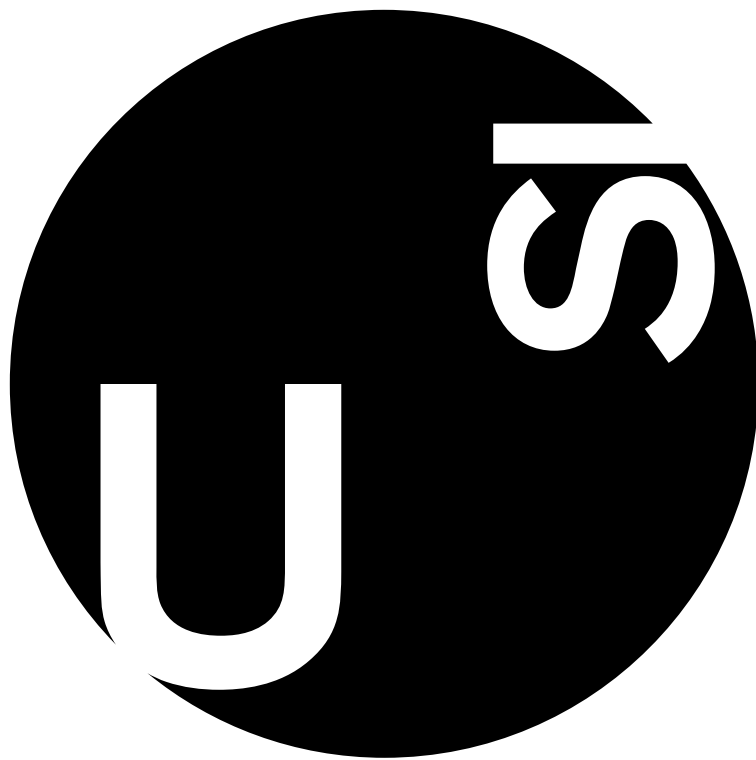


Università della Svizzera italiana

Team Notebook

illUSIon

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1 Template

Write it before the contest begins and then reuse.

```
#include <iostream>
#include <fstream>
#include <vector>
#include <string>
#include <algorithm>
#include <set>
#include <map>
#include <cmath>
#include <cstring>
#include <ctime>
#include <unordered_map>
#include <iomanip>
#include <complex>
#include <cassert>
using namespace std;

#define fi first
#define se second
#define pb push_back
#define all(v) (v).begin(),(v).end()
#define deb(a) cerr<< #a << " = " << (a)<<"\n";

typedef long long ll;
typedef unsigned long long ull;
typedef unsigned int uint;
typedef long double ld;
typedef vector<int> vi;
typedef pair<int,int> pii;

template<class T>
ostream& operator<<(ostream& stream, const vector<T> v
){ stream << "["; for (int i=0; i<(int)v.size();
i++) stream << v[i] << " "; stream << "]; return
stream; }

ll fpow(ll x, ll p, ll m){ll r=1; for (;p>=1){ if (
p&1) r=r*x%m; x=x*x%m; } return r;}
int gcd(int a, int b){ if (!b) return a; return gcd(b,
a%b);}
ll gcd(ll a, ll b){ if (!b) return a; return gcd(b,a%b
);}

int main(){
ios::sync_with_stdio(0);
cin.tie(0);

return 0;
}
```

2 Mathematics

2.1 Number Theory

2.1.1 Mobius inversion formula

$$g = f \star 1 \Leftrightarrow f = \mu \star g$$

Example:

$$\sum_{d|n} \phi(d) = n \Leftrightarrow \phi(n) = \sum_{d|n} \mu(d) \cdot \frac{n}{d}$$

2.1.2 Cipolla's Algorithm

Computes the square root of an elements in a field of prime order.

Let $p \geq 3$ prime and $n \in F_p$. To compute x with $x^2 = n$, take $a \in F_p$ such that $a^2 - n$ is not a square. Take the extension $F_{p^2} = F_p(\sqrt{a^2 - n})$ and compute $x = (a + \sqrt{a^2 - n})^{(p+1)/2}$

2.1.3 Extendend Euclidean Algorithm

```
void gcdext(LL a, LL b, LL &x, LL &y){
    if (b==0){
        x=1, y=0, gcd=a;
        return;
    }

    LL x0, y0;
    gcdext(b, a%b, x0, y0);
    x=y0;
    y=x0-y0*(a/b);
}
```

2.1.4 Lucas's theorem

$$\binom{m}{n} \equiv \prod_{i=0}^k \binom{m_i}{n_i} \pmod{p}$$

m_i and n_i are the i^{th} digits in the base p representation.

2.1.5 Bezout identity

Let $a \neq 0, b \neq 0$. $d = \gcd(a, b)$ is the smallest positive integer for which there integer solutions to:

$$xa + yb = d$$

If (x, y) is one solution, the other solutions are given by:

$$\left(x + \frac{kb}{\gcd(a, b)}, y - \frac{kb}{\gcd(a, b)} \right), k \in \mathbb{Z}$$

2.1.6 Chinese remainder theorem

Let m_1 and m_2 such that $\gcd(m_1, m_2) = 1$. Then the following system:

$$x \equiv a_1 \pmod{m_1}$$

$$x \equiv a_2 \pmod{m_2}$$

has an unique solution $\pmod{m_1 m_2}$, given by:

$$a_1 \cdot m_2 \cdot m_2^{-1} + a_2 \cdot m_1 \cdot m_1^{-1}$$

2.1.7 Miller-Rabin

To check wether n is prime, write $n = 2^s d$. Now if:

$$a^d \not\equiv 1 \pmod{n}$$

and

$$a^{2^r d} \not\equiv -1 \pmod{n}$$

for $0 \leq r \leq s-1$ then n is not prime. Use around 10 iterations for accuracy.

```
def checkPrime(n):
    if n==2: return 1
    if n%2==0: return 0

    d=n-1
    s=0

    while d%2==0:
        s+=1
        d//=2

    for a in range(2, min(17, n-1)):
        if not a==2 and a%2==0: continue

        v=fpow(a, d, n)
        if v==1 or v==n-1: continue

        i=1
        while i<=s-1:
            v=v*v%n
            if v==1: return 0
            if v==n-1: break
            i+=1

        if i==s:
            return 0

    return 1
```

2.2 Equations

2.2.1 Cramer's Rule

$$\begin{cases} ax + by = e \\ cx + dy = f \end{cases} \Rightarrow \begin{cases} x = \frac{ed-bf}{ad-bc} \\ y = \frac{af-ec}{ad-bc} \end{cases}$$

For a general system of equations, use Cramer's rule. If we have a system $Ax = b$, then the solutions are:

$$x_i = \frac{\det A'_i}{\det A}$$

where A'_i is formed by replacing the i^{th} column with b .

2.2.2 Gaussian Elimination

```
for (j=1; j<=min(N,M); j++){
    int ind=-1;
    for (i=j; i<=N; i++){
        if (a[i][j]<-EPS || a[i][j]>EPS){
            ind=i;
            break;
        }
    }

    if (ind==-1) continue;

    LD val=a[ind][j];
    for (k=j; k<=M+1; k++){
        swap(a[j][k], a[ind][k]);
```

```

    a[ind][k]/=val;
}

ind = j;
for (i=ind+1; i<=N; i++){
    LD coef=-a[i][j]/a[ind][j];
    for (k=j; k<=M+1; k++)
        a[i][k]+=coef*a[ind][k];
}
}

for (i=N; i>0; i--){
    for (j=1; j<=M+1; j++)
        if (a[i][j]<-EPS || a[i][j]>EPS){
            if (j==M+1){
                printf("Impossibil\n");
                return 0;
            }

            x[j]=a[i][M+1];
            for (k=j+1; k<=M; k++)
                x[j]-=x[k]*a[i][k];
            x[j]/=a[i][j];
            break;
        }
}
}

```

2.3 Combinatorics

2.3.1 Catalan numbers:

$$C_n = \frac{1}{n+1} \binom{2n}{n}$$

Number of lattice paths that don't go above $x = y$, number of rooted ordered trees, number of full binary trees with $n + 1$ leaves, etc.

