

Module

11

# Session Hijacking

# Learning Objectives

01

Summarize Session Hijacking Concepts

03

Explain Network-Level Session Hijacking

02

Explain Application-Level Session Hijacking

04

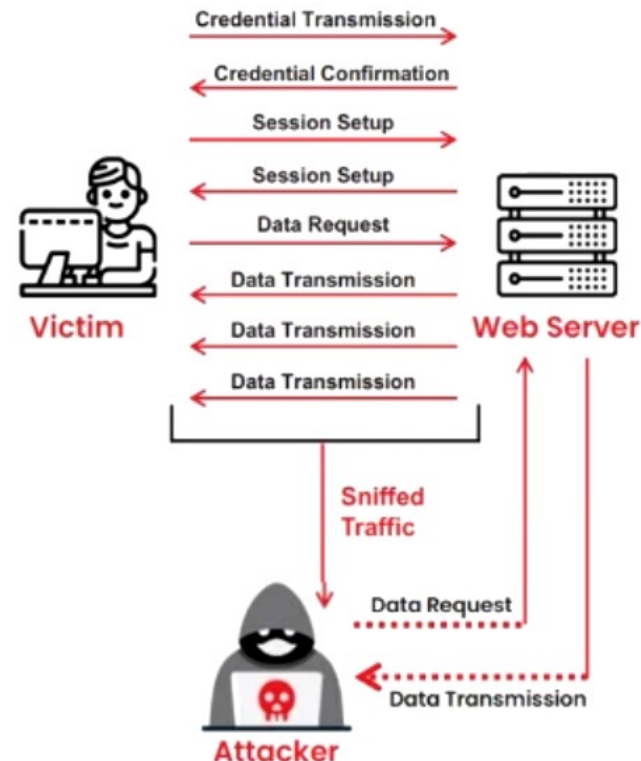
Explain Session Hijacking Countermeasures

## Objective 01

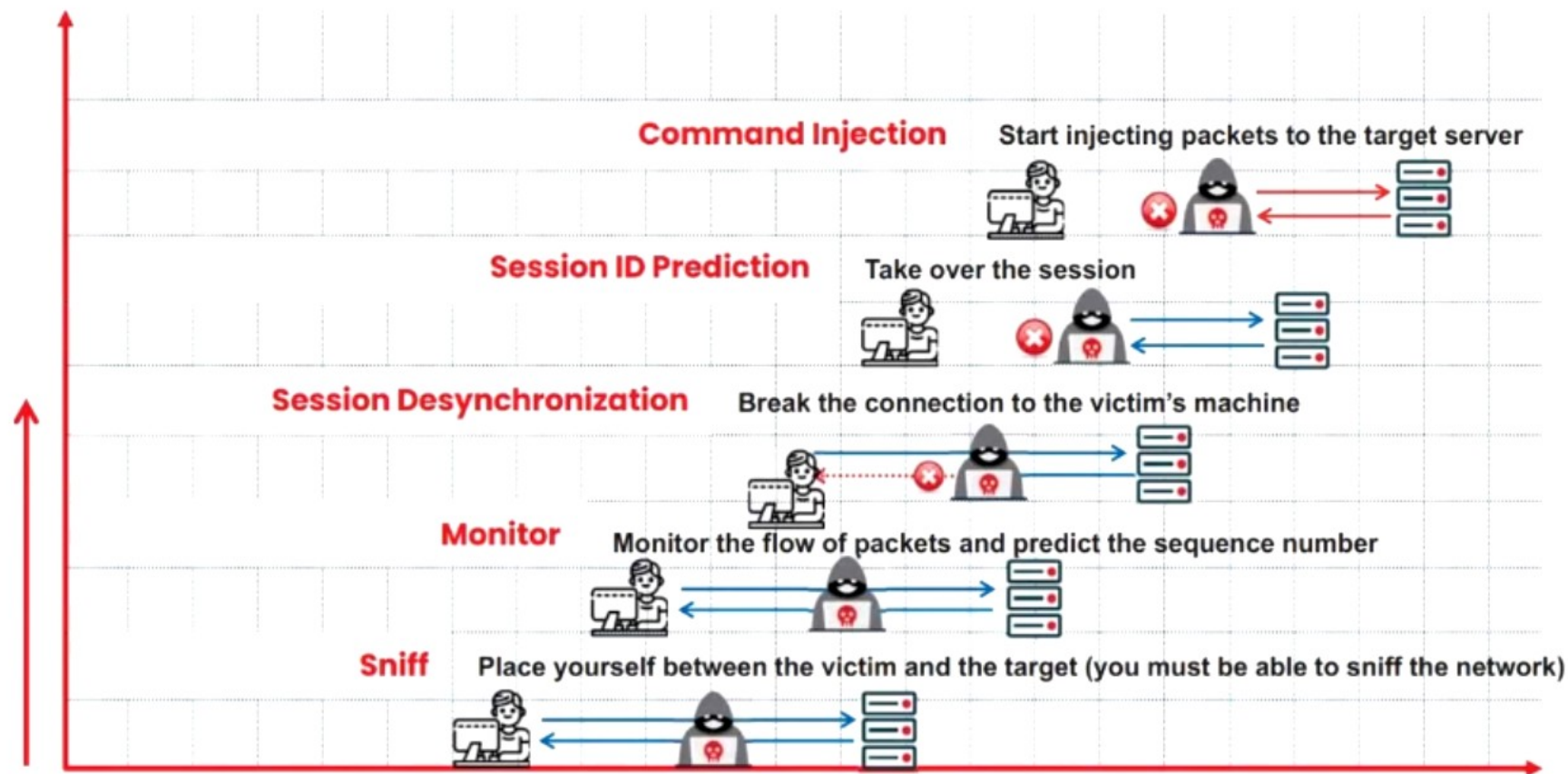
# Summarize Session Hijacking Concepts

# What is Session Hijacking?

- Session hijacking refers to an attack in which an attacker seizes control of a **valid TCP communication session** between two computers
- As most **authentications only occur at the start of a TCP session**, this allows the attacker to gain access to a machine
- Attackers can sniff all the traffic from the established TCP sessions and perform **identity theft, information theft, fraud**, etc.
- The attacker steals a valid session ID and uses it to **authenticate himself with the server**



# Session Hijacking Process



# Session Hijacking in OSI Model

## Network-Level Hijacking

Network-level hijacking can be defined as the **interception of packets** during the transmission between a client and the server in a TCP or UDP session

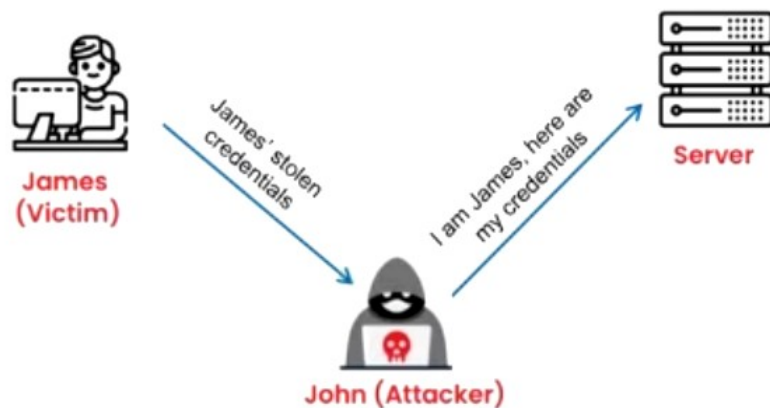
## Application-Level Hijacking

Application-level hijacking refers to **gaining control over the HTTP's user session** by obtaining the session IDs

