

Mastering Linked Lists – From Basics to Advanced

What is a Linked List?

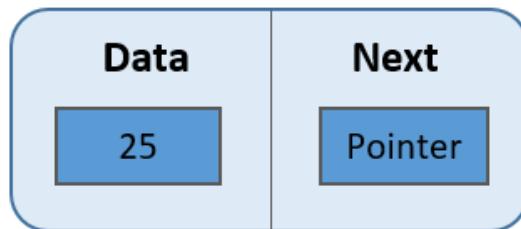
A Linked List is a way of organizing data where each item, called a node, holds two things — the actual data and a reference to the next item in the list.

Think of it like a chain of people, where each person is holding something (the data) and also pointing to the next person in line. The important part is: unlike arrays, these people (or nodes) aren't standing next to each other — they could be scattered all around, but as long as each one knows who's next, the chain stays connected.

So, each node has:

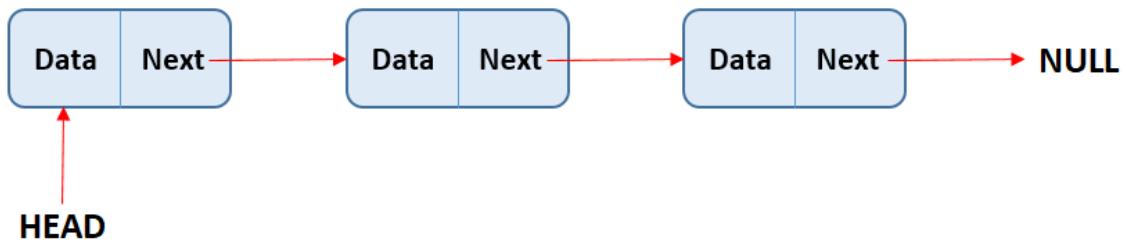
- A part that stores the data.
- A part that stores a pointer (or reference) to the next node.

This setup allows the linked list to grow or shrink easily, without needing a large block of continuous memory.



This image shows how a Linked List works in a very simple way.

- It all starts with the HEAD, which is just a reference to the first node in the list. Think of it like a starting point — without the HEAD, you can't access the list.
- Each node in the list has two parts:
 1. Data – which stores the actual information (like a number, name, etc.).
 2. Next – which holds the address (or pointer) of the next node in the chain.
- The red arrows in the image represent the connections between the nodes — each node points to the next one.
- This continues until the last node, which doesn't point to anything further. Instead, its "Next" part contains NULL, meaning “end of the list.”



AB CD Node Structure

Each node in a singly linked list contains:

- **Data** – The value stored.
- **Next** – A pointer/reference to the next node.

```
class Node {
    int data;
    Node next;

    Node(int data) {
        this.data = data;
        this.next = null;
    }
}
```

🔗 Types of Linked Lists

1. Singly Linked List – Each node points to the next. (\rightarrow)



- ◆ Insert at Beginning - Time Complexity: O(1)

```
newNode.next = head;
head = newNode;
```

- ◆ Insert at End - Time Complexity: O(n)*

```
tail.next = newNode;
tail = newNode;
```

- ◆ Delete from Beginning - Time Complexity: O(1)

```
head = head.next;
```

- ◆ Delete from End - Time Complexity: O(n)

```
while (temp.next.next != null) temp = temp.next;
temp.next = null;
```

- ◆ Search - Time Complexity: O(n)

```
while (temp != null) {
    if (temp.data == key) return true;
    temp = temp.next;
}
```

- ◆ Traverse - Time Complexity: O(n)

```
Node temp = head;
while (temp != null) {
    System.out.print(temp.data + " ");
    temp = temp.next;
}
```

* O(1) if using tail pointer

2. Doubly Linked List – Each node has both next and prev pointers. (\leftrightarrow)



- ◆ Insert at Beginning - Time Complexity: O(1)

```
newNode.next = head;
head.prev = newNode;
head = newNode;
```

- ◆ Insert at End - Time Complexity: O(1)*

```
tail.next = newNode;
newNode.prev = tail;
tail = newNode;
```

- ◆ Delete from Beginning - Time Complexity: O(1)

```
head = head.next;
head.prev = null;
```

- ◆ Delete from End - Time Complexity: O(1)

```
tail = tail.prev;
tail.next = null;
```

- ◆ Search - Time Complexity: O(n)

```
while (temp != null) {
    if (temp.data == key) return true;
    temp = temp.next;
}
```

- ◆ Traverse Forward - Time Complexity: O(n)

