CIS*2030 - Lab#5

Writing and Assembling 68000 Programs

The purpose of this lab is to gain a better understanding of how to write and assemble 68000 assembly-language programs.

Upon completion of this lab, students will be able to:

- Read and interpret a *listing* file.
- Read and interpret a Motorola S-Record,
- Understand how the assembler's main directives (EQU, ORG, DC, DS, and END) affect the assembly process,
- Illustrate the effects of assembler directives by creating a memory map,
- Understand the difference between an assembler directive and a machine instruction, and
- Write simple programs using both the assembler directives and arithmetic and data movement instructions.

During this lab you will make use of the following 68000 assembly-language programs which are available from course website:

•	L5_1.X68	provided
•	L5_2.X68	provided
•	L5_3.X68	provided
•	L5_4.X68	To be written by you
•	L5 5.X68	To be written by you

Part A: Listing Files

Recall that in Lab 4 we learned how to take pre-written assembly-language programs and automatically convert them into machine-language programs using a program called an *assembler*. (Strictly speaking, the assembler that we used is called a *cross-assembler*. This is because the assembler program generates machine language for a different type of computer than the one that it is running on. In our case, the assembler is running on an Intel (x86) architecture, but producing machine language for the Motorola 68000.) Using the assembler is easy. Any ASCII text editor can be used to create a source file; that is, an assembly language program. This source file is then assembled to produce a *binary* file (the 68000 machine language program to be executed) and a *listing* file. The listing file contains the source program together with the assembled code in hexadecimal format and any error messages. If errors are found during the assembly phase, you must return to the source file and edit it to correct them. When the program has been assembled and no syntax errors have been made, only then will the assembler produce the binary file that can be run on the 68KMB

To see how the assembler creates a listing file, we will work with the source program illustrated in Fig. 1 (and also given in class).

```
* L5 1.x68: Program given in class
                         ;program starts at $8000
   ORG $8000
      MOVE.B VALUE1,D0 ;load first byte from $9000
START
       MOVE.B VALUE2,D1 ;load second byte from $9001
              DO, D1 ; sum bytes
       ADD.B
       MOVE.B D1, RESULT ; store sum at $9002
               #14
       TRAP
                          ;return to MON68K
   ORG $9000
                  :data starts at $9000
VALUE1 DC.B
              11 ; initialize RAM with
                                         11
VALUE2 DC.B
             22 ; initialize RAM with
                                         22
RESULT DS.B
             1 ; reserve one byte of
                                         RAM
       END START
```

Figure 1: Simple program presented in class (L5 1.x68).

The previous program loads two bytes stored in memory (with predefined values 11 and 22, respectively) into data registers D0 and D1, respectively. The (byte) contents of both registers are then added together, and the resulting sum is stored back in memory at hexadecimal address 9002. The program terminates by passing control back to the monitor program.

To assemble and generate a *listing file* for the previous program, follow the two steps:

- 1. After logging in, create the directory structure $H: \langle 2030 \rangle L5$ under the H: directory. Download the L5_1.x68 file from the course website and save it as L5_1.x68 under the H: directory. Note: The text in the file should be left justified; that is, it should start at column 0. Also, be sure to explicitly type the .x68 extension when saving!
- 2. To create a listing file, simply assemble the source program; that is, first open a DOSBox window from **Start Menu>All Programs>DOSBox...**, then change into the directory H:\2030\L5 and enter the following command:

```
H:\2030\L5>asm68 L5 1
```

The previous command causes the assembler (asm68) to take the ASCII encoded source file (L5_1.x68) and convert it into a binary (or object) file (L5_1.bin) that contains the binary machine-language code produced by the assembly process; but it also causes the assembler to generate a listing file that can be viewed using a text editor (L5_1.lis). **Note:** the assembler will only produce a binary file if there are no syntax errors in the original source file.

Now use a text editor to examine the listing file. Your listing file should look like the one given below in Fig. 2. Notice that the listing file shows the original source program before (right-hand side) and after (left-hand side) assembly.

```
Source file: L5_1.X6
2
     L5_1.x68
3
                                                   ;program starts at
;load first byte from
;load second byte from
  0008000
                                ORG
                                          $8000
5
  00008000 103900009000
                                MOVE.B
  0000800C D200
                                MOVE.B
                                ADD.B
                                          D0,D1
                                                    ;sum bytes
  0000800E 13C100009002
                                                    ;store sum at $9002
8
                                MOVE.B
  00008014 4E4E
                                          #14
                                TRAP
                                                    ;return to MON68K
1
1
1
                                                   ;data starts at $9000 ;initialize RAM with
   00009000
                                ORG
   00009000 OB
                                DC.B
                                            11
   00009001 16
00009002 00000001
                                DC.B
DS.B
                                                    ;initialize RAM with
                                                   ;reserve one byte of
1
1
            000080
                                END START
   Lines: 16, Errors: 0, Warnings: 0.
```

Figure 2: List file produced by assembler (L5_1.LIS).