# Spoofing vs. Hijacking

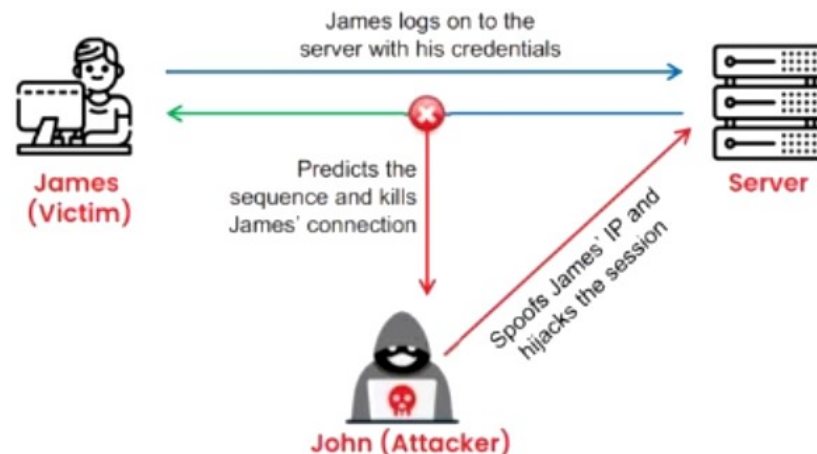
## Spoofing Attack

- An attacker **pretends to be another user** or machine (victim) to gain access
- The attacker does not seize control of an existing active session; instead, he or she initiates a new session using the victim's **stolen credentials**



## Hijacking

- Session hijacking is the process of seizing control of an **existing active session**
- The attacker relies on the **legitimate user** to create a connection and authenticate





## Objective 02

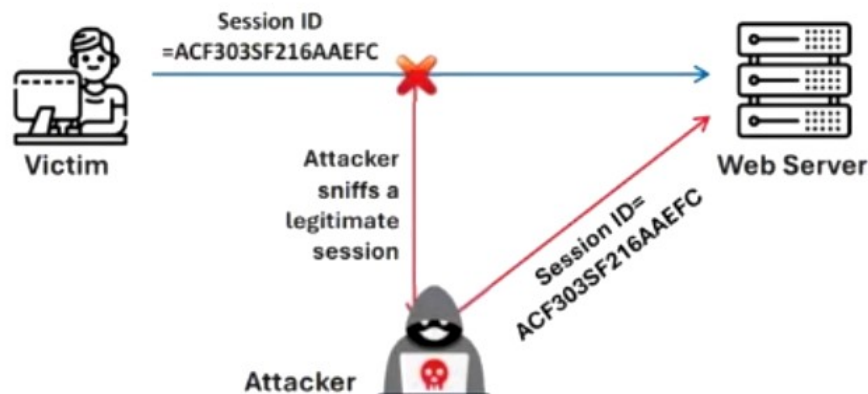
# Explain Application-Level Session Hijacking



# Compromising Session IDs using Sniffing and by Predicting Session Token

## Compromising Session IDs using Sniffing

- An attacker uses a sniffer to **capture a valid session token** or **session ID**
- The attacker then uses the valid token session to **gain unauthorized access** to the web server



## Compromising Session IDs by Predicting Session Token

- Attackers can **predict session IDs** generated by weak algorithms and **impersonate a website user**
- Attackers analyze variable sections of session IDs to **determine a pattern**
- The analysis is performed **manually** or **using various cryptanalytic tools**
- Attackers **collect a high number of simultaneous session IDs** to gather samples in the same time window and keep the variable constant

# How to Predict a Session Token

## Analyzing Token Patterns

### Sequential Tokens

Tokens can be predicted by attackers if they follow an identifiable pattern

<http://www.certifiedhacker.com/view/JBEX1001>  
<http://www.certifiedhacker.com/view/JBEX1002>  
<http://www.certifiedhacker.com/view/JBEX1003>

↓ ↓  
Constant Sequential

### Timestamp-based Tokens

Tokens are easier to predict if they include a timestamp

<http://www.certifiedhacker.com/view/JBEX20240611T1234>  
<http://www.certifiedhacker.com/view/JBEX20240611T1236>  
<http://www.certifiedhacker.com/view/JBEX20240611T1238>

↓ ↓ ↓  
Constant Date Time

## Brute Force Attacks

### Small Token Space

A small token space allows attackers to use brute-force attacks to guess all possible tokens

<http://www.certifiedhacker.com/view/0011>  
<http://www.certifiedhacker.com/view/0033>  
<http://www.certifiedhacker.com/view/0055> } → Small token size

### Lack of Rate Limiting

Without **rate limiting**, attackers can make numerous token guesses without being blocked

## Weak Random Number Generators

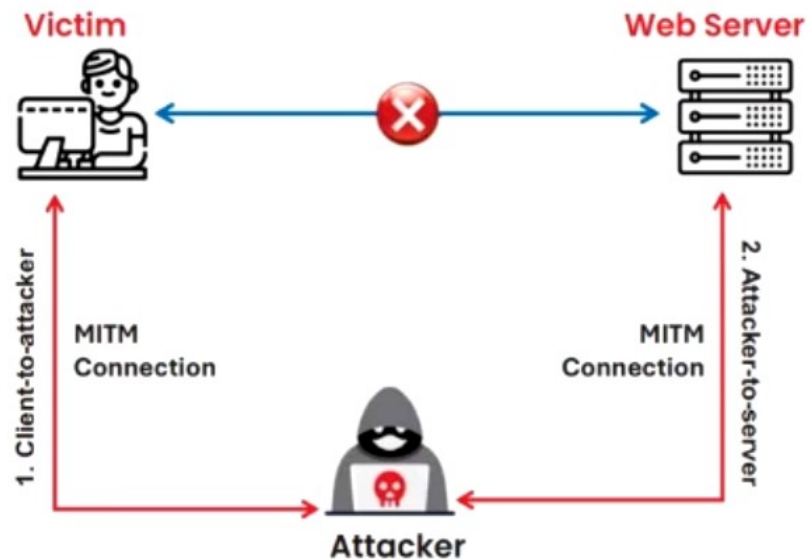
### Predictable PRNG

Predictable PRNGs can produce token sequences that attackers can guess if they know the seed or algorithm

# Compromising Session IDs Using Man-in-the-Middle/ Manipulator-in-the-Middle Attack

The man-in-the-middle/manipulator-in-the-middle attack is used to **intrude into an existing connection** between systems and intercept the messages being exchanged

- Attackers use different techniques and **split the TCP connection** into two connections:
  - Client-to-attacker connection
  - Attacker-to-server connection
- After the interception of the TCP connection, an attacker can read, modify, and insert fraudulent data into the **intercepted communication**
- In the case of an **http transaction**, the TCP connection between the client and the server becomes the target



# Compromising Session IDs Using Man-in-the-Browser /Manipulator-in-the-Browser **Attack**

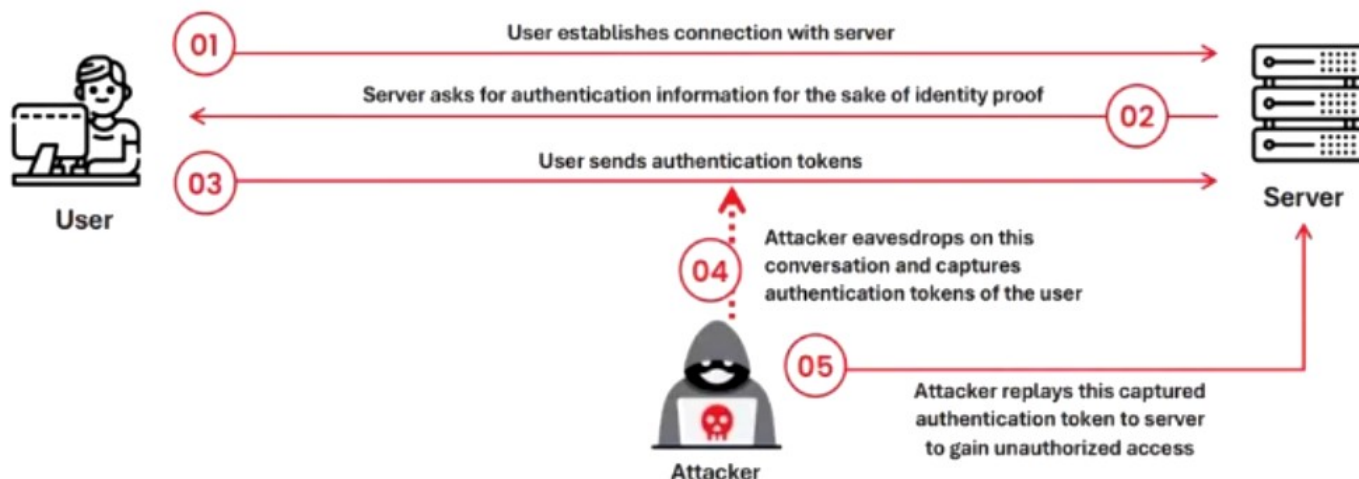
The man-in-the-browser/manipulator-in-the- browser attack **uses a Trojan horse** to intercept the calls between the browser and its security mechanisms or libraries

