Improve Your Software Quality Using Optional Type Systems



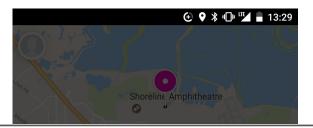
https://eisop.github.io/
Live demo: http://eisop.uwaterloo.ca/live/

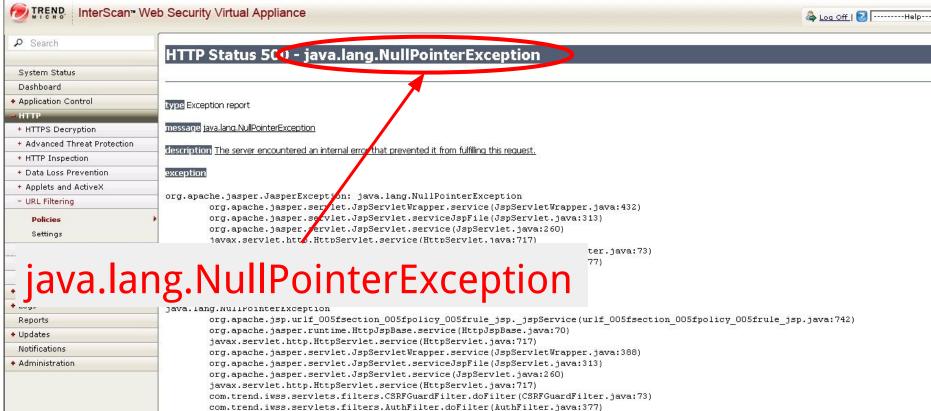
Werner Dietl, University of Waterloo



https://ece.uwaterloo.ca/~wdietl/

Motivation





Cost of software failures

\$312 billion per year global cost of software bugs (2013) **\$300 billion** dealing with the Y2K problem

\$440 million loss by Knight Capital Group Inc. in 30 minutes in August 2012

\$650 million loss by NASA Mars missions in 1999; unit conversion bug

\$500 million Ariane 5 maiden flight in 1996; 64-bit to 16-bit conversion bug



Software bugs can cost lives

1985-2000: >8 deaths: Radiation therapy

1991: 28 deaths: Patriot missile guidance system

1997: **225** deaths: jet crash caused by radar software

2003: **11 deaths**: blackout

2011: Software caused 25% of all medical device recalls



Java's type system is too weak

```
Type checking prevents many errors
int i = "hello";
```

Type checking doesn't prevent enough errors

```
System.console().readLine();
```

Collections.emptyList().add("one");



Java's type system is too weak

```
Type checking prevents many errors
  int i = "hello";
Type checking of NullPointerException
  System.console().readLine();
  Collections.emptyList().add("one");
```



Java's type system is too weak

```
Type checking prevents many errors
int i = "hello";
```

Type checking doesn't prevent enough errors

```
System UnsupportedOperationException
Collections.emptyList().add("one");
```



```
Date date = new Date();
myMap.put(date, "now");
date.setSeconds(0); // round to minute
myMap.get(date);
```



Corrupted map

```
Date date = new Date();
myMap.put(date, "now");
date.setSeconds(0); // round to minute
myMap.get(date);
```



dbStatement.executeQuery(userInput);



dbStatement.executeQuery(userInput);

SQL injection attack

Initialization, data formatting, equality tests, ...



Solution: Optional Type Checking

- 1. Design a type system to solve a specific problem
- 2. Write type qualifiers in code (or, use type inference)

```
void foo (@Nullable Date date) {
  date.setSeconds(0); // compile-time error
```

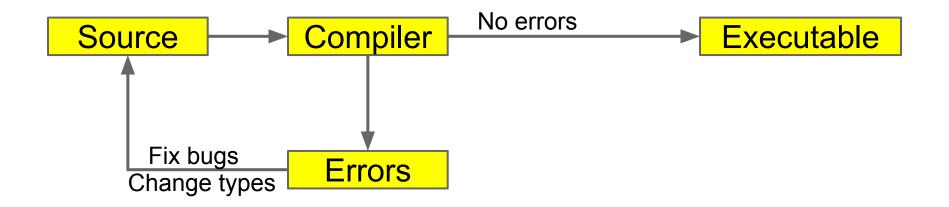
3. Type checker warns about violations (bugs)

```
% javac -processor NullnessChecker MyFile.java
```

