#### **ELICITING AWARENESS**

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<b>\ \</b>	A decision maker must choose a 'plan-of-action;' what action to take
	provided the future resolution of uncertainty

♦ He is unaware of some relevant contingencies and *knows this is possible* 

♦ He can seek the council of an expert who is more aware than himself

# Why is the interesting?

- When preferences are not aligned, the expert might strategically conceal her awareness
- Can the dm do anything to incentivize revelation?
- Importantly, even with full/complete contracting, the dm cannot articulate what he wants
- A(n unaware) designer may not be able to solve the problem, if mechanisms depend on the unknowns

- A politician (the decision maker) is trying to write environmental legislation that
  - can be contingent on the future realized environmental state-of-affairs, but
  - can depend only on those contingencies he is aware of.

 He can enlist the help of an environmental scientist (the expert) who may reveal what she is aware of

 $\diamond$  The true state-space is  $\Omega = \{\omega, \nu\}$ ; equally likely

 $\diamond$  Set of actions  $\mathcal{A} = \{a, b, c\}$ 

 $\diamond$  The politician must choose legislation  $\mathfrak{c}:\Omega \to \mathcal{A}$ 

The expert can tell distinguish the states, but the politician cannot.

$$\mathcal{P}_{\mathbf{e}} = \big\{ \{\omega\}, \{\nu\} \big\} \qquad \qquad \mathcal{P}_{d} = \big\{ \{\omega, \nu\} \big\}$$

$$u_{d} = \begin{cases} \begin{array}{c|c|c|c} a & b & c \\ \hline \omega & 4 & 0 & 2 \\ \hline \nu & 0 & 6 & 0 \end{array} & u_{e} = \begin{cases} \begin{array}{c|c|c} a & b & c \\ \hline \omega & 0 & 2 & 4 \\ \hline \nu & 0 & 2 & 4 \end{array} \\ u_{d} = \begin{cases} \begin{array}{c|c|c} a & b & c \\ \hline \omega, \nu \} & 2 & 3 & 1 \end{array} & u_{e} = \begin{cases} \begin{array}{c|c|c} a & b & c \\ \hline \omega, \nu \} & 0 & 2 & 4 \end{array} \end{cases}$$

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- Without revelation c = b
- $\bullet \mathbb{E}[u_d] = 3, \mathbb{E}[u_e] = 2$

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$$\diamond$$
 With revelation:  $\mathfrak{c}': \left\{ egin{array}{l} \omega \mapsto a \\ \nu \mapsto b \end{array} \right.$ 

$$\bullet \mathbb{E}[u_d] = 5, \mathbb{E}[u_e] = 1$$
; So the expert won't reveal.

$$a_d = \left\{ \begin{array}{c|ccc} a & b & c \\ \hline \omega & 4 & 0 & 2 \end{array} \right.$$

$$\diamond$$
 But,  $\mathfrak{c}^{\star}: \left\{ egin{array}{l} \omega \mapsto c \\ \nu \mapsto b \end{array} 
ight.$  is a Pareto improvement over no revelation

$$\diamond \ \mathbb{E}[u_d] = 4, \mathbb{E}[u_e] = 3$$

- ♦ The Pareto improvement c\*, requires revelation
- But revealing allows the politician to exploit the expert
- What is the politician could commit:
  - ⋄ Propose  $\mathfrak{c} = \mathfrak{b}$  (his outside option)
  - $\diamond$  After the expert reveals, propose some other contract  $\mathfrak{c}^{\dagger}$
  - $\diamond$   $c^{\dagger}$  only get implemented if the expert agrees

Internalizing this, the politician solves:

$$\max_{\mathfrak{c}^\dagger:\Omega\to\mathcal{A}}\mathbb{E}[u_d(\mathfrak{c}^\dagger)] \hspace{1cm} \text{subject to} \hspace{1cm} \mathbb{E}[u_{\pmb{e}}(\mathfrak{c}^\dagger)] \geq \mathbb{E}[u_{\pmb{e}}(\mathfrak{c})] \hspace{1cm} \text{(IC)}$$

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 The solution is  $\mathfrak{c}^{\star}: \left\{ \begin{array}{l} \omega \mapsto c \\ \nu \mapsto b \end{array} \right.$ 



- full revelation
- ♦ an efficient contract

So a two stage game with commitment to never revoke prior proposals resulted in

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Does this always work?

