

Elapsed Time: 0.045s

Application execution time is too short. Metrics data may be unreliable. Consider reducing the sampling interval or increasing your application execution time.

Clockticks: 153,360,000
Instructions Retired: 140,220,000
CPI Rate: 1.094

The CPI may be too high. This could be caused by issues such as memory stalls, instruction starvation, branch misprediction or long latency instructions. Explore the other hardware-related metrics to identify what is causing high CPI.

MUX Reliability: 0.986
Retiring: 24.9% of Pipeline Slots
 Light Operations: 26.9% of Pipeline Slots
 FP Arithmetic: 0.0% of uOps
 FP x87: 0.0% of uOps
 FP Scalar: 0.0% of uOps
 FP Vector: 0.0% of uOps
 Other: 100.0% of uOps
 Heavy Operations: 0.0% of Pipeline Slots
 Microcode Sequencer: 7.4% of Pipeline Slots
 Assists: 0.0% of Pipeline Slots
Front-End Bound: 21.7% of Pipeline Slots

Issue: A significant portion of Pipeline Slots is remaining empty due to issues in the Front-End.

Tips: Make sure the code working size is not too large, the code layout does not require too many memory accesses per cycle to get enough instructions for filling four pipeline slots, or check for microcode assists.

Front-End Latency: 24.9% of Pipeline Slots

This metric represents a fraction of slots during which CPU was stalled due to front-end latency issues, such as instruction-cache misses, ITLB misses or fetch stalls after a branch misprediction. In such cases, the front-end delivers no uOps.

ICache Misses:	0.0% of Clockticks
ITLB Overhead:	1.1% of Clockticks
Branch Resteers:	0.0% of Clockticks
Mispredicts Resteers:	0.0% of Clockticks
Clears Resteers:	0.0% of Clockticks
Unknown Branches:	0.0% of Clockticks
DSB Switches:	0.0% of Clockticks
Length Changing Prefixes:	0.0% of Clockticks
MS Switches:	0.0% of Clockticks

Issue: A significant fraction of cycles was stalled due to switches of uOp delivery to the Microcode Sequencer (MS). Commonly used instructions are optimized for delivery by the DSB or MITE pipelines. Certain operations cannot be handled natively by the execution pipeline, and must be performed by microcode (small programs injected into the execution stream). Switching to the MS too often can negatively impact performance. The MS is designated to deliver long uOp flows required by CISC instructions like CPUID, or uncommon conditions like Floating Point Assists when dealing with Denormals. Note that this metric value may be highlighted due to Microcode Sequencer issue.

Front-End Bandwidth:	0.0% of Pipeline Slots
Front-End Bandwidth MITE:	20.7% of Clockticks
Front-End Bandwidth DSB:	4.1% of Clockticks
(Info) DSB Coverage:	31.6%
Bad Speculation:	3.1% of Pipeline Slots
Branch Mispredict:	0.0% of Pipeline Slots
Machine Clears:	3.1% of Pipeline Slots
Back-End Bound:	50.3% of Pipeline Slots

A significant portion of pipeline slots are remaining empty. When operations take too long in the back-end, they introduce bubbles in the pipeline that ultimately cause fewer pipeline slots containing useful work to be retired per cycle than the machine is capable to support. This opportunity cost results in slower execution. Long-latency operations like divides and memory operations can cause this, as can too many operations being directed to a single execution port (for example, more multiply operations arriving in the back-end per cycle than the execution unit can support).

Memory Bound:	33.5% of Pipeline Slots
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The metric value is high. This can indicate that the significant fraction of execution pipeline slots could be stalled due to demand memory load and stores. Use Memory Access analysis to have the metric breakdown by memory hierarchy, memory bandwidth information, correlation by memory objects.

L1 Bound:

14.1% of Clockticks

This metric shows how often machine was stalled without missing the L1 data cache. The L1 cache typically has the shortest latency. However, in certain cases like loads blocked on older stores, a load might suffer a high latency even though it is being satisfied by the L1. Note that this metric value may be highlighted due to DTLB Overhead or Cycles of 1 Port Utilized issues.

DTLB Overhead:

0.7% of Clockticks

Load STLB Hit:

0.0% of Clockticks

Load STLB Miss:

0.7% of Clockticks

Loads Blocked by Store Forwarding:

0.0% of Clockticks

Lock Latency:

0.0% of Clockticks

A significant fraction of CPU cycles spent handling cache misses due to lock operations. Due to the microarchitecture handling of locks, they are classified as L1 Bound regardless of what memory source satisfied them. Note that this metric value may be highlighted due to Store Latency issue.