2.3.2 Stirling numbers

$$S(n, k) = kS(n-1, k) + S(n, k-1)$$

$$B_n = \sum_{k=0}^n S(n, k)$$

Number of ways to partition n elements into k sets.

2.3.3 Inclusion-Exclusion

$$\begin{aligned}
 \left| \bigcup_{1 \leq i \leq n} A_i \right| &= \sum_{1 \leq i_1 \leq n} |A_{i_1}| - \sum_{1 \leq i_1 < i_2 \leq n} |A_{i_1} \cap A_{i_2}| \\
 &+ \sum_{1 \leq i_1 < i_2 < i_3 \leq n} |A_{i_1} \cap A_{i_2} \cap A_{i_3}| - \dots + (-1)^{n+1} \left| \bigcap_{i=1}^n A_i \right|.
 \end{aligned}$$

2.3.4 Burnside lemma

Let G be a finite group which acts on X . X^g - number of elements fixed by g . We have:

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

Polya enumeration theorem follows from Burnside - look at the number of cycles.

2.3.5 Cayley's formula

Number of labeled trees on n vertices:

$$n^{n-2}$$

2.3.6 Generalized Cayley formula

If there are k connected components, each of size s_i , the number of possible trees on these components is:

$$\begin{aligned}
 \sum_{\substack{d_i \geq 1, \\ \sum_{i=1}^k d_i = 2k-2}} s_1^{d_1} \dots s_k^{d_k} \binom{k-2}{d_1-1, \dots, d_k-1} &= \\
 &= s_1 \dots s_k \cdot n^{k-2}
 \end{aligned}$$

2.3.7 Matrix tree theorem

D - degree matrix, A adjacency matrix. $L = D - A$ - laplacian. The number of spanning trees is calculated as the determinant of L after removing the last row and column.

2.4 Geometry

2.4.1 Pick's theorem

$$A = i + \frac{b}{2} - 1$$

i - number of lattice points in the interior, b - number of lattice points on the boundary. Works for polygons with integer coordinates.

2.4.2 Lattice points on segment

Number of points on a lattice segment from (a, b) to (c, d) is:

$$\gcd(c-a, d-b) + 1$$

2.4.3 Triangles

Side lengths: a, b, c

Semiperimeter: $p = \frac{a+b+c}{2}$

Area: $A = \sqrt{p(p-a)(p-b)(p-c)}$

Circumradius: $R = \frac{abc}{4A}$

Inradius: $r = \frac{A}{p}$

Law of sines: $\frac{\sin \alpha}{a} = \frac{1}{2R}$

Law of cosines: $a^2 = b^2 + c^2 - 2bc \cos \alpha$

2.5 Numerical methods

2.5.1 Simpson's rule

We have:

$$\int_a^b f(x) dx \approx$$

$$(f(x_0) + 4f(x_1) + 2f(x_2) + 4f(x_3) + \dots + 4f(x_{2n-1}) + f(x_{2n})) \frac{h}{3}$$

Where we break up the interval $[a, b]$ into $2n$ parts: $x_i = a + ih$, $h = \frac{b-a}{2n}$. The error is:

$$\leq \frac{1}{180} h^4 (b-a) \max_{x \in [a, b]} |f^{(4)}(x)|$$

3 Graph Algorithms

3.1 LCA in $O(\log n)$

```
const int lg=18;

void dfs(int nod, int t, int ad){
    f[nod][0]=t,d[nod]=ad;
    for (int i=1; i<=lg; i++)
        f[nod][i]=f[f[nod][i-1]][i-1];

    for (int nxt : g[nod])
        if (nxt!=t)
            dfs(nxt,nod,ad+1);
}

int lca(int x, int y){
    if (d[x]>d[y]) swap(x,y);
    for (int i=0; i<=lg; i++)
        if ((d[y]-d[x])&(1<<i))
            y=f[y][i];

    if (x==y) return x;
    for (int i=lg; i>=0; i--)
        if (f[x][i]!=f[y][i])
            x=f[x][i],y=f[y][i];

    return f[x][0];
}
```

3.2 Max-Flow - Edmonds-Karp - $O(N \cdot M^2)$

```
struct MaxFlow{
    vector<vi> c,f,g;
    vi tt;
    int s,t,N;

    MaxFlow(int N, int s, int t){
        this->N = N, this->s = s, this->t = t;

        c.resize(N),f.resize(N),g.resize(N);
        for (int i=0; i<N; i++)
            c[i].resize(N),f[i].resize(N);
    }

    int getmaxflow(){
        int res=0;
        for (int i=0; i<N; i++)
            fill(all(f[i]),0);

        while (bfs()){
            int nod;
            for (nod=0; nod<N; nod++){
                if (c[nod][t]-f[nod][t]==0 || tt[nod]==-1)
                    continue;

                tt[t]=nod;
                int cr,fmin=(1<<30);
                for (cr=t; cr!=s; cr=tt[cr])
                    fmin=min(fmin,c[tt[cr]][cr]-f[tt[cr]][cr]);

                res+=fmin;

                for (cr=t; cr!=s; cr=tt[cr]){
                    f[tt[cr]][cr]+=fmin;
                    f[cr][tt[cr]]-=fmin;
                }
            }
        }

        return res;
    }

    void add_uedge(int from, int to, int cap){
```

```
        g[from].pb(to);
        g[to].pb(from);
        c[from][to]+=cap;
    }

    void add_bedge(int from, int to, int cap){
        g[from].pb(to);
        g[to].pb(from);

        c[from][to]+=cap;
        c[to][from]+=cap;
    }

    void print(){
        cout << N << "\n";
        for (int i=0; i<N; i++, cout << "\n")
            for (int j=0; j<N; j++)
                cout << c[i][j] << " ";
    }

    bool bfs(){
        tt.clear(),tt.resize(N,-1);
        int *q = new int[N+10],K=0,i;
        bool *v = new bool[N+10];

        for (i=0; i<N; i++) v[i]=0;

        q[K++]=s; v[s]=1;
        bool ok=0;
        for (i=0; i<K; i++){
            int nod = q[i];

            for (int nxt : g[nod]){
                if (c[nod][nxt]-f[nod][nxt]>0 && !v[nxt]){
                    if (nxt==t){
                        ok=1;
                        continue;
                    }

                    tt[nxt]=nod;
                    q[K++]=nxt;
                    v[nxt]=1;
                }
            }
        }

        delete[] q;
        delete[] v;
        return ok;
    }
};
```