So a two stage game with commitment to never revoke prior proposals resulted in

- ♦ full revelation
- ♦ an efficient contract

Does this always work? No

- ♦ Take  $\Omega = \{\omega, \nu, \upsilon\}$ ; equally likely
- $\diamond$  Set of actions  $\mathcal{A} = \{a, b, c, d\}$
- ♦ The expert can tell distinguish the states, but the decision maker cannot.

$$\mathscr{P}_{e} = \big\{\{\omega\}, \{\nu\}, \{v\}\big\} \qquad \qquad \mathscr{P}_{d} = \big\{\{\omega, \nu, v\}\big\}$$

- Without revelation c = a
- $\Rightarrow \mathbb{E}[u_d] = \frac{6}{2}, \mathbb{E}[u_e] = \frac{3}{2}$

$$\diamond$$
 With full revelation:  $\mathbf{c}': egin{cases} \omega \mapsto a \\ \nu \mapsto b \\ v \mapsto c \end{cases}$ 

$$\bullet \ \mathbb{E}[u_d] = \frac{9}{3}, \mathbb{E}[u_e] = \frac{4}{3}$$

$$\diamond \ \ \mathsf{Revealing} \ \big\{ \{\omega\}, \{\nu, \upsilon\} \big\} : \mathbf{c''} : \left\{ \begin{array}{ll} \omega & \mapsto a \\ \{\nu, \upsilon\} \mapsto b \end{array} \right.$$

$$\bullet \ \mathbb{E}[u_d] = \frac{7}{3}, \mathbb{E}[u_e] = \frac{5}{3}$$

So, even with commitment, the expert does not fully reveal

$$\diamond$$
 Again, this is inefficient:  $\mathfrak{c}^\star: \left\{ egin{array}{l} \omega \mapsto a \\ \nu \mapsto b \\ v \mapsto d \end{array} \right.$ 

$$\bullet \ \mathbb{E}[u_d] = \frac{8}{3}, \mathbb{E}[u_e] = \frac{6}{3}$$



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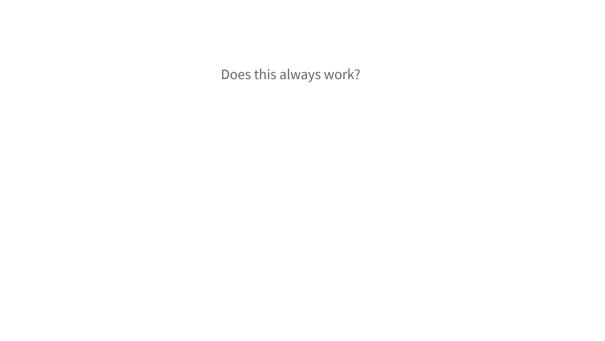
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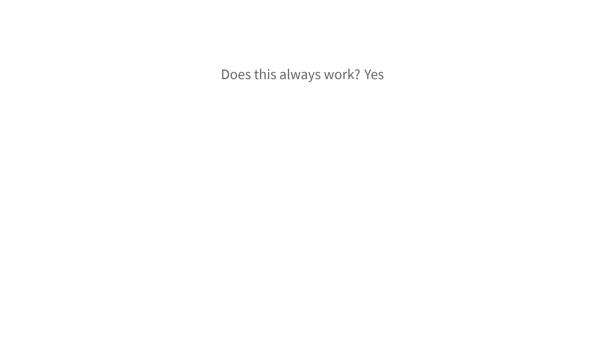
$$\max_{\mathfrak{c}^{\dagger}:\Omega\to\mathcal{A}}\mathbb{E}[u_{\boldsymbol{d}}(\mathfrak{c}^{\dagger})] \qquad \text{subject to} \qquad \mathbb{E}[u_{\boldsymbol{e}}(\mathfrak{c}^{\dagger})] \geq \mathbb{E}[u_{\boldsymbol{e}}(\mathfrak{c}'')] \qquad \text{(IC)}$$

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  - ♦ This becomes the new IC constraint
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⋄ c\* is the solution





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- This iterated procedure (ending when nothing novel is revealed) always results in:
  - full revelation
  - an efficient contract
- More remarkably: any procedure yielding full revelation is equivalent to this one

The latter point can be seen as an impossibility result:	
<ul> <li>Without commitment to leave proposed contracts on the table, full revelation cannot be obtained.</li> </ul>	

## Each proposed contract specifies:

- (1) The outcome should the game end
  - dm wants to maximize his own payoff
- (2) The incentive constraint should the game continue
  - dm wants to minimize the expert's payoff

In the examples, contracts solved (1) ignoring (2)

- The dm cannot conceive of what the expert is aware of
- ♦ It seems prudent, therefore, to consider *robust* strategies:
  - A strategy is *robust* when it maximizes the worst case outcome (over what set?)
- Robust strategies turn out to be exactly those that follow the principle of myopic optimization

## Model

How does the dm reason about the expert?