Using your knowledge of listing files and instruction formats from class, answer the following questions related to the listing file above.

iatec	ited to the listing file above.			
a)	At what hexadecimal address in memory does the instruction ADD.B D0,D1 start?			
b)	At what hexadecimal address in memory does the MOVE.B D1,RESULT end?			
c)	What is the hexadecimal value of the <i>operation word</i> for the MOVE.B D1,RESULT instruction?			
d)	What are the hexadecimal addresses of the first and second <i>extension words</i> for the instruction MOVE.B D1,RESULT?			
e)	Using the instruction format for the ADD instruction given on page 267 of your textbook and the listing file in Fig. 2, identify which bits of the instruction ADD.B D0,D1 represent the following components of the assembly-language instruction:			
	ADDB Do Now Consider the following S-Record:			

S1138000103900009000123900009001D20013C111

Identify the following five fields in the S-Record:			
a) Record Type:			
b) Byte Count:			
c) Load Address:			
d) Data Bytes:			
e) Checksum:			

Part B: Motorola S-Records

You will recall from lab 4 that the (binary) machine-language program produced by the assembler must be converted into Motorola S-Records before it can be downloaded to the 68KMB. The previous S-Records are a common format for transferring binary data between devices. The binary data are encoded as ASCII characters which can then be transferred between two devices using a simple serial protocol, like the one that exists between the SunRay and the 68KMB.

1. To convert the binary file created in step 2 of Part A (L5_1.bin) to Motorola S-Records (L5_1.txt) **open a Windows' Command Prompt** from **Start Menu>All Programs>Accessories**, change into directory H:\2030\L5 and enter the following command:

H:\2030\L5>srec L5 1

This will cause the file "L5_1.txt" to be created. This file can be downloaded to the 68KMB and is illustrated in Fig. 3. The file contains a total of five records: one of type S0, one of type S9, and three of type S1. The latter type of record is known as a data record. In general, data records can appear in any order in the file, and each record's subfields are composed of ASCII bytes whose associated characters, when paired, represent one-byte hexadecimal values. **Review the description of Motorola S-Records on page 50 of your textbook**.

```
S1138000103900009000123900009001D20013C111
S1098010000090024E4E38
S10590000B1649
S90380007C
```

Figure 3: Machine-language program L5_1.BIN encoded using Motorola S-Records.

Part C: Assembler Directives

As discussed in class, an assembly language is made up of two types of statements: executable instructions and assembler directives. An executable instruction is one of the processor's valid instructions and is translated into the appropriate machine-language code by the assembler. We have already encountered a number of typical instructions (e.g., ADD, MOVE, SUB, DIV, etc.). Assembler directives cannot be translated into machine code; they simply tell the assembler things it needs to know about the program and its environment. Basically, assembler directives link symbolic names to actual values, allocate storage for data in RAM, set up pre-defined constants, and control the assembly process. The assembler directives to be covered in this section are: EQU, DC, DS, ORG, and END. Before proceeding review each of the previous directives by reading Sections 4.6.1 through 4.6.5 of your textbook.

In Fig. 4 (below) you will find a 68000 assembly-language program that employs all five of the previous assembler directives. Download the program (called L5_2.x68) from the CourseLink and load it into the 68000's memory. If you forget how to do this, review the procedure in Part A of Lab 4.

```
*L5 2 sample program
       EQU $8000
CODE
       EQU $9000
DATA
        ORG CODE
       MOVE.W $9000,D0
START
                #7,D0
        MULU
        DIVU
                #3,D0
        ADD.W
                $9004,D0
        MOVE.W D0,$9002
        TRAP
                #14
        END CODE
ORG DATA
   DC.W
            $10
   DS.W
            1
            10
   DC.W
   END DATA
START
```

Figure 4: Sample program - L5 2.X68.

1.	Execute the program in trace mode and answer the questions below.	
	a) What does the program do?	
	b) What is the effect of the assembler directive 'ORG CODE'?	
	c) What is the effect of the assembler directive 'DS.W 1'?	1
	d) What is the effect of the assembler directive 'DC.W \$10?	
	e) What is the effect of the assembler directive 'DC.W 10'?	

f) What is the effect of the assemble directive 'CODE EQU \$8000'?

2. Given the DC and DS directives below, fill in the contents of the memory map (on the next page) and give the starting address that each symbol represents.

	ORG	\$9000
list	DC.B	\$A, 2, %110
val1	DC.W	\$CB
val2	DS.B	3
val3	DC.L	list + 2
team	DC.B	'Jays'

Symbol	Address
list	
val l	
val2	
val3	
team	

	KAM
9011	
9010	
900F	
900E	
900D	
900C	
900B	
900A	
9009	
9008	
9007	
9006	
9005	
9004	
9003	
9002	
9001	
9000	

RAM.

