Intel[®] oneAPI VTune[™] Profiler 2021.1.1 Gold

Elapsed Time: 0.047s

Application execution time is too short. Metrics data may be unreliable. Consider reducing the sampling interval or increasing your application execution time.

Clockticks: 43,020,000 **Instructions Retired:** 66,960,000

CPI Rate: 0.642 MUX Reliability: 0.996

Retiring: 46.6% of Pipeline Slots 44.5% of Pipeline Slots

FP Arithmetic:
FP x87:
0.0% of uOps
100.0% of uOps
100.0% of uOps
2.2% of Pipeline Slots

Microcode Sequencer: 3.1% of Pipeline Slots
Assists: 0.0% of Pipeline Slots

Front-End Bound: 25.1% of Pipeline Slots

Issue: A significant portion of Pipeline Slots is remaining empty due to issues in the Front-End.

Tips: Make sure the code working size is not too large, the code layout does not require too many memory accesses per cycle to get enough instructions for filling four pipeline slots, or check for microcode assists.

14.3% of Pipeline Slots Front-End Latency: ICache Misses: 0.0% of Clockticks 0.0% of Clockticks ITLB Overhead: **Branch Resteers:** 0.0% of Clockticks 0.0% of Clockticks **Mispredicts Resteers:** 0.0% of Clockticks **Clears Resteers: Unknown Branches:** 0.0% of Clockticks **DSB Switches:** 0.0% of Clockticks **Length Changing Prefixes:** 0.0% of Clockticks **MS Switches:** 0.0% of Clockticks Front-End Bandwidth: 10.8% of Pipeline Slots

This metric represents a fraction of slots during which CPU was stalled due to front-end bandwidth issues, such as inefficiencies in the instruction decoders or code restrictions for caching in the DSB (decoded uOps cache). In such cases, the front-end typically delivers a non-optimal amount of uOps to the back-end.

Front-End Bandwidth MITE: 43.0% of Clockticks

This metric represents a fraction of cycles during which CPU was stalled due to the MITE fetch pipeline issues, such as inefficiencies in the instruction decoders.

Front-End Bandwidth DSB: 14.3% of Clockticks (Info) DSB Coverage: 33.3%

Issue: A significant fraction of uOps was not delivered by the DSB (known as Decoded ICache or uOp Cache). This may happen if a hot code region is too large to fit into the DSB.

Tips: Consider changing the code layout (for example, via profile-guided optimization) to help your hot regions fit into the DSB.

See the "Optimization for Decoded ICache" section in the Intel 64 and IA-32 Architectures Optimization Reference Manual.

Bad Speculation:
Branch Mispredict:
Machine Clears:

Back-End Bound:

7.2% of Pipeline Slots
0.0% of Pipeline Slots
7.2% of Pipeline Slots

A significant portion of pipeline slots are remaining empty. When operations take too long in the back-end, they introduce bubbles in the pipeline that ultimately cause fewer pipeline slots containing useful work to be retired per cycle than the machine is capable to support. This opportunity cost results in slower execution. Long-latency operations like divides and memory operations can cause this, as can too many operations being directed to a single execution port (for example, more multiply operations arriving in the back-end per cycle than the execution unit can support).

Memory Bound: 8.6% of Pipeline Slots 12.6% of Clockticks L1 Bound: **DTLB Overhead:** 1.3% of Clockticks **Load STLB Hit:** 0.0% of Clockticks **Load STLB Miss:** 1.3% of Clockticks **Loads Blocked by Store Forwarding:** 0.0% of Clockticks 0.0% of Clockticks **Lock Latency: Split Loads:** 0.0% of Clockticks 4K Aliasing: 0.0% of Clockticks FB Full: 0.0% of Clockticks 0.0% of Clockticks L2 Bound: L3 Bound: 0.0% of Clockticks **Contested Accesses:** 0.0% of Clockticks 0.0% of Clockticks **Data Sharing:** 0.0% of Clockticks L3 Latency: SO Full: 0.0% of Clockticks **DRAM Bound:** 0.0% of Clockticks **Memory Bandwidth:** 0.0% of Clockticks **Memory Latency:** 12.6% of Clockticks **Store Bound:** 0.0% of Clockticks **Store Latency:** 0.0% of Clockticks False Sharing: 0.0% of Clockticks 0.0% of Clockticks **Split Stores:**

Core Bound: 12.5% of Pipeline Slots

DTLB Store Overhead:

Store STLB Hit:

Store STLB Hit:

This metric represents how much Core non-memory issues were of a bottleneck. Shortage in hardware compute resources, or dependencies software's instructions are both categorized under Core Bound. Hence it may indicate the machine ran out of an 000 resources, certain execution units are overloaded or dependencies in program's data- or instruction- flow are limiting the performance (e.g. FP-chained long-latency arithmetic operations).