- 01 The Trojan first infects the **OS** or **application**
- 02 The Trojan installs malicious code (extension files) and saves it in the **browser configuration**
- 03 When the user restarts the browser, it loads the **malicious extension** files
- 04 The extension files register a **handler** for every visit to a webpage
- 05 When a page is loaded, the extension matches the **URL** with a list of **targeted sites**
- 06 The user **logs in securely** to the website
- 07 The extension registers a **button event handler** for specific page loads
- 08 When the user clicks the button; the extension uses **DOM** to extract and modify **form data**
- 09 The browser sends the **form** and **modified values** to the server
- 10 The server **receives modified values** but cannot distinguish from the original
- 11 After the server performs the transaction, a **receipt is generated**
- 12 Now, the browser receives the receipt for the **modified transaction**
- 13 The browser displays the receipt with the **original details**
- 14 The user believes the original transaction was processed **without interception**



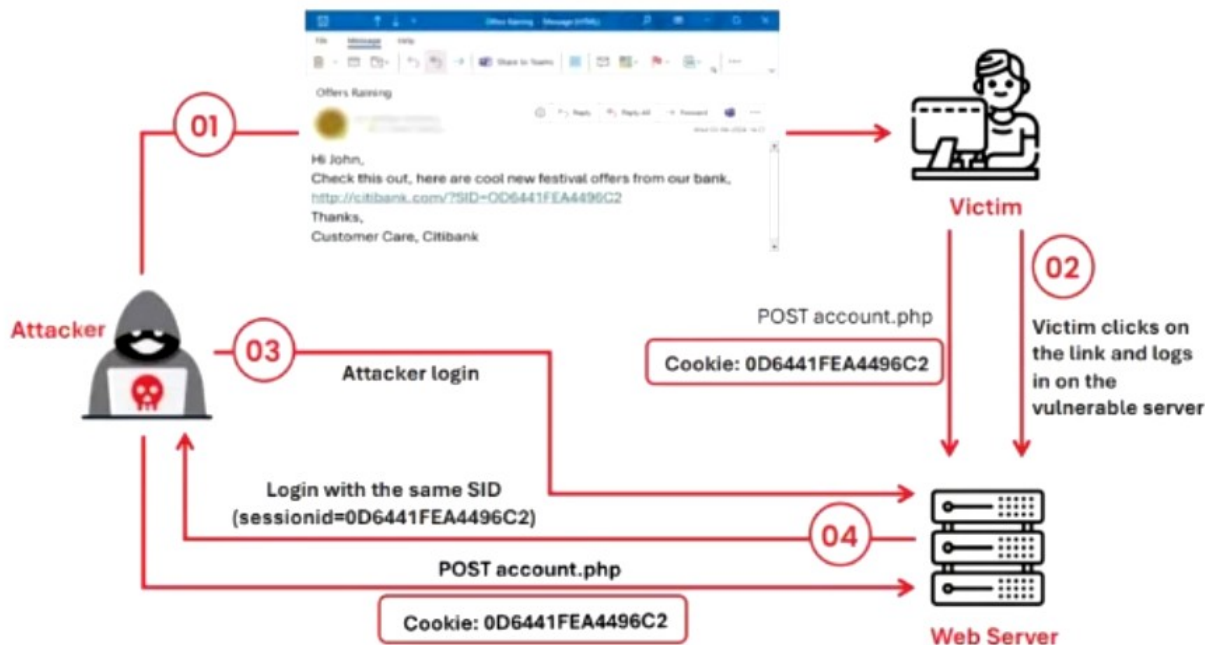
# Compromising Session IDs Using Session Replay Attacks

- In a session replay attack, the attacker listens to the conversation between the **user and the server** and captures the **authentication token** of the user
- Once the authentication token is captured, the attacker **replays the request to the server** with the captured **authentication token** and gains **unauthorized access** to the server



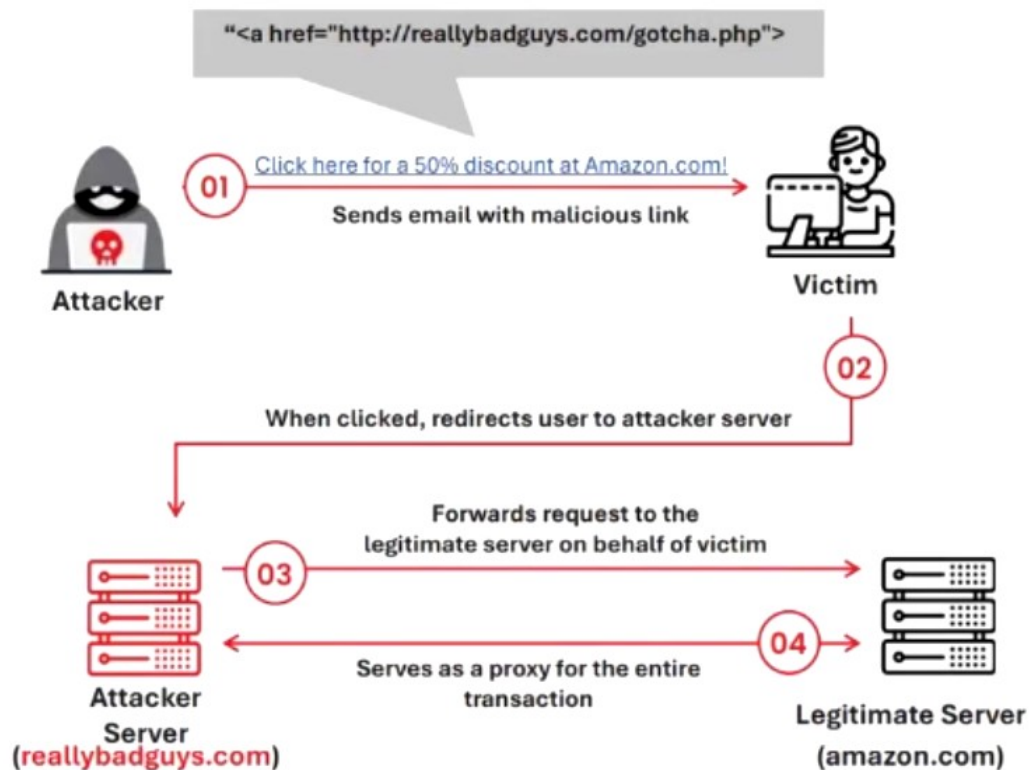
# Compromising Session IDs Using Session **Fixation**

- Session fixation is an attack that allows an attacker to hijack a **valid user session**
- The attacker exploits the **vulnerability of a server** that allows a user to use a fixed SID
- The attacker provides a **valid SID** to a victim and lures him or her to **authenticate** using that SID
- An attacker attempts to lure a user to authenticate himself or herself with a known session ID and then hijacks the **user-validated session** with the knowledge of the used session ID
- The attacker has to provide a **legitimate web application session ID** and attempt to lure the victim's browser to use it
- Some techniques for executing session fixation attacks are as follows:
  - Session token in the **URL argument**
  - Session token in a **hidden form field**
  - Session ID in a **cookie**