3.3 Max Flow - Dinic - $O(N^2 \cdot M)$

```
struct MaxFlow{
    vector<vi> c,f,g;
    ll s,t,N;
    vi ptr,d;
    MaxFlow(int N, int s, int t){
        this->N = N, this->s = s, this->t = t;

        c.resize(N),f.resize(N),g.resize(N);
        for (int i=0; i<N; i++)
            c[i].resize(N),f[i].resize(N);
    }

    ll getmaxflow(){
        ll res=0;
        for (int i=0; i<N; i++)
            fill(all(f[i]),0);

        while (bfs()){
            ptr.clear();
            ptr.resize(N,0);

            while (ll pushed=dfs(s,(1LL<<60)))
                res+=pushed;
        }
    }
};
```

```

    }

    return res;
}

ll dfs(int nod, ll flow){
    if (!flow) return 0;
    if (nod==t) return flow;

    for (ll &i=ptr[nod]; i<g[nod].size(); i++){
        int nxt=g[nod][i];
        if (d[nxt]!=d[nod]+1) continue;

        ll pushed=dfs(nxt,min(flow,c[nod][nxt]-f[nod][nxt]));
        if (pushed){
            f[nod][nxt]+=pushed;
            f[nxt][nod]-=pushed;
            return pushed;
        }
    }

    return 0;
}

void add_uedge(int from, int to, ll cap){
    g[from].pb(to);
    g[to].pb(from);
    c[from][to]+=cap;
}

void add_bedge(int from, int to, ll cap){
    c[from][to]+=cap;
    c[to][from]+=cap;
    g[from].pb(to);
    g[to].pb(from);
}

void print(){
    cout << N << "\n";
    for (int i=0; i<N; i++, cout << "\n")
        for (int j=0; j<N; j++)
            cout << c[i][j] << " ";
}

bool bfs(){
    d.clear(),d.resize(N,-1);
    int *q = new int[N+10],K=0,i;

    q[K++]=s; d[s]=0;
    for (i=0; i<K; i++){
        int nod = q[i];

        for (int nxt : g[nod]){
            if (c[nod][nxt]-f[nod][nxt]>0 && d[nxt]==-1){
                d[nxt]=d[nod]+1;
                q[K++]=nxt;
            }
        }
    }

    delete[] q;

    return d[t]!=-1;
}
};

```

3.4 Min Cost Max Flow

```

const int inf = (1<<29);

struct MCMFlow{
    struct edge{
        int to,c,cst;
        edge(int to=0, int c=0, int cst=0) : to(to), c(c),
            cst(cst) {}
        edge(){}
    };

```

```

};

vector<edge> e;
set<pii> st;
vi di;
vector<vi> g;

int n,S,F,flow,cflow;

MCMFlow(int n, int S, int F){
    this->n=n;
    g.resize(n);
    di.resize(n);
    this->S=S,this->F=F;
}

void add_edge(int x, int y, int cp, int cst){
    edge e1(y,cp,cst);
    edge e2(x,0,-cst);
    g[x].pb(e.size());
    e.pb(e1);
    g[y].pb(e.size());
    e.pb(e2);
}

bool dijkstra(){
    vi d(n),p(n),pe(n);
    fill(all(d),inf);
    d[S]=0;
    st.clear();
    st.insert({0,S});
    vi real_d(n);

    while (!st.empty()){
        pii cr = *st.begin();
        st.erase(st.begin());

        int nod=cr.se;
        real_d[nod]=d[nod]-di[nod];

        for (int id : g[nod]){
            int nxt = e[id].to;
            int cst=d[nod]+e[id].cst-di[nod]+di[nxt];

            if (e[id].c && cst<d[nxt]){
                if (d[nxt]!=inf) st.erase({d[nxt],nxt});
                d[nxt] = cst;
                p[nxt]=nod;
                pe[nxt]=id;
                st.insert({d[nxt],nxt});
            }
        }
    }

    di = real_d;
    if (d[F] == inf) return 0;

    int minv = inf,cr=0;
    for (cr=F; cr!=S; cr=p[cr])
        minv=min(minv,e[pe[cr]].c);

    flow+=minv;
    cflow += minv*real_d[F];

    for (cr = F; cr!=S; cr=p[cr]){
        e[pe[cr]].c-=minv;
        e[pe[cr]^1].c+=minv;
    }

    return 1;
}

int get_flow(){
    flow=cflow=0;
    bellman();

    while (dijkstra()) ;
    return flow;
}

```

```

}

void bellman(){
    fill(all(di),inf);
    di[S]=0;

    vi q,nq;
    vi v(n,0);
    q.pb(S);

    for (int i=1; i<=n; i++){
        nq.clear();
        for (int x : q){
            for (int id : g[x]){
                int nxt = e[id].to;
                if (e[id].c && e[id].cst+di[x]<di[nxt]){
                    di[nxt]=e[id].cst+di[x];

                    if (!v[nxt]) nq.pb(nxt);
                    v[nxt]=1;
                }
            }
        }

        q.clear();
        for (int x : nq){
            v[x]=0;
            q.pb(x);
        }
    }
};

```

3.5 Heavy-Light Decomposition

```

int N,M,a[100100],sz[100100],L,d[100100],lid[100100],
    pl[100100]; vi g[100100];
int pos[100100],off[100100];

int t[800100];

vi lant[100100];

void bdfs(int nod, int f, int h){
    d[nod]=h;
    sz[nod]=1;

    if ((nod!=1 && g[nod].size()==1) || N==1){
        lid[nod]=++L;
        lant[L].pb(nod);
        return;
    }

    int bst=0;
    for (int nxt : g[nod]){
        if (nxt==f) continue;

        bdfs(nxt,nod,h+1);
        if (sz[bst]<sz[nxt]) bst=nxt;
        sz[nod]+=sz[nxt];
    }

    lid[nod]=lid[bst];
    lant[lid[nod]].pb(nod);

    for (int nxt : g[nod])
        if (nxt!=f && nxt!=bst) pl[lid[nxt]]=nod;
}

void upd(int nod, int l, int r, int off, int p, int
    val){
    if (l==r){
        t[nod+off]=val;
        return ;
    }
}

```

```

    int mid=(l+r)/2;
    if (p<=mid) upd(nod*2,l,mid,off,p,val);
    else upd(nod*2+1,mid+1,r,off,p,val);
    t[nod+off]=max(t[nod*2+off],t[nod*2+1+off]);
}

void upd(int x, int val){
    upd(1,1,lant[lid[x]].size(),off[lid[x]],pos[x],val);
}

int query(int nod, int l, int r, int ql, int qr, int
    off){
    if (ql<=l && r<=qr) return t[nod+off];

    int mid=(l+r)/2,lv=0,rv=0;
    if (ql<=mid) lv=query(nod*2,l,mid,ql,qr,off);
    if (mid<qr) rv=query(nod*2+1,mid+1,r,ql,qr,off);

    return max(lv,rv);
}

void build(){
    int i,j;

    bdfs(1,0,1);
    int crsz=0;
    for (i=1; i<=L; i++){
        reverse(all(lant[i]));
        off[i]=crsz;
        for (j=1; j<=(int)lant[i].size(); j++){
            pos[lant[i][j-1]]=j;
            upd(lant[i][j-1],a[lant[i][j-1]]);
        }

        crsz+=4*lant[i].size();
    }
}

int query(int x, int y){
    int res=0;
    while (x!=-1){
        if (d[x]>d[y]) swap(x,y);

        if (lid[x]==lid[y]){
            res=max(res,query(1,1,lant[lid[x]].size(),pos[x],
                pos[y],off[lid[x]]));
            x=-1;
        }
        else {
            if (d[pl[lid[x]]]<d[pl[lid[y]]]) swap(x,y);

            res=max(res,query(1,1,lant[lid[x]].size(),1,pos[x],
                off[lid[x]]));
            x=pl[lid[x]];
        }
    }

    return res;
}

```

3.6 Max Bipartite matching - $O(N\sqrt{N})$

```

int N,M,E,l[10100],r[10100]; vector<int> g[10100];
bool v[10100];

bool try_match(int nod){
    if (v[nod]) return 0;
    v[nod]=1;

    int i,nd;
    for (i=0; i<g[nod].size(); i++){
        nd=g[nod][i];
        if (!r[nd] || try_match(r[nd])){
            l[nod]=nd;
            r[nd]=nod;
            return 1;
        }
    }
}

```

```

    }
}

return 0;
}

//inside code
bool ex=1;
while (ex){
    ex=0;
    memset(v,0,sizeof(v));
    for (i=1; i<=N; i++)
        if (!l[i])
            ex|=try_match(i);
}