Split Loads:

0.0% of Clockticks

4K Aliasing:

0.0% of Clockticks

FB Full:

0.0% of Clockticks

This metric does a rough estimation of how often L1D Fill Buffer unavailability limited additional L1D miss memory access requests to proceed. The higher the metric value, the deeper the memory hierarchy level the misses are satisfied from. Often it hints on approaching bandwidth limits (to L2 cache, L3 cache or external memory). Avoid adding software prefetches if indeed memory BW limited.

L2 Bound:	0.0% of Clockticks
L3 Bound:	3.5% of Clockticks
Contested Accesses:	0.0% of Clockticks
Data Sharing:	4.5% of Clockticks
L3 Latency:	0.0% of Clockticks
SQ Full:	0.0% of Clockticks
DRAM Bound:	0.0% of Clockticks
Memory Bandwidth:	14.1% of Clockticks
Memory Latency:	21.1% of Clockticks
Store Bound:	0.0% of Clockticks
Store Latency:	31.7% of Clockticks
False Sharing:	0.0% of Clockticks
Split Stores:	0.0% of Clockticks
DTLB Store Overhead:	1.2% of Clockticks
Store STLB Hit:	0.0% of Clockticks
Store STLB Hit:	1.2% of Clockticks
Core Bound:	16.8% of Pipeline Slots

This metric represents how much Core non-memory issues were of a bottleneck. Shortage in hardware compute resources, or dependencies software's instructions are both categorized under Core Bound. Hence it may indicate the machine ran out of an OOO resources, certain execution units are overloaded or dependencies in program's data- or instruction- flow are limiting the performance (e.g. FP-chained long-latency arithmetic operations).

Divider:	0.0% of Clockticks
Port Utilization:	8.8% of Clockticks
Cycles of 0 Ports Utilized:	22.8% of Clockticks
Serializing Operations:	21.1% of Clockticks
Mixing Vectors:	0.0% of uOps
Cycles of 1 Port Utilized:	4.1% of Clockticks
Cycles of 2 Ports Utilized:	14.5% of Clockticks
Cycles of 3+ Ports Utilized:	16.6% of Clockticks
ALU Operation Utilization:	13.5% of Clockticks
Port 0:	8.3% of Clockticks
Port 1:	8.3% of Clockticks
Port 5:	8.3% of Clockticks
Port 6:	29.0% of Clockticks
Load Operation Utilization:	4.1% of Clockticks
Port 2:	8.3% of Clockticks
Port 3:	8.3% of Clockticks
Store Operation Utilization:	8.3% of Clockticks
Port 4:	8.3% of Clockticks
Port 7:	0.0% of Clockticks
Vector Capacity Usage (FPU):	0.0%
Average CPU Frequency:	1.198 GHz
Total Thread Count:	9
Paused Time:	0s

Effective Physical Core Utilization: 59.5% (2.381 out of 4)

The metric value is low, which may signal a poor physical CPU cores utilization caused by:

- load imbalance
- threading runtime overhead
- contended synchronization
- thread/process underutilization
- incorrect affinity that utilizes logical cores instead of physical cores

Explore sub-metrics to estimate the efficiency of MPI and OpenMP parallelism or run the Locks and Waits analysis to identify parallel bottlenecks for other parallel runtimes.

Effective Logical Core Utilization: 35.0% (2.801 out of 8)

The metric value is low, which may signal a poor logical CPU cores utilization. Consider improving physical core utilization as the first step and then look at opportunities to utilize logical cores, which in some cases can improve processor throughput and overall performance of multi-threaded applications.

Collection and Platform Info:

Application Command Line: ./codecs/HHI-VVC/decoder/vvdecapp "-b"
"./bin/HHI-VVC/randomaccess_fast.cfg/CLASS_C/
RaceHorses_416x240_30_QP_27_HHI-VVC.bin"

User Name: root

Operating System: 5.4.0-72-generic DISTRIB_ID=Ubuntu
DISTRIB_RELEASE=18.04 DISTRIB_CODENAME=bionic
DISTRIB_DESCRIPTION="Ubuntu 18.04.5 LTS"

Computer Name: eimon

Result Size: 14.5 MB

Collection start time: 22:29:07 18/04/2021 UTC

Collection stop time: 22:29:07 18/04/2021 UTC

Collector Type: Event-based sampling driver

CPU:

Name: Intel(R) Processor code named Kabylake
ULX

Frequency: 1.992 GHz

Logical CPU Count: 8

Cache Allocation Technology:

Level 2 capability: not detected

Level 3 capability: not detected

Intel® oneAPI VTune™ Profiler 2021.1.1 Gold

Elapsed Time: 0.025s

Application execution time is too short. Metrics data may be unreliable. Consider reducing the sampling interval or increasing your application execution time.

