Algorithmic Game Theory

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October 7, 2023

About the Course

Class

- Simultaneous and sequential decision making
- Solution concepts and optimal policies
- Practical algorithm for finding the optimal policies

Homeworks

• You will get to implement the games and algorithms!

Understand These!







aStar (c) DeepStack

Game Theory - Reinforcement Learning

Reinforcement Learning

- Single agent settings
- Maximize reward
- Scalable practical algorithms

Game Theory

- Multi agent settings
- Analyzes agent interaction, incentives
- Optimal solution concepts
- Algorithms (historically) tabular and not scalable

Terminology

Reinforcement Learning

- 1. Environment
- 2. Agent
- 3. Policy
- 4. Reward

Game Theory

- 1. Game
- 2. Player
- 3. Strategy
- 4. Utility



Normal Form Games

The normal form games is a model in which each player chooses his strategy, and then all players play simultaneously. The outcome depends on the actions chosen by the players.

Definition: Normal Form Game

is a tuple $\langle N, (A_i), (u_i) \rangle$, where

- *N* is the **finite** set of players
- A_i is the nonempty set of actions available to the player i
- u_i is a **payoff/utility** function for the player i. Let $A = \times_{i \in N} A_i$. $u_i : A \to \mathbb{R}$

Normal Form Games

- If there are only two players (|N| = 2), we can conveniently described the game using a table
- Rows/columns correspond to actions of player one/two
- The cell (i,j) contains the players' payoffs $u_1(i,j)$ and $u_2(i,j)$

Normal Form Games

| | Rock | Paper | Scissors |
|----------|---------|---------|----------|
| Rock | (0, 0) | (-1, 1) | (1, -1) |
| Paper | (1, -1) | (0, 0) | (-1, 1) |
| Scissors | (-1, 1) | (1, -1) | (0, 0) |

Table: Rock-Paper-Scissors

| | Confess | Be Quiet |
|----------|---------|----------|
| Confess | (8, 8) | (0, 10) |
| Be Quiet | (10, 0) | (2, 2) |

Table: Prisoner's dillema

Constant Sum Games

- Constant-sum game is a game for which $u_1 + u_2 = c$
- Zero-sum game is a constant-sum game for c=0, so $u_1=-u_2$
- Critical implications!

Zero Sum Games

| | Rock | Paper | Scissors |
|----------|---------|---------|----------|
| Rock | (0, 0) | (-1, 1) | (1, -1) |
| Paper | (1, -1) | (0, 0) | (-1, 1) |
| Scissors | (-1, 1) | (1, -1) | (0, 0) |

| | Rock | Paper | Scissors |
|----------|------|-------|----------|
| Rock | 0 | -1 | 1 |
| Paper | 1 | 0 | -1 |
| Scissors | -1 | 1 | 0 |

Normal Form Game Strategies

Definition: Pure Strategy

 $a_i \in A_i$ is player i's pure strategy. This strategy is referred to as pure, because there's no probability involved. For example, the player can always play Scissors.

Definition: Mixed Strategy

is a probability measure over the player's pure strategies. The set of player i's mixed strategies is denoted as Π_i . Given $\pi_i \in \Pi_i$, we denote the probability that the player chooses the action $a_j \in A_i$ as $\pi_i(a_j)$ Mixed strategies allow a player to probabilistically choose actions.

Normal Form Game Strategies II

Definition: Support

For a strategy π_i , support is the set of actions with non-zero probability $\{a \in A | \pi_i(a) > 0\}$.

Definition: Strategy profile

Is the set of all players' strategies, denoted as $\pi = (\pi_0, \pi_1 \cdots \pi_n)$. Finally, π_{-i} refers to all the strategies in π except π_i .

