

# G6021: Comparative Programming

Some more exercises on the  $\lambda$ -calculus

1. Insert all the missing parentheses and  $\lambda$ 's into the following abbreviated  $\lambda$ -terms, where  $I = \lambda x.x$ .

- $\lambda xy.x$
- $\lambda xy.xyy$
- $\lambda xyz.xz(yz)$
- $\lambda xyz.x(yz)$
- $\lambda xyz.xzy$
- $\lambda xy.x(xy)$
- $\lambda x.xI$
- $\lambda x.xII$
- $\lambda x.xIII$
- $(\lambda xy.y(xxy))(\lambda xy.y(xxy))$
- $\lambda f.(\lambda x.f(xx))(\lambda x.f(xx))$

2. Let  $I = \lambda x.x$ ,  $B = \lambda xyz.x(yz)$ ,  $C = \lambda xyz.xzy$ ,  $T = \lambda xy.x(xy)$  and  $Y = (\lambda xy.y(xxy))(\lambda xy.y(xxy))$ .

Reduce the following terms to normal form, if they have one. You do not need to draw a reduction graph.

- $CI$
- $C(\lambda x.xII)I$
- $Y$
- $TBCI$
- $TI$
- $TT$
- $TTT$  (Just joking - this takes 42 beta reductions...)