# Session Hijacking Using Proxy Servers

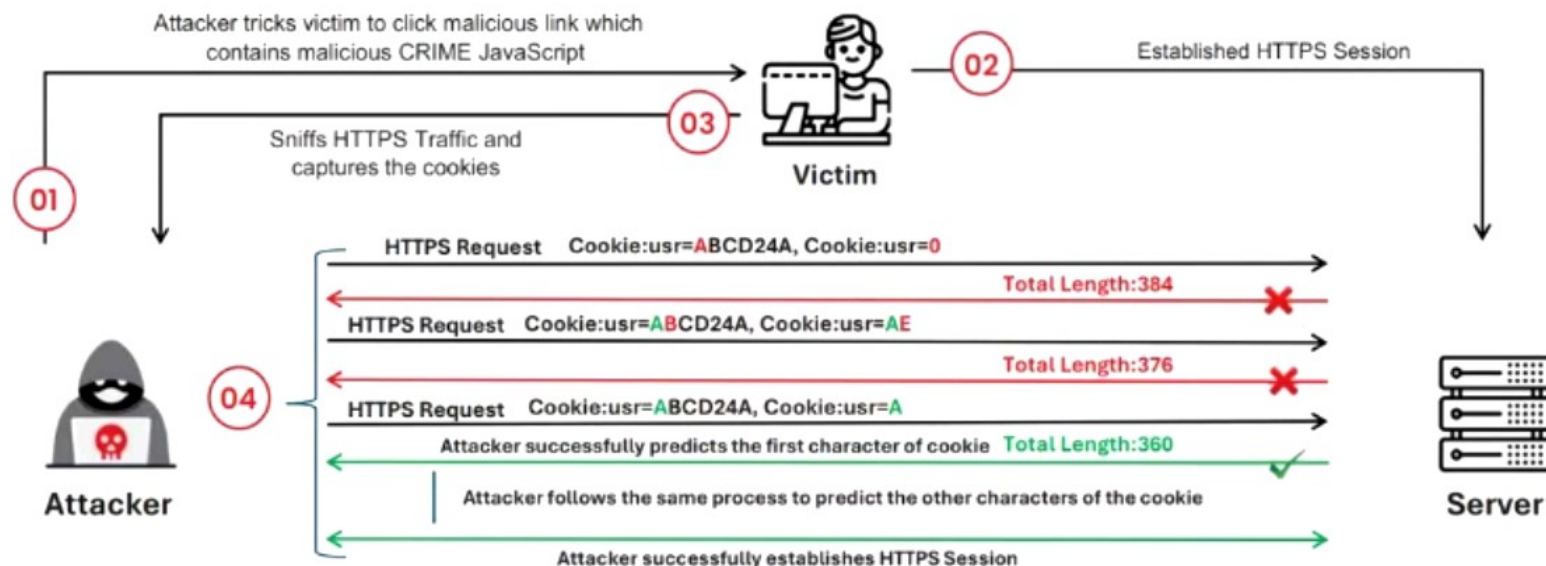
- An attacker lures the victim to **click on a bogus link**, which looks legitimate but redirects the user to the attacker server
- The attacker forwards the request to the legitimate server on behalf of the victim and **serves as a proxy** for the entire transaction
- The attacker then **captures the session's information** during the interaction of the legitimate server and user





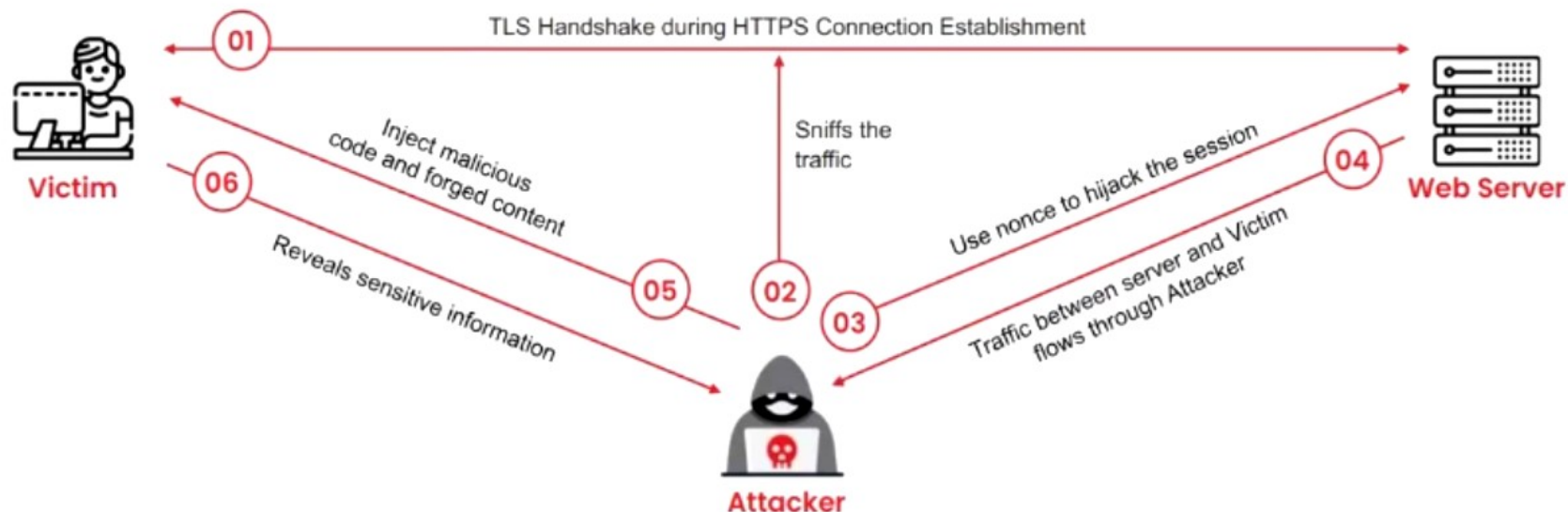
# Session Hijacking Using CRIME Attack

- Compression Ratio Info-Leak Made Easy (CRIME) is a client-side attack that exploits the vulnerabilities present in the **data compression** feature of protocols, such as SSL/TLS, SPDY, and HTTPS
- Attackers hijack the session by decrypting secret **session cookies**
- The authentication information obtained from the session cookies is used to establish a **new session** with the web application



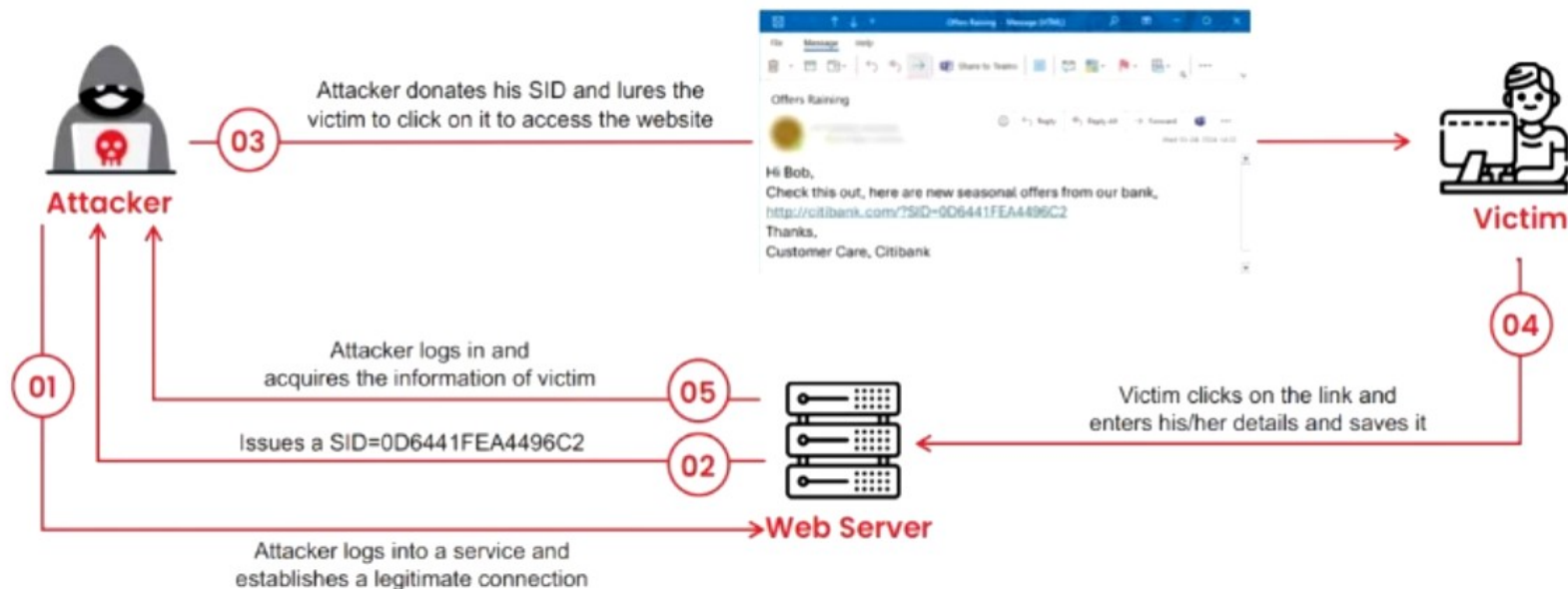
# Session Hijacking Using Forbidden Attack