Clockticks:	173,160,000
Instructions Retired:	151,740,000
CPI Rate:	1.141

The CPI may be too high. This could be caused by issues such as memory stalls, instruction starvation, branch misprediction or long latency instructions. Explore the other hardware-related metrics to identify what is causing high CPI.

MUX Reliability:	0.967
Retiring:	46.8% of Pipeline Slots

A high fraction of pipeline slots was utilized by useful work. While the goal is to make this metric value as big as possible, a high Retiring value for non-vectorized code could prompt you to consider code vectorization. Vectorization enables doing more computations without significantly increasing the number of instructions, thus improving the performance. Note that this metric value may be highlighted due to Microcode Sequencer (MS) issue, so the performance can be improved by avoiding using the MS.

Light Operations:	32.9% of Pipeline Slots
FP Arithmetic:	0.0% of uOps
FP x87:	0.0% of uOps
FP Scalar:	0.0% of uOps
FP Vector:	0.0% of uOps
Other:	100.0% of uOps
Heavy Operations:	13.9% of Pipeline Slots

CPU retired heavy-weight operations (instructions that required 2+ uops) in a significant fraction of cycles.

Microcode Sequencer:	1.1% of Pipeline Slots
Assists:	0.0% of Pipeline Slots
Front-End Bound:	36.3% of Pipeline Slots

Issue: A significant portion of Pipeline Slots is remaining empty due to issues in the Front-End.

Tips: Make sure the code working size is not too large, the code layout does not require too many memory accesses per cycle to get enough instructions for filling four pipeline slots, or check for microcode assists.

Front-End Latency: 28.1% of Pipeline Slots

This metric represents a fraction of slots during which CPU was stalled due to front-end latency issues, such as instruction-cache misses, ITLB misses or fetch stalls after a branch misprediction. In such cases, the front-end delivers no uOps.

ICache Misses: 6.2% of Clockticks

Issue: A significant portion of instruction fetches is missing in the instruction cache.

Tips:

1. Use profile-guided optimization to reduce the size of hot code regions.
2. Consider compiler options to reorder functions so that hot functions are located together.
3. If your application makes significant use of macros, try to reduce this by either converting the relevant macros to functions or using linker options to eliminate repeated code.
4. Consider the Os/O1 optimization level or the following subset of optimizations to decrease your code footprint:
 - use inlining only when it decreases the footprint
 - disable loop unrolling
 - disable intrinsic inlining

Optimization examples:

Instruction Cache Misses recipe from Intel VTune Profiler Performance Analysis Cookbook

ITLB Overhead:	1.2% of Clockticks
Branch Resteers:	1.4% of Clockticks
Mispredicts Resteers:	0.0% of Clockticks
Clears Resteers:	0.0% of Clockticks
Unknown Branches:	1.4% of Clockticks
DSB Switches:	0.0% of Clockticks
Length Changing Prefixes:	0.0% of Clockticks
MS Switches:	0.0% of Clockticks

Issue: A significant fraction of cycles was stalled due to switches of uOp delivery to the Microcode Sequencer (MS). Commonly used instructions are optimized for delivery by the DSB or MITE pipelines. Certain operations cannot be handled natively by the execution pipeline, and must be performed by microcode (small programs injected into the execution stream). Switching to the MS too often can negatively impact performance. The MS is designated to deliver long uOp flows required by CISC instructions like CPUID, or uncommon conditions like Floating Point Assists when dealing with Denormals. Note that this metric value may be highlighted due to Microcode Sequencer issue.

