

Problem 1

QP:

$$\min_{\underline{x}} \frac{1}{2} \underline{x}^T G \underline{x} + c^T \underline{x}$$

$$\text{s.t. } a_i^T \underline{x} = b_i \quad i \in \mathbb{E}$$

$$a_i^T \underline{x} \geq b_i \quad i \in \mathbb{I}$$

- a) For the QP to be convex the matrix G must be positive semi-definite, $\underline{\underline{G \geq 0}}$

Convexity is important because it guarantees that local solutions are in fact global.

- b) If $\underline{z}^T G \underline{z} \geq 0$, we would have $g(\underline{z}) \geq g(\underline{z}^*)$ meaning \underline{z}^* would not be a strict local minimizer, only a local minimizer and \underline{z}^* is only one of possibly many global solutions.

$$c) \min_{\mathbf{x}} (x_1 - 1)^2 + (x_2 - 2,5)^2$$

$$= x_1^2 - 2x_1 + 1 + x_2^2 - 5x_2 + 2,5^2$$

$$\mathbf{G} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = \mathbf{I}$$

$$\mathbf{c}^T = [-2 \quad -5]$$

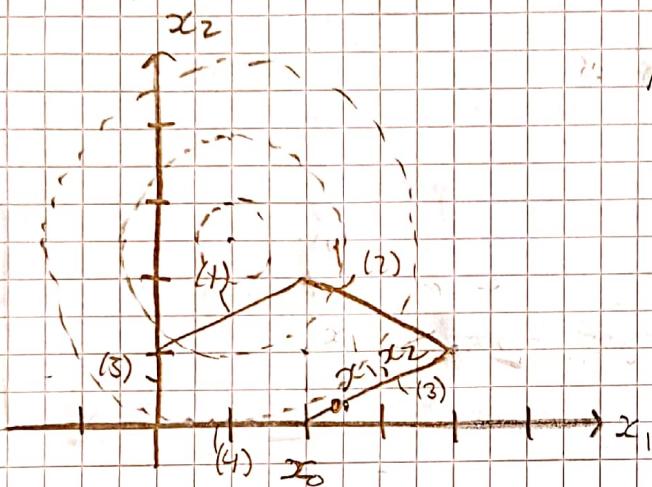
$$\text{s.t. } x_1 - 2x_2 + 2 \geq 0 \quad (1)$$

$$-x_1 - 2x_2 + 6 \geq 0 \quad (2)$$

$$-x_1 + 2x_2 + 2 \geq 0 \quad (3)$$

$$x_1 \geq 0 \quad (4)$$

$$x_2 \geq 0 \quad (5)$$



$$A = \begin{bmatrix} 1 & -2 \\ -1 & -2 \\ -1 & 2 \\ 1 & 0 \\ 0 & 1 \end{bmatrix} \quad \mathbf{x} \geq \mathbf{b} = \begin{bmatrix} -2 \\ 6 \\ 2 \\ 0 \\ 0 \end{bmatrix}$$

$$(0) \mathbf{x}_0 = [2 \quad 0]^T, \quad \mathcal{W}_0 = \{3\}$$

$$\min_p \frac{1}{2} \mathbf{p}_0^T \mathbf{G} \mathbf{p}_0 + \mathbf{g}_0^T \mathbf{p}_0 \quad \text{s.t.} \quad [-1 \quad 2] \mathbf{p} = 0$$

$$\rightarrow \mathbf{p}_0 = [0, 2 \quad 0, 1]^T \neq 0$$

$$\rightarrow \alpha_0 = \min_i \left(1, \frac{b_i - \mathbf{a}_i^T \mathbf{x}_0}{\mathbf{a}_i^T \mathbf{p}_0} \quad (i \notin \mathcal{W}_0) \right) = 1$$

$$\Rightarrow \mathbf{x}_1 = \mathbf{x}_0 + \alpha_0 \mathbf{p}_0 = [2, 2 \quad 0, 1]^T$$

$$(1) \quad x_1 = [2, 2, 0, 1]^T \quad \mathcal{W}_1 = \{3\}$$

$$\rightarrow p_1 \approx 0$$

$$\rightarrow [-1 \quad 2] \hat{\lambda} = g_1 = Gx_1 + c = \begin{bmatrix} 4 \\ 1 \end{bmatrix} + \begin{bmatrix} -2 \\ -5 \end{bmatrix} = \begin{bmatrix} 2 \\ -4 \end{bmatrix}$$

$$\Leftrightarrow \begin{bmatrix} -1 \\ 2 \end{bmatrix} \hat{\lambda} = \begin{bmatrix} 2 \\ -4 \end{bmatrix} \rightarrow \hat{\lambda} = -2, 4 \neq 0$$