- A forbidden attack is a type of man-in-the-middle attack used to **hijack HTTPS sessions**
- It exploits the reuse of **cryptographic nonce** during the TLS handshake
- After hijacking the HTTPS session, the attackers **inject malicious code** and **forged content** that prompts the victim to disclose sensitive information, such as bank account numbers, passwords, and social security numbers



# Session Hijacking Using Session Donation Attack

- In a session donation attack, an attacker **donates his/her own session identifier (SID)** to the target user
- The attacker first **obtains a valid SID** by logging into a service and later feeds the same SID to the target user
- This SID **links a target user back to the attacker's account page** without any information to the victim



## Objective 03

# Explain Network-Level Session Hijacking

# Network-Level Session Hijacking

- The network-level hijacking relies on hijacking **transport** and **Internet protocols** used by web applications in the application layer
- By attacking the network-level sessions, the attacker gathers some **critical information**, which are used to **attack the application-level sessions**

Network-level hijacking includes:

01 Blind hijacking

02 UDP hijacking

03 TCP/IP hijacking

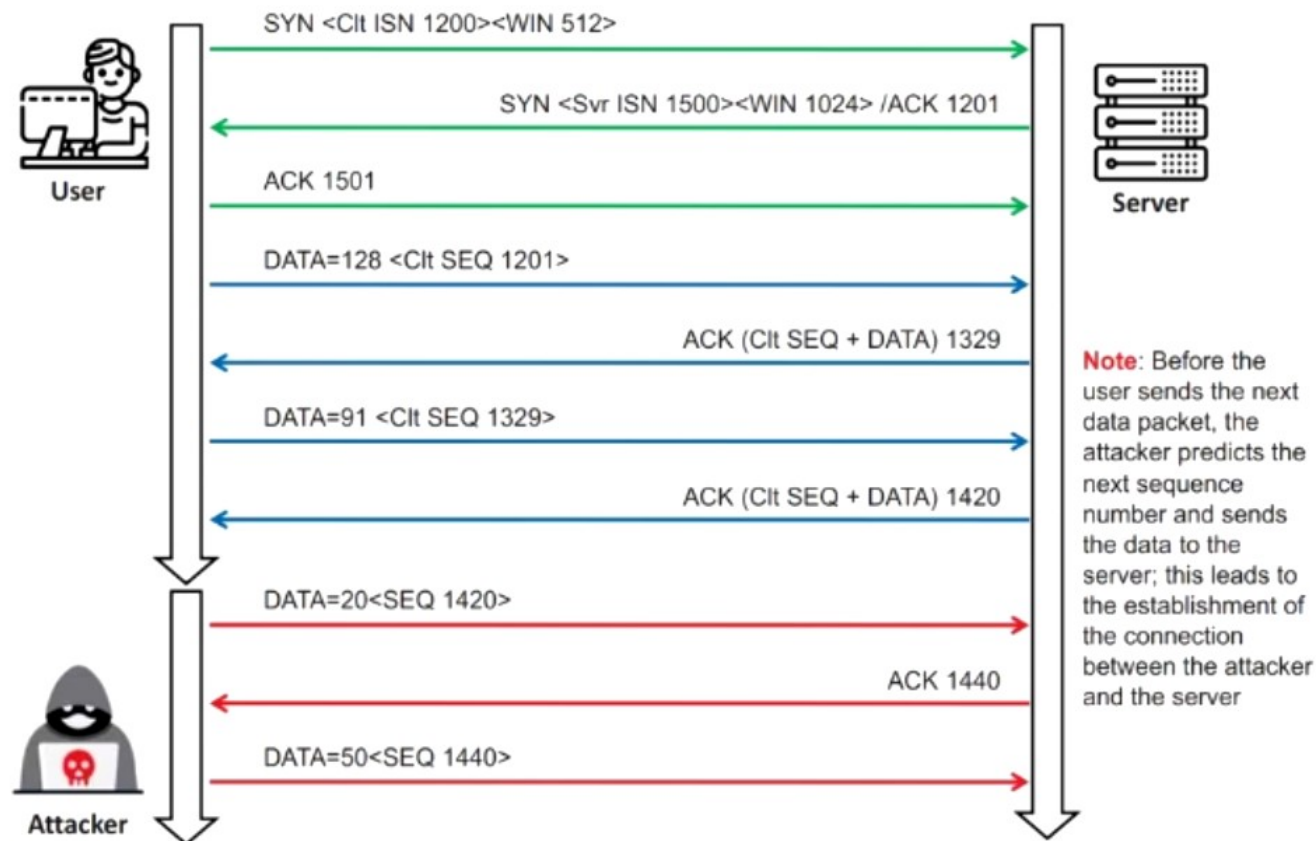
04 RST hijacking

05 Man-in-the-middle: Packet sniffer

06 IP spoofing: Source routed packets

# TCP/IP Hijacking

- TCP/IP hijacking involves using **spoofed packets** to seize control of a connection between a victim and target machine
- A victim's connection hangs, and an attacker is then able to **communicate with the host's machine** as if the attacker is the victim
- To launch a TCP/IP hijacking attack, the **attacker must be on the same network as the victim**
- The target server and the victim machines can be located anywhere





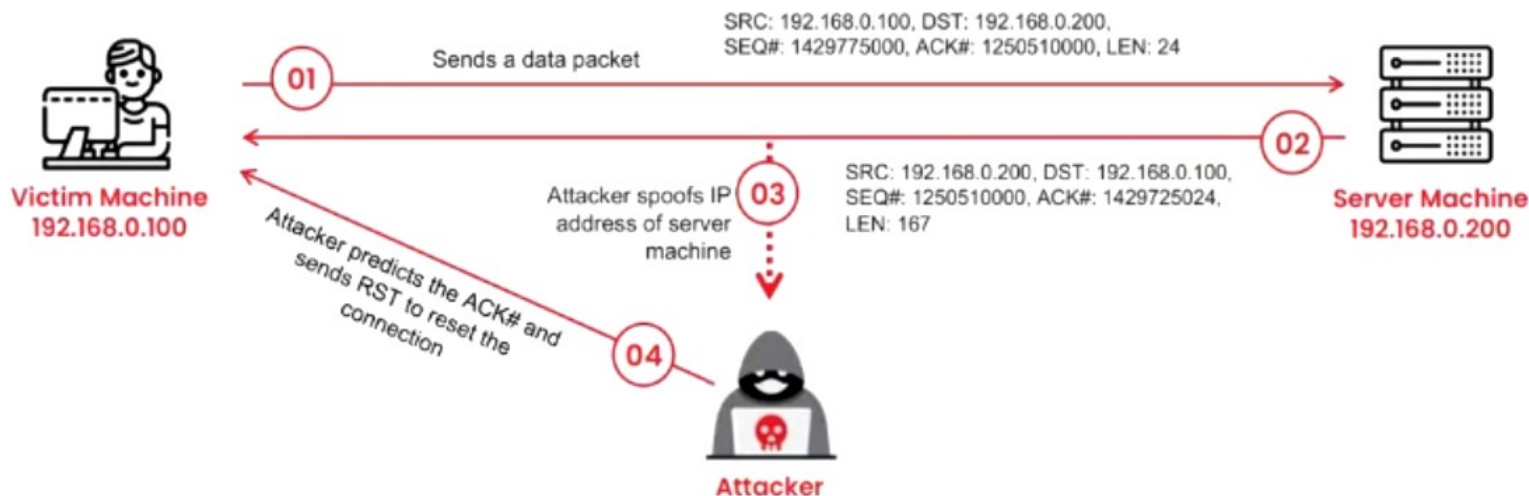
## IP Spoofing: Source Routed **Packets**

