

On the number of crossing-free perfect matchings with designated left and right endpoints

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Abstract

[TODO main results: the spc bound, the bound dependent on the number of interior points. Do I mention the expressions for the Catalan triangle (are they known)?]

Introduction

In the study of plane embeddings of planar graphs, also known as plane graphs, embeddings with non-crossing straight edges (called crossing-free geometric graphs) are of particular interest. For instance, it is known [10] that any optimal solution to the travelling salesman problem in the Euclidean plane has non-crossing straight edges. In 1936, K. Wagner proved¹ [CITE] that any planar graph has such an embedding.

A natural enumerative question is then, given a vertex set of n points in the Euclidean plane, to bound the number of plane graphs of a certain sort (cycles, matchings, etc.) which can be drawn with non-crossing straight edges. For spanning cycles, the question was introduced by M. Newborn and W. O. J. Moser in 1976 [6], and a super-exponential upper bound was given. In 1980, M. Ajtai, V. Chvátal, M. Newborn, and E. Szemerédi [1] proved that there are only exponentially many crossing-free geometric graphs on a given point set.

Exponential bounds for various sorts of graphs were then improved, in particular for triangulations—the maximal crossing-free geometric graphs. Indeed, finding a bound for them yields a bound for the overall number of crossing-free geometric graphs. [TODO citation] In addition, bounds on the number of triangulations are of interest in geometric modeling [CITE Studies in computational geometry motivated by mesh generation].

Some special configurations of points are better understood; in particular, for point sets in convex position, the numbers of crossing-free perfect matchings, triangulations, and spanning cycles are known. In fact, the study of point sets in convex position vastly predates the general problem, with the number of triangulations given by Euler in 1751 [4].

For perfect matchings, an optimal lower bound of $C_{\frac{n}{2}}$ was obtained as early as 1995 by A. García, M. Noy, and J. Tejel [5]; upper bounds have been harder to come by, with the best one yet being asymptotically 10.05^n , given by M. Sharir and E. Welzl in 2005 [9] In that same work, an upper bound of 5.83^n is given for the number of crossing-free perfect matchings on a point set whose points are designated as left or right endpoints; moreover, an upper bound of 4^n is given when all the left endpoints are to the left of the right endpoints (matchings across a line).

In this thesis, we start by giving an overview of some of the existing results, in particular a way to classify arrangements of designated left and right endpoints, and we restate the proof of the bound for matchings across a line by M. Sharir and E. Welzl (section 1). We then study in detail the inequalities in that bound and the objects it actually counts, yielding criteria that will allow us to tighten the bound in specific

¹This result was independently proved by I. Fáry in 1948 [CITE], and by S. K. Stein in 1951 [CITE]; it is generally known as Fáry's theorem.

cases (section 2). Using these criteria, we give a bound for matchings across a line that depends on the number of points in the interior of the convex hull (section 3). We then give new bounds on the number of crossing-free perfect matchings on a point set with designated left and right endpoints for a specific family of left-to-right orderings of the left and right endpoints (4).

1 Background

We first give some definitions, as well as existing proofs whose ideas we will use to derive our results.

1.1 Geometric graphs

Definition (geometric graph). Given a set of points P in the Euclidean plane \mathbb{R}^2 , a *geometric graph* is a collection of straight line segments (edges) whose endpoints are elements of P .

It can be described as a simple graph (in the combinatorial sense) on the vertices P , where the edge $\{v, w\}$ corresponds to the segment joining v and w . \diamond

Definition (crossing-free). A geometric graph is *crossing-free* if no two edges share points other than their endpoints; it is called *crossing* otherwise. \diamond

Note that being crossing-free implies that the corresponding simple graph is planar, and that the geometric graph is a plane embedding.

Definition (triangulation). A *triangulation* is a maximal crossing-free geometric graph, that is, geometric graph such that for all v and w in P that are not joined by a segment, adding the segment joining v and w would result in a crossing geometric graph. \diamond

Note that the faces (in the sense of plane graphs) formed by a triangulation are all triangles, with the possible exception of the outer face (thus this definition is *not* equivalent to that of a triangulation of the 2-sphere).

Since a geometric graph corresponds to a simple graph on the underlying point set, we can also look at geometric graphs that belong to special classes of simple graphs. As we will only be focusing on geometric graphs, we simply call *perfect matching* a geometric graph which corresponds to a perfect matching in the combinatorial sense.

As outlined in the introduction, we are interested in statements regarding the number of geometric graphs, possibly restricted to a certain type, on a given point set. Evidently, that number would depend on the choice of the point set P , so instead we seek bounds on that number depending on the cardinality $|P|$, and possibly restricting P so that it satisfies certain properties.

An important property that P can have is general position. We say that a point set in the plane is *in general position* if no three points are collinear. In this we will only study point sets in general position.

In general, if $g(P)$ is the number of geometric graphs of a certain sort on the point set P , we will look for lower bounds l and upper bounds u of the form

$$\forall n \in \mathbb{N}, \forall P \text{ such that } |P| = n, l(n) \leq g(P) \leq u(n),$$

where P runs over all point sets that satisfy the relevant properties.

