

Randomised Algorithms

Winter term 2022/2023, Exercise Sheet No. 3

Hand-out: Mon, 31. Oct.

Hand-in: Sun, 6. Nov.

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Exercise 1

[4 points]

Consider running the random QUICKSORT algorithm on the following input:

[11, 19, 1, 8, 7, 10, 15, 9, 4, 18]

- (a) Explain what will happen to the comparison between 8 and 15, in the case the pivot is 7? Same question for the pivot being 15, and being 10?
- (b) Overall, what is the probability that 8 and 15 are compared?

Exercise 2

[6 points]

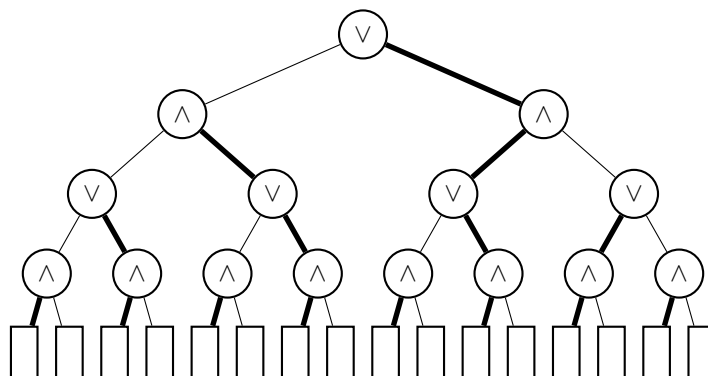
Let \mathcal{A} be a Monte Carlo algorithm with two-sided errors that runs in polynomial time $T(n)$ and outputs a correct result for a decision problem L with probability at least $1/2 + 1/\sqrt{n}$, where n is size of the input.

- (a) Consider the probability amplification with \mathcal{A}_t algorithm like in the lecture, ie. run \mathcal{A} t times then output the decision based on the majority vote, give an upper bound for the error probability of \mathcal{A}_t . You can use the result from the lecture without repeating its arguments in answering this question.
- (b) Find t such that the running time of \mathcal{A}_t remains polynomial while its error probability is reduced to no more than $1/2 - \varepsilon$ for some constant $\varepsilon > 0$, and use this to prove that $L \in \text{BPP}$.
Hint: The inequality $1 + x \leq e^x$ for all $x \in \mathbb{R}$ might be useful.
- (c) Suppose now that our original algorithm \mathcal{A} only has probability $1/2 + 1/n$ to output a correct answer, is L still in BPP? Give a short justification for your answer.

Exercise 3

[6 points]

- (a) Outline an algorithm that allows the construction of a worst input for a deterministic algorithm, denoted $\mathcal{A}(T_k)$, to evaluate a complete OR-AND binary tree T_k (the root is an OR) with k levels of \vee operators and k levels of \wedge operators.
- (b) Execute your algorithm on the following tree, on which each bold edge indicates the first child that will be examined by $\mathcal{A}(T_k)$ during the evaluation of a parent node.



Exercise 4

[4 points]

Prove the following inequality for all matrices M of real numbers:

$$\max_i \min_j M_{ij} \leq \min_j \max_i M_{ij}.$$