MyFile.java:149: dereference of possibly-null reference bb2
 allVars = bb2.vars;

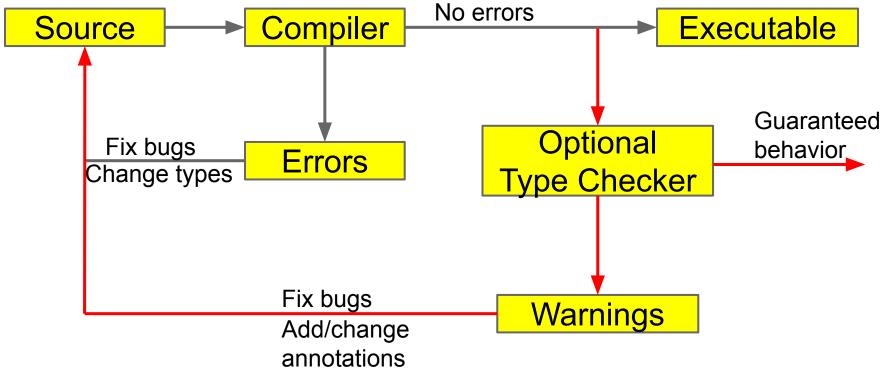


Type Checking



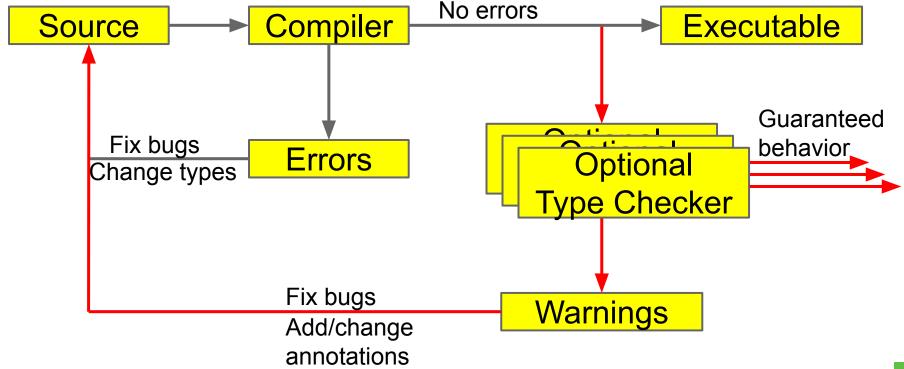


Optional Type Checking





Optional Type Checking



The Checker Framework

A framework for pluggable type checkers "Plugs" into the OpenJDK or OracleJDK compiler

javac -processor MyChecker ...

Standard error format allows tool integration



Ant, Maven, Gradle integration

```
cpresetdef name="jsr308.javac">
  <javac fork="yes"</pre>
   executable="${checkerframework}/checker/bin/${cfJavac}" >
    <!-- JSR-308-related compiler arguments -->
    <compilerarg value="-version"/>
    <compilerarg value="-implicit:class"/>
 </javac>
                                         <dependencies>
</presetdef>
                                           ... existing <dependency> items ...
                                           <!-- annotations from the Checker Framework:
                                                  nullness, interning, locking, ... -->
                                             <dependency>
                                               <groupId>io.github.eisop/groupId>
                                               <artifactId>checker-qual</artifactId>
                                               <version>3.34.0-eisop1/version>
                                             </dependency>
```

</dependencies>

Live demo: http://eisop.uwaterloo.ca/live/

Checker Framework Live Demo

Write Java code here:

```
import org.checkerframework.checker.nullness.qual.Nullable;
class YourClassNameHere {
   void foo(Object nn, @Nullable Object nbl) {
        nn.toString(); // OK
        nbl.toString(); // Error
   }
}
```

Choose a type system: Nullness Checker ▼

Check

Examples:

Nullness: NullnessExample | NullnessExampleWithWarnings

MapKey: <u>MapKeyExampleWithWarnings</u>

Interning: InterningExample | InterningExampleWithWarnings



Prevent null pointer exceptions

Java 8 introduces the Optional<T> type

- Wrapper; content may be present or absent
- Constructor: of(T value)
- Methods: boolean isPresent(), T get()

```
Optional<String> maidenName;
```