Part D: Assemble-Time Expressions

Most 68000 assemblers permit the programmer to employ an expression as a valid name. The assembler, to yield a numeric value, evaluates expressions. For example, you can write the following:

FRAME	EQU	128
FRAME2	EQU	FRAME+
FRAME3	EQU	FRAME2

In this example, the first assembler directive equates the value 128 to the symbolic name FRAME. The next directive equates the symbolic name FRAME2 to the value FRAME+16. Since FRAME has already been equated to 128, the assembler equates FRAME2 to 128+16 or

144. As explained in class, the following code is illegal.

```
FRAME EQU 128
FRAME2 EQU FRAME3+
FRAME3 EQU FRAME*4
```

The code is illegal because the second line attempts to equate FRAME2 to the FRAME3 plus 16, but the value of FRAME3 has not yet been evaluated. The following listing file shows how the assembler reports the error.

```
1
  0000900
                              OR $9000
2
           0800000
                              EQ 128
                      FRAME:
3
           0000000
                              EQ FRAME3+
                      FRAME2
**** Illegal forward reference - operand
                      FRAME3 EQU FRAME*4
4
           00000200
5
           00009000
                              END $9000
```

It is almost natural to regard arithmetic operations (e.g., FRAME+16) as part of the executable program. They are not. Arithmetic expressions are evaluated by the assembler and the resulting numbers used in the assembly. Since the expressions are evaluated when the program is assembled, they are called assemble-time expressions. There are two reasons to use assemble- time expressions. The first is to make your code easier to maintain. The second is to make your code easier to read. Review Section 4.5 of your textbook to learn more about the assemble- time expressions supported by the assembler used in this course.

In Fig. 5 (below) you will find a 68000 assembly-language program that employs a single assemble-time expression. Download the program (called L5_3.x68) from the course website and load it into the 68000's memory. If you forget how to do this, review the procedure in Part A of Lab 4.

```
*L5 3.x68 example of program with assemble-time expression
        ORG
              $8000
START MOVE.B $9001,D0
                        ; numeric memory address
       MOVE.B VAR1+1,D0 ;assemble-time expression
                  VAR2,D0
       MOVE.B
                           ;label
        TRAP #14
      ORG $9000
      DC.B
              $01
VAR1
      DC.B
              $02
VAR2
VAR3
      DC.B
              $03
VAR4
      DC.B
              $04
              START
        END
```

Figure 5: Example of assemble-time expression.

a)	Examine the listing file for the previous program. Is there any difference between the machine-language code for the three MOVE instructions? Explain your findings.
b)	Trace each MOVE instruction. What value is being placed into D0 each time a move operation is performed?

c) If the instruction "MOVE.B VAR1+1,D0" was replaced with "MOVE.B VAR3-1,D0", how this effect the result in D0?			
d) If the instruction "MOVE.B VAR2,D0" was changed to "MOVE.B #VAR2,D0", what would happen?			
Part E: Write Your Own Program			
Write a program called evaluate to compute the following algebraic expression: $Z = 3Y + 2X$ Assume that variables X, Y and Z are treated as unsigned 16-bit values stored in memory starting hexadecimal address 9000. Use the assembler directives DS and DC to create the variables in memory you can initialize these variables to any values that you consider appropriate. However, you must a labels in your program to access the variables. Save your program as H:\2030\L5\L5_4.X68 . Verify the your program runs correctly by running it on the 68KMB.	ry; ise		
Now re-write the previous program (save it as H:\2030\L5\L5_5.X68), but this time assume that only the first variable (e.g., X) stored in memory starting at hexadecimal address 9000 is identified by a label. Use an assemble-time expression to compute the addresses of the remaining variables. Examine the isting files for the programs in questions 12 through 13. Is there a difference when using labels <i>and</i> assemble-time expressions? Explain.			

As a result of completing this lab you should be able to read and interpret a listing file, read and interpret a Motorola S-Record, understand how the assembler's main directives (EQU, ORG, DC, DS, and END) affect the assembly process, be able to illustrate the effects of assembler directives by drawing a memory map, understand the difference between an assembler directive and a machine instruction, and write (simple) programs using both the assembler directives and arithmetic and data movement instructions discussed thus far in the course.

Looking ahead to Lab 6

Next time we will introduce a new type of register, called an address register. Up to now we have assumed that addresses are encoded as part of the instruction itself. Address registers give us another, more efficient option. Assuming that our data is in memory, we can store the memory address of the data that we wish to access in and address register, then, when we want to access the data, we can tell the instruction to look inside the address register to find the address of the data it needs to fetch. If this reminds you of pointers in C, good, because address registers are used for just that – implementing pointers in high-level languages.