- 01 The packet source routing technique is used for **gaining unauthorized access** to a computer with the help of a trusted host's IP address
- 02 An attackers spoofs the host's IP address so that the server **managing a session** with the host accepts the packets from the attacker
- 03 When the session is established, the attacker **injects forged packets** before the host responds to the server
- 04 The original packet from the host is lost as the server receives the packet with a **sequence number** already used by the attacker
- 05 The packets from the attacker are source-routed through the host with the **destination IP** specified by the attacker



# RST Hijacking

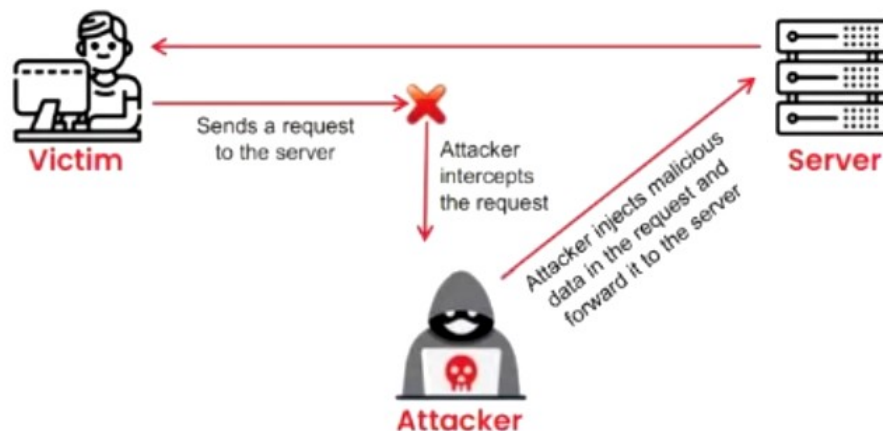
- RST hijacking involves injecting an **authentic-looking reset (RST) packet** using a spoofed source address and predicting the acknowledgment number
- A hacker can reset a victim's connection if it uses an **accurate acknowledgment number**
- The victim would believe that the source sent the **reset packet**, and **reset the connection**
- RST Hijacking can be performed using a **packet crafting tool**, such as Colasoft Packet Builder, and TCP/IP analysis tools, such as tcpdump



# Blind and UDP Hijacking

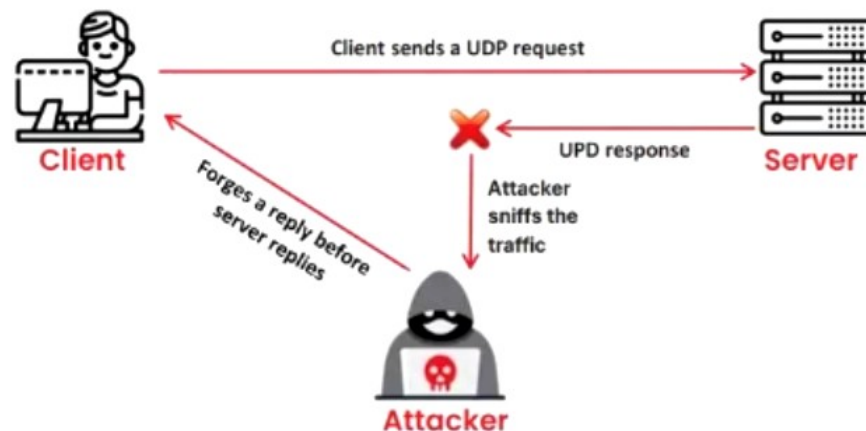
## Blind Hijacking

- An attacker can inject **malicious data or commands** into the intercepted communications in the TCP session even if the source-routing is disabled
- The attacker can send the data or commands but has no **access to see the response**



## UDP Hijacking

- A network-level session hijacking where the attacker sends **forged server reply** to a victim's UDP request before the intended server replies to it
- The attacker uses a **man-in-the-middle** attack to intercept the server's response to the client and sends a forged reply



# MiTM Attack Using Forged ICMP and ARP Spoofing

- In this attack, the packet sniffer is **used as an interface** between the client and server
- An attacker changes the **default gateway** of the client's machine and attempts to reroute packets
- The packets between the client and server are routed through the **hijacker's host** using two techniques, as shown below:

## Forged Internet Control Message Protocol (ICMP)

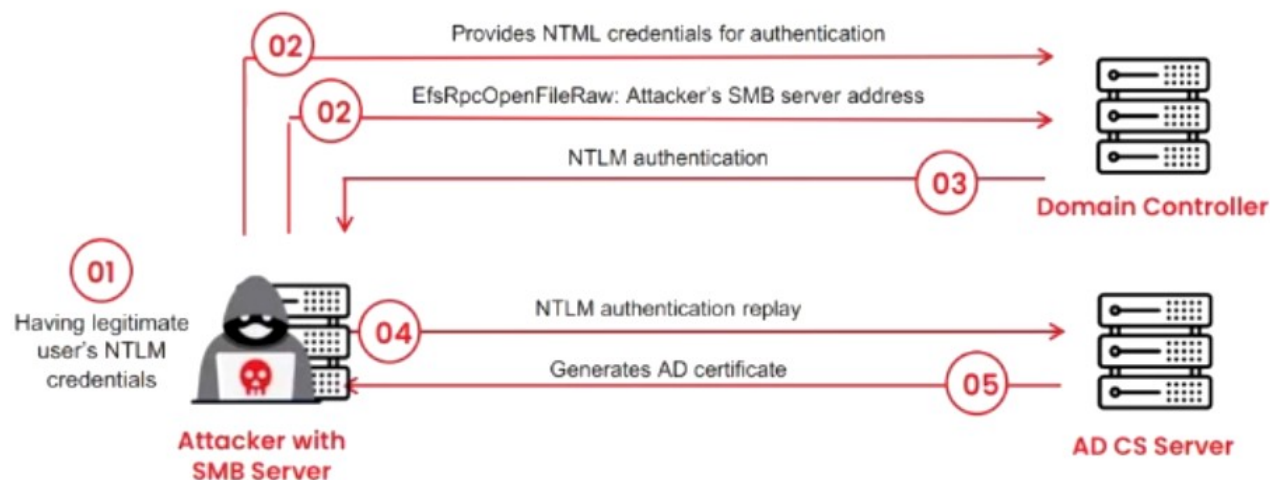
It is an extension of IP to **send error messages** where the attacker can send messages to **fool the client and server**

## Address Resolution Protocol (ARP) Spoofing

ARP is used to map the **network layer addresses** (IP address) to **link layer addresses** (MAC address)