```

3.7 Centroid Decomposition

```

int N,K,sz[100010],cnt[100010]; vi g[100010];
bool v[100010]; ll res;

void dfs_sz(int nod, int f){
    sz[nod]=1;
    for (int nxt : g[nod]){
        if (v[nxt] || f==nxt) continue;
        dfs_sz(nxt,nod);
        sz[nod]+=sz[nxt];
    }
}

void dfs(int nod, int f, int d, bool add){
    if (d>K) return;

    if (add) cnt[d]++;
    if (!add) res+=cnt[K-d];
    for (int nxt : g[nod])
        if (!v[nxt] && nxt!=f) dfs(nxt,nod,d+1,add);
}

void calc(int nod){
    int i;

    cnt[0]=1;
    for (int nxt : g[nod])
        if (!v[nxt]){
            dfs(nxt,nod,1,0);
            dfs(nxt,nod,1,1);
        }

    for (i=0; i<=sz[nod]; i++)
        cnt[i]=0;
}

void decomp(int nod){
    dfs_sz(nod,-1);

    int tot=sz[nod];
    bool f=0;
    while (!f){
        f=1;

        for (int nxt : g[nod])
            if (!v[nxt] && 2*sz[nxt]>tot){
                f=0;
                sz[nod]-=sz[nxt];
                sz[nxt]+=sz[nod];
                nod=nxt;
                break;
            }
    }

    calc(nod);
    v[nod]=1;

    for (int nxt : g[nod])
        if (!v[nxt])
            decomp(nxt);
}

```

3.8 Euler tour

```

void euler(){
    int nod;
    T=1,st[1]=1;

    while (T){
        nod=st[T];
        while (cr[nod]<g[nod].size())
            if (!ve[g[nod]][cr[nod]].ind) break;
            else cr[nod]++;

        if (cr[nod]<g[nod].size()){
            ve[g[nod]][cr[nod]].ind=1;
            st[++T]=g[nod][cr[nod]].d;
        }
        else{
            ans[++K]=nod;
            T--;
        }
    }
}

```

3.9 Biconnected components

```

void cache_bc(const int x, const int y)
{
    vector <int> con; int tx, ty;
    do {
        tx = stk.top().first, ty = stk.top().second;
        stk.pop();
        con.push_back(tx), con.push_back(ty);
    }
    while (tx != x || ty != y);
    C.push_back(con);
}

void DF(const int n, const int fn, int number)
{
    vector <int>::iterator it;

    dfn[n] = low[n] = number;
    for (it = adj[n].begin(); it != adj[n].end(); ++
it) {
        if (*it == fn) continue;
        if (dfn[*it] == -1) {
            stk.push( make_pair(n, *it) );
            DF(*it, n, number + 1);
            low[n] = Min(low[n], low[*it]);
            if (low[*it] >= dfn[n])
                cache_bc(n, *it);
        }
        else
            low[n] = Min(low[n], dfn[*it]);
    }
}

int main(void)
{
    int n;
    read_in(adj, n);
    dfn.resize(n + 1), dfn.assign(n + 1, -1);
    low.resize(n + 1);
    DF(1, 0, 0);

    ofstream out(oname);
    out << C.size() << "\n";
    for (size_t i = 0; i < C.size(); ++ i) {
        sort(C[i].begin(), C[i].end());
        C[i].erase(unique(C[i].begin(), C[i].end()), C[i]
].end());
        for (size_t j = 0; j < C[i].size(); ++ j)
            out << C[i][j] << " ";
        out << "\n";
    }
}

```


4 Data structures

4.1 Treap

```

struct Treap{
    ll x,p,sum,sumi;
    int sz;
    Treap *l,*r;

    Treap(ll x, ll p, Treap* l, Treap* r, int sz);
    void upd();

} *nil=new Treap(0,0,nullptr,nullptr,0);

typedef Treap* tp;
tp root = nil;

void Treap::upd(){
    if (this==nil) return;

    sum=x+l->sum+r->sum;
    sz=1+l->sz+r->sz;
    sumi=l->sumi+1LL*(l->sz+1)*x+r->sumi+(l->sz+1)*r->sum;
}

Treap::Treap(ll x, ll p, tp l, tp r, int sz=1){
    this->x = x, this->p = p;
    this->sum=this->sumi=x;
    this->l = l, this->r = r;
    this->sz=sz;
}

void split(tp root, int x, tp &L, tp &R){
    if (root==nil){
        L=R=nil;
        return;
    }

    root->upd();

    if (root->x<=x){
        split(root->r,x,root->r,R);
        L=root;
        L->upd();
    }
    else {
        split(root->l,x,L,root->l);
        R=root;
        R->upd();
    }
}

tp merge(tp l, tp r){
    if (l==nil || r==nil)
        return (l!=nil ? l : r);

    if (l->p>r->p){
        l->r=merge(l->r,r);
        l->upd();
        return l;
    }
    else {
        r->l=merge(l,r->l);
        r->upd();
        return r;
    }
}

void insert(tp &root, tp nod){
    if (root==nil){
        root = nod;
        nod->upd();
        return;
    }
}

```

```

    if (root->p<nod->p){
        split(root,nod->x,nod->l,nod->r);
        root=nod;
    }
    else if (root->x>nod->x) insert(root->l,nod);
    else insert(root->r,nod);

    root->upd();
}

void del(tp &root, ll x){
    if (root==nil) return;

    if (root->x==x){
        tp cr = root;
        root = merge(root->l, root->r);

        delete cr;
        return;
    }
    else if (root->x>x) del(root->l,x);
    else del(root->r,x);

    root->upd();
}

```

4.2 Persistent Segment Tree

```

struct tree{
    int l,r,s;
    tree (int l=0, int r=0, int s=0) : l(l), r(r), s(s) {}
};

int N,Q,a[300010],T,rt[300010]; tree t[7000000];

int upd(int nod, int l, int r, int p){
    tree &cr=t[++T];
    cr=t[nod];

    if (l==r){
        cr.s++;
        cr.l=cr.r=0;
        return T;
    }

    int mid=(l+r)/2,id=T;
    if (p<=mid) cr.l=upd(t[nod].l,l,mid,p);
    else cr.r=upd(t[nod].r,mid+1,r,p);
    cr.s=t[cr.l].s+t[cr.r].s;
    return id;
}

int query(int nod, int l, int r, int ql, int qr){
    if (nod==0) return 0;
    if (ql<=l && r<=qr) return t[nod].s;

    int mid=(l+r)/2,res=0;
    if (ql<=mid) res+=query(t[nod].l,l,mid,ql,qr);
    if (mid<qr) res+=query(t[nod].r,mid+1,r,ql,qr);
    return res;
}

int find_kth(int p, int q, int l, int r, int k){
    if (l==r) return l;

    int mid=(l+r)/2;
    int totl=(t[p].l).s-t[q].l.s;
    if (totl>=k) return find_kth(t[p].l,t[q].l,l,mid,k);
    else return find_kth(t[p].r,t[q].r,mid+1,r,k-totl);
}