1.4% of Clockticks

0.0% of Clockticks

1.4% of Clockticks

Divider: 0.0% of Clockticks **Port Utilization:** 18.4% of Clockticks

Cycles of 0 Ports Utilized: 14.3% of Clockticks
Serializing Operations: 0.0% of Clockticks

Mixing Vectors: 0.0% of uOps

Cycles of 1 Port Utilized: 7.2% of Clockticks

Cycles of 2 Ports Utilized: 7.2% of Clockticks

Cycles of 3+ Ports Utilized: 14.3% of Clockticks

ALU Operation Utilization: 21.5% of Clockticks

Port 0: 14.3% of Clockticks
Port 5: 28.7% of Clockticks
Port 6: 28.7% of Clockticks
Load Operation Utilization: 14.3% of Clockticks
Port 2: 28.7% of Clockticks
28.7% of Clockticks

Port 3: 28.7% of Clockticks
Store Operation Utilization: 28.7% of Clockticks
Port 4: 28.7% of Clockticks
28.7% of Clockticks
0.0% of Clockticks

Vector Capacity Usage (FPU): 0.0%

Average CPU Frequency: 1.035 GHz

Total Thread Count: 1 Paused Time: 0s

Effective Physical Core Utilization: 19.5% (0.781 out of 4)

The metric value is low, which may signal a poor physical CPU cores utilization caused by:

- load imbalance
- threading runtime overhead
- contended synchronization
- thread/process underutilization
- incorrect affinity that utilizes logical cores instead of physical cores

Explore sub-metrics to estimate the efficiency of MPI and OpenMP parallelism or run the Locks and Waits analysis to identify parallel bottlenecks for other parallel runtimes.

Effective Logical Core Utilization: 11.2% (0.892 out of 8)

The metric value is low, which may signal a poor logical CPU cores utilization. Consider improving physical core utilization as the first step and then look at opportunities to utilize logical cores, which in some cases can improve processor throughput and overall performance of multi-threaded applications.

Collection and Platform Info:

Application Command Line: ./codecs/HM/decoder/TAppDecoderStatic

"-b" "./bin/HM/encoder_randomaccess_main.cfg/CLASS_C/

RaceHorses 416x240 30 QP 37 HM.bin"

User Name: root

Operating System: 5.4.0-72-generic DISTRIB_ID=Ubuntu

DISTRIB RELEASE=18.04 DISTRIB CODENAME=bionic

DISTRIB DESCRIPTION="Ubuntu 18.04.5 LTS"

Computer Name: eimon

Result Size: 12.5 MB

Collection start time: 22:26:58 18/04/2021 UTC

Collection stop time: 22:26:58 18/04/2021 UTC

Collector Type: Event-based sampling driver

CPU:

Name: Intel(R) Processor code named Kabylake

ULX

Frequency: 1.992 GHz

Logical CPU Count: 8

Cache Allocation Technology:

Level 2 capability: not detected

Level 3 capability: not detected

Intel[®] oneAPI VTune[™] Profiler 2021.1.1 Gold

Elapsed Time: 0.042s

Application execution time is too short. Metrics data may be unreliable. Consider reducing the sampling interval or increasing your application execution time.

Clockticks: 42,300,000 **Instructions Retired:** 66,960,000

CPI Rate: 0.632 MUX Reliability: 0.894

Retiring: 49.2% of Pipeline Slots

A high fraction of pipeline slots was utilized by useful work. While the goal is to make this metric value as big as possible, a high Retiring value for non-vectorized code could prompt you to consider code vectorization. Vectorization enables doing more computations without significantly increasing the number of instructions, thus improving the performance. Note that this metric value may be highlighted due to Microcode Sequencer (MS) issue, so the performance can be improved by avoiding using the MS.

Light Operations: 33.9% of Pipeline Slots

 FP Arithmetic:
 0.0% of uOps

 FP x87:
 0.0% of uOps

 FP Scalar:
 0.0% of uOps

 FP Vector:
 0.0% of uOps

 Other:
 100.0% of uOps

Heavy Operations: 15.3% of Pipeline Slots

CPU retired heavy-weight operations (instructions that required 2+ uops) in a significant fraction of cycles.

Microcode Sequencer: 2.7% of Pipeline Slots 0.0% of Pipeline Slots

Front-End Bound: 16.4% of Pipeline Slots Front-End Latency: 10.9% of Pipeline Slots ICache Misses: 12.8% of Clockticks 1.3% of Clockticks ITLB Overhead: 0.0% of Clockticks **Branch Resteers: Mispredicts Resteers:** 0.0% of Clockticks **Clears Resteers:** 0.0% of Clockticks Unknown Branches: 0.0% of Clockticks 0.0% of Clockticks **DSB Switches: Length Changing Prefixes:** 0.0% of Clockticks **MS Switches:** 0.0% of Clockticks Front-End Bandwidth: 5.5% of Pipeline Slots

Front-End Bandwidth MITE: 32.8% of Clockticks **Front-End Bandwidth DSB:** 0.0% of Clockticks

(Info) DSB Coverage: 33.3%

Bad Speculation:
Branch Mispredict:
Machine Clears:

Back-End Bound:

0.0% of Pipeline Slots
0.0% of Pipeline Slots
34.3% of Pipeline Slots

A significant portion of pipeline slots are remaining empty. When operations take too long in the back-end, they introduce bubbles in the pipeline that ultimately cause fewer pipeline slots containing useful work to be retired per cycle than the machine is capable to support. This opportunity cost results in slower execution. Long-latency operations like divides and memory operations can cause this, as can too many operations being directed to a single execution port (for example, more multiply operations arriving in the back-end per cycle than the execution unit can support).