# PetitPotam Hijacking

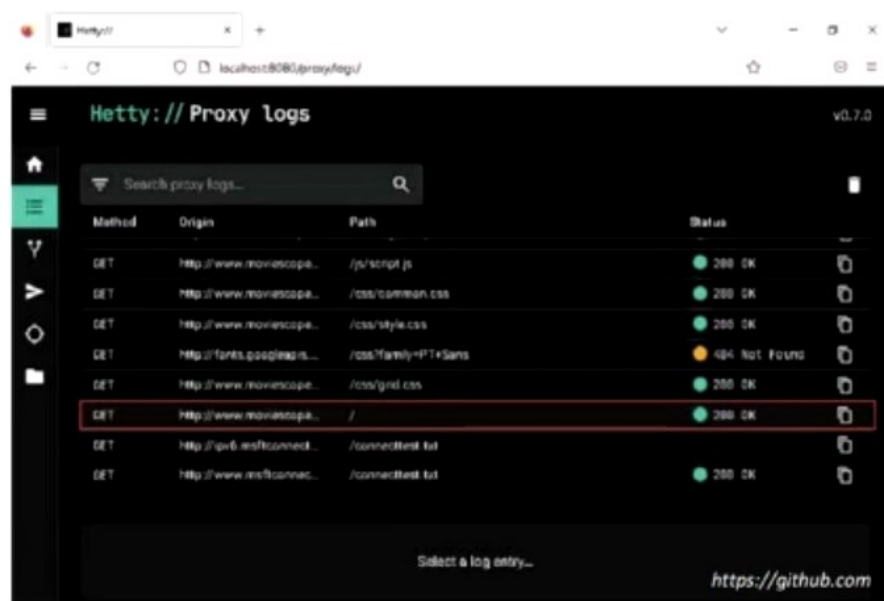
- In a PetitPotam attack, a **domain controller (DC)** is forced by an attacker to initiate authentication to the attacker's server
- The attacker uses **Microsoft's Encrypting File System Remote Protocol (MS-EFSRPC) API** for authentication session hijacking
- The attacker relays the **NTLM authentication** shared by the domain controller to the **Active Directory Certificate Services (AD CS)** server and generates a certificate to acquire **admin-level privileges**



# Session Hijacking Tools

## Hetty

Hetty is an HTTP toolkit that allows attackers to perform **machine-in-the-middle (MITM) HTTP proxy** attack using logs and advanced search



## Caido

Caido is a web **security auditing** toolkit that security professionals can use to intercept and view HTTP requests in **real-time** while browsing



Some more tools are:

**bettercap**  
<https://www.bettercap.org>

**OWASP ZAP**  
<https://www.zaproxy.org>

**WebSploit Framework**  
<https://sourceforge.net>

**sslstrip**  
<https://pypi.org>

**Burp Suite**  
<https://portswigger.net>

## Objective 04

# Explain Session Hijacking Countermeasures



# Protecting against Session Hijacking

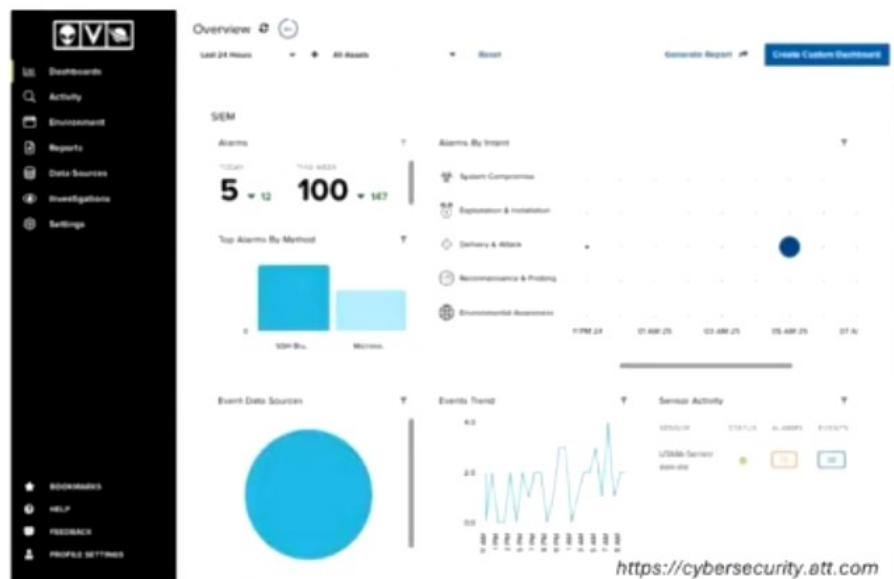
- 01 Use **Secure Shell (SSH)** or **OpenSSH** to create a secure communication channel
- 02 Implement the **log-out functionality** for the user to end the session
- 03 Generate a **session ID** after a successful login and accept only session IDs generated by the server only
- 04 Ensure that data in transit is **encrypted** and implement the **defense-in-depth** mechanism
- 05 Use **string** or **long random numbers** as session keys
- 06 Switch from a hub network to a **switch network** to reduce the risk of session hijacking attacks
- 07 Implement **timeout()** to destroy sessions when expired
- 08 Avoid including the session ID in the **URL** or **query string**
- 09 Ensure that **client-side** and **server-side** protection software are in the active state and up-to-date
- 10 Use **strong authentication** (such as Kerberos) or peer-to-peer virtual private networks (VPNs)
- 11 Configure appropriate **internal** and **external spoof rules** on gateways
- 12 Use **IDS products** or **ARPwatch** for monitoring ARP cache poisoning
- 13 Enable browsers to **verify website authenticity** using network notary servers
- 14 Use SFTP, AS2 managed file transfer, or FTPS to send data using **encryption** and **digital certificates**



# Session Hijacking Detection Tools

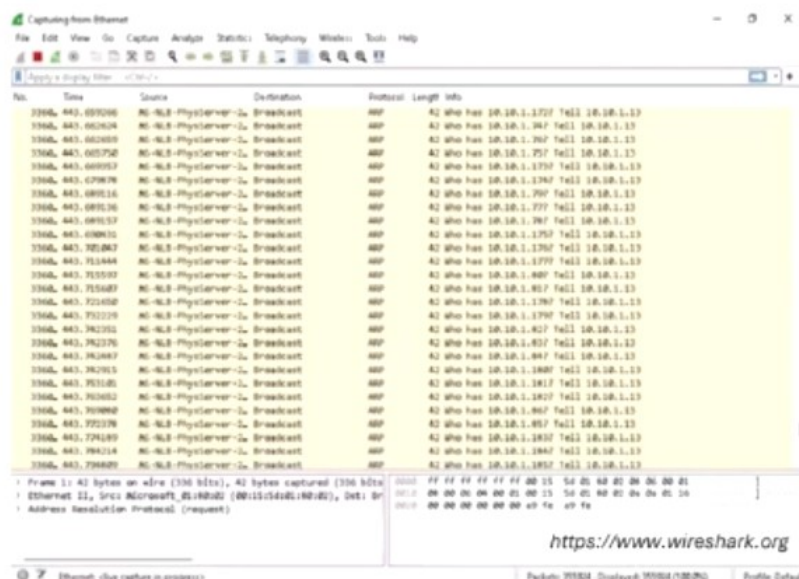
USM  
Anywhere

USM Anywhere delivers **threat detection**, incident response, and compliance management across cloud, on-premises, etc.



Wireshark

Wireshark allows you to **capture and interactively browse the traffic** running on a computer network



Session Hijacking  
Detection Tools:

Quantum Intrusion Prevention System  
(IPS) (<https://www.checkpoint.com>)

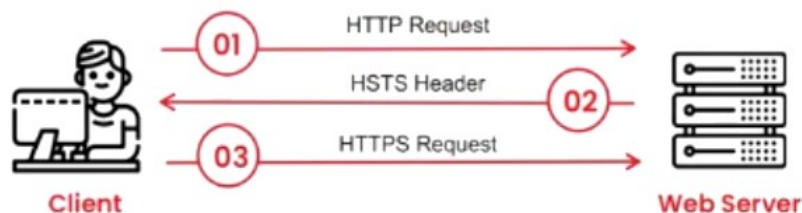
SolarWinds Security Event Manager  
(<https://www.solarwinds.com>)

IBM Security Network Intrusion Prevention  
System (<https://www.ibm.com>)

# Approaches to Prevent Session Hijacking

## HTTP Strict Transport Security (HSTS)

- HTTP Strict Transport Security (HSTS) is a **web security policy** that protects HTTPS websites against MITM attacks
- It allows web servers to **enforce web browsers** to interact with it using secure HTTPS protocol