In addition, we will want asymptotics on such bounds l and u , to compare them with existing results; since the bounds are often exponential, we tend to ignore polynomial factors; we will thus write that

$$f(n) \preccurlyeq u(n)$$

if $f(n) \leq p(n)u(n)$ for some polynomially-bounded p .

1.2 Triangulations of convex point sets and the Catalan numbers

A set of points is said to be *in convex position* if it is in general position and all points are in the boundary of its convex hull.

For points in convex position, the numbers of triangulations and perfect matchings are known. The following recurrence, given by J. A. Segner in 1758 [8], defines the *Catalan numbers*.

Theorem (Segner). *Let C_n be the number of triangulations of a set of $n + 2$ points in convex position. Then $C_0 = 1$, and for $n > 0$,*

$$C_n = \sum_{i=0}^{n-1} C_i C_{n-i-1}. \quad (1.1)$$

Proof. Consider a set of $n + 2$ points in convex position. For $n = 0$ the result is trivial, the only triangulation is an edge.

Let $n > 0$, and number the points along the boundary of the convex hull, from 0 to $n + 1$. The edge e joining points $n + 1$ and (n) is in the boundary of the convex hull, and thus must be in any triangulation, since no other edge can cross it. Let i be the index of the third point of the triangle containing edge e , which may be any of the others, from 0 to $n - 1$.

Then, restricting the triangulation to the $i + 2$ points $\{0, \dots, i, n + 1\}$ yields a triangulation of those points. Similarly, restricting the triangulation of the whole point set to points i through n yields a triangulation of $n - i + 1$ points. It follows that $C_n \leq \sum_{i=0}^{n-1} C_i C_{n-i-1}$.

Moreover, given a triangulation of $\{0, \dots, i, n + 1\}$ and a triangulation of $\{i, \dots, n\}$, adding the edge e forms a triangulation of the whole point set, proving equality. \square

An explicit expression for C_n was given earlier by Euler [4], along with the generating function. A convenient form is

$$C_n = \frac{1}{n+1} \binom{2n}{n}, \quad (1.2)$$

yielding asymptotically

$$C_n \sim \frac{4^n}{n^{\frac{3}{2}} \sqrt{\pi}} \leq 4^n. \quad (1.3)$$

1.3 Crossing-free perfect matchings

We will call PM_P the set of perfect matchings on the point set P , and CFPM_P the set of crossing-free perfect matchings on the point set P .

[TODO talk about the number of crossing-free perfect matchings in convex position, give the proof; find the Alfred Errera thing from the octavo of the académie royale de belgique, reference 10 in Sharir–Welzl; this might require a trip to the library?]

For a set P of n points in general position, it was proved by Micha Sharir and Emo Welzl in 2005 [9] that $|\text{CFPM}_P| \leq 10.05^n$.

When talking about a perfect matching, we will denote the unique edge incident to point p by e_p .

1.4 Bracket expressions and their relations to perfect matchings

We say that a point set is *in general position with respect to the horizontal* if it is in general position and no two points lie on a vertical line. Note that any point set in general position can be put in general position with respect to the horizontal by an arbitrarily small rotation. Moreover, note that points in general position with respect to the horizontal are ordered from left to right.

This ordering can be used to classify perfect matchings on the point set; to this end, we introduce bracket expressions. A *bracket expression* of length n is a sequence

Theorem (Ruiz-Vargas–Welzl). *Let P be a point set of size $n = 2k$ in general position with respect to the horizontal, and let B be a well-formed bracket expression of length n . Then there exists a crossing-free perfect matching consistent with B ; in other words, $v_P(B) \geq 1$.* \square

Proof. Let μ_0 be a perfect matching on P consistent with B . This is always possible, for instance, parsing the bracket expression, match the point corresponding to an opening bracket and the point corresponding to the matching closing bracket.

Define $l(m)$ for a perfect matching m on P to be the sum of the lengths of the edges of m .

Then, repeat the following procedure, starting at $i = 0$. If there is no crossing in μ_i , we have found a perfect matching with the desired properties. If there is a crossing, let a, b, c , and d be the points involved, so that the edge ab crosses the edge cd . Remove these edges, and replace them by ad and cb (thus untangling them). This yields another perfect matching μ_{i+1} . By the triangle inequality (see figure 2), $l(\mu_{i+1}) < l(\mu_i)$.

If this procedure did not terminate, it would yield a sequence $(\mu_i)_{i \in \mathbb{N}}$ of crossing perfect matchings on P on which l is strictly decreasing, thus an infinite sequence of distinct graphs on P . Since there are only finitely many graphs on P , this is a contradiction, so we eventually⁴ find a crossing-free perfect matching. \square

This immediately yields a lower bound for the number of crossing-free perfect matchings, since there are $C_{\frac{n}{2}}$ well-formed bracket expressions of length n . Note that this lower bound had already been proved by A. García, M. Noy, and J. Tejel in 1995, by a recurrence that split the point set, rather than via well-formed bracket expressions [5].

