

A: Zoom In

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- 1 A: Zoom In
- 2 C: Escher Sequence
- 3 F: Cow's Landing
- 4 E: Color Rur
- 5 B: Building for Days
- 6 D: Attend Dance



D: Attend Dance



Problem A: Zoom In

Author: Ryan Hechenberger Just do IT!



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Author: Shizhe Zhao

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■ Brute-Force: try all possible length from n to 2n: $O(n^2)$





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- Brute-Force: try all possible length from n to 2n: $O(n^2)$
- Observation: let s2=flip(reverse(s)), then concated string s s2 must be an Escher Sequence.
- Minimize the length by removing their maximum common part: maximize sufix(s) == prefix(s2).
- Can be done by KMP. (Or other string matching algorithms)





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Let given 2^k by 2^k grid map, there are exactly 2^{2k-2} number of cows in each quadrant.





Problem F: Cow's Landing

Author: Shizhe Zhao

- Let given 2^k by 2^k grid map, there are exactly 2^{2k-2} number of cows in each quadrant.
- Find which quadrant the cow belongs to, and we will deal with a same problem on 2^{k-1} by 2^{k-1} grid map.





Problem F: Cow's Landing

Author: Shizhe Zhao

- Let given 2^k by 2^k grid map, there are exactly 2^{2k-2} number of cows in each quadrant.
- Find which quadrant the cow belongs to, and we will deal with a same problem on 2^{k-1} by 2^{k-1} grid map.
- Repeat such process until k=0.





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 - Let the first color segment be [1, i], maximizing i doesn't affect global minimum.
 - \blacksquare So just choose a color that maximize i.
- It is the compression algorithm in Compress Path Database.





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Problem B: Building for Days

Author: Ryan Hechenberger Determine:

- whether a cell is traversable
- the type of a traversable cell: vertical hallway, horizontal hallway, hallway intersection, room.
- the the type of a non-traversable cell





How to determine the typ of a non-traversable cell?







How to determine the typ of a non-traversable cell? There must be in same connected component!





Problem B: Building for Days

How to determine the typ of a non-traversable cell? There must be in same connected component! FloodFill:

```
def floodfill(r, c):
    # over border
    if not (0 <= r and r < h and 0 <= c and c < w):
        return

g[r][c] = '#'
floodfill(r-1, c)
floodfill(r, c-1)
floodfill(r, c+1)</pre>
```



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Problem D: Attend Dance

Author: Shizhe Zhao If no-cow is lying:



Problem D: Attend Dance

Author: Shizhe Zhao If no-cow is lying:

■ then $|A| = max(A_i|0 \le i < n)$ (same for |B| and |C|)





Author: Shizhe Zhao If no-cow is lying:

- then $|A| = max(A_i|0 \le i < n)$ (same for |B| and |C|)
- let C_a be the number of $A_i = |A|$, then $C_a|A|$





Problem D: Attend Dance

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Problem D: Attend Dance

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- then $|A| = max(A_i|0 \le i < n)$ (same for |B| and |C|)
- let C_a be the number of $A_i = |A|$, then $C_a|A|$
- ...
- identify these constraints, and all of them must be true.