Optional reminds you to check

```
Without Optional:
                               With Optional:
         possible
                                        possible
         NullPointerException
                                        NoSuchElementException
                               Optional < St omName;
String mNz , e;
                               omName.get().equals(...);
mName.equals(...);
                               if (omName.: Present()) {
if (mName != null) {
                                               guals(...);
                                 omName.ge
  mName.equals(...);
```

possible NullPointerException

How <u>not</u> to use Optional

Stuart Marks's rules:

- Never, ever, use null for an Optional variable or return value.
- Never use Optional.get() unless you can prove that the Optional is present.
- Prefer alternative APIs over Optional.isPresent() and Optional.get().
- It's generally Let's enforce the
- If an Optiona rules with a tool. result of Opt

nal for the specific purpose of

chain, or has an intermediate

- Avoid using Optional in fields, method parameters, and collections.
- Don't use an Optional to wrap any collection type (List, Set, Map). Instead, use an empty collection to represent the absence of values.



Which rules to enforce with a tool

Stuart Marks's rules:

- 1. **Never**, ever, use null for an Optional variable or return value.
- 2. **Never** use Optional.get() unless you can prove that the Optional is present.
- 3. *Prefer* alternative APIs over Optional.isPresent() and Optional.get().
- 4. It's *generally a bad idea* to create an Optional for the specific purpose of chaining methods from it to get a value.
- 5. If an Optional chain has a nested Optional chain, or has an intermediate result of Optional, it's *probably too complex*.
- 6. Avoid using Optional in fields, method parameters, and collections.
- 7. **Don't** use an Optional to wrap any collection type (List, Set, Map). Instead, use an empty collection to represent the absence of values.



Which rules to enforce with a tool

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specific purpose of

Define a type system

```
Addr
Obj
                                                                                                                                                                                                                                                                                     = Set of Addresses \cup \{\text{null}_a\}
                                                                                                                                                                                                                                                                                     = rType, Fields
                                                                                                                                                                                                              \in <sup>r</sup>Type
                                                                                                                                                                                                                                                                                     = OwnerAddr ClassId<\(\bar{r}\)Type>
       P \in Program ::= \overline{Class}, ClassId, \underline{Expr} Fs
                                                                                                                                                                                                             \in
                                                                                                                                                                                                                              Fields
                                                                                                                                                                                                                                                                                     = FieldId \rightarrow Addr
                                     Class
                                                                                                   class ClassId<TVarId
Cls ∈
                                                                                                                                                                                                           \in
                                                                                                                                                                                                                             OwnerAddr
                                                                                                                                                                                                                                                                                     = Addr \cup \{any_a\}
                                                                                                    extends ClassId< Typ
                                                                                                                                                                                                                                                                                      = TVarId rType; ParId Addr
                                                                                                                                                                                                                              <sup>r</sup>Env
                                                                                                    { FieldId SType; Met
                                     <sup>s</sup>Type
                                                                                                    SNType | TVarId
                                                                                                                                                                                                                                                                                                                         h, {}^{r}\Gamma, e_0 \rightsquigarrow h', \iota_0
                                                                                                    OM ClassId < Type >
                                      <sup>s</sup>NType
                                                                                                                                                                                                                                                                                                                                     \iota_0 \neq \mathtt{null}_a
                                      OM
                                                                                                                                                         h, {}^{r}\Gamma, e_0 \rightsquigarrow h_0, \iota_0
                                                                                                                                                                                                                                                                    OS-Read \frac{\iota = h'(\iota_0) \downarrow_2 (f)}{h, {}^{\mathbf{r}}\Gamma, e_0.f \leadsto h', \iota}
                                     Meth
                                                                                                                                                                     \iota_0 \neq \mathtt{null}_a
                                      MethSig
                                                                                                                                                        h_0, {}^{\mathbf{r}}\Gamma, e_2 \rightsquigarrow h_2, \iota
                                                                                                 OS-Upd \frac{h' = h_2[\iota_0.f := \iota]}{h, {}^{r}\Gamma, e_0.f = e_2 \leadsto h',}
                      ∈ Purity
                   ∈ Expr
                                                                                                                                                                                                                                                                                    \Gamma \vdash e_0 : N_0 \qquad N_0 = u_0 C_0 < >
                                                                                                    Expr.MethId<sType>(Expr) |
                                                                                                                                                                                                                                                                                                        \mathtt{T}_1 = fType(\mathtt{C}_0,\mathtt{f})
                                                                                                    new SType (SType) Expr
                                                                                                                                                                                                                                                                                                        \Gamma \vdash e_2 : N_0 \triangleright T_1
                                                                             \text{GT-Read} \frac{\Gamma \vdash \mathsf{e}_0 : \mathsf{N}_0 \quad \mathsf{N}_0 = \_(\text{GT-Upd} \frac{\mathsf{u}_0 \neq \mathsf{any} \quad rp(\mathsf{u}_0, \mathsf{T}_1)}{\Gamma \vdash \mathsf{e}_0.\mathsf{f} : \mathsf{N}_0 \triangleright fType(\mathsf{C}_0, \mathsf{f})} 
    εГ
                                      Env
                                                                              ::= TVarId sNType; ParId sType
   h \vdash {}^{\mathbf{r}}\Gamma : {}^{\mathbf{s}}\Gamma
   h \vdash \iota_1 : dyn({}^{\mathfrak s}N, h, {}^{\mathfrak l}_{1,1})
   h \vdash \iota_2 : dyn(^{s}T, \iota_1, h(\iota_1)\downarrow_1)
                                                                                                                                             \implies h \vdash \iota_2 : dyn({}^{\mathtt{s}}\mathtt{N} \triangleright {}^{\mathtt{s}}\mathtt{T},\mathtt{h},{}^{\mathtt{r}}\Gamma)
   {}^{\mathtt{s}}\mathtt{N} = \mathtt{u}_N \; \mathtt{C}_N < >
                                                                                                                               u_N = this_u \Rightarrow {}^{\mathbf{r}}\Gamma(this)
                                                                                                                                                                     dom(C) = \overline{X}
                                                                                                                                                                                                                                                                                                      free(^{s}T) \subseteq \overline{X} \circ \overline{X'}
   free(^{s}T) \subseteq dom(C_N)
                                                                                                   DYN-
                                                                                                                             dyn(^{s}T, \iota, ^{r}T, (X' ^{r}T'; \bot)) = {^{s}T[\iota'/this, \iota'/peer, \iota/rep, any, \iota/rep, any,
```

