CS 221 Fall 2005

# **Programming Languages**

**Handout 3 Nov 23, 2005** 

# Landin's SECD Machine

**Initial State** (for expression e): (nil, empty, e :: nil) :: nil

**Final States**: (v :: nil, empty, nil) :: nil, where v is the result value.

## Analyze

$$(S, E, op(e_1, e_2) :: C) :: D \quad \mapsto \quad (S, E, e_1 :: e_2 :: Prim(op) :: C) :: D \tag{1}$$

$$(S, E, \mathtt{apply}(e_1, e_2) :: C) :: D \quad \mapsto \quad (S, E, e_1 :: e_2 :: \mathtt{App} :: C) :: D \tag{2}$$

$$(S, E, if(e, e_1, e_2) :: C) :: D \mapsto (S, E, e :: If :: e_1 :: e_2 :: C) :: D$$

$$(3)$$

#### **Value Expressions**

$$(S, E, n :: C) :: D \mapsto (n :: S, E, C) :: D$$
 (4)

$$(S, E, b :: C) :: D \mapsto (b :: S, E, C) :: D$$
 (5)

$$(S, E, x :: C) :: D \quad \mapsto \quad (E(x) :: S, E, C) :: D \tag{6}$$

$$(S, E, (\operatorname{fun} f(x) \operatorname{is} e) :: C) :: D \mapsto ([\operatorname{fun} f(x) \operatorname{is} e; E] :: S, E, C) :: D \tag{7}$$

## **Function Call**

$$(v_2 :: v_1 :: S, E, \mathsf{App} :: C) :: D \mapsto$$

$$(nil, E'[f \mapsto v_1][x \mapsto v_2], e :: nil) :: (S, E, C) :: D$$

$$(8)$$

where  $v_1 = [\operatorname{fun} f(x) \operatorname{is} e; E'].$ 

## **Function Return**

$$(v :: S, E, nil) :: (S', E', C') :: D \quad \mapsto \quad (v :: S', E', C') :: D \tag{9}$$

#### **Consume values**

$$(v_2 :: v_1 :: S, E, Prim(op) :: C) :: D \mapsto ((v_1 op v_2) :: S, E, C) :: D$$

$$(10)$$

$$(true :: S, E, If :: e_1 :: e_2 :: C) :: D \mapsto (S, E, e_1 :: C) :: D$$
 (11)

$$(false :: S, E, \mathtt{If} :: e_1 :: e_2 :: C) :: D \quad \mapsto \quad (S, E, e_2 :: C) :: D \tag{12}$$