Corollary. *Let P be a point set of even size n in general position. There are at least $C_{\frac{n}{2}}$ distinct crossing-free perfect matchings on P , i.e., $|\text{CFPM}_P| \geq C_{\frac{n}{2}}$.* \square

Moreover, this lower bound is optimal, since it is attained if P is in convex position (see section 1.3).

Now that we have a tight uniform lower bound for v_P , we will start looking at upper bounds dependent on the bracket expression.

1.6 Matchings across a line

Again we consider $n = 2k$ points in general position with respect to the horizontal. The matchings consistent with bracket expressions with k opening brackets followed by k closing brackets, $\langle^k \rangle^k$, are called *matchings across a line*. Indeed, any edge in such a matching will cross any vertical line that separates the left-points from the right-points.

The following result, and its proof, were given by Micha Sharir and Emo Welzl in 2005 [9].

Theorem (Sharir–Welzl). *Let P be a set of $n = 2k$ points in general position with respect to the horizontal. Then there are at most $C_{\frac{n}{2}}^2$ crossing-free perfect matchings across a line on P , i.e., $v_P(\langle^k \rangle^k) \leq C_k^2$.*

The idea of the proof is as follows. First, pick a vertical line that separates the left-points from the right-points; we will call it *the vertical line*. Further, let us call the set of left-points L and the set of right-points R .

A perfect matching consistent with a given well-formed bracket expression (and in particular, a perfect matching across a line) is uniquely defined by a bijection $\mu : L \rightarrow R$ from the left-points to the right-points. Consider such a perfect matching across a line, and let E be the set of its edges. Now, number the intersections between the edges of the perfect matching and the vertical line from top to bottom. This yields a numbering $\iota : E \rightarrow [k]$ of the edges.

⁴The untangling procedure itself predates this proof, and was shown to terminate in $\mathcal{O}(n^3)$ steps by J. van Leeuwen and A. A. Schoone in 1980 [10].

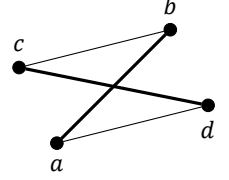


Figure 2. Untangling in a perfect matching. Replacing the thick edges by the thin ones reduces the total edge length, while preserving left and right endpoints.

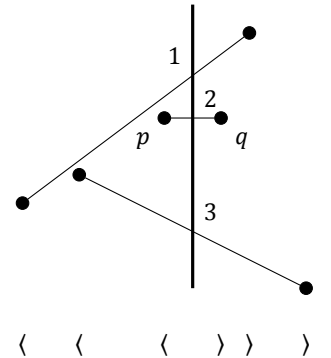


Figure 3. Constructing the numberings λ and ρ of the left- and right-points from a crossing-free perfect matching across a line; here $\iota(\{p, q\}) = 2$, and thus $\lambda(p) = \rho(q) = 2$.

Define $\lambda(l) := \iota(e_l)$ mapping a left-point to the intersection number of its edge, and similarly $\rho(r) := \iota(e_r)$ for the right-points. We have $\mu = \rho^{-1} \circ \lambda$.

The bijection λ (respectively ρ) determines the order in which the left points (respectively right points) reach the vertical line.

If the matching is crossing-free, we will show that λ and ρ have to be in sets of size C_k , thus that there are at most C_k^2 functions $\mu = \rho^{-1} \circ \lambda$, i.e., that there can be at most C_k^2 crossing-free perfect matchings across a line.

Since we are going to reuse these concepts in subsequent sections, we will formalize and name the properties of λ and ρ that we will consider.

Definition (crossing-free alignment). Let Λ be a set of k points in general position. Let V be a vertical line to the right of Λ . A bijection λ from Λ to $[k]$ is a *crossing-free left-alignment of Λ on V* if there exists a set $A = \{a_p \mid p \in \Lambda\}$ of k line segments such that the following hold:

1. for every $p \in \Lambda$, the left endpoint of a_p is p , and its right endpoint is on V ;
2. no two segments cross, and no two segments share any endpoints;
3. let $\iota : A \rightarrow [k]$ the numbering of the right endpoints of the segments in A from top to bottom, then $\lambda(p) = \iota(a_p)$.

Such a set of segments is said to *realize λ* .

Correspondingly, we define a *crossing-free right-alignment*, where “right” and “left” are swapped in the definition. \diamond

A crossing-free left-alignment of five points, with a realizing set of segments—drawn as arrows—is shown in figure 4

Lemma. Let Λ be a set of k points in general position, and let V be a vertical line to the right of Λ . There are at most C_k crossing-free left-alignments of Λ on V .

Proof. Let l be the leftmost point of Λ , and let λ be a crossing-free left-alignment of Λ on V . Let $A = \{a_p \mid p \in \Lambda\}$ be a set of non-crossing line segments realizing λ . For a point $p \in \Lambda$, we will call a_p *the segment of p* .

$\lambda(l)$, being the index on the vertical line V of the right endpoint of the segment a_l of the leftmost point, is equal to one plus the number of points of Λ that are above a_l . Indeed, the segments of points above a_l must themselves reach the vertical line V above a_l , otherwise they would cross a_l , and correspondingly for points below a_l , so that there are as many edges reaching the vertical line V below a_l as there are points below a_l .