Heap

= Addr \rightarrow Obj



Define a type system

- 1. **Type hierarchy** (subtyping)
- 2. **Type rules** (what operations are illegal)
- 3. **Type introduction** (what types for literals, ...)
- 4. **Dataflow** (run-time tests)

We will define two type systems: Nullness and Optional

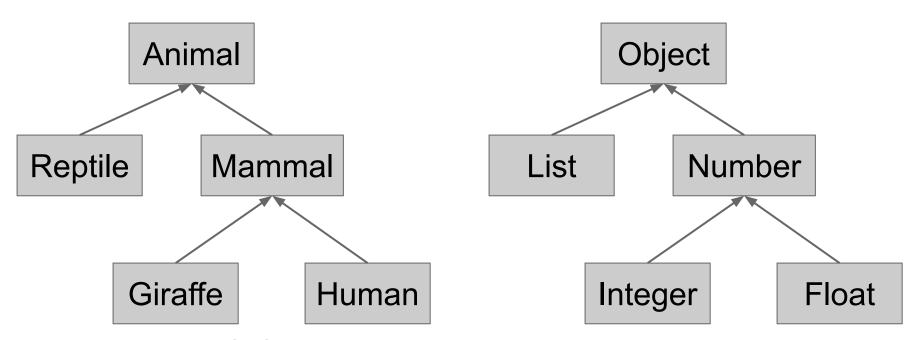


Define a type system

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1. Type hierarchy

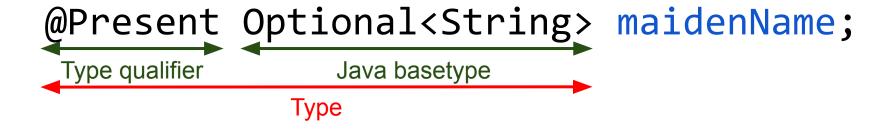


2 pieces of information:

- the types
- their relationships (lower = fewer values, more properties)



Type = type qualifier + Java basetype





```
String op(Data in) {
  return "transform: " + in.getF();
}
...
String s = op(null);
```



Where is the defect?

```
String op(Data in) {
  return "transform: " + in.getF();
}
...
String s = op(null);
```