```

4.3 2D Segment Tree

Basically use treaps for the second dimension.

```

/* insert treap implementation here */

struct tree2d{
    tree2d *l, *r;
    Treap* root;

    tree2d(){
        root=nil;
    }
} *root;

ll getmxv(tp root, int x1, int x2, int l, int r, int k){
    ll res=-inf;
    if (root==nil || r<x1 || l>x2) return res;

    if (x1<=l && r<=x2) return root->mv[k];

    if (root->x<x1) res=getmxv(root->r,x1,x2,root->x+1,r,k);
    else if (root->x>x2) res=getmxv(root->l,x1,x2,l,root->x-1,k);
    else{
        res=root->v[k];

        ll lres=getmxv(root->l,x1,x2,l,root->x-1,k);
        ll rres=getmxv(root->r,x1,x2,root->x+1,r,k);

        res=max(res,max(lres,rres));
    }

    return res;
}

void upd2d(tree2d *&nod, int l, int r, int x, int y)
{
    if (nod==NULL) nod=new tree2d();

    //change this in case the value can be deleted
    ins(nod->root,new Treap(y,rnd()+1,nil,nil,x,y));

    if (l==r)
        return;

    int mid=(l+r)/2;
    if (x<=mid) upd2d(nod->l,l,mid,x,y);
    else upd2d(nod->r,mid+1,r,x,y);
}

ll query2d(tree2d *nod, int l, int r, int ql, int qr, int y1, int y2, int k){
    if (nod==NULL) return -inf;
    if (ql<=l && r<=qr) return getmxv(nod->root,y1,y2,0,C, k);

    int mid=(l+r)/2; ll res=-inf;
    if (ql<=mid) res=max(res,query2d(nod->l,l,mid,ql,qr,y1,y2,k));
    if (mid<qr) res=max(res,query2d(nod->r,mid+1,r,ql,qr,y1,y2,k));
    return res;
}

```

4.4 Merging Segment Trees

```

struct stree{
    stree *l,*r;
    int sz;

    stree();
} *nil = new stree();
typedef stree* st;

```

```

stree::stree(){
    l=r=nil;
    sz=0;
}

void insert(st &nod, int l, int r, int p){
    cnt++;
    if (nod==nil) nod=new stree();

    if (l!=r){
        int mid=(l+r)/2;
        if (p<=mid) insert(nod->l,l,mid,p);
        else insert(nod->r,mid+1,r,p);
    }

    nod->sz++;
}

st merge(st &a, st b){
    cnt++;
    if (a==nil || b==nil)
        return (a!=nil ? a : b);

    a->sz+=b->sz;
    a->l=merge(a->l,b->l);
    a->r=merge(a->r,b->r);

    delete b;
    return a;
}

void split(st &nod, int l, int r, int p, st &rn){
    cnt++;
    if (nod==nil || l==r){
        rn=nil;
        return;
    }

    int mid=(l+r)/2;
    rn=new stree();

    if (p<=mid){
        split(nod->l,l,mid,p,rn->l);
        rn->r=nod->r;
        nod->r=nil;
    }
    else
        split(nod->r,mid+1,r,p,rn->r);

    rn->sz=rn->l->sz+rn->r->sz;
    nod->sz-=rn->sz;
}

int query(st nod, int l, int r, int ql, int qr){
    cnt++;
    if (nod==nil) return 0;
    if (ql<=l && r<=qr) return nod->sz;

    int mid=(l+r)/2,res=0;
    if (ql<=mid) res+=query(nod->l,l,mid,ql,qr);
    if (mid<qr) res+=query(nod->r,mid+1,r,ql,qr);
    return res;
}

```

4.5 Binary Indexed Tree

```

int N,M,tree[100010];

int sum(int r){
    int res=0;
    while (r){
        res+=tree[r];
        r=(r&-r);
    }
    return res;
}

int get_sum(int l,int r){

```

```

    return sum(r)-sum(l-1);
}

void update(int ind, int val){
    while (ind<=N){
        tree[ind]+=val;
        ind+=(ind&-ind);
    }
}

int find_min(int sum){
    int lg=0,p=0;
    for (lg=0; (1<<(lg+1))<=N; lg++) ;

    for (; lg>=0; lg--){
        if (p+(1<<lg)<=N)
            if (sum>=tree[p+(1<<lg)]){
                p+=(1<<lg), sum-=tree[p];
                if (!sum) return p;
            }
    }
    return -1;
}

```

4.6 SQRT Decomposition

```

struct block{
    vector<ll> arr,nxt,nrpasi,fnxt,nxtb;
    int lzadd=0,id;

    int getp(int p){
        if (p==-1) return -1;
        return sz*id+p;
    }

    void init(){
        int sz = arr.size();
        //the elements should already be pushed to array
        //do necessary initialization here for any
        additional operations
    }

    void unlz(){
        if (lzadd==0) return;
        int i;
        for (i=0; i<(int)arr.size(); i++)
            arr[i]+=lzadd;
        lzadd=0;
    }

    //augment the structure with additional operations
    //call unlz() when needed
}

int bid(int x){
    return x/sz;
}

int bpos(int x){
    return x%sz;
}

cin >> N >> Q;

int i,j;
sz=102;
for (i=0; i<N; i++) cin >> a[i];

if (sz*sz<N){
    while (sz*sz<=N) sz++;
    sz--;
}

for (i=0; i<N; i+=sz){
    int ind = i/sz;
    for (j=i; j<min(i+sz,N); j++)
        b[ind].arr.pb(a[j]);
}

```

```

b[ind].init();
b[ind].id = ind;
b[ind].calc_nxt();
if (ind) b[ind-1].calc_nxtb(b[ind]);
}

//an example of how an update to a range should look
like
//NOTE: this is a rough sketch
//the upd function updates a range of elements inside
a block
int l,r; ll x;

cin >> l >> r >> x;
l--,r--;
int idl=bid(l),idr=bid(r);

if (idl==idr){
    b[idl].upd(l,r,x);
}
else {
    b[idl].upd(l,(idl+1)*sz-1,x);
    b[idr].upd(idr*sz,r,x);

    for (i=idl+1; i<idr; i++)
        b[i].lzadd+=x;
}

```

4.7 Mo's Algorithm

```

void mvr(int &r, int pos){
    while (r!=pos){
        if (r<pos){
            r++;
            cnte[p[r-1]]++;
            cnt[p[r]]++;
            res+=cnte[p2[r][0]];
        }
        else{
            res-=cnte[p2[r][0]];
            cnt[p[r]]--;
            cnte[p[r-1]]--;
            r--;
        }
    }
}

void mvl(int &l, int pos){
    while (l!=pos){
        if (l>pos){
            l--;
            cnte[p[l-1]]++;
            cnt[p[l]]++;
            res+=cnt[p2[l-1][1]];
        }
        else{
            res-=cnt[p2[l-1][1]];
            cnte[p[l-1]]--;
            cnt[p[l]]--;
            l++;
        }
    }
}

//main code

while (sr*sr<=N)
    sr++;
sr--;
cin >> Q;
for (i=1; i<=Q; i++){
    cin >> q[i].l >> q[i].r;
    q[i].id=i;
}

```

```

sort(q+1,q+Q+1,[sr](query A, query B) { return (A.l/
    sr==B.l/sr ? A.r<B.r : A.l/sr<B.l/sr);});

int l=q[1].l,r=q[1].r;
for (i=l; i<=r; i++){
    cnte[p[i-1]]++;
    cnt[p[i]]++;
    res+=cnte[p2[i][0]];
}

for (i=1; i<=Q; i++){
    if (r<q[i].r) mvr(r,q[i].r),mvl(1,q[i].l);
    else mvl(1,q[i].l),mvr(r,q[i].r);

    ans[q[i].id]=res;
}

```

5 Transformations

5.1 Fast Fourier Transform - $O(n \log n)$

```

const ld pi = acos(-1);
typedef complex<double> base;

int rev(int x, int lg){
    int res=0;
    for (int i=0; i<lg; i++)
        if (x&(1<<i)) res+=(1<<(lg-1-i));
    return res;
}

void fft(vector<base> &a, bool inv){
    int lg=1,sz=a.size(),i,j;
    while ((1<<lg)<sz) lg++;

    for (i=0; i<sz; i++){
        if (i<rev(i,lg)) swap(a[i],a[rev(i,lg)]);

        for (int len=2; len<=sz; len<=1){
            ld ang = 2*pi/len * (inv ? -1 : 1);
            base wlen(cos(ang),sin(ang));

            for (i=0; i<sz; i+=len){
                base w(1,0);
                for (j=0; j<len/2; j++){
                    base u=a[i+j],v=w*a[i+len/2+j];
                    a[i+j]=u+v;
                    a[i+len/2+j]=u-v;

                    w*=wlen;
                }
            }
        }

        if (inv)
            for (i=0; i<sz; i++)
                a[i]/=sz;
    }

    vi conv(vi &a, vi &b){
        vector<base> na,nb,nc;
        for (int x : a)
            na.pb(base(x));
        for (int x : b)
            nb.pb(base(x));

        int sz=2*max(a.size(),b.size());
        int lg=1;
        while ((1<<lg)<sz) lg++;
        sz=1<<lg;

        na.resize(sz),nb.resize(sz),nc.resize(sz);
        fft(na,0),fft(nb,0);

        for (int i=0; i<sz; i++)
            nc[i]=na[i]*nb[i];
        fft(nc,1);
    }
}