Moreover, as the oriented angle between a_l and the horizontal increases, points are only added to the set of points below a_l , so that choosing the number of points below a_l determines the sets of points below and above a_l .

Further, since points above a_l must reach the vertical line V above a_l , and correspondingly for points below, if the point p is above a_l , then $\lambda(p) < \lambda(l)$, and if it is below, $\lambda(p) > \lambda(l)$. Thus λ restricted to the points above a_l is a crossing-free left-alignment of the $\lambda(l) - 1$ points above a_l , and $\lambda - \lambda(l)$ restricted to the points below a_l is a crossing-free left-alignment of the $k - \lambda(l)$ points below a_l , where the appropriate subsets of A provide the set of segments required in the definition.

It follows that λ is determined by the choice of $\lambda(l)$ and crossing-free left-alignments of $\lambda(l) - 1$ and $k - \lambda(l)$ points. Letting $i := \lambda(l) - 1$, these are crossing-free alignments of i and $k - i - 1$ points (see figure 4). Therefore, if ϖ_j is an upper bound for the number of crossing-free left-alignments of j points when $j < k$, we can give an upper bound on the number of crossing-free left-alignments of k points,

$$\varpi_k := \sum_{i=0}^{k-1} \varpi_i \varpi_{k-i-1}.$$

We can start the recurrence with $\varpi_0 = 1$; this is the recurrence for the Catalan numbers, thus $\varpi_k = C_k$. \square

The same results holds for crossing-free right-alignments, simply exchange “left” and “right” in the above lemma and its proof.

Proof (of the theorem). If we have a crossing-free perfect matching across a line, then λ as constructed above is a crossing-free left-alignment of L on the vertical line; indeed, letting a_p be the portion of the edge e_p to the left of the vertical line yields segments that realize λ . Similarly, ρ is a crossing-free right-alignment of R on the vertical line. \square

2 Analysing the overcounting in the upper bound for matchings across a line

The bound given above for the number of matchings across a line is not optimal; an example is given by Sharir and Welzl of a crossing-free left-alignment λ and a crossing-free right-alignment ρ such that $\rho^{-1} \circ \lambda$ is not crossing-free (see figure 5).

There is however more to the overcounting than just counting some crossing matchings. To characterize that, it helps to name some functions and sets. Again, are considering a set P of $2k$ points in general position with respect to the horizontal and the bracket expression $\langle^k \rangle^k$. We let L be the set of left-points, and R be the set of right-points. Moreover, we pick a vertical line separating L and R , and call it V .

In the previous section, given a crossing-free perfect matching μ , we constructed crossing-free left- and right-alignments λ and ρ (see figure 3). Let us call that construction ω , thus $\omega(\mu) = (\lambda, \rho)$. Moreover, let us call $\text{CFLA}_{L,V}$ the set of crossing-free left-alignments of L on V , and $\text{CFRA}_{R,V}$ the set of crossing-free left-alignments of R on V . Given sets of equal sizes L and R , $\lambda \in \text{CFLA}_{L,V}$, and $\rho \in \text{CFRA}_{R,V}$, let $c(\lambda, \rho)$ be the (possibly crossing) perfect matching $\rho^{-1} \circ \lambda$. Finally, let us name the sets of perfect matchings across a line and of crossing-free perfect matchings across a line,

$$\begin{aligned} \text{PMAL}_P &:= \beta_P^{-1}(\langle^k \rangle^k) \\ \text{CFPMAL}_P &:= \text{PMAL}_P \cap \text{CFPM}_P. \end{aligned}$$

Then the following diagram commutes, *i.e.*, constructing crossing-free alignments from a matching and then composing them yields the initial matching.

$$\begin{array}{ccc} \text{CFPMAL}_P & \xrightarrow{\omega} & \text{CFLA}_{L,V} \times \text{CFRA}_{R,V} \\ \downarrow \subseteq & \swarrow c & \\ \text{PMAL}_P & & \end{array}$$

The result of Sharir and Welzl consists in deducing from the commutative diagram that ω must be injective, and so that

$$|\text{CFPMAL}_P| \leq |\text{CFLA}_{L,V} \times \text{CFRA}_{R,V}|, \quad (2.1)$$

and then a bound for the number of crossing-free alignments yields a bound for the number of crossing-free matchings across a line.

The inequality (2.1) is strict, as shown in figure 5. We would thus like to study the pairs of crossing-free alignments which are *not* in the image of ω , since they are responsible for that excess; we will call them *overcounted*. Overcounted pairs of crossing-free alignments fall into one of two categories;

1. evidently, if $c(\lambda, \rho)$ is crossing, then (λ, ρ) is not in the image of ω , since the diagram commutes;
2. on the other hand, figure 6 shows a pair of crossing-free alignments (λ, ρ) which composes to a crossing-free alignment $\mu = c(\lambda, \rho)$, but which is not equal to $\omega(\mu)$.

In addition, we will call a crossing-free left alignment λ overcounted if any pair (λ, ρ) is overcounted.

TODO

Figure 5. Composing crossing-free alignments, only to get a crossing perfect matching.

