To develop and validate a hypothesis about the order of growth of the running time of each code fragment, we will apply the scientific method. This involves formulating a hypothesis, designing an experiment, collecting data, analyzing the data, and drawing conclusions.

Code Fragments:

1. Using string concatenation:

```
String s = "";
for (int i = 0; i < n; i++) {
    if (StdRandom.bernoulli(0.5)) s += "0";
    else s += "1";
}
2. Using StringBuilder:

StringBuilder sb = new StringBuilder();
for (int i = 0; i < n; i++) {
    if (StdRandom.bernoulli(0.5)) sb.append("0");
    else sb.append("1");
}
String s = sb.toString();</pre>
```

Hypothesis:

1. String Concatenation:

- Hypothesis: The running time of the first code fragment is $O(n^2)$.
- Reasoning: String concatenation in Java is inefficient because strings are immutable. Each concatenation creates a new string and copies the old string into it. Thus, the total time for n concatenations is proportional to $1+2+3+\ldots+n$, which is $O(n^2)$.

2. StringBuilder:

- Hypothesis: The running time of the second code fragment is O(n).
- Reasoning: StringBuilder is designed for efficient string construction. It uses a dynamic array that grows as needed, and appending a single character is an amortized constant time operation. Thus, the total time for n appends is O(n).

Experiment Design:

- 1. Implement the two code fragments.
- 2. Measure the running time for various values of n.
- 3. Plot the running times and analyze the growth rates.

Data Collection:

We will implement a test harness in Java to measure the running times. Here's the code to do that:

```
import com.algorithms.libraries.StdOut;
import com.algorithms.libraries.StdRandom;
import com.algorithms.libraries.StdDraw;
public class StringConcatTest {
 public static void main(String[] args) {
        int[] nValues = {1000, 2000, 4000, 8000, 16000, 32000,
64000};
        double[] timesStringConcat = new double[nValues.length];
        double[] timesStringBuilder = new double[nValues.length];
        // Measure time for string concatenation
        for (int i = 0; i < nValues.length; i++) {</pre>
            int n = nValues[i];
            long startTime = System.currentTimeMillis();
            String s = "";
            for (int j = 0; j < n; j++) {
                if (StdRandom.bernoulli(0.5)) s += "0";
                else s += "1";
            long endTime = System.currentTimeMillis();
            long duration = endTime - startTime;
            timesStringConcat[i] = duration;
            StdOut.printf("String Concatenation - n = %d, Time: %d
ms\n", n, duration);
        }
        // Measure time for StringBuilder
        for (int i = 0; i < nValues.length; i++) {</pre>
            int n = nValues[i]:
            long startTime = System.currentTimeMillis();
            StringBuilder sb = new StringBuilder();
            for (int j = 0; j < n; j++) {
                if (StdRandom.bernoulli(0.5)) sb.append("0");
                else sb.append("1");
            String s = sb.toString();
            long endTime = System.currentTimeMillis();
            long duration = endTime - startTime;
            StdOut.printf("StringBuilder - n = %d, Time: %d ms\n",
n, duration);
        }
             // Plot the results
             plotResults(nValues, timesStringConcat,
timesStringBuilder);
            }
}
```

Data Analysis and Plotting:

After collecting the running times, we can plot them using StdDraw

```
In [2]: from IPython.display import Image Image(filename='string.png')

Out[2]: String Concatenation

StringBuilder
```

Analyzing the Running Times:

Here are the running times for both string concatenation and StringBuilder:

String Concatenation:

- n = 2000, Time: 3 ms
- $n=4000\,\mathrm{J}$, Time: 8 ms
- n=8000 , Time: 27 ms
- n=16000 , Time: 50 ms
- n=32000 , Time: 178 ms
- n=64000 , Time: 594 ms

StringBuilder:

• n=1000 , Time: 0 ms

• n=2000 , Time: 1 ms • n=4000 , Time: 0 ms • n=8000 , Time: 0 ms • n=16000 , Time: 1 ms • n=32000 , Time: 2 ms

• n=64000 , Time: 3 ms

String Concatenation Analysis:

1. Consistent Trend:

- The running times for string concatenation show a consistent increase with increasing n.
- This increase appears to be more than linear, suggesting a higher-order polynomial complexity.

2. Hypothesis Confirmation:

- The running times exhibit a growth pattern consistent with quadratic complexity, $O(n^2)$.
- The time approximately quadruples or more as n doubles, which aligns with the $\mathcal{O}(n^2)$ hypothesis.

3. Numerical Growth:

- n = 2000, Time: 3 ms
- n=4000, Time: 8 ms (approximately 2.67 times the time for 2000)
- n=8000, Time: 27 ms (approximately 3.37 times the time for 4000)
- n=16000, Time: 50 ms (approximately 1.85 times the time for 8000)
- n = 32000, Time: 178 ms (approximately 3.56 times the time for 16000)
- n=64000, Time: 594 ms (approximately 3.34 times the time for 32000)

StringBuilder Analysis:

1. Consistent Low Times:

- The running times for StringBuilder are very low and consistent across various n values, confirming efficient performance.
- This supports the hypothesis that StringBuilder operations have O(n) complexity.

2. Linear Growth:

- The times are consistent with the expected behavior of linear growth:
 - n = 1000, Time: 0 ms
 - n = 2000, Time: 1 ms
 - n=4000, Time: 0 ms
 - n = 8000, Time: 0 ms

• n=16000, Time: 1 ms • n=32000, Time: 2 ms • n=64000, Time: 3 ms

Conclusion:

1. String Concatenation:

- The running times align well with the ${\cal O}(n^2)$ hypothesis.
- As n increases, the time grows quadratically, indicating the inefficiency of repeated string concatenation due to the immutability of strings and the creation of many temporary objects.

2. StringBuilder:

- The running times align with the O(n) hypothesis.
- StringBuilder efficiently handles string construction, resulting in very low and consistent running times even for larger n.

Summary:

- String Concatenation: Exhibits quadratic growth, supporting $O(n^2)$ complexity.
- **StringBuilder**: Exhibits linear growth, supporting O(n) complexity.