```

```

vi c(a.size()+b.size()-1);
for (int i=0; i<(int)c.size(); i++)
    c[i]=(int)(nc[i].real()+0.5); //be careful

return c;
}

```

5.2 FFT with big modulo

```

namespace fft {
typedef double dbl;

struct num {
    dbl x, y;
    num() { x = y = 0; }
    num(dbl x, dbl y) : x(x), y(y) { }
};

inline num operator+(num a, num b) { return num(a.x +
    b.x, a.y + b.y); }
inline num operator-(num a, num b) { return num(a.x -
    b.x, a.y - b.y); }
inline num operator*(num a, num b) { return num(a.x *
    b.x - a.y * b.y, a.x * b.y + a.y * b.x); }
inline num conj(num a) { return num(a.x, -a.y); }

int base = 1;
vector<num> roots = {{0, 0}, {1, 0}};
vector<int> rev = {0, 1};

const dbl PI = acosl(-1.0);

void ensure_base(int nbase) {
    if (nbase <= base) {
        return;
    }
    rev.resize(1 << nbase);
    for (int i = 0; i < (1 << nbase); i++) {
        rev[i] = (rev[i >> 1] >> 1) + ((i & 1) << (nbase -
            1));
    }
    roots.resize(1 << nbase);
    while (base < nbase) {
        dbl angle = 2 * PI / (1 << (base + 1));
        // num z(cos(angle), sin(angle));
        for (int i = 1 << (base - 1); i < (1 << base); i++)
            {
                roots[i << 1] = roots[i];
                // roots[(i << 1) + 1] = roots[i] * z;
                dbl angle_i = angle * (2 * i + 1 - (1 << base));
                roots[(i << 1) + 1] = num(cos(angle_i), sin(
                    angle_i));
            }
        base++;
    }
}

void fft(vector<num> &a, int n = -1) {
    if (n == -1) {
        n = a.size();
    }
    assert((n & (n - 1)) == 0);
    int zeros = __builtin_ctz(n);
    ensure_base(zeros);
    int shift = base - zeros;
    for (int i = 0; i < n; i++) {
        if (i < (rev[i] >> shift)) {
            swap(a[i], a[rev[i] >> shift]);
        }
    }
    for (int k = 1; k < n; k <= 1) {
        for (int i = 0; i < n; i <= 2 * k) {
            for (int j = 0; j < k; j++) {
                num z = a[i + j + k] * roots[j + k];
                a[i + j + k] = a[i + j] - z;
                a[i + j] = a[i + j] + z;
            }
        }
    }
}

```

```

    }
}

vector<num> fa, fb;

vector<int> multiply(vector<int> &a, vector<int> &b)
{
    int need = a.size() + b.size() - 1;
    int nbase = 0;
    while ((1 << nbase) < need) nbase++;
    ensure_base(nbase);
    int sz = 1 << nbase;
    if (sz > (int) fa.size()) {
        fa.resize(sz);
    }
    for (int i = 0; i < sz; i++) {
        int x = (i < (int) a.size() ? a[i] : 0);
        int y = (i < (int) b.size() ? b[i] : 0);
        fa[i] = num(x, y);
    }
    fft(fa, sz);
    num r(0, -0.25 / sz);
    for (int i = 0; i <= (sz >> 1); i++) {
        int j = (sz - i) & (sz - 1);
        num z = (fa[j] * fa[j] - conj(fa[i] * fa[i])) * r;
        if (i != j) {
            fa[j] = (fa[i] * fa[i] - conj(fa[j] * fa[j])) * r;
        }
        fa[i] = z;
    }
    fft(fa, sz);
    vector<int> res(need);
    for (int i = 0; i < need; i++) {
        res[i] = fa[i].x + 0.5;
    }
    return res;
}

vector<int> multiply_mod(vector<int> &a, vector<int>
    &b, int m, int eq = 0) {
    int need = a.size() + b.size() - 1;
    int nbase = 0;
    while ((1 << nbase) < need) nbase++;
    ensure_base(nbase);
    int sz = 1 << nbase;
    if (sz > (int) fa.size()) {
        fa.resize(sz);
    }
    for (int i = 0; i < (int) a.size(); i++) {
        int x = (a[i] % m + m) % m;
        fa[i] = num(x & ((1 << 15) - 1), x >> 15);
    }
    fill(fa.begin() + a.size(), fa.begin() + sz, num {0,
        0});
    fft(fa, sz);
    if (sz > (int) fb.size()) {
        fb.resize(sz);
    }
    if (eq) {
        copy(fa.begin(), fa.begin() + sz, fb.begin());
    } else {
        for (int i = 0; i < (int) b.size(); i++) {
            int x = (b[i] % m + m) % m;
            fb[i] = num(x & ((1 << 15) - 1), x >> 15);
        }
        fill(fb.begin() + b.size(), fb.begin() + sz, num
            {0, 0});
        fft(fb, sz);
    }
    dbl ratio = 0.25 / sz;
    num r2(0, -1);
    num r3(ratio, 0);
    num r4(0, -ratio);
    num r5(0, 1);
    for (int i = 0; i <= (sz >> 1); i++) {
        int j = (sz - i) & (sz - 1);
        num a1 = (fa[i] + conj(fa[j]));

```

```

        num a2 = (fa[i] - conj(fa[j])) * r2;
        num b1 = (fb[i] + conj(fb[j])) * r3;
        num b2 = (fb[i] - conj(fb[j])) * r4;
        if (i != j) {
            num c1 = (fa[j] + conj(fa[i]));
            num c2 = (fa[j] - conj(fa[i])) * r2;
            num d1 = (fb[j] + conj(fb[i])) * r3;
            num d2 = (fb[j] - conj(fb[i])) * r4;
            fa[i] = c1 * d1 + c2 * d2 * r5;
            fb[i] = c1 * d2 + c2 * d1;
        }
        fa[j] = a1 * b1 + a2 * b2 * r5;
        fb[j] = a1 * b2 + a2 * b1;
    }
    fft(fa, sz);
    fft(fb, sz);
    vector<int> res(need);
    for (int i = 0; i < need; i++) {
        long long aa = fa[i].x + 0.5;
        long long bb = fb[i].x + 0.5;
        long long cc = fa[i].y + 0.5;
        res[i] = (aa + ((bb % m) << 15) + ((cc % m) << 30))
            % m;
    }
    return res;
}

vector<int> square_mod(vector<int> &a, int m) {
    return multiply_mod(a, a, m, 1);
}
};