TODO

Figure 6. Composing crossing-free alignments, and getting a crossing-free perfect matching, but one that does not decompose to the original alignments.

To hunt down overcounted crossing-free alignments, we will examine in more detail the anatomy of a crossing-free left-alignment of L and the proof of the Sharir–Welzl bound for the number of crossing-free left-alignments, with some fresh definitions.

First, for a point set L , note that we can partition the half-plane to the right of $\text{lm}(L)$ in $|L|$ regions, separated by rays emanating from $\text{lm}(L)$ and going through other points of L . We will call those regions *left L -cones*, and number them from top to bottom from 1 to $|L|$, see figure ?? . Note that there are i points of L above the interior of left L -cone number $i + 1$, and $|L| - i - 1$ points below.

Given a crossing-free left-alignment λ of L on V , for any set of segments $A = \{a_p \mid p \in L\}$ realizing λ , the segment $a_{\text{lm}(L)}$ of the leftmost point then has to be contained in left L -cone number $\lambda(\text{lm}(L))$. As we remarked in the previous section, for a crossing-free perfect matching μ , the edges of μ restricted to the left side of V realize the left-alignment in $\omega(\mu)$. Thus if (λ, ρ) is in the image of ω , $\text{lm}(L)$ has to be matched to a right-point in left L -cone number $\lambda(\text{lm}(L))$, see figure ?? .

To extend this reasoning beyond the leftmost point of L , we define *left conifers of L* . A left conifer κ of a point set L is a map from points in L to left S -cones for subsets S of L , satisfying the following properties:

1. $\kappa(\text{lm}(L))$ is a left L -cone;
2. κ restricted to the points above the interior of $\kappa(\text{lm}(L))$ is a left conifer of these points;
3. κ restricted to the points below the interior of $\kappa(\text{lm}(L))$ is a left conifer of these points.

See figure ?? for an example. We call LC_L the set of left conifers of L . From the observations above on the number of cones and the size of the sets of points above and below their interiors, if ϖ_k is the number of left conifers of k points, we have

$$\varpi_k = \sum_{i=0}^{k-1} \varpi_i \varpi_{k-i-1} \text{ and } \varpi_0 = 1,$$

so that $|\text{LC}_L| = C_{|L|}$.

Recall that a crossing-free left-alignment λ is determined by the value of $\lambda(\text{lm}(L))$ and by crossing-free left alignments of the points above and below any edge realizing λ , *i.e.*, above and below cone number $\lambda(\text{lm}(L))$. This is equivalent to choosing left L -cone number $\lambda(\text{lm}(L))$, and recursively doing so in the point subsets above and below the interior of that cone, see figure ?? . This gives an injective map $\mathfrak{K} : \lambda \mapsto \mathfrak{K}_\lambda$ from crossing-free left-alignments of L on any line to left conifers. Note that it is not bijective, see figure ?? . In turn, this yields $|\text{CFLA}_{L,V}| \leq |\text{LC}_L| = C_{|L|}$, which is exactly the Sharir–Welzl bound.

Now, note that the map \mathfrak{K} has the property that for any set of segments $A = \{a_p \mid p \in L\}$ realizing λ , a_p lies inside the cone $\mathfrak{K}_\lambda(p)$; since the edges of a matching μ realize the crossing-free alignments $\omega(\mu)$, $\mathfrak{K}_\lambda(p)$ must contain the right-point to which p is matched.

3 Highly convex matchings across a line

In particular, this means that if there is no right-point in $\mathfrak{K}_\lambda(p)$, then λ is overcounted. We will call *vacuous* a left cone that does not contain any right-points, and we will call *vacuous* a left conifer that has vacuous cones; thus λ is overcounted if \mathfrak{K}_λ is vacuous.

We will now use convexity considerations to give an upper bound on the number of non-vacuous left conifers. Let us first define some concepts related to convex hulls. Given a set of points P in general position, the *extreme points of P* are the points of P that are in the boundary of the convex hull of P . Given a set of points L in general position with respect to the horizontal and a set of points R in general position with respect to the horizontal to the right of L , the *R -liminary points of L* are the points

of L which are extreme points of $L \cup R$, see figure ?? . We will call $\text{lm}(L)$ the leftmost point of L ; note that it is an R -liminary point of L .

Note that in polar order around $\text{lm}(L)$, the R -liminary points of L other than $\text{lm}(L)$ are separated by R , that is, in polar order, one subset of the R -liminary points of L appears first, followed by all points of R , followed by the rest of the R -liminary points of L (except $\text{lm}(L)$). We will call those R -liminary points of L which appear below R the *lower R -liminary points of L* , and those that appear above R the *upper R -liminary points of L* .