Where is the defect?

```
String op(Data in) {
  return "transform: " + in.getF();
}
...
String s = op(null);
```



```
Where is the defect?
```

```
String op(Data in) {
   return "transform: " + in.getF();
}
Can't decide without specification!
...
String s = op(null);
```



Specification 1: non-null parameter

```
String op(@NonNull Data in) {
  return "transform: " + in.getF();
}
...
String s = op(null);
```



Specification 1: non-null parameter

```
String op(@NonNull Data in) {
  return "transform: " + in.getF();
}
...
String s = op(null); // error
```



Specification 2: nullable parameter

```
String op(@Nullable Data in) {
  return "transform: " + in.getF();
}
...
String s = op(null);
```

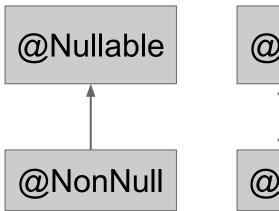


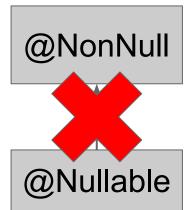
Specification 2: nullable parameter

```
String op(@Nullable Data in) {
  return "transform: " + in.getF();
}
  // error
...
String s = op(null);
```



Type hierarchy for nullness





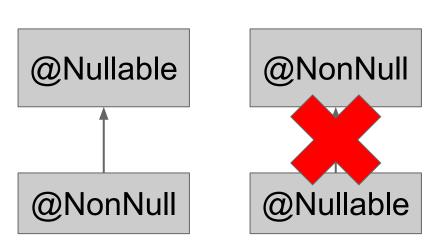


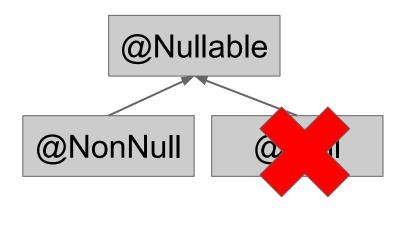
2 pieces of information:

- the types
- their relationships



Type hierarchy for nullness





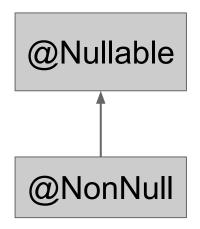
2 pieces of information:

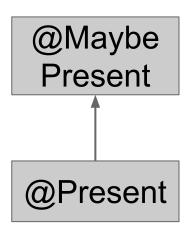
- the types
- their relationships



Type hierarchy for Optional

"Never use Optional.get() unless you can prove that the Optional is present."





2 pieces of information:

- the types
- their relationships



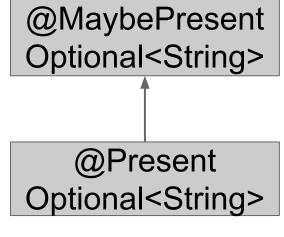
Type = type qualifier + Java basetype



Default qualifier = @MaybePresent

Type

- @MaybePresent Optional<String>
- Optional<String>



equivalent



Define a type system

- 1. Type hierarchy (subtyping)
- 2. Type rules (what operations are illegal)
- 3. Type introduction (what types for literals, ...)
- 4. Dataflow (run-time tests)

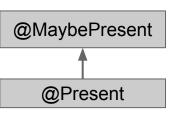


2. Type rules

To prevent <u>null pointer exceptions</u>:

- expr.field
 expr.getValue()
 receiver must be non-null
- synchronized (expr) { ... }
 monitor must be non-null
- ...





"Never use Optional.get() unless you can prove that the Optional is present."

Only call Optional.get() on a receiver of type

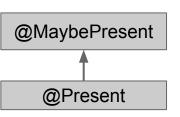
@Present Optional.

```
class Optional<T> {
   T get() { ... }
}
```

```
example call:
myOptional.get()
```

```
example call:
a.equals(b)
```

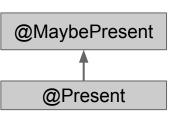




"Never use Optional.get() unless you can prove that the Optional is present."

```
Only call Optional.get() on a receiver of type
@Present Optional.
                           example call:
                           myOptional.get()
class Optional<T> {
  T get(Optional<T> this) { ... }
```

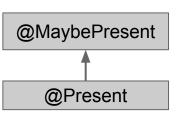




"Never use Optional.get() unless you can prove that the Optional is present."

```
Only call Optional.get() on a receiver of type
@Present Optional.
                          example call:
                          myOptional.get()
class Optional<T> {
  T get(@Present Optional<T> this) {...}
```





