

Introduction to High Performance Scientific Computing

Autumn, 2016

Lecture 14

Vectorizing code

In general (Python, Fortran, Matlab,...), avoid for loops and *vectorize* calculations involving arrays.

Example:

```
In [27]: x=np.linspace(0,1,101)
```

```
In [28]: f = np.empty_like(x)
```

```
In [29]: for i in range(size(x)):
.....:         f[i] = cos(x[i])**2
.....:
```

```
In [30]: f = cos(x)**2      Vectorized version of loop
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- **Vectorized code will usually be faster, sometimes *much faster***
- **Exception: parallelizing Fortran code with OpenMp:**
vectorized code → loops → parallel loops

Today

Programming example: from PDE → algorithm → serial code → parallel code

Programming example

Task: Compute temperature distribution in a room

Programming example

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Governing equation: Heat equation (diffusion equation):

$$\frac{\partial T}{\partial t} = \alpha \nabla^2 T + S(\mathbf{x}, t)$$

$$T(\mathbf{x}, t = 0) = f(\mathbf{x}) \quad \text{Initial condition}$$

Here, S is a *heat source*. Boundary conditions should also be specified as appropriate.

Problem: given the source, initial condition, and boundary conditions, solve for the temperature distribution, $T(\mathbf{x}, t)$

Programming example

Today: 1-D problem

$$\frac{\partial T}{\partial t} = \frac{\partial^2 T}{\partial x^2} + S(x, t)$$

$$T(x, t = 0) = f(x)$$

Initial condition

$$T(x = 0, t) = a(t), \quad T(x = 1, t) = b(t)$$

Boundary conditions

$$0 \leq x \leq 1$$

Programming example

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First consider steady problem, e.g., $S = S(x)$, a and b are constants:

$$\frac{\partial^2 T}{\partial x^2} + S(x, t) = 0$$

Poisson equation

Programming example

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$$\frac{\partial^2 T}{\partial x^2} + S(x) = 0$$

Notes:

- 1. This is an extremely simple problem, easy to write down the analytical solution**
- 2. No need to use compiled language**
- 3. Certainly no need to parallelize**
- 4. But what about two-dimensional or three-dimensional problems?**
 - Then, the picture changes considerably!**
- 5. We are just considering the 1-D problem for illustrative purposes**

Programming example

First consider steady problem, e.g., $S = S(x)$, a and b are constants:

$$\frac{\partial^2 T}{\partial x^2} + S(x) = 0 \quad \text{Poisson equation}$$

Numerical method:

1. Discretize the derivative:

$$\frac{\partial^2 T}{\partial x^2} \approx \frac{T_{i+1} - 2T_i + T_{i-1}}{\Delta x^2} \quad \text{2nd-order, centered scheme}$$

$$x_i = i * \Delta x, \quad i = 1, 2, \dots, N$$

$$(N + 1) * \Delta x = 1$$