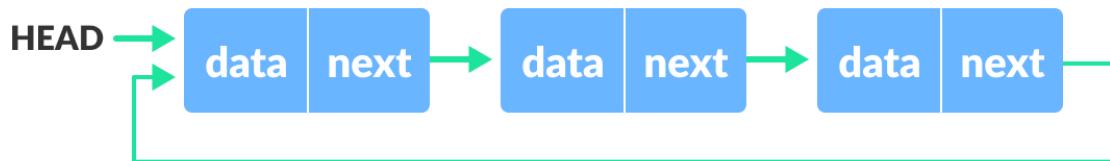
```
for (DNode temp = head; temp != null; temp = temp.next)
    System.out.print(temp.data + " ");
```

- ◆ Traverse Backward - Time Complexity: O(n)

```
for (DNode temp = tail; temp != null; temp = temp.prev)
    System.out.print(temp.data + " ");
```

* O(1) if tail maintained

3. Circular Linked List



- Singly Circular: Last node points to the head. ()

- ◆ Insert at End - Time Complexity: O(1)

```
tail.next = newNode;
newNode.next = head;
tail = newNode;
```

- ◆ Delete from Beginning - Time Complexity: O(1)

```
head = head.next;
tail.next = head;
```

- ◆ Traverse - Time Complexity: O(n)

```
Node temp = head;
do {
    System.out.print(temp.data + " ");
    temp = temp.next;
} while (temp != head);
```

- Doubly Circular: Both ends connected forming a circle.

- ◆ Insert at End - Time Complexity: O(1)

```
tail.next = newNode;
newNode.prev = tail;
newNode.next = head;
head.prev = newNode;
tail = newNode;
```

- ◆ Delete from Beginning - Time Complexity: O(1)

```
head = head.next;
head.prev = tail;
tail.next = head;
```

- ◆ Delete from End - Time Complexity: O(1)

```
tail = tail.prev;
tail.next = head;
head.prev = tail;
```

- ◆ Traverse Forward - Time Complexity: O(n)

```
DCNode temp = head;
do {
    System.out.print(temp.data + " ");
    temp = temp.next;
} while (temp != head);
```

Terminology

- **Head** – First node in the list.
- **Tail** – Last node, typically points to null (or head in circular list).
- **Traversal** – Visiting each node.
- **Insertion** – Adding nodes.
- **Deletion** – Removing nodes.
- **Search** – Finding nodes with a specific value.

Advanced Understanding of Linked Lists

This section helps you understand how linked lists work in real life — when to use them, where they perform better than arrays, and what common mistakes to avoid. It's about seeing the why behind the code.

1 Linked List vs Array – A Deep Dive

Arrays and linked lists are both linear data structures, but they function very differently under the hood.

Feature	Array	Linked List
Memory Allocation	Contiguous	Non-contiguous
Access Time	$O(1)$ (Random)	$O(n)$ (Sequential)
Insertion/Deletion	Costly ($O(n)$)	Efficient ($O(1)$ at head)
Cache Performance	High	Low
Resizing	Fixed/Resize cost	Dynamic (flexible)

 **Use Case Tip:** Use arrays when you need frequent indexed access; use linked lists when insert/delete operations dominate.

2 Real-World Applications of Linked Lists

Linked lists are everywhere — here's how they're used:

Use Case	Type of List	Description
Music/Video Playlists	Doubly Linked List	Navigate forward/backward
Undo/Redo in Editors	Doubly Linked List	Reversible actions
Operating Systems (Memory)	Singly Linked List	Manage free blocks
LRU Cache Implementation	Doubly Linked List + HashMap	Fast eviction
Round-Robin Task Scheduling	Circular Linked List	Tasks run in cycles

3 Why Linked Lists Are a Favorite in Interviews

- Pointer management:** Interviewers test how well you handle references and memory.
- Edge cases:** Think nulls, empty lists, one-node lists.
- Recursion vs iteration:** Tests logical flow handling.
- Follow-ups:** Easy to build on—merge, reverse, detect cycles.

 *Pro Tip:* Always clarify whether the list is singly or doubly linked, and whether it's circular, before solving.

4 Types of Linked Lists: Comparison and Trade-offs

Type	Forward	Backward	Circular	Memory Cost	Use Case Example
Singly Linked List	✓	✗	✗	Low	Stack
Doubly Linked List	✓	✓	✗	Medium	Undo/Redo, LRU
Circular SLL	✓	✗	✓	Low	Round Robin
Circular DLL	✓	✓	✓	High	Cyclic Navigation

5 Pointer Manipulation Concepts

Pointers (or references) are the heart of a linked list.

