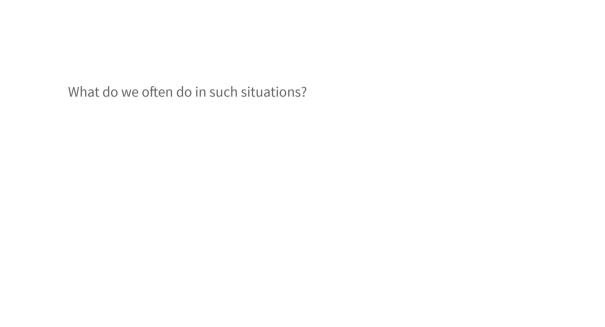
#### **ELICITING AWARENESS**

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- A decision maker must choose a 'plan-of-action;' what action to take provided the future resolution of uncertainty
- ♦ He is unaware of some relevant contingencies and *knows this is possible* 
  - Purchasing a home
  - Hiring a new AP
  - Writing legislation



What do we often do in such situations? seek the council of an expert who is more aware

- Purchasing a home: surveyor
- Hiring a new AP: letter writers, search committee
- Writing legislation: technical advisors

# Why is the interesting?

- When preferences are not aligned, the expert might strategically conceal her awareness
- Can the dm do anything to incentivize revelation?
- Importantly, even with full/complete contracting, the dm cannot articulate what he wants
- A(n unaware) designer may not be able to solve the problem, if mechanisms depend on the unknowns

#### Literature

- Contracting under unawareness
  - Tirole (2009); Filiz-Ozbay (2012); Von Thadden and Zhao (2012); Auster (2013) Auster and Pavoni (2021); Piermont (2017); Lei and Zhao (2021); Francetich and Schipper (2021)
- Modeling unawareness and awareness of unawareness
  - Halpern and Rêgo (2006, 2013); Heifetz et al. (2013); Karni and Vierø (2013, 2017), Halpern and Piermont (2020); Piermont, (2021)
- Robust Mechanism Design
  - ♦ Bergemann and Morris (2005); Jehiel et al., (2006); Carroll (2015, 2019).

- A politician (the decision maker) is trying to write environmental legislation that
  - can be contingent on the future realized environmental state-of-affairs, but
  - can depend only on those contingencies he is aware of.

 He can enlist the help of an environmental scientist (the expert) who may reveal what she is aware of

 $\diamond$  The true state-space is  $\Omega = \{\omega, \nu\}$ ; equally likely

 $\diamond$  Set of actions  $\mathcal{A} = \{a, b, c\}$ 

 $\diamond$  The politician must choose legislation  $\mathfrak{c}:\Omega \to \mathcal{A}$ 

The expert can tell distinguish the states, but the politician cannot.

$$\mathcal{P}_{\mathbf{e}} = \big\{ \{\omega\}, \{\nu\} \big\} \qquad \qquad \mathcal{P}_{d} = \big\{ \{\omega, \nu\} \big\}$$

How does the politician view payoffs in coarse states?
♦ Assume it is aggregated via expectations
⋄ As if he correctly assesses randomness, but cannot explain what causes it

$$u_{d} = \begin{cases} \begin{array}{c|c|c|c} a & b & c \\ \hline \omega & 4 & 0 & 2 \\ \hline \nu & 0 & 6 & 0 \end{array} & u_{e} = \begin{cases} \begin{array}{c|c|c} a & b & c \\ \hline \omega & 0 & 2 & 4 \\ \hline \nu & 0 & 2 & 4 \end{array} \\ \\ u_{d} = \begin{cases} \begin{array}{c|c|c} a & b & c \\ \hline \omega, \nu \} & 2 & 3 & 1 \end{array} & u_{e} = \begin{cases} \begin{array}{c|c|c} a & b & c \\ \hline \omega, \nu \} & 0 & 2 & 4 \end{array} \end{cases}$$

$$u_{d} = \begin{cases} \frac{a}{\{\omega, \nu\}} & 2 & 3 & 1 \end{cases} \qquad u_{e} = \begin{cases} \frac{a}{\{\omega, \nu\}} & 0 & 2 & 4 \end{cases}$$

- Without revelation c = b
- $\bullet \mathbb{E}[u_d] = 3, \mathbb{E}[u_e] = 2$

$$\diamond$$
 With revelation:  $\mathfrak{c}': \left\{ egin{array}{l} \omega \mapsto a \\ \nu \mapsto b \end{array} \right.$ 

$$\bullet \mathbb{E}[u_d] = 5, \mathbb{E}[u_e] = 1$$
; So the expert won't reveal.