If there are δ_1 upper R -liminary points of L , then the first δ_1 left L -cones are vacuous. Similarly, if there are δ_2 lower R -liminary points of L , then the last δ_2 left L -cones are vacuous. Moreover, for any non-vacuous left L -cone ℓ , any upper R -liminary points of L are R -liminary points of the subset of L above ℓ , and any lower R -liminary points of L are R -liminary points of the subset of L below ℓ , so that we get the recurrence

$$\varpi_{k,\delta} := \max_{\delta_1 + \delta_2 = \delta - 1} \sum_{i=\delta_1}^{k-\delta_2-1} \varpi_{i,\delta_1} \varpi_{k-i-1,\delta_2},$$

where $\varpi_{k,\delta}$ is an upper bound on the number of crossing-free non-vacuous left conifers of k points of which at least δ are R -liminary. Moreover, we can start that recurrence with

$$\begin{aligned} \varpi_{k,0} &:= \varpi_{k,1}, \\ \varpi_{k,k} &:= 1, \end{aligned}$$

since there is at least one R -liminary point of L (the leftmost one), and if all points are R -liminary, only one cone is non-vacuous.

4 Three changes of bracket direction

We now consider left-right matchings corresponding to bracket expressions which have three changes of bracket directions, i.e., bracket expressions of the form $\langle^k \rangle^l \langle^q \rangle^p$, k opening brackets, l closing brackets, q opening brackets, p closing brackets, where $k - l = p - q \geq 0$ and $k + l + p + q = n$.

In term of points, this means that four sets can be separated by vertical lines, from left to right, k left-points forming the set K , l right-points forming L , q left-points forming Q , and p right-points forming P . We pick a vertical line separating K and L and call it *the left line*, and we pick a vertical line separating P and Q and call it *the right line*.

Given a crossing-free perfect matching on those points, numbering from top to bottom the intersections between edges incident to the points of K and the left line, we get a crossing-free left-alignment of K . $k - l$ of the k edges on this vertical line are incident to points in P ; the other l are incident to points in L . Numbering those l edges yields a crossing-free right-alignment λ of L .

Similarly on the right side, we get a crossing-free right-alignment of π of P , numbering the subset of edges incident to a point in P and a point in Q , we get a left-alignment θ of q .

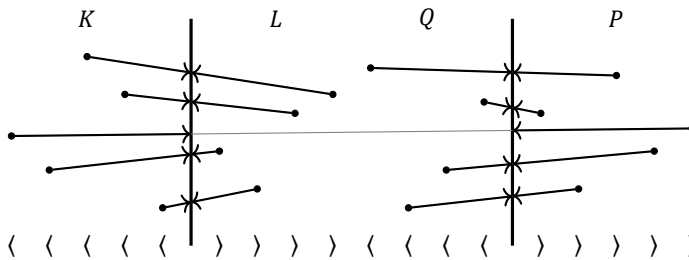


Figure 7. Four crossing-free alignments.

The matching is uniquely determined by κ , π , the choice of the $k - l$ among k points and $k - l = p - q$ among p points that get matched to each other, and by λ and θ , which gives the following bound for the number of these matchings:

$$v_P(\langle^k \rangle^l \langle^q \rangle^p) \leq C_k \binom{k}{l} C_l C_q \binom{p}{q} C_p.$$

Asymptotically, the factor involving k and l is

$$C_k \binom{k}{l} C_l \leq 4^{k+l} \binom{k}{l} \leq 4^{k+l} \left(\left(\frac{(1-\alpha)^{1-\frac{1}{\alpha}}}{\alpha} \right)^{\left(\frac{\alpha}{\alpha+1} \right)^{k+l}} \right),$$

where $l = \alpha k$. The base of the exponential bound for the binomial coefficient is maximal when $\alpha = \frac{3-\sqrt{5}}{2}$, where it is $\phi = \frac{1+\sqrt{5}}{2}$. This gives the overall bound

$$v_P(\langle^k \rangle^l \langle^q \rangle^p) \leq 4^{k+l} \phi^{k+l} 4^{p+q} \phi^{p+q} = (4\phi)^n \approx 6.472^n.$$

4.1 Improving the binomial bound

We can however improve upon that bound: indeed, once κ , π , and the set of points of K matched to P are fixed, the edges of the matching that cross both the left line and the right line—let us call these *long edges*—are determined. As a result, the region between the left line and the right line is partitioned in trapezoidal cells, and the portion of any edge from K to L and from P to Q that lies between the left line and the right line is confined to a single of those cells. It follows that λ is composed of crossing-free right-alignments of subsets of L separated by the long edges, and similarly for θ with subsets of Q (in figure 7, there is one long edge, and thus two cells).

Let us look at the edges crossing the left line (the same argument applies to the right line), numbered from top to bottom: $k - l$ of those are long edges, let $S \subseteq [k]$ be their numbers; in between two long edges, above the first long edge, and below the last one, we have the edges that define the crossing-free right-alignments that make up λ . It follows that a crossing-free right-alignment of m points that makes up λ corresponds to a maximal sequence of m consecutive elements of $[k] \setminus S$.

We will call the set of maximal sequences of consecutive elements of S' the *cells* of S' , written $\text{cells}(S')$.