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With boundary conditions: $T_0 = T_a, \quad T_N = T_b$

Programming example

Equation for T_i :
$$\frac{T_{i+1} - 2T_i + T_{i-1}}{\Delta x^2} = -S_i$$

In matrix form: $AT = b$

$$A = \begin{bmatrix} -2 & 1 & 0 & 0 & 0 & \dots & 0 \\ 1 & -2 & 1 & 0 & 0 & \dots & 0 \\ 0 & 1 & -2 & 1 & 0 & \dots & 0 \\ \vdots & \ddots & \ddots & \ddots & \ddots & \vdots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 1 & -2 & 1 \\ 0 & \dots & 0 & 0 & 0 & 1 & -2 \end{bmatrix}, \quad b = \frac{1}{\Delta x^2} \begin{bmatrix} -\Delta x^2 T_a - S_1 \\ -S_2 \\ \vdots \\ -S_i \\ \vdots \\ -S_{N-1} \\ -\Delta x^2 T_b - S_N \end{bmatrix}$$

- In 1-D, this is just a tridiagonal system of equations
- Easy to solve directly (with, say, *DGTSV*)

Programming example

- In two or three dimensions, A loses its simple banded structure
- Then, direct solution becomes very expensive for large N
- *Iterative* methods are a popular alternative

Programming example

- In two or three dimensions, A loses its simple banded structure
- Then, direct solution becomes very expensive for large N
- *Iterative* methods are a popular alternative
- Basic idea: rewrite $Ax=b$ as $A_1x = A_2x + b$
- Choose A_1 so that it is easy to invert, then solve iterative system:
- $A_1x^{k+1} = A_2x^k + b$
 - Requires guess, x^0

Jacobi iteration

- **Basic idea:** rewrite $Ax=b$ as $A_1x = A_2x + b$
- **Choose A_1** so that it is easy to invert, then solve iterative system:
- $A_1x^{k+1} = A_2x^k + b$
 - Requires guess, x^0
- **Jacobi iteration:** Choose A_1 to be diagonal matrix (main diagonal of A):

$$\frac{T_{i+1}^{k-1} - 2T_i^k + T_{i-1}^{k-1}}{\Delta x^2} = -S_i$$

$$T_i^k = \frac{\Delta x^2}{2} S_i + \frac{1}{2} (T_{i+1}^{k-1} + T_{i-1}^{k-1})$$

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Main algorithm, easy to code!

Jacobi iteration in Fortran

- **Plan:**
 - **Set parameters:** a, b, n, tol
 - **Construct grid** x_i
 - **Construct source function, $S(x)$, initialize $T=T(x,t=0)$**
 - **Iterate using formula below**
 - **Each iteration check if $|T_k - T_{k-1}| < \text{tol}$**

$$T_i^k = \frac{\Delta x^2}{2} S_i + \frac{1}{2} (T_{i+1}^{k-1} + T_{i-1}^{k-1})$$

Jacobi iteration in Fortran

- **One Fortran trick: set variables to be dimension(0:N+1)**
 - $x(0)=0$, $x(N+1)=1$, $T(0)=a$, $T(N+1)=b$
 - **Then, easy to compute T_1 using:**

$$T_i^k = \frac{\Delta x^2}{2} S_i + \frac{1}{2} (T_{i+1}^{k-1} + T_{i-1}^{k-1})$$

Jacobi iteration in Fortran

- **One Fortran trick: set variables to be dimension(0:N+1)**
 - $x(0)=0$, $x(N+1)=1$, $T(0)=a$, $T(N+1)=b$
 - **Then, easy to compute T_1 using:**

$$T_i^k = \frac{\Delta x^2}{2} S_i + \frac{1}{2} (T_{i+1}^{k-1} + T_{i-1}^{k-1})$$

Core part of code (see *jacobi1s.f90*):