Outcome

- Given a pure strategies of all players, we can easily compute the utilities/reward. Player i's utility is $u_i(a)$.
- How to compute the outcome if the players use mixed strategy (they randomize among the pure strategies)? We simply compute the expected value given the probability measure.
- Since the players choose the actions simultaneously, the events are independent and consequently $\pi(a) = \prod \pi_i(a_i)$
- Using this fact, computing the expected value is easy $u_i(\pi) = \sum_{a \in A} \pi(a) u_i(a)$

$$u_i(\pi) = \sum_{a \in A} \pi(a) u_i(a)$$

Outcome Examples

| | Rock | Paper | Scissors |
|----------|---------|---------|----------|
| Rock | (0, 0) | (-1, 1) | (1, -1) |
| Paper | (1, -1) | (0, 0) | (-1, 1) |
| Scissors | (-1, 1) | (1, -1) | (0, 0) |

- $\pi_1 = (0.2, 0.2, 0.6), \pi_2 = (0.2, 0.2, 0.6)$
- $\pi_1 = (0.6, 0.2, 0.2), \pi_2 = (0.2, 0.2, 0.6)$

| | Confess | Be Quiet |
|----------|---------|----------|
| Confess | (8, 8) | (0, 10) |
| Be Quiet | (10, 0) | (2, 2) |

- $\pi_1 = (0.4, 0.6), \pi_2 = (0.4, 0.6)$
- $\pi_1 = (0.6, 0.4), \pi_2 = (0.4, 0.6)$

Best Response

- One of the key concepts, that you will see throughout the class
- Given the strategies π_{-i} of the opponents, the **best response** is the strategy that maximizes the utility for the player.

Definition: Best Response

Best response against a policy π_i is:

$$\underset{\pi_{-i} \in \Pi_{-i}}{\operatorname{arg max}} \, R_{-i}(\pi_i, \pi_{-i})$$

We use $\mathbb{BR}(\pi_i)$ to denote the set of best response policies against the policy π_i .

Best Response

Note that for zero-sum games, opponent maximizing their reward is equivalent to opponent minimizing our reward.

$$\underset{\pi_{-i}}{\operatorname{arg\,max}} R_{-i}(\pi_i, \pi_{-i}) = \underset{\pi_{-i}}{\operatorname{arg\,min}} R_i(\pi_i, \pi_{-i})$$

As this means the player's value against any best-response strategy is unique, we denote this unique value as $BRV_i(\pi_i)$.

$$BRV_i(\pi_i) = \min_{\pi_{-i}} R_i(\pi_i, \pi_{-i}) = -\max_{\pi_{-i}} R_{-i}(\pi_i, \pi_{-i})$$

Best Response

Lemma

For any best response strategy $\pi_i \in \mathbb{BR}_i(\pi_{-i})$, all the actions in the support have the same expected value.

Lemma

The best response set $\mathbb{BR}(\pi_{-i})$ is convex.

Dominated Strategies

- Some actions can be clearly poor choises, and it makes no sense for a rational player to take.
- Strategy π_i^a strictly dominates π_i^b iff for any π_{-i}

$$u_i(\pi_i^a, \pi_{-i}) > u_i(\pi_i^b, \pi_{-i})$$

• Strategy $\pi_i^{\it a}$ weakly dominates $\pi_i^{\it b}$ iff for any $\pi_{-\it i}$

$$u_i(\pi_i^a,\pi_{-i}) \geq u_i(\pi_i^b,\pi_{-i})$$

- Strategy is **strictly/weakly** dominated if there's a strategy that strictly/weakly dominates it.
- Strategies π_i^a, π_i^b are **intransitive** iff one neither dominates nor is dominated by the other.

Examples

Can a weakly/strictly dominated strategy be a best response?

Elimination of Dominated Strategies

• A rational player does not play dominated strategy

| | Left | Center | Right |
|--------|---------|--------|---------|
| Тор | (13, 3) | (1, 4) | (7, 3) |
| Middle | (4, 1) | (3, 3) | (6, 2) |
| Up | (-1, 9) | (2, 8) | (8, -1) |

- A rational player does not play dominated strategy
- Iterated elimination of dominated strategies

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Dominated Strategies and Best Response

Examples

Can a weakly/strictly dominated strategy that we found during the iterated elimination be a best response in the original game?

Week 1 Homework

- 1. Python and notebooks
- 2. Strategy pair evaluation for a matrix game
- 3. Best response calculation
- 4. Strategy evaluation against a best response
- 5. Iterated removal of dominated strategies