- next → points to the next node
- prev → in DLL, points to previous node
- null → denotes end (or start in some circular types)

⚠ Common Mistakes:

- Skipping a node by misplacing `next = next.next`
- Losing head reference → causes memory leaks
- Infinite loops from improper circular handling

6 Linked Lists in System Design

Linked lists are often embedded inside more complex data structures:

Data Structure	Role of Linked List
Stack (LL-based)	Push/Pop at head ($O(1)$)
Queue (LL-based)	Enqueue at tail, Dequeue at head
HashMap (Chaining)	Handle collisions via separate chaining
Graph (Adjacency List)	Maintain neighbor nodes

📌 Bonus: LRU Cache uses DLL + HashMap for $O(1)$ access & eviction.

7 Time and Space Complexity Table

Operation	Singly LL	Doubly LL	Array
Insert at Head	$O(1)$	$O(1)$	$O(n)$
Insert at Tail	$O(n)*$	$O(1)*$	$O(n)$
Delete at Head	$O(1)$	$O(1)$	$O(n)$
Delete at Tail	$O(n)$	$O(1)$	$O(n)$
Search	$O(n)$	$O(n)$	$O(1)$
Access by Index	$O(n)$	$O(n)$	$O(1)$

*With maintained tail pointer

8 Static vs Dynamic Memory

- **Arrays:** Stored in stack (static memory)
- **Linked Lists:** Stored in heap (dynamic memory)
- In C/C++, manual memory management with malloc()/free()
- In Java/Python, handled by Garbage Collector

 **Tip:** Linked lists give more control but can lead to memory leaks if not handled carefully (especially in C/C++).

9 When Not to Use a Linked List

Avoid linked lists if:

- You need random indexed access (arrays are O(1))
- Memory is tight (linked lists use extra space for pointers)
- Cache performance is critical (linked lists are scattered in memory)

10 Common Pitfalls and How to Avoid Them

Mistake	Solution
Infinite loop in circular list	Always set termination condition
Losing head reference	Use a dummy node or temp variable
Skipping node on delete	Ensure correct next linkage
Modifying while traversing	Use two-pointer or dummy technique
Forgetting to handle null pointers	Always null-check before access

Advanced Linked List Operations (with Code Snippets)

These operations are frequently asked in interviews and competitive programming. Below are concise Java code snippets with time and space complexities.

1 Reverse a Linked List

```
Node reverse(Node head) {  
    Node prev = null, curr = head;  
    while (curr != null) {  
        Node next = curr.next;  
        curr.next = prev;  
        prev = curr;  
        curr = next;  
    }  
    return prev;  
}  
// Time: O(n), Space: O(1)
```

2 Detect Cycle in Linked List (Floyd's Algorithm)

```
boolean hasCycle(Node head) {  
    Node slow = head, fast = head;  
    while (fast != null && fast.next != null) {  
        slow = slow.next;  
        fast = fast.next.next;  
        if (slow == fast) return true;  
    }  
    return false;  
}  
// Time: O(n), Space: O(1)
```

3 Find the Middle Node

```
Node findMiddle(Node head) {  
    Node slow = head, fast = head;  
    while (fast != null && fast.next != null) {  
        slow = slow.next;  
        fast = fast.next.next;  
    }  
    return slow;  
}  
// Time: O(n), Space: O(1)
```

4 Check if a Linked List is a Palindrome

```
boolean isPalindrome(Node head) {  
    Node slow = head, fast = head;  
    while (fast != null && fast.next != null) {  
        slow = slow.next;  
        fast = fast.next.next;  
    }  
    Node secondHalf = reverse(slow);  
    while (secondHalf != null) {  
        if (head.data != secondHalf.data) return false;  
        head = head.next;  
        secondHalf = secondHalf.next;  
    }  
    return true;  
}  
// Time: O(n), Space: O(1)
```

5 Merge Two Sorted Linked Lists

```
Node merge(Node l1, Node l2) {
    if (l1 == null) return l2;
    if (l2 == null) return l1;
    if (l1.data < l2.data) {
        l1.next = merge(l1.next, l2);
        return l1;
    } else {
        l2.next = merge(l1, l2.next);
        return l2;
    }
}
// Time: O(n + m), Space: O(n + m) - recursive stack
```

6 Remove N-th Node from End

```
Node removeNthFromEnd(Node head, int n) {
    Node dummy = new Node(0);
    dummy.next = head;
    Node fast = dummy, slow = dummy;
    for (int i = 0; i <= n; i++) fast = fast.next;
    while (fast != null) {
        fast = fast.next;
        slow = slow.next;
    }
    slow.next = slow.next.next;
    return dummy.next;
}
// Time: O(n), Space: O(1)
```