$$\diamond \; \mathsf{But}, \mathfrak{c}^\star : \left\{ egin{array}{l} \omega \mapsto c \\ \nu \mapsto b \end{array} 
ight. ext{is a Pareto improvement over no revelation} 
ight.$$

$$\diamond \mathbb{E}[u_d] = 4, \mathbb{E}[u_e] = 3$$

- ♦ The Pareto improvement c\*, requires revelation
- But revealing allows the politician to exploit the expert
- What is the politician could commit:
  - ♦ Propose  $\mathfrak{c} = \mathfrak{b}$  (his outside option)
  - $\diamond$  After the expert reveals, propose some other contract  $c^{\dagger}$
  - $\diamond c^{\dagger}$  only get implemented if the expert agrees; else c = b

Internalizing this, the politician solves:

$$\max_{\mathfrak{c}^\dagger:\Omega\to\mathcal{A}}\mathbb{E}[u_d(\mathfrak{c}^\dagger)] \hspace{1cm} \text{subject to} \hspace{1cm} \mathbb{E}[u_{\pmb{e}}(\mathfrak{c}^\dagger)] \geq \mathbb{E}[u_{\pmb{e}}(\mathfrak{c})] \hspace{1cm} \text{(IC)}$$

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$$\diamond$$
 The solution is  $\mathfrak{c}^{\star}: \left\{ \begin{array}{l} \omega \mapsto c \\ \nu \mapsto b \end{array} \right.$ 



- full revelation
- ♦ an efficient contract

So a two stage game with commitment to not revoke the prior proposal resulted in

- ♦ full revelation
- ⋄ an efficient contract

Does this always work?

So a two stage game with commitment to not revoke the prior proposal resulted in

- ♦ full revelation
- ♦ an efficient contract

Does this always work? No

What if the politician as initially unaware of action b?

$$u_e = \begin{cases} & \begin{array}{c|cccc} a & b & c \\ \hline \omega & 0 & 2 & 4 \\ \hline \nu & 0 & 2 & 4 \end{array} \end{cases}$$

- Without revelation  $c^* = a$
- $\bullet \mathbb{E}[u_d] = 2, \mathbb{E}[u_e] = 0$

$$\left\{ \begin{array}{c|ccccc}
 & a & b & c \\
\hline
 & \omega & 0 & 2 & 4 \\
\hline
 & \nu & 0 & 2 & 4
\end{array} \right.$$

- $\diamond$  Under full revelation (same as before):  $\mathbf{c}': \left\{ \begin{array}{l} \omega \mapsto a \\ v \mapsto b \end{array} \right.$
- $\bullet \mathbb{E}[u_d] = 5, \mathbb{E}[u_e] = 1$ ; this satisfies the incentive constraint.

$$d = \begin{cases} \begin{array}{c|cccc} a & b & c \\ \hline \omega & 4 & 0 & 2 \\ \hline \end{array}$$

- But, revealing only b leads to  $\mathfrak{c} = b$
- $\bullet \mathbb{E}[u_d] = 3, \mathbb{E}[u_e] = 2$ , partial revelation is preferred

$$u_d = \left\{ \begin{array}{c|ccc} & a & b & c \\ \hline \omega & 4 & 0 & 2 \end{array} \right.$$

$$\diamond$$
 As before,  $\mathfrak{c}^{\star}: \left\{ egin{array}{l} \omega \mapsto c \\ \nu \mapsto b \end{array} \right.$  is a Pareto improvement over  $\mathfrak{c}=b$ 



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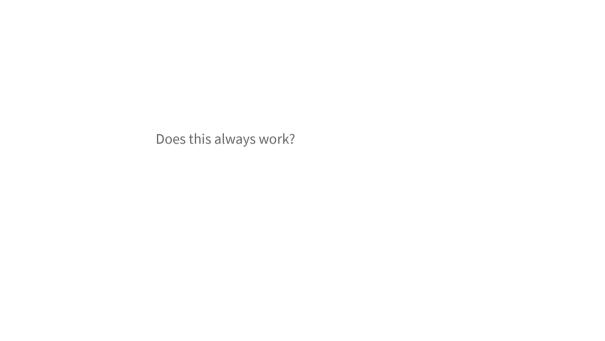
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⋄ c\* is the solution



Does this always work? Yes, but what is 'always'?

# Model

# **Hypothetical State-Space**

Call  $h = (W, A, (v_d, v_e), p)$  a hypothetical states-space,

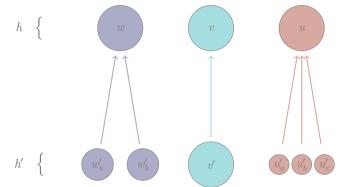
- $\diamond$  W is a finite set of states
- $\diamond$  A is a set of actions
- $v_i: W \times A \to \mathbb{R}$ , for  $i \in \{d, e\}$ , determines payoffs, and,
- $\diamond \ p \in \Delta(W)$  is a probability over states

Let  ${\mathscr H}$  collect all hss;  ${\mathscr H}$  are the possible types of expert

Say that  $h' = (W', A', (v'_d, v'_e), p')$  refines  $h = (W, A, (v_d, v_e), p)$ :

- $\diamond A \subset A'$
- $\diamond$  Each state in h corresponds to an event in h' such that
  - probabilities aggregate
  - expected utilities of h 'measurable' acts are invariant

Then write:  $h \leq h'$ .