## Token Binding

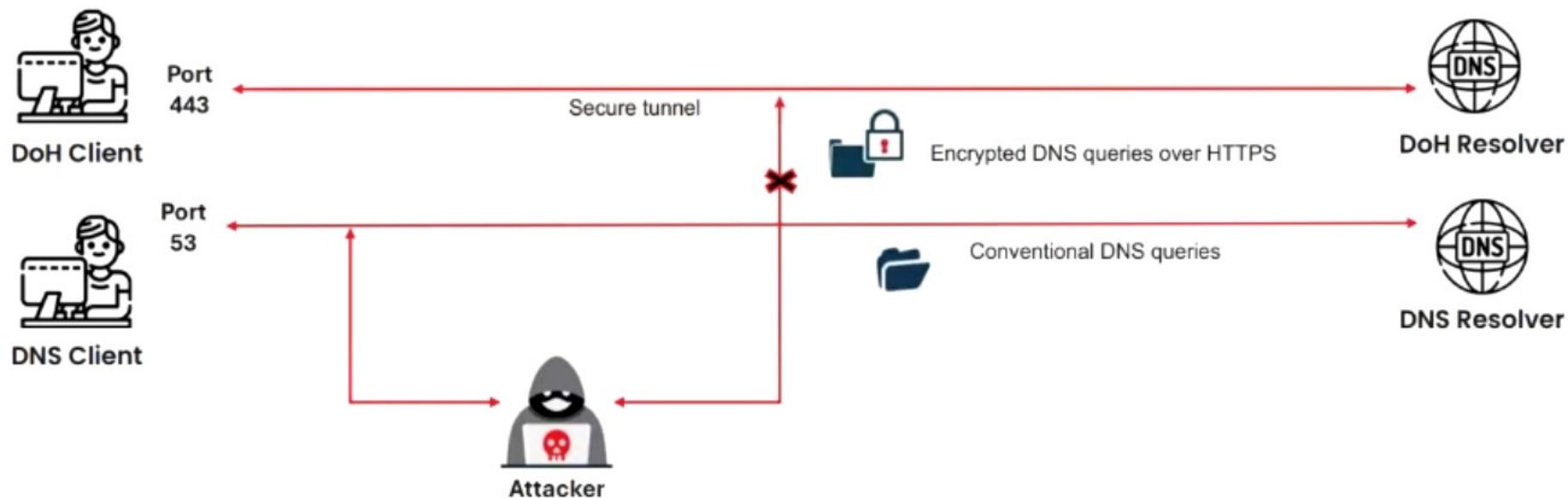
- When a user logs on to a web application, it generates a cookie with an **SID**, called a **token**
- Token binding **protects client-server communications** against session hijacking attacks



# Approaches to Prevent MITM Attacks

## DNS over HTTPS

- DNS over HTTPS (DoH) is an enhanced version of DNS protocol, which is used to **prevent snooping** of user's web activities or DNS queries during the DNS lookup process
- The web queries and traffic are sent through **encrypted HTTPS** via port 443



# IPsec

- IPsec is a protocol suite developed by the IETF for **securing IP communications** by **authenticating** and **encrypting** each IP packet of a communication session
- It is deployed widely to implement **VPNs** and for **remote user access** through dial-up connection to private networks

## IPsec Authentication and Confidentiality

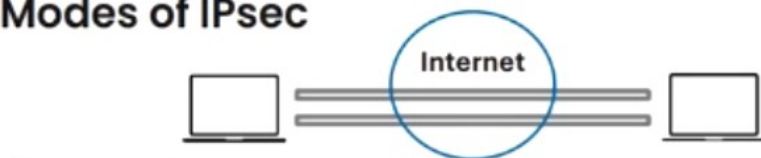
- IPsec uses two different security services for authentication and confidentiality
  - **Authentication Header (AH)**: Provides the data authentication of the sender
  - **Encapsulation Security Payload (ESP)**: Provides both the data authentication and encryption (confidentiality) of the sender

## Benefits of IPsec

- Network-level peer authentication
- Data origin authentication
- Data integrity
- Data confidentiality (encryption)
- Replay protection

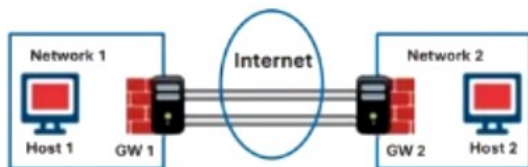
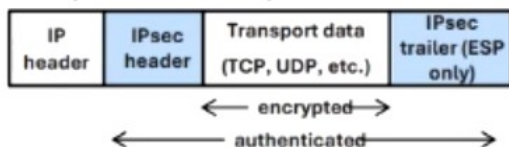
# IPsec (Cont'd)

## Modes of IPsec



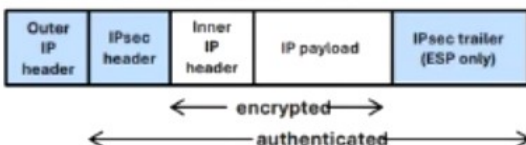
### Transport Mode

Transport – mode encapsulation

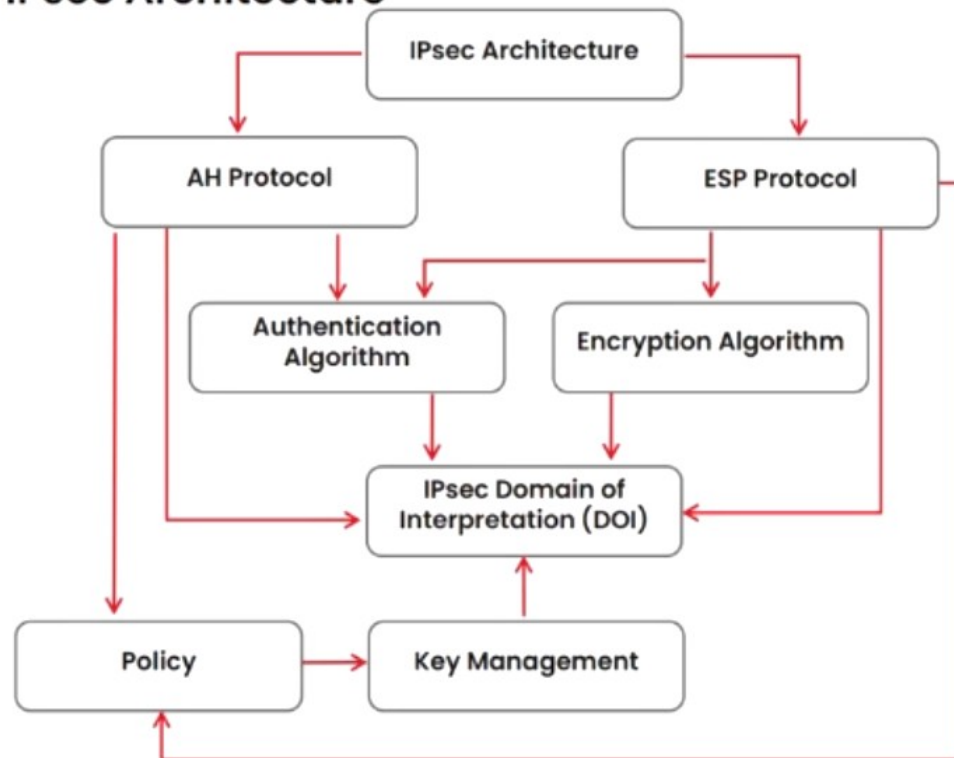


### Tunnel Mode

Tunnel – mode encapsulation



## IPsec Architecture





# Module Summary



- In this module, we have discussed the following:
  - ✓ Session hijacking concepts and different types of session hijacking
  - ✓ Application-level and network-level session hijacking attacks
  - ✓ Various session hijacking tools
  - ✓ How to detect, protect, and defend against session hijacking attacks, as well as various session hijacking detection and prevention tools
  - ✓ We concluded with a detailed discussion on various countermeasures to be employed to prevent session hijacking attempts by threat actors
- In the next module, we will discuss in detail how attackers, as well as ethical hackers and pen-testers, evade network security components such as IDSs and firewalls to compromise the infrastructure