Then, we can improve the $\binom{k}{l} C_l$ factor in the bound (in which the binomial comes from the choice of the long edges amongst the k edges on the left line, and the Catalan number comes from the choice of λ), summing over the choices of the long edges (and thus of S above). The improved factor becomes

$$\text{spc}(k, l) := \sum_{S \in \binom{[k]}{k-l}} \prod_{c \in \text{cells}([k] \setminus S)} C_{|c|} = \sum_{S' \in \binom{[k]}{l}} \prod_{c \in \text{cells}(S')} C_{|c|}, \quad (4.1)$$

and the overall bound becomes

$$v_P(\langle^k \rangle^l \langle^q \rangle^p) \leq C_k \text{spc}(k, l) \text{spc}(p, q) C_p. \quad (4.2)$$

4.2 A recurrence

In order to compute spc efficiently, and eventually, get its asymptotics, it is useful to get rid of the cells function. We can express spc as a recurrence instead. First, we note that $\text{spc}(k, k) = C_k$: there is only one summand, S' is the whole set, so it has only one cell, namely itself. Otherwise, $k - l \geq 1$; in the sum over the S , consider the greatest element j of S , which is at least $k - l$, and split the sum over that,

$$\text{spc}(k, l) = \sum_{j=k-l}^k \sum_{S \in \binom{[k]}{k-l}} \prod_{c \in \text{cells}([k] \setminus S)} C_{|c|}.$$

For fixed j , all summands (of the sum over S) will have a factor with $c = \{j+1, \dots, k\}$, and thus a factor of C_{k-j} . Factoring out this C_{k-j} , we get

$$\text{spc}(k, l) = \sum_{j=k-l}^k C_{k-j} \sum_{\substack{S \subseteq \binom{[k]}{k-l} \\ j = \max S}} \prod_{\substack{c \in \text{cells}([k] \setminus S) \\ c \neq \{j+1, \dots, k\}}} C_{|c|}.$$

Now, note that choosing a subset S of $[k]$ of size $k-l$ whose maximum is j is equivalent to choosing a subset s of $[j-1]$ of size $k-l-1$, where $S = s \cup \{j\}$. Moreover, the cells of $[k] \setminus S$ other than $\{j+1, \dots, k\}$ are exactly the cells of $[j-1] \setminus s$, thus

$$\text{spc}(k, l) = \sum_{j=k-l}^k C_{k-j} \sum_{\substack{s \subseteq \binom{[j-1]}{k-l-1} \\ j = \max s}} \prod_{c \in \text{cells}([j-1] \setminus s)} C_{|c|}.$$

By definition of spc , this means

$$\text{spc}(k, l) = \sum_{j=k-l}^k C_{k-j} \text{spc}(j-1, l+j-k).$$

Rewriting this as a sum over $i := k-j$, this gives us the following recurrence for spc :

$$\text{spc}(k, l) = \sum_{i=0}^l C_i \text{spc}(k-i-1, l-i) \quad \text{for } l < k, \quad (4.3)$$

$$\text{spc}(k, k) = C_k. \quad (4.4)$$

4.3 A better recurrence

We can now turn this recurrence into a simpler recurrence, which we will prove by recurrence.

Since spc has not been formally defined for negative arguments, we extend the definition with $\text{spc}(k, l) = 0$ for $l < 0$; this is consistent with the definition, since it yields a sum over subsets $S \subseteq [k]$ bigger than k , and it yields an empty sum in the recurrence we just derived. In addition to that, we also let $\text{spc}(k, l) = 0$ when $l > k$.

Note that for $k-1 = l \geq 0$, we have

$$\begin{aligned} \text{spc}(k, l) &= \sum_{i=0}^l C_i \text{spc}(k-i-1, l-i) = \sum_{i=0}^l C_i \text{spc}(l-i, l-i) = \sum_{i=0}^l C_i C_{l-i} \\ &= C_{l+1}, \end{aligned} \quad (4.5)$$

and thus

$$\text{spc}(k, l) = \text{spc}(k-1, l) + \text{spc}(k, l-1) = C_k \text{ for } k = l > 0. \quad (4.6)$$

Further, for $k-1 = l$, we get

$$\text{spc}(k, l) = \sum_{i=0}^l C_i C_{l-i} = C_l + \sum_{i=0}^{l-1} C_i C_{((l-1)-i)+1}$$

applying (4.4) on the left and (4.5) on the right,

$$\begin{aligned} &= \text{spc}(k-1, l) + \sum_{i=0}^{l-1} C_i \text{spc}(k-i-1, (l-1)-i) \\ &= \text{spc}(k-1, l) + \text{spc}(k, l-1). \end{aligned} \quad (4.7)$$

1					
1	1				
1	2	2			
1	3	5	5		
1	4	9	14	14	
1	5	14	28	42	42

Figure 8. The first few values of spc ; k vertically from 0 to 5, l horizontally from 0 to k . The $l+1$ summands in the recurrence (4.3) with $k = 6$ and $l = 3$ are highlighted.