## **Outcome Profiles**

An outcome profile is a function from types to contracts:

$$\begin{array}{cccc} f \colon h & \mapsto & \mathfrak{c} \\ & & & & & \\ & & & & & \\ \mathscr{H} & & & \mathbb{R}^W \end{array}$$

Call and outcome profile **fully revealing** if the experts payoff is monotone in her type

$$\mathbb{E}_{p'}\left[v'_{\boldsymbol{e}}(w, f(h'))\right] \ge \mathbb{E}_{p}\left[v_{\boldsymbol{e}}(w, f(h))\right]$$

for any h' such that  $h \leq h'$ 

### Iterated Revision Mechanism

An *iterated revelation mechanism* (IRM), is parameterized by a function from sequences of types to contracts:

$$\alpha:(h_0\ldots h_n)\mapsto (\mathfrak{c}:W_n\to\mathbb{R})$$

- (1) Set n = 0. The decision maker proposes  $\mathfrak{c}_0 = \alpha(h_0)$ .
- (2) The expert reveals  $h_{n+1}$ .
  - $\diamond$  If  $h_n \prec h_{n+1}$ , the decision maker proposes  $\mathfrak{c}_{n+1} = \alpha(h_{n+1})$ . Set n = n+1 and repeat step 2.
  - ♦ Otherwise, continue to step 3.
- (3) The mechanism is over and the expert selects from  $\{c_0, \dots c_n\}$ .

# Full Revelation

#### Theorem

The following are equivalent for an outcome profile f

- (1) f can be implemented by an IRM
- (2) f is a fully revealing

The can be seen as an impossibility result:	
<ul> <li>Without commitment to leave proposed contracts on the table, full revelation cannot be obtained.</li> </ul>	

### Each proposed contract in an IRM specifies:

- (1) The outcome should the game end
  - ⋄ dm wants to maximize his own payoff
- (2) The implicit incentive constraint should the game continue
  - dm wants to minimize the expert's payoff

In the examples, contracts solved (1) ignoring (2)

- The dm cannot conceive of what the expert is aware of
- ♦ It seems prudent, therefore, to consider *robust* strategies: those that maximize the worst case outcome
- Robust strategies turn out to be exactly those that follow the principle of myopic optimization

## Robustness

Call an IRM,  $\alpha$ , robust if at every sequence of (partial) revelations  $\sigma = (h_0, \dots h_n)$ :

- $\diamond \ \alpha$  maximizes the worst case payoff over all best responses that extend  $\sigma$ .
- $\diamond$  relative to any other  $\alpha'$  that coincides with  $\alpha$  over  $\sigma$

#### Robustness

#### Theorem

The following are equivalent:

- (1)  $\alpha$  is robust
- 2)  $\alpha$  is myopically optimal: at each sequence  $(h_0, \ldots, h_n)$ ,

$$lpha(h_0,\dots,h_n)\in \mathop{\sf argmax}_{{\mathfrak c}\colon W o {\mathbb R}} V_d(h,{\mathfrak c}) \qquad \qquad {\sf subject\ to} \ V_e(h_n,{\mathfrak c})>V_e(h_n,lpha(h_0\dots h_m)) \qquad \qquad {\sf for\ } 0\le m< n$$

# The Designers Problem

More generally, often awareness is decentralized:

- A designer wants the decision maker to take some action
- The designer does not know the dm's or the expert's awareness
- ♦ A mechanism elicits awareness and returns an action recommendation

# Mechanism

Mechanis

$$\mathcal{M}:(h^d,h^{m{e}})\mapsto \mathcal{M}(h^d,h^{m{e}})$$

where  $\mathcal{M}(h^d,h^e): \mathit{W}(h^d \wedge h^e) \to \mathcal{A}$ 

#### Desiderata:

INDIVIDUAL RATIONALITY: the dm can not do better alone (there is no constraint for the expert)

INCENTIVE COMPATIBILITY: i prefers to report  $h^i$  than any  $h \prec h^i$ 

PARETO OPTIMALITY: there is no feasible contract that dominates the outcome of the mechanism

These are all **ex-post** restrictions — they must hold for all type realizations

Consider the mechanism,  $\mathcal{M}^{IRM}$ , that implements an myopically optimal IRM.

#### **Theorem**

The mechanism  $\mathcal{M}^{IRM}$ 

- ⋄ is individually rational, incentive compatible, and Pareto optimal, and,
- $\diamond\ V_d$ -dominates any other such mechanism (point-wise over the typespace)