7 Find Intersection Point of Two Linked Lists

```
Node getIntersection(Node a, Node b) {  
    Node p1 = a, p2 = b;  
    while (p1 != p2) {  
        p1 = (p1 == null) ? b : p1.next;  
        p2 = (p2 == null) ? a : p2.next;  
    }  
    return p1;  
}  
// Time: O(n + m), Space: O(1)
```

8 Clone Linked List with Random Pointer

```
Node copyRandomList(Node head) {  
    if (head == null) return null;  
    Node curr = head;  
    while (curr != null) {  
        Node copy = new Node(curr.data);  
        copy.next = curr.next;  
        curr.next = copy;  
        curr = copy.next;  
    }  
    curr = head;  
    while (curr != null) {  
        if (curr.random != null)  
            curr.next.random = curr.random.next;  
        curr = curr.next.next;  
    }  
    Node dummy = new Node(0), copyCurr = dummy;  
    curr = head;  
    while (curr != null) {  
        copyCurr.next = curr.next;  
        curr.next = curr.next.next;  
        curr = curr.next;  
        copyCurr = copyCurr.next;  
    }  
    return dummy.next;  
}  
// Time: O(n), Space: O(1)
```

9 Rotate Linked List by K

```
Node rotateRight(Node head, int k) {  
    if (head == null) return null;  
    int len = 1;  
    Node tail = head;  
    while (tail.next != null) {  
        tail = tail.next;  
        len++;  
    }  
    k %= len;  
    if (k == 0) return head;  
    tail.next = head; // make circular  
    for (int i = 0; i < len - k; i++) tail = tail.next;  
    Node newHead = tail.next;  
    tail.next = null;  
    return newHead;  
}  
// Time: O(n), Space: O(1)
```

10 Flatten Multilevel Linked List

```
Node flatten(Node head) {  
    Node curr = head;  
    while (curr != null) {  
        if (curr.child != null) {  
            Node temp = curr.next;  
            curr.next = curr.child;  
            Node tail = curr.child;  
            while (tail.next != null) tail = tail.next;  
            tail.next = temp;  
            curr.child = null;  
        }  
        curr = curr.next;  
    }  
    return head;  
}  
// Time: O(n), Space: O(1)
```

Basic Linked List Problems

#	Problem Title	Platform	Link	Key Concepts
1	Reverse Linked List	LeetCode – Easy	 Link	Iteration, Recursion
2	Delete Node in a Linked List	LeetCode – Medium	 Link	Deletion
3	Middle of the Linked List	LeetCode – Easy	 Link	Fast/Slow Pointer
4	Merge Two Sorted Lists	LeetCode – Easy	 Link	Recursion
5	Remove Linked List Elements	LeetCode – Easy	 Link	Deletion
6	Palindrome Linked List	LeetCode – Easy	 Link	Reverse, Compare
7	Convert Binary Number in LL to Integer	LeetCode – Easy	 Link	Traversal
8	Linked List Cycle	LeetCode – Easy	 Link	Floyd's Algorithm
9	Remove Duplicates from Sorted List	LeetCode – Easy	 Link	Basic Traversal
10	Search in Linked List	GFG – Easy	 Link	Traversal

Advanced Linked List Problems

#	Problem Title	Platform	Link	Key Concepts
1	Linked List Cycle II	LeetCode – Medium	 Link	Detect Cycle Start
2	Intersection of Two Linked Lists	LeetCode – Easy	 Link	Length Diff, Pointers
3	Rotate List	LeetCode – Medium	 Link	Modulo, Tail Handling
4	Remove N-th Node from End	LeetCode – Medium	 Link	Two Pointers
5	Add Two Numbers	LeetCode – Medium	 Link	Math in Linked List
6	Copy List with Random Pointer	LeetCode – Medium	 Link	Cloning, Hashing
7	Flatten a Multilevel Doubly Linked List	LeetCode – Medium	 Link	Flattening, Recursion
8	Reverse Nodes in k-Group	LeetCode – Hard	 Link	Recursion, k-Group
9	Merge k Sorted Lists	LeetCode – Hard	 Link	Min Heap
10	LRU Cache	LeetCode – Medium	 Link	DLL + HashMap

GitHub DSA Repository (by Harsh Raj)

-  [BasicLinkedListOperations.java](#)
-  [LinkedListProblems.java](#)
-  <https://github.com/eh-harsh02/dsa>