```

5.3 Number Theoretic Transform

```

const int mod = 7340033;
const int root = 5;
const int root_1 = 4404020;
const int root_pw = 1<<20;

void fft (vector<int> &a, bool invert) {
    int n = (int) a.size();

    for (int i=1, j=0; i<n; ++i) {
        int bit = n >> 1;
        for (; j>=bit; bit>>=1)
            j -= bit;
        j += bit;
        if (i < j)
            swap (a[i], a[j]);
    }

    for (int len=2; len<=n; len<=1) {
        int wlen = invert ? root_1 : root;
        for (int i=len; i<root_pw; i<=1)
            wlen = int (wlen * 111 * wlen % mod);
        for (int i=0; i<n; i+=len) {
            int w = 1;
            for (int j=0; j<len/2; ++j) {
                int u = a[i+j], v = int (a[i+j+len/2] * 111
                    * w % mod);
                a[i+j] = u+v < mod ? u+v : u+v-mod;
                a[i+j+len/2] = u-v >= 0 ? u-v : u-v+mod;
                w = int (w * 111 * wlen % mod);
            }
        }
    }
    if (invert) {
        int nrev = reverse (n, mod);
        for (int i=0; i<n; ++i)
            a[i] = int (a[i] * 111 * nrev % mod);
    }
}

```

5.4 XOR convolutin (\oplus) - $O(n \log n)$

```

const int lg = 16, sz=(1<<16), mod=30011, i2=15006;
vi a(sz);

```

```

vi conv(const vi &a, const vi &b){
    int len,i,k;
    vi c;
    c.insert(c.end(),all(a));
    c.insert(c.end(),all(b));

    for (len=sz; len>1; len>>=1){
        for (i=0; i<2*sz; i+=2*len){
            for (k=0; k<len/2; k++){
                int a=c[i+k],na=c[i+len/2+k];
                int b=c[i+len+k],nb=c[i+3*len/2+k];

                c[i+k]=(a+na)%mod;
                c[i+len/2+k]=(b+nb)%mod;
                c[i+len+k]=(a-na+mod)%mod,c[i+3*len/2+k]=(b-
                    nb+mod)%mod;

                // for OR (a+na), a
                // for AND (a+na),na
            }
        }
    }

    for (i=0; i<2*sz; i+=2)
        c[i]=(c[i]*c[i+1])%mod;

    for (len=2; len<=sz; len<<=1){
        for (i=0; i<2*sz; i+=2*len){
            for (k=0; k<len/2; k++){
                int a=c[i+k],b=c[i+len+k];
                c[i+k]=(a+b)*i2%mod;
                c[i+len/2+k]=(a-b+mod)*i2%mod;

                //for OR b, a-b
                //for AND a-b, b
            }
        }
    }

    c.resize(sz);
    return c;
}

```

6 String Algorithms

6.1 Suffix automaton - $O(n)$

Every state represents an equivalence class of substrings. To get the number of different substrings, compute the sum of differences between the len of node, and node.fail.

```

struct state {
    int len, link;
    map<char,int> next;
};

const int MAXLEN = 100000;
state st[MAXLEN*2];
int sz, last;

void sa_init() {
    sz = last = 0;
    st[0].len = 0;
    st[0].link = -1;
    ++sz;
    //clear next in case of multiple automaton
}

void sa_extend (char c) {
    int cur = sz++;
    st[cur].len = st[last].len + 1;
    int p;
    for (p=last; p!=-1 && !st[p].next.count(c); p=st[p].link)
        st[p].next[c] = cur;

```

```

    if (p == -1)
        st[cur].link = 0;

    else {
        int q = st[p].next[c];
        if (st[p].len + 1 == st[q].len)
            st[cur].link = q;

        else {
            int clone = sz++;
            st[clone].len = st[p].len + 1;
            st[clone].next = st[q].next;
            st[clone].link = st[q].link;

            for (; p!=-1 && st[p].next[c]==q; p=st[p].link)
                st[p].next[c] = clone;

            st[q].link = st[cur].link = clone;
        }
    }
    last = cur;
}

```

6.2 Aho-Corasick

```

int to[1000010][26],K=1,cnt[1000010],T,fail[1000010],pnod[120],root=1;
int q[1000010];

int ins(int &nod, int p, string &s){
    if (nod==0) nod=++K;

    if (p==s.length())
        return nod;

    return ins(to[nod][s[p]-'a'],p+1,s);
}

void calc_fail(){
    T=1,q[1]=root;
    fail[root]=root;

    int i,j;
    for (i=1; i<=T; i++){
        int cr=q[i];

        for (j=0; j<26; j++){
            if (to[cr][j]!=0){
                int nxt=to[cr][j];

                int f=fail[cr];
                while (f!=1 && to[f][j]==0)
                    f=fail[f];

                fail[nxt]=f;
                if (to[f][j]!=0 && to[f][j]!=nxt)
                    fail[nxt]=to[f][j];

                q[++T]=nxt;
            }
        }
    }

    void prop(){
        for (int i=K; i>1; i--){
            cnt[fail[q[i]]]+=cnt[q[i]];
        }
    }
}

```

6.3 Suffix Array

```
const int MAXN = 65536;
const int MAXLG = 17;

char A[MAXN];
struct entry {
    int nr[2], p;
} L[MAXN];
int P[MAXLG][MAXN], N, i, stp, cnt;

bool cmp(const entry &a, const entry &b) {
    return a.nr[0] == b.nr[0] ? (a.nr[1] < b.nr[1]) :
        (a.nr[0] < b.nr[0]);
}

int main() {
    gets(A);
    for (N = strlen(A), i = 0; i < N; ++i)
        P[0][i] = A[i] - 'a';
    for (stp = 1, cnt = 1; cnt >> 1 < N; ++stp, cnt
        <= 1) {
        for (i = 0; i < N; ++i) {
            L[i].nr[0] = P[stp - 1][i];
            L[i].nr[1] = i + cnt < N ? P[stp - 1][i + cnt]
                : -1;
            L[i].p = i;
        }
        sort(L, L + N, cmp);
        for (i = 0; i < N; ++i)
            P[stp][L[i].p] = i > 0 && L[i].nr[0] == L[i - 1].nr[0]
                && L[i].nr[1] == L[i - 1].nr[1] ? P[stp - 1][L[i - 1].p] : i;
    }
    return 0;
}
```

6.4 KMP

Find all occurrences of A in B :

```
int i, q=0;
for (i=2; i<=N; i++){
    while (q && A[q+1]!=A[i])
        q--;

    if (A[q+1]==A[i]) q++;
    P[i]=q;
}

q=0;
for (i=1; i<=M; i++){
    while (q && A[q+1]!=B[i])
        q=P[q];

    if (A[q+1]==B[i]) q++;

    if (q==N){
        K++;
        q=P[N];
        if (K<=1000)
            pos[K]=i-N;
    }
}
```

6.5 Manacher's Algorithm

```
void get_pal(string &s, int *q){
    int i, N=s.length(), hr=-1, l, p;

    for (i=0; i<N; i++){
        l=0;
        if (hr<i) l=0, hr=i;
        else l=q[2*p-i];
```

```
        if (i+l<hr){
            q[i]=l;
            continue;
        }
        else l=hr-i;

        while (i+l<N && i-l>=0 && s[i+l]==s[i-l])
            l++;

        l--;
        q[i]=l;
        hr=i+l;
        p=i;
    }
}
```

7 Geometry

7.1 Common structures

```
struct point{
    ld x, y;
    point (ld x=0, ld y=0) : x(x), y(y) {}
};

struct line{
    ld A, B, C;
    line (ld A=0, ld B=0, ld C=0) : A(A), B(B), C(C) {}
};

ld norm(point P){
    return sqrt(P.x*P.x+P.y*P.y);
}
```

7.2 Angle between vectors

$$\cos \alpha = \frac{u \cdot v}{|u||v|}$$

7.3 Line from 2 points

```
line getline(point u, point v){
    line r;
    r.A=u.y-v.y;
    r.B=v.x-u.x;
    r.C=u.x*(v.y-u.y)-u.y*(v.x-u.x);
    return r;
}
```