Now let $k - 1 > l \geq 0$. Assume $\text{spc}(k', l') = \text{spc}(k' - 1, l') + \text{spc}(k', l' - 1)$ for $0 \leq l' < k' < k$. Then we can apply this assumption to the summands of $\text{spc}(k, l)$:

$$\begin{aligned}
 \text{spc}(k, l) &= \sum_{i=0}^l C_i \text{spc}(k - i - 1, l - i) \\
 &= \sum_{i=0}^l C_i (\text{spc}(k - i - 2, l - i) + \text{spc}(k - i - 1, l - i - 1)) \\
 &= \sum_{i=0}^l C_i \text{spc}(k - i - 2, l - i) + \sum_{i=0}^l C_i \text{spc}(k - i - 1, l - i - 1) \\
 &= \sum_{i=0}^l C_i \text{spc}((k - 1) - i - 1, l - i) + \sum_{i=0}^{l-1} C_i \text{spc}(k - i - 1, (l - 1) - i) \quad (4.8) \\
 &\quad + C_l \text{spc}(k - l - 1, -1),
 \end{aligned}$$

so, substituting the recurrence for spc ,

$$= \text{spc}(k - 1, l) + \text{spc}(k, l - 1). \quad (4.9)$$

We thus have the following recurrence for spc :

$$\text{spc}(k, l) = \text{spc}(k - 1, l) + \text{spc}(k, l - 1) \quad \text{for } k \geq 1, 0 \leq l \leq k, \quad (4.10)$$

$$\text{spc}(k, l) = 0 \quad \text{for } l < 0 \text{ or } l > k, \quad (4.11)$$

$$\text{spc}(0, 0) = 1. \quad (4.12)$$

This is the recurrence defining the Catalan triangle⁵, and its solution is known, namely

$$\text{spc}(k, l) = \frac{(k + l)! (k - l + 1)}{l! (k + 1)!} = \frac{k - l + 1}{k + 1} \binom{k + l}{k}, \quad (4.13)$$

as shown by L. F. A. Arbogast in 1800 [2, p. 214–217].

4.3.1 A combinatorial interpretation of spc

The Catalan triangle counts well-formed prefix bracket expressions [CITATION]. This can in fact be seen from all the definitions of spc above, thus providing a more combinatorial solution.

In the original definition (4.1) of spc , the sum is over the choices for the positions of the unmatched $k - l$ opening brackets; between those, well-formed bracket expressions (counted by Catalan numbers) are inserted.

The first recurrence (4.3) corresponds to the following grammatical definition⁶ of well-formed prefixes, where again wfbe denotes a well-formed bracket expression:

$$\begin{aligned}
 \text{wfprefix} ::= & \text{wfbe} \langle \text{wfprefix} \\
 & | \text{wfbe},
 \end{aligned}$$

where the sum ranges over the length of the well-formed bracket expression preceding the first unmatched bracket.

Finally, (4.10) reflects the fact that a well-formed prefix with k opening brackets and l closing brackets either ends with an opening bracket (preceded by a well-formed prefix with $k - 1$ opening and l closing brackets), or with a closing bracket (preceded by a well-formed prefix with k opening and $l - 1$ closing brackets).

⁵[TODO cite something that actually uses that name here]

⁶This grammar requires infinite look-ahead to parse, but doing a first pass to mark the unmatched opening brackets resolves that.

1					
1	1				
1	2	2			
1	3	5	5		
1	4	9	14	14	
1	5	14	28	42	42

Figure 9. The recurrence (4.10).

4.4 Asymptotics

We can now use the expression (4.13) for spc to study the asymptotics of the bound (4.2). We are interested in the asymptotics as a function of the length of the bracket expression $n = k + l + p + q$. Since the bound is a product of two identical two-parameter factors, we study one of them,

$$C_k \text{spc}(k, l),$$

as a function of $n_1 := k + l$. Let $l = \alpha k$, thus $0 \leq \alpha \leq 1$, the above expression becomes

$$C_{\frac{n_1}{\alpha+1}} \left(\frac{n_1}{\frac{n_1}{\alpha+1}} \right) \frac{(1-\alpha)n + \alpha + 1}{n + \alpha + 1}.$$

Asymptotically as $n_1 \rightarrow \infty$, this yields

$$\frac{1}{n_1 \sqrt{\pi n_1}} 4^{\frac{n_1}{\alpha+1}} \frac{\sqrt{1+\alpha}}{\sqrt{2\pi n_1 \alpha}} \left(\frac{1+\alpha}{\alpha^{\frac{\alpha}{\alpha+1}}} \right)^{n_1} (1-\alpha),$$

or, up to a polynomially-bounded factor,

$$4^{\frac{n_1}{\alpha+1}} \left(\frac{1+\alpha}{\alpha^{\frac{\alpha}{\alpha+1}}} \right)^{n_1} = \left(4^{\alpha+1} \frac{1+\alpha}{\alpha^{\frac{\alpha}{\alpha+1}}} \right)^{n_1}.$$

As α ranges from 0 to 1, the base of that exponential reaches a maximum of 5 at $\alpha = \frac{1}{4}$. We thus have the following asymptotic bound.

$$v_P(\langle^k \rangle^l \langle^q \rangle^p) \leq C_k \text{spc}(k, l) \text{spc}(p, q) C_p \leq 5^{k+l} 5^{p+q} = 5^n. \quad (4.14)$$

[TODO: would it be interesting or feasible to average over all $\langle^k \rangle^l \langle^q \rangle^p$, for fixed n ?

[TODO: here a section about spc generalizing to arbitrary WFBs, and the applicability of overcounting analyses]

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