 + 0) + 0) = 0$

Sample Tree 4:

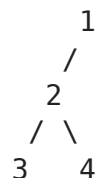
Consider the following tree:



- The function call starts with node 1.
- It recursively calls on node 2 (left child of 1) and node 3 (right child of 1).
- Node 2 then recursively calls on its left child 4, whose children are both null. So, $\text{mystery}(4)$ returns 0.
- Similarly, $\text{mystery}(3)$ returns 0 since node 3 is a leaf node with no children.

The return value of $\text{mystery}(1)$ would be $0 + 0 = 0$.

Sample Tree 5:



- Here, node 1 calls $\text{mystery}(2)$, and 2 calls on its children 3 and 4.
- Both 3 and 4 return 0 because they are leaf nodes.
- The return value of $\text{mystery}(1)$ would be $0 + 0 = 0$.

3. Proving the hypothesis

We can prove this hypothesis using mathematical induction.

Base case:

For a null node (empty tree), the function explicitly returns 0.

Inductive step:

Assume the hypothesis holds for all trees of height h or less. We'll prove it holds for a tree of height $h+1$.

For a tree of height $h+1$:

$$\text{mystery}(\text{root}) = \text{mystery}(\text{root.left}) + \text{mystery}(\text{root.right})$$

By our inductive hypothesis, both `mystery(root.left)` and `mystery(root.right)` return 0, as they are trees of height h or less.

Therefore:

$$\text{mystery}(\text{root}) = 0 + 0 = 0$$

This proves that the hypothesis holds for a tree of height $h+1$ if it holds for all trees of height h or less.

Conclusion:

By the principle of mathematical induction, we have proved that the `mystery()` function always returns 0 for any binary tree.

The function `mystery()` simply counts the non-leaf nodes in a tree recursively, but since it never adds any values for these nodes, it will always return 0. This is because the function only counts non-null nodes by recursively summing results from left and right subtrees, but never accumulates any positive values itself. Since the base case (null node) returns 0, all subsequent additions will only add up zeros, resulting in a final sum of 0 for any tree. To fix the function to count non-leaf nodes correctly, we would need to modify it to return `1 + mystery(x.left) + mystery(x.right)` for non-leaf nodes.