Memory Bound: 13.8% of Pipeline Slots 0.0% of Clockticks L1 Bound: **DTLB Overhead:** 5.1% of Clockticks **Load STLB Hit:** 0.0% of Clockticks **Load STLB Miss:** 5.1% of Clockticks **Loads Blocked by Store Forwarding:** 0.0% of Clockticks 0.0% of Clockticks **Lock Latency: Split Loads:** 0.0% of Clockticks 4K Aliasing: 0.0% of Clockticks FB Full: 0.0% of Clockticks 0.0% of Clockticks L2 Bound: L3 Bound: 12.8% of Clockticks **Contested Accesses:** 0.0% of Clockticks 0.0% of Clockticks **Data Sharing:** 0.0% of Clockticks L3 Latency: SO Full: 0.0% of Clockticks **DRAM Bound:** 0.0% of Clockticks **Memory Bandwidth:** 0.0% of Clockticks **Memory Latency:** 12.8% of Clockticks **Store Bound:** 0.0% of Clockticks **Store Latency:** 0.0% of Clockticks False Sharing: 0.0% of Clockticks 0.0% of Clockticks **Split Stores:**

Core Bound: 20.6% of Pipeline Slots

DTLB Store Overhead:

Store STLB Hit:

Store STLB Hit:

This metric represents how much Core non-memory issues were of a bottleneck. Shortage in hardware compute resources, or dependencies software's instructions are both categorized under Core Bound. Hence it may indicate the machine ran out of an 000 resources, certain execution units are overloaded or dependencies in program's data- or instruction- flow are limiting the performance (e.g. FP-chained long-latency arithmetic operations).

0.5% of Clockticks

0.0% of Clockticks

0.5% of Clockticks

Divider: 0.0% of Clockticks **Port Utilization:** 19.1% of Clockticks

Cycles of 0 Ports Utilized: 10.9% of Clockticks
Serializing Operations: 0.0% of Clockticks
Mixing Vectors: 0.0% of uOps

Cycles of 1 Port Utilized: 10.9% of Clockticks
Cycles of 2 Ports Utilized: 5.5% of Clockticks
Cycles of 3+ Ports Utilized: 16.4% of Clockticks
ALU Operation Utilization: 21.9% of Clockticks

Port 0: 21.9% of Clockticks
Port 5: 21.9% of Clockticks
Port 6: 21.9% of Clockticks

Port 3: 21.9% of Clockticks
Store Operation Utilization: 10.9% of Clockticks
Port 4: 10.9% of Clockticks
10.9% of Clockticks
10.9% of Clockticks

Vector Capacity Usage (FPU): 0.0%

Average CPU Frequency: 1.139 GHz

Total Thread Count: 1 Paused Time: 0s

Effective Physical Core Utilization: 26.1% (1.044 out of 4)

The metric value is low, which may signal a poor physical CPU cores utilization caused by:

- load imbalance
- threading runtime overhead
- contended synchronization
- thread/process underutilization
- incorrect affinity that utilizes logical cores instead of physical cores

Explore sub-metrics to estimate the efficiency of MPI and OpenMP parallelism or run the Locks and Waits analysis to identify parallel bottlenecks for other parallel runtimes.

Effective Logical Core Utilization: 11.2% (0.895 out of 8)

The metric value is low, which may signal a poor logical CPU cores utilization. Consider improving physical core utilization as the first step and then look at opportunities to utilize logical cores, which in some cases can improve processor throughput and overall performance of multi-threaded applications.

Collection and Platform Info:

Application Command Line: ./codecs/HM/decoder/TAppDecoderStatic

"-b" "./bin/HM/encoder_randomaccess_main.cfg/CLASS_C/

RaceHorses 416x240 30 QP 37 HM.bin"

User Name: root

Operating System: 5.4.0-72-generic DISTRIB ID=Ubuntu

DISTRIB RELEASE=18.04 DISTRIB CODENAME=bionic

DISTRIB DESCRIPTION="Ubuntu 18.04.5 LTS"

Computer Name: eimon

Result Size: 12.3 MB

Collection start time: 07:47:24 19/04/2021 UTC

Collection stop time: 07:47:24 19/04/2021 UTC

Collector Type: Event-based sampling driver

CPU:

Name: Intel(R) Processor code named Kabylake

ULX

Frequency: 1.992 GHz

Logical CPU Count: 8

Cache Allocation Technology:

Level 2 capability: not detected

Level 3 capability: not detected