"Never use Optional.get() unless you can prove that the Optional is present."

```
Only call Optional.get() on a receiver of type
@Present Optional.
                           example call:
                           myOptional.get()
class Optional<T> {
  T get(@Present Optional<T> this) {...}
  T orElseThrow(@Present O... this, ...) {...}
```



Define a type system

- 1. Type hierarchy (subtyping)
- 2. Type rules (what operations are illegal)
- 3. Type introduction (what types for literals...)
- 4. Dataflow (run-time tests)



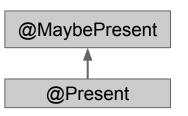
Type introduction rules

For Nullness type system:

- null : @Nullable
- "Hello World" : @NonNull



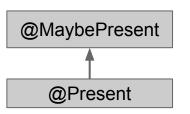
Type introduction for Optional



```
Optional<T> of(T value) {...}
Optional<T> ofNullable(T value){...}
```



Type introduction for Optional



```
@Present Optional<T> of(T value) {...}
Optional<T> ofNullable(@Nullable T value){...}
```



Define a type system

- 1. Type hierarchy (subtyping)
- 2. Type rules (what operations are illegal)
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Flow-sensitive type refinement

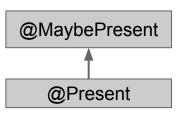
After an operation, give an expression a more specific type

```
@Nullable Object x;
if (x != null) {
    ...    x is @NonNull here
}

y = new SomeType();
    ...    y is @NonNull here

y = unknownValue;
    ...    y is @Nullable again
...    y is @Nullable again
```

Type refinement for Optional



"Never use Optional.get() unless you can prove that the Optional is present."

After receiver.isPresent() returns true, the receiver's type is @Present

```
@MaybePresent Optional<String> x;
if (x.isPresent()) {
    ...     x is @Present here
}
...     x is @MaybePresent again
```



Let's look at the Optional Checker

https://github.com/eisop/checker-framework/tr ee/master/checker/src/main/java/org/checkerf ramework/checker/optional

https://github.com/eisop/checker-framework/tr ee/master/checker-qual/src/main/java/org/che ckerframework/checker/optional/qual



Design the type system first

Before you start coding, first write the user manual.

What problem are you solving?

What qualifiers will you need?

What rules do you need to enforce?



Implement type qualifiers and hierarchy

```
@Documented
@Retention(RetentionPolicy.RUNTIME)
@Target({ElementType.TYPE USE,
     ElementType.TYPE PARAMETER})
@SubtypeOf({MaybePresent.class})
public @interface Present {}
```

https://github.com/eisop/checker-framework/tree/master/checker-qual/src/main/java/org/checkerframework/checker/optional/qual



Use the Subtyping Checker

Experiment with the type qualifiers using the Subtyping Checker:

```
javac -processor ...SubtypingChecker \
  -Aquals=...Present,...MaybePresent \
  SubtypeCheck.java
```



Name your checker

Simplify usage by adding a main class:

```
public class OptionalChecker
   extends BaseTypeChecker {}
```



Annotate parts of the JDK

Add jdk.astub JDK specifications

```
class Optional<T extends Object> {
    static <T extends Object> @Present Optional<T> of(T value) ...
    T get(@Present Optional<T> this) ...
    ...
}
```

https://github.com/eisop/jdk/blob/master/src/java.base/share/classes/java/util/Optional.java



Type rules and type introductions

No additional rules needed for the Optional Checker — everything can be specified declaratively and use Checker Framework rules.



Implement dataflow refinement

Declarative specification possible:

```
@EnsuresQualifierIf(result = true,
  expression = "this",
  qualifier = Present.class)
public boolean isPresent() {
    return value != null;
```