```
do k1=1, kmax
```

```
  Tnew(1:n) = S(1:n)*dx2f + 0.5d0*(T(0:n-1) + T(2:n+1)) !Jacobi
```

```
  deltaT(k1) = maxval(abs(Tnew(1:n)-T(1:n))) !compute relative error
```

```
  T(1:n)=Tnew(1:n) !update variable
```

```
  if (deltaT(k1)<tol) exit !check convergence criterion
```

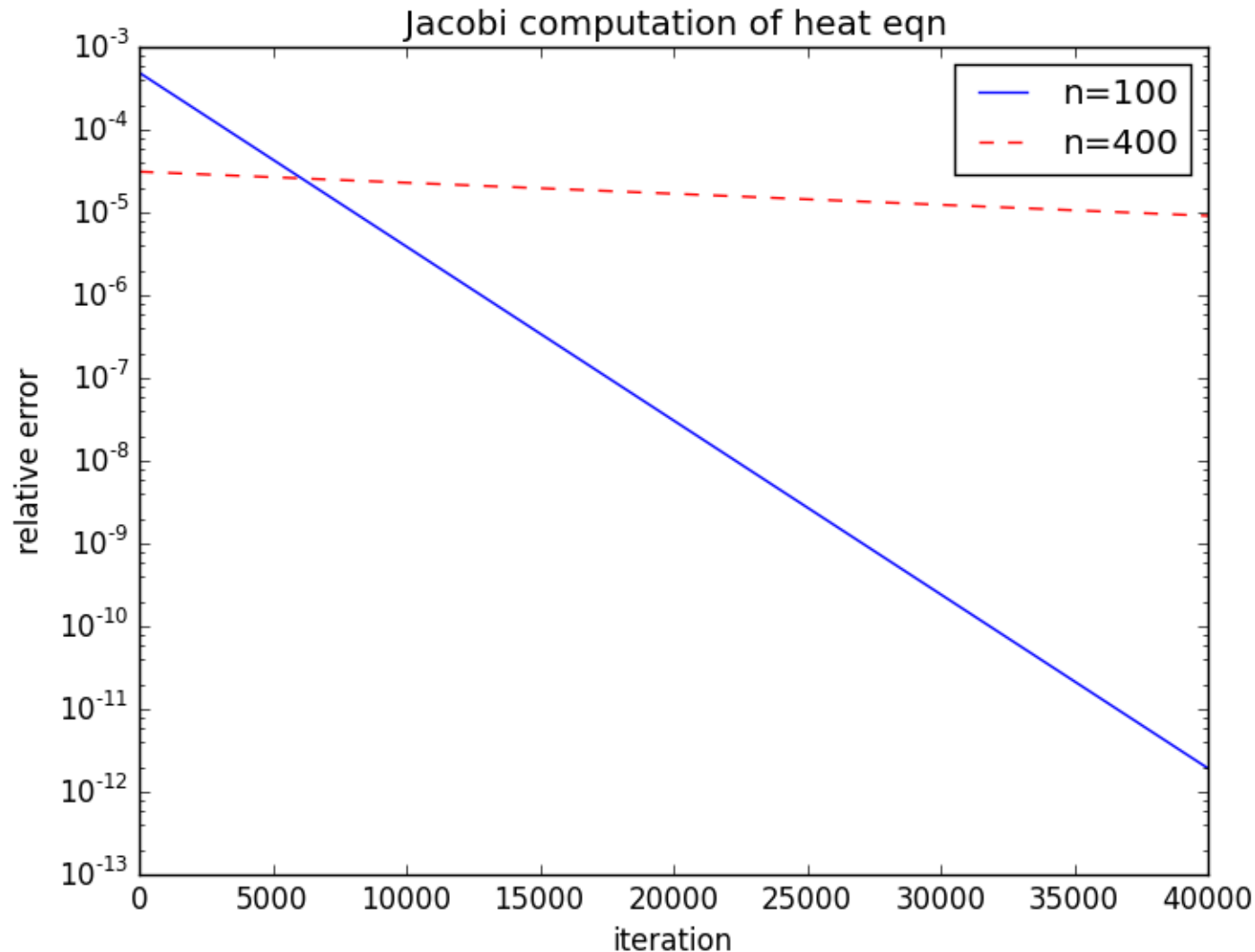
```
end do
```

Jacobi results

1. Does the solution converge?

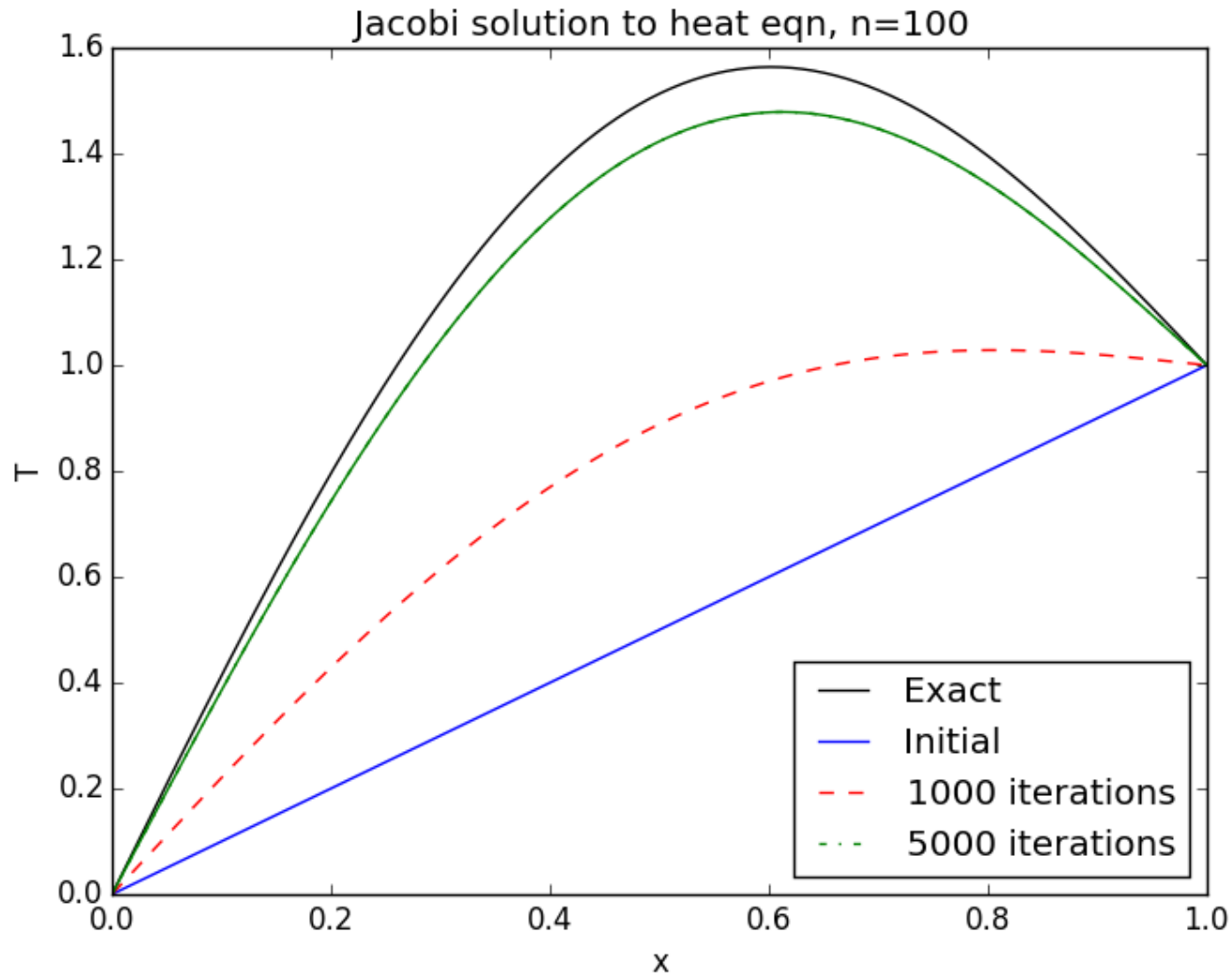
Jacobi results

1. Does the solution converge? Yes, but very slowly for large n



Jacobi results

1. Does it converge to the correct solution? Yes, can check that error $\sim \Delta x^2$



Jacobi results

- Jacobi is simplest, but *most inefficient* iterative solver
- Good illustration of basic ideas
- Better methods: Gauss-Seidel, SOR, conjugate gradient, multigrid

Parallel Jacobi

Let's now parallelize the solver with OpenMP

- Look for loops that can be parallelized
- Look for vectorized operations that can be converted to loops that can be parallelized

Serial:

```
Tnew(1:n) = S(1:n)*dx2f + 0.5d0*(T(0:n-1) + T(2:n+1)) !Jacobi
```

```
deltaT(k1) = maxval(abs(Tnew(1:n)-T(1:n))) !compute relative error
```


Parallel Jacobi

Let's now parallelize the solver with OpenMP

- Look for loops that can be parallelized
- Look for vectorized operations that can be converted to loops that can be parallelized

Parallel:

```
    dmax=0.d0
    !$omp parallel do reduction(max:dmax)
    do i1=1,n
        Tnew(i1) = S(i1)*dx2f + 0.5d0*(T(i1-1) + T(i1+1))
        dmax = max(dmax,abs(Tnew(i1)-T(i1)))
    end do
    !$omp end parallel do
    deltaT(k1) = dmax
```

Parallel Jacobi

Let's now parallelize the solver with OpenMP

- Look for loops that can be parallelized
- Look for vectorized operations that can be converted to loops that can be parallelized

Serial:

```
do i1=0,n+1
    x(i1) = i1*dx
end do
!-----

!set initial condition
T = (b-a)*x + a

!set source function
S = S0*sin(pi*x)
```

Parallel Jacobi

Let's now parallelize the solver with OpenMP

- Look for loops that can be parallelized
- Look for vectorized operations that can be converted to loops that can be parallelized

Parallel:

```
!$omp parallel do
do i1=0,n+1
    x(i1) = i1*dx
    T(i1) = (b-a)*x(i1) + a !set initial condition
    S(i1) = S0*sin(pi*x(i1)) !set source function
end do
!$omp end parallel do
```

Parallel Jacobi notes

- Will only see speedup with $n > \sim 20000$ (commonly seen in 2D problems)
- See *`jacobi1s_omp.f90`*, *`jacobi1_omp.py`*
- **f2py and OpenMP:** `f2py --f90flags='-fopenmp' -lgomp -c jacobi1s_omp.f90 -m j1`
-

Time-dependent problem

Today: 1-D problem

$$\frac{\partial T}{\partial t} = \frac{\partial^2 T}{\partial x^2} + S(x, t)$$

$$T(x, t = 0) = f(x)$$

$$T(x = 0, t) = a(t), \quad T(x = 1, t) = b(t)$$

$$0 \leq x \leq 1$$

Simple (inefficient) approach: *method of lines*

- 1. Discretize spatial variable \rightarrow $N+2$ points between 0 and 1**
- 2. Solve resulting N ODEs with solver of choice (*odeint, ode15s, ...*)**

Time-dependent problem

Again, we discretize the derivative:

$$\frac{\partial^2 T}{\partial x^2} \approx \frac{T_{i+1} - 2T_i + T_{i-1}}{\Delta x^2}$$

$$x_i = i * \Delta x, \quad i = 0, 1, 2, \dots, N + 1$$

$$(N + 1) * \Delta x = 1$$

Time-dependent problem

Again, we discretize the derivative as:

$$\frac{\partial^2 T}{\partial x^2} \approx \frac{T_{i+1} - 2T_i + T_{i-1}}{\Delta x^2}$$

$$x_i = i * \Delta x, \quad i = 0, 1, 2, \dots, N + 1$$

$$(N + 1) * \Delta x = 1$$

So, we have **N ODEs**:

$$\frac{dT_i}{dt} = S_i(t) + \frac{T_{i+1} - 2T_i + T_{i-1}}{\Delta x^2}, \quad i = 1, 2, \dots, N$$

with the boundary conditions substituted in the RHS when needed.

Solving single ODE in python (lecture 6)

- Use *odeint* from *scipy.integrate* module to solve:

$$\frac{dy}{dt} = -ay$$

- **Basic idea:** discretize time, $t = 0, dt, \dots, N*dt$, and starting from $y(0)$ march forward in time and compute $y(dt), \dots y(N*dt)$
- *odeint* chooses the stepsize, dt , so that error tolerances are satisfied
- Need to specify:
 - Initial condition
 - Timespan for integration
 - A Python function which provides RHS of the ODE to *odeint*
- Look at *ode_example.py* and lab 4

Time-dependent problem

Solving N ODEs:

$$\frac{dT_i}{dt} = S_i(t) + \frac{T_{i+1} - 2T_i + T_{i-1}}{\Delta x^2}, \quad i = 1, 2, \dots, N$$

- Will need to provide N initial conditions when calling *odeint*.
- The python function which provides RHS to *odeint* will:
 - Take t and T_1, \dots, T_N and any other needed parameters as input
 - Return N values for dT/dt as output
- No need for Fortran for 1D problems, but may be faster for two and three dimensions.