7.4 Intersection of 2 lines

```
point intersect(line l1, line l2){
    if (l1.A*l2.B==l2.A*l1.B){
        //parallel, or completely intersecting
        return {0,0};
    }
    else{
        ld x=(-l1.C*l2.B+l1.B*l2.C)/(l1.A*l2.B-l1.B*l2.A);
        ld y=(-l1.A*l2.C+l1.C*l2.A)/(l1.A*l2.B-l1.B*l2.A);
        return point(x,y);
    }
}
```

7.5 Distance from point to line

$$d = \frac{|Ax + By + C|}{\sqrt{A^2 + B^2}}$$

If line is given as a pair of points, or distance to ray/segment is required, use:

```
ld dist(point A, point B, point X){
    B.x-=A.x, B.y-=A.y;
    X.x-=A.x, X.y-=A.y;

    ld t=(B.x*X.x+B.y*X.y)/(B.x*B.x+B.y*B.y);
```

```
//in case of line/ray intersection put conditions on
t
point P(t*B.x,t*B.y);
return sqrt(pow(P.x-X.x,2)+pow(P.y-X.y,2));
}
```

7.6 Bisector

```
point bisector(point X, point A, point B){
    A.x-=X.x,A.y-=X.y;
    B.x-=X.x,B.y-=X.y;

    ld v;
    v=norm(A);
    A.x/=v,A.y/=v;

    v=norm(B);
    B.x/=v,B.y/=v;

    if (abs(A.x*B.x+A.y*B.y+1)<eps) return point(X.x-A.y
        ,X.y+A.x);
    //returns vector
    return point(X.x+(A.x+B.x)/2,X.y+(A.y+B.y)/2);
}
```

7.7 Convex hull

```
int N,K,st[120100]; point a[120100];

double area(point A, point B, point C){
    return (A.x*B.y+B.x*C.y+C.x*A.y-B.y*C.x-C.y*A.x-A.y*
        B.x);
}

double sdist(point A, point B){
    return (A.x-B.x)*(A.x-B.x)+(A.y-B.y)*(A.y-B.y);
}

inline bool cmp(point A, point B){
    return (area(a[1],A,B)==0 ? sdist(a[1],A)<sdist(a
        [1],B) : area(a[1],A,B)>0);
}

int i,ind=1;
for (i=1; i<=N; i++){
    scanf("%lf %lf",&a[i].x,&a[i].y);
    if (a[i].x<a[ind].x) ind=i;
}

swap(a[1],a[ind]);
sort(a+2,a+N+1,cmp);
a[N+1]=a[1];

K=1,st[1]=1;
for (i=2; i<=N; i++){
    st[++K]=i;
    while (area(a[st[K-1]],a[st[K]],a[i+1])<0) K--;
}
```

7.8 Testing if point is inside convex polygon - $O(\log n)$

```
bool is_between(point A, point B, point X){
    return (min(A.x,B.x)<=X.x && X.x<=max(A.x,B.x) &&
        min(A.y,B.y)<=X.y && X.y<=max(A.y,B.y));
}

bool inside_triangle(point A, point B, point C, point
    X){
    if (sign(ccw(X,A,B))*sign(ccw(X,B,C))>=0 && sign(ccw
        (X,B,C))*sign(ccw(X,C,A))>=0) return 1;
    return 0;
}
```

```
}

bool is_inside(point &A, vector<point> &poly){
    int l=1,r=poly.size()-1,mid;
    if (ccw(poly[0],poly[1],A)<0) return 0;
    while (l<r){
        mid=(l+r+1)/2;
        if (ccw(poly[0],poly[mid],A)>=0) l=mid;
        else r=mid-1;
    }

    if (r==(int)poly.size()-1) return (ccw(poly[0],poly[
        r],A)==0 && is_between(poly[0],poly[r],A));
    return inside_triangle(poly[0],poly[r],poly[r+1],A);
}
```

7.9 Signed area of a polygon

$$A = \frac{1}{2} \sum_{i=1}^n (x_i \cdot y_{i+1} - x_{i+1} \cdot y_i)$$

7.10 Closest pair of points

```
int N; point a[100100];
point v[100100]; int K;

double mind(int l, int r){
    if (r-l+1<=1) return (1LL<<30);
    if (r-l+1==2){
        if (a[l].y>a[r].y) swap(a[l],a[r]);
        return dist(a[l],a[r]);
    }

    double ans=(1LL<<30);
    int mid=(l+r)/2,mx=a[mid].x;
    ans=min(ans,mind(l,mid));
    ans=min(ans,mind(mid+1,r));

    int i,j;
    merge(a+l,a+mid+1,a+mid+1,a+r+1,v,[](point A,
        point B){return A.y<B.y; });
    memcpy(a+l,v,(r-l+1)*sizeof(point));

    vector<point> med;
    for (i=1; i<=r; i++){
        if (abs(a[i].x-mx)<=ans)
            med.pb(a[i]);

    for (i=0; i<(int)med.size(); i++){
        for (j=i+1; j<(int)med.size() && med[j].y-med[i
            ].y<=ans; j++){
            ans=min(ans,dist(med[i],med[j]));
        }

    return ans;
}
```

8 Misc

8.1 2-SAT

```
const int MAXN=100010;
int N,M,st[2*MAXN+10],K;
vector<int> g[2*MAXN+10],gc[2*MAXN+10];
int comp[2*MAXN+10],nrC;

int id(int x){
    if (x>0) return 2*x-1;
    return 2*(-x);
}
```



```

void dfs1(int nod){
    comp[nod]=1;
    for (int nxt : g[nod])
        if (!comp[nxt]) dfs1(nxt);
    st[++K]=nod;
}

void dfs2(int nod){
    comp[nod]=nrC;
    for (int nxt : gc[nod])
        if (!comp[nxt]) dfs2(nxt);
}

int main(){
    // N - number of variables, M-clauses

    cin >> N >> M;

    int i,x,y;
    for (i=1; i<=M; i++){
        fcn >> x >> y;
        g[id(-x)].push_back(id(y));
        g[id(-y)].push_back(id(x));

        gc[id(y)].push_back(id(-x));
        gc[id(x)].push_back(id(-y));
    }

    for (i=1; i<=2*N; i++)
        if (!comp[i]) dfs1(i);

    memset(comp,0,sizeof(comp));
    for (i=2*N; i>0; i--)
        if (!comp[st[i]]) nrC++,dfs2(st[i]);

    for (i=1; i<=N; i++)
        if (comp[id(i)]==comp[id(-i)]){
            fout << -1 << "\n";
            return 0;
        }

    for (i=1; i<=N; i++)
        fout << (comp[id(i)]>comp[id(-i)]) << " ";
    fout << "\n";
    return 0;
}

```

8.2 Custom set/multiset comparator

```

struct lex_compare {
    bool operator() (const int64_t& lhs, const int64_t
        & rhs) const {
        stringstream s1, s2;
        s1 << lhs;
        s2 << rhs;
        return s1.str() < s2.str();
    }
};

//use like this
set<int64_t, lex_compare> s;

```

8.3 Parsing expressions

```

const long MAX = 100010;
char S[MAX], *p=S;

long termen();
long factor();

long eval() {
    long r = termen();
    while ( *p=='+' || *p=='-' ) {
        switch ( *p ) {
            case '+':
                ++p;

```

```

                r += termen();
                break;
            case '-':
                ++p;
                r -= termen();
                break;
        }
    }
    return r;
}

long termen() {
    long r = factor();
    while ( *p=='*' || *p=='/' ) {
        switch ( *p ) {
            case '*':
                ++p;
                r *= factor();
                break;
            case '/':
                ++p;
                r /= factor();
                break;
        }
    }
    return r;
}

long factor() {
    long r=0;
    if ( *p == '(' ) {
        ++p;
        r = eval();
        ++p;
    } else {
        while ( *p>='0' && *p<='9' ) {
            r = r*10 + *p - '0';
            ++p;
        }
    }
    return r;
}

```