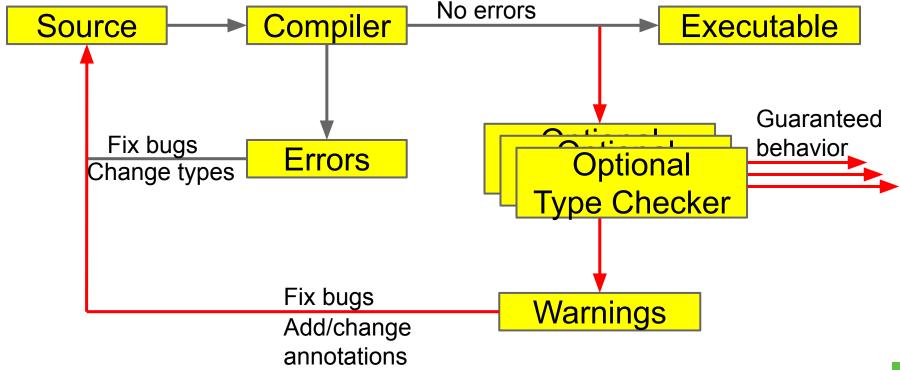
You can use the Optional Checker

Distributed with the Checker Framework Checks 6 of the 7 rules for using Optional

https://eisop.github.io/cf/manual/#optional-checker



Optional Type Checking





Benefits of type systems

- Find bugs in programs
 - Guarantee the absence of errors
- Improve documentation
 - Improve code structure & maintainability
- Aid compilers, optimizers, and analysis tools
 - E.g., could reduce number of run-time checks
- Possible negatives:
 - Must write the types (or use type inference)
 - False positives are possible (can be suppressed)



Checker Framework facilities

- Full type systems: inheritance, overriding, ...
- Generics (type polymorphism)
 - Also qualifier polymorphism
- Qualifier defaults
- Pre-/post-conditions
- Warning suppression



Example type systems

problems in Swing, JabRef

```
Null dereferences (@Nullable)
   >200 errors in Google Collections, javac, ...
Equality tests (@Interned)
   >200 problems in Xerces, Lucene, ...
Concurrency / locking (@GuardedBy)
   >500 errors in BitcoinJ, Derby, Guava, Tomcat, ...
Fake enumerations / typedefs (@Fenum)
```



String type systems

```
Regular expression syntax (@Regex)
   56 errors in Apache, etc.; 200 annos required
printf format strings (@Format)
   104 errors, only 107 annotations required
Method signature format (@FullyQualified)
   28 errors in OpenJDK, ASM, AFU
Compiler messages (@CompilerMessageKey)
   8 wrong keys in Checker Framework
```



Security type systems

Command injection vulnerabilities (@OsTrusted)

5 missing validations in Hadoop

Information flow privacy (@Source)

SPARTA detected malware in Android apps



It's easy to write your own type system!



What a checker guarantees

The program satisfies the type property. There are:

- no bugs (of particular varieties)
- no wrong annotations
- Caveat 1: only for code that is checked
 - Native methods (handles reflection!)
 - Code compiled without the pluggable type checker
 - Suppressed warnings
 - Indicates what code a human should analyze
 - Checking part of a program is still useful
 - Caveat 2: The checker itself might contain an error



Verification

- Goal: prove that no bug exists
- Specifications: user provides
- False negatives: none
- False positives: user suppresses warnings
- Downside: user burden

Bug-finding

- **Goal**: find some bugs at low cost
- Specifications: infer likely specs
- False negatives: acceptable
- False positives: heuristics focus on most important bugs
- Downside: missed bugs



Neither is "better"; each is appropriate in certain circumstances.

Checkers are usable

- Type-checking is familiar to programmers
- Modular: fast, incremental, partial programs
- Annotations are not too verbose
 - @Nullable: 1 per 75 lines
 - **@Interned**: 124 annotations in 220 KLOC revealed 11 bugs
 - @Format: 107 annotations in 2.8 MLOC revealed 104 bugs
 - Possible to annotate part of program
 - Fewer annotations in new code
- Few false positives
- First-year CS majors preferred using checkers to not
- Practical: in use in Silicon Valley, on Wall Street, etc.



Tips

- Start by type-checking part of your code
- · Only type-check properties that matter to you
- Use subclasses (not type qualifiers) if possible
- Write the spec first (and think of it as a spec)
- Avoid warning suppressions when possible
- Avoid raw types such as List; use List<String>



More at JAX 2023

Null Pointer Exceptions: the Problems, Current Approaches, and Ongoing Efforts

Donnerstag, 11. Mai 2023

16:45 - 17:45

Raum: Kongresssaal



Pluggable type-checking improves code

Checker Framework for creating type checkers

• Featureful, effective, easy to use, scalable

Prevent bugs at compile time

Create custom type-checkers

Improve your code!

https://eisop.github.io/