$\Rightarrow x_1$ not optimal point

$$(2) \quad x_2 = [2, 2, 0, 1]^T \quad \mathcal{W}_2 = \emptyset$$

$$\rightarrow p_2 = [-1, 2, 2, 4]^T \neq 0$$

$$\rightarrow x_2 = \min(1, \dots) = 0, 607, \quad i=1$$

$$(3) \quad x_3 = [1, 4, 1, 7]^T \quad \mathcal{W}_3 = \{1\}$$

$$\rightarrow p_3 \approx 0$$

$$\rightarrow \hat{\lambda}_1 = 0, 8 \geq 0$$

$$\Rightarrow x^* = x_3 = \underline{\underline{\begin{bmatrix} 1, 4 \\ 1, 7 \end{bmatrix}}}$$

d)

$$q(x) = \frac{1}{2} x^T G x + c^T x \quad \text{s.t. } Ax \geq b$$

$$\mathcal{L}(x, \lambda) = \frac{1}{2} x^T G x + c^T x - \lambda^T (Ax - b)$$

$$\nabla_x \mathcal{L}(x, \lambda) = Gx + c - A^T \lambda = 0$$

$$\rightarrow x = G^{-1}(A^T \lambda - c)$$

\Rightarrow Dual objective function:

$$f(\lambda) = \inf_x \mathcal{L}(x, \lambda)$$

$$= \frac{1}{2} (A^T \lambda - c)^T G^{-1} (A^T \lambda - c) + c^T G^{-1} (A^T \lambda - c) \\ - \lambda^T (A G^{-1} (A^T \lambda - c) - b)$$

$$= -\frac{1}{2} (A^T \lambda - c)^T G^{-1} (A^T \lambda - c) + \lambda^T b$$

\Rightarrow Dual problem:

$$\max_{\lambda} -\frac{1}{2} (A^T \lambda - c)^T G^{-1} (A^T \lambda - c) + \lambda^T b \quad \text{s.t. } \lambda \geq 0.$$

e) theorem D.11: $f(\bar{\lambda}) \leq q(\bar{x})$, \bar{x} feasible, $\bar{\lambda} \geq 0$

$$, q(x^*) - q(\bar{x}) \leq q(x^*) - f(\bar{\lambda})$$

As we can see, $f(\bar{\lambda})$ will be an upper bound for the error $q(x^*) - q(\bar{x})$ when $q(x^*)$ is unknown.

Problem 2

a)

$$R_1: \begin{cases} 2A + 1B \leq 8 \\ 1A + 3B \leq 15 \\ A, B \geq 0 \end{cases} \quad \underline{x} = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} A \\ B \end{bmatrix}$$

$$\underline{x}^T \begin{bmatrix} 2 & 1 \\ 1 & 3 \end{bmatrix} \underline{x} \leq \begin{bmatrix} 8 \\ 15 \end{bmatrix}$$

$$\text{Profit: } (3 - 0.4x_1)x_1 + (2 - 0.2x_2)x_2$$

$$= 3x_1 - 0.4x_1^2 + 2x_2 - 0.2x_2^2$$

$$= \underline{x}^T \begin{bmatrix} -0.4 & 0 \\ 0 & -0.2 \end{bmatrix} \underline{x} + \begin{bmatrix} 3 & 2 \end{bmatrix} \underline{x}$$

$$= \frac{1}{2} \underline{x}^T \begin{bmatrix} -0.8 & 0 \\ 0 & -0.4 \end{bmatrix} \underline{x} + \begin{bmatrix} 3 & 2 \end{bmatrix} \underline{x}$$

\Rightarrow

$$\min_{\underline{x}} \frac{1}{2} \underline{x}^T \underbrace{\begin{bmatrix} -0.8 & 0 \\ 0 & -0.4 \end{bmatrix}}_G \underline{x} + \underbrace{\begin{bmatrix} 3 & 2 \end{bmatrix}}_{c^T} \underline{x}$$

$$\text{s.t. } \underbrace{\begin{bmatrix} -2 & -1 \\ -1 & -3 \\ 1 & 0 \\ 0 & 1 \end{bmatrix}}_A \underline{x} \geq \underbrace{\begin{bmatrix} -8 \\ -15 \\ 0 \\ 0 \end{bmatrix}}_b \quad \left. \begin{array}{l} \text{From } R_1, R_2 \\ \text{From } A, B \geq 0 \end{array} \right.$$

b) See below

c) $x^* = \begin{bmatrix} 2,3 \\ 3,5 \end{bmatrix}$ $x_{LP}^* = \begin{bmatrix} 1,8 \\ 4,4 \end{bmatrix}$

The optimal point of the QP problem x^* is not at the intersection point of the constraints. It's instead on one of the constraints, making only one constraint active.

d) From x_0 the algorithm "jumps" onto one constraint. There, it follows the constraint until it reaches the optimal point.

e) As we can see, the QP solution can be anywhere in the feasible area as opposed to LP where it must be at the intersection of two or more constraints.

TTK4135 Assignment 4: Problem 2
Contour with constraints (in red) and iterations