## **Hypothetical State-Space**

Call  $h = (W, (v_d, v_e), p)$  a hypothetical states-space,

- $\diamond$  W is a finite set
- $\diamond \ v_i: W \times \mathcal{A} \rightarrow \mathbb{R}$ , for  $i \in \{d, e\}$ , and,
- $\diamond p \in \Delta(W)$

Let  $\mathcal{H}$  collect all hss;  $\mathcal{H}$  are the possible types of expert

Say that  $h' = (W', (v'_d, v'_e), p')$  refines  $h = (W, (v_d, v_e), p)$ , written  $h \leq h'$ , if there exists a surjection  $g: W' \to W$  such that

there exists a surjection 
$$q:W'\to W$$
 such that

$$\sum_{w' \in q^{-1}(w)} p'(w') = p(w) \qquad \qquad \text{for all } w \in W \tag{H1}$$
 and

and

$$\sum_{w'\in g^{-1}(s)}w_i'(w',\cdot)p'(w')=v_i(w,\cdot)\qquad \text{ for all }w\in W \text{ and }i\in\{d,e\} \quad \text{(H2)}$$

# Strategies

#### Let $h_0$ be the awareness of the dm:

- ♦ The actions of the expert (of type h) are all sequences of revelations  $H = (h_0, h_1, \dots h_n)$  where  $h_i \prec h_{i+1} \preccurlyeq h$ .
- A strategy of the dm is a mapping from sequences of revelations to contracts

$$\sigma:(h_0,h_1,\ldots h_n)\mapsto\mathfrak{c}$$

where  $\mathfrak{c}:W_n\to\mathcal{A}$ .

## Strategies

Let  $h_0$  the type of the dm and h the type of the expert with  $h = (W, (v_d, v_e), p)$ .

#### **Best Response**

Call  $H = (h_0, h_1, \dots, h_n)$  a best response to strategy  $\sigma$  if

$$\mathbb{E}(v_{e}(\sigma(H)) \ge \mathbb{E}(v_{e}(\sigma(H')))$$

for any other feasible sequence of revelations, H'.

### Full Revelation

#### **Theorem**

The following are equivalent:

(1)  $\sigma$  is equivalent<sup>1</sup> to some  $\bar{\sigma}$  that globally satisfies the IC constraint:

$$\mathbb{E}(v_e(\bar{\sigma}(H)) \geq \mathbb{E}(v_e(\bar{\sigma}(H')))$$
 whenever  $H$  extends  $H'$ . (1

(2)  $\sigma$  is fully revealing: for every type h, it is a best response to reveal h

<sup>&</sup>lt;sup>1</sup>I.e., for all types the strategies yield the same set of best responses / revelations and same implemented contract

#### Robustness

The dm cannot properly envision hypothetical state spaces:

- It seems unreasonable to assume a probability over them
- ♦ Instead, the dm could maximize the worst case outcome
- ⋄ Follows the literature on robustness mechanism design

### Robustness

#### Robustness

Call an strategy,  $\sigma$ , robust if it is IC and for all (reachable) H

$$\inf_{\substack{H' \text{ extends } H \\ H' \text{ best response to } \sigma}} \mathbb{E}(v_d(\sigma(H'))) \geq \inf_{\substack{H' \text{ extends } H}} \mathbb{E}(v_d(\sigma'(H')))$$

for other  $\sigma'$  that coincides with  $\sigma$  up-to H.

### Robustness

#### Theorem

The following are equivalent:

- (1)  $\sigma$  is robust
- (2)  $\sigma$  is equivalent to some  $\bar{\sigma}$  that is myopically optimal:  $\bar{\sigma}(H)$  maximizes  $\mathbb{E}(v_d)$  subject to the IC constraint.

# The Designers Problem

More generally, often awareness is decentralized:

- A designer wants the decision maker to take some action
- The designer does not know the dm's or the expert's awareness
- ♦ A mechanism elicits awareness and returns an action recommendation

## Mechanism

#### mechanis

 $\mathcal{M}: (h^d, h^e) \mapsto \mathcal{M}(h^d, h^e)$ 

where  $\mathcal{M}(h^d,h^e): \mathit{W}(h^d \wedge h^e) \to \mathcal{A}$ 

#### Desiderata:

INDIVIDUAL RATIONALITY: the dm can not do better alone (there is no constraint for the expert)

INCENTIVE COMPATIBILITY: i prefers to report  $h^i$  than any  $h \prec h^i$ 

PARETO OPTIMALITY: there is no feasible contract that dominates the outcome of the mechanism

These are all **ex-post** restrictions — they must hold for all type realizations

Consider the mechanism,  $\mathcal{M}^*$ , that enforces round-by-round commitment then implements the game described above.

#### Theorem

The mechanism  $\mathcal{M}^*$ 

- ⋄ is individually rational, incentive compatible, and Pareto optimal, and,
- $\diamond$   $\mathbb{E}(v_d)$ -dominates any other such mechanism (point-wise over the typespace)

- there is a 'dual' mechanism that is expert-optimal:
- it reverses the proposer and acceptor roles.
- requires only one round