Front-End Bandwidth:	8.2% of Pipeline Slots
Front-End Bandwidth MITE:	28.1% of Clockticks
Front-End Bandwidth DSB:	4.7% of Clockticks
(Info) DSB Coverage:	24.0%
Bad Speculation:	3.5% of Pipeline Slots
Branch Mispredict:	0.0% of Pipeline Slots
Machine Clears:	3.5% of Pipeline Slots

Back-End Bound:	13.5% of Pipeline Slots
Memory Bound:	9.5% of Pipeline Slots
L1 Bound:	9.4% of Clockticks
DTLB Overhead:	0.6% of Clockticks
Load STLB Hit:	0.0% of Clockticks
Load STLB Miss:	0.6% of Clockticks
Loads Blocked by Store Forwarding:	0.0% of Clockticks
Lock Latency:	0.0% of Clockticks
Split Loads:	0.0% of Clockticks
4K Aliasing:	0.0% of Clockticks
FB Full:	0.0% of Clockticks
L2 Bound:	3.1% of Clockticks
L3 Bound:	6.2% of Clockticks
Contested Accesses:	0.0% of Clockticks
Data Sharing:	0.0% of Clockticks
L3 Latency:	0.0% of Clockticks
SQ Full:	0.0% of Clockticks
DRAM Bound:	0.0% of Clockticks
Memory Bandwidth:	0.0% of Clockticks
Memory Latency:	18.7% of Clockticks
Store Bound:	0.0% of Clockticks
Store Latency:	12.5% of Clockticks
False Sharing:	0.0% of Clockticks
Split Stores:	0.0% of Clockticks
DTLB Store Overhead:	0.7% of Clockticks
Store STLB Hit:	0.0% of Clockticks
Store STLB Hit:	0.7% of Clockticks
Core Bound:	3.9% of Pipeline Slots
Divider:	0.0% of Clockticks
Port Utilization:	7.7% of Clockticks
Cycles of 0 Ports Utilized:	25.7% of Clockticks
Serializing Operations:	12.5% of Clockticks
Mixing Vectors:	0.0% of uOps
Cycles of 1 Port Utilized:	9.4% of Clockticks
Cycles of 2 Ports Utilized:	14.0% of Clockticks
Cycles of 3+ Ports Utilized:	18.7% of Clockticks
ALU Operation Utilization:	15.2% of Clockticks
Port 0:	9.4% of Clockticks
Port 1:	9.4% of Clockticks
Port 5:	9.4% of Clockticks
Port 6:	32.7% of Clockticks
Load Operation Utilization:	9.4% of Clockticks
Port 2:	14.0% of Clockticks
Port 3:	9.4% of Clockticks
Store Operation Utilization:	9.4% of Clockticks
Port 4:	9.4% of Clockticks
Port 7:	4.7% of Clockticks
Vector Capacity Usage (FPU):	0.0%
Average CPU Frequency:	2.244 GHz
Total Thread Count:	9

Paused Time: 0s

Effective Physical Core Utilization: 49.6% (1.984 out of 4)

The metric value is low, which may signal a poor physical CPU cores utilization caused by:

- load imbalance
- threading runtime overhead
- contended synchronization
- thread/process underutilization
- incorrect affinity that utilizes logical cores instead

of physical cores

Explore sub-metrics to estimate the efficiency of MPI and OpenMP parallelism or run the Locks and Waits analysis to identify parallel bottlenecks for other parallel runtimes.

Effective Logical Core Utilization: 37.2% (2.975 out of 8)

The metric value is low, which may signal a poor logical CPU cores utilization. Consider improving physical core utilization as the first step and then look at opportunities to utilize logical cores, which in some cases can improve processor throughput and overall performance of multi-threaded applications.

Collection and Platform Info:

Application Command Line: ./codecs/HHI-VVC/decoder/vvdecapp "-b" ".bin/HHI-VVC/randomaccess_fast.cfg/CLASS_C/RaceHorses_416x240_30_QP_27_HHI-VVC.bin"

User Name: root

Operating System: 5.4.0-72-generic DISTRIB_ID=Ubuntu
DISTRIB_RELEASE=18.04 DISTRIB_CODENAME=bionic
DISTRIB_DESCRIPTION="Ubuntu 18.04.5 LTS"

Computer Name: eimon

Result Size: 14.1 MB

Collection start time: 07:49:28 19/04/2021 UTC

Collection stop time: 07:49:28 19/04/2021 UTC

Collector Type: Event-based sampling driver

CPU:	
Name: ULX	Intel(R) Processor code named Kabylake
Frequency:	1.992 GHz
Logical CPU Count:	8
Cache Allocation Technology:	
Level 2 capability:	not detected
Level 3 capability:	not detected