Implement your own Type System, today! The Checker Framework



http://CheckerFramework.org/

Twitter: @CheckerFrmwrk

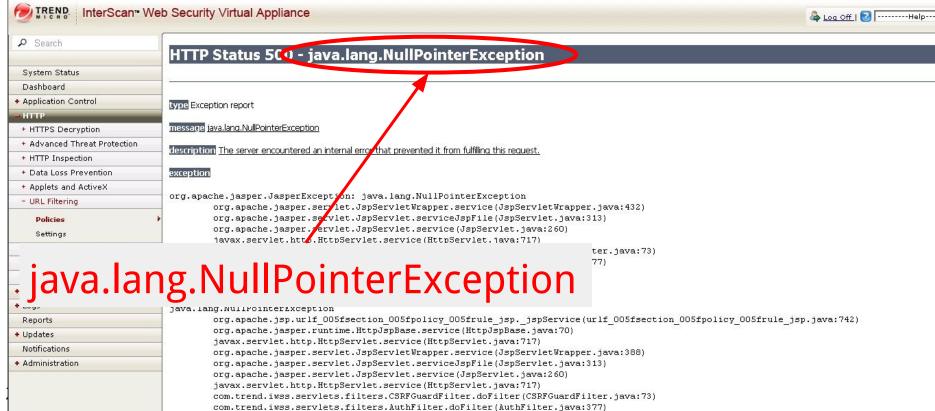
Live demo: http://CheckerFramework.org/live/

Werner Dietl, University of Waterloo



Motivation





Java's type system is too weak

Type checking prevents many errors
int i = "hello";

Type checking doesn't prevent enough errors

```
System.console().readLine();
```



Java's type system is too weak

```
Type checking prevents many errors
int i = "hello";
```

```
Type checking (NullPointerException Orself). System.console().readLine();
```



Prevent null pointer exceptions

Java 8 introduces the Optional<T> type

- Wrapper; content may be present or absent
- Constructor: of(T value)
- Methods: boolean isPresent(), T get()

Optional<String> maidenName;



Optional reminds you to check

```
Without Optional:
         possible
         NullPointerException
String mNz , e;
mName.equals(...);
if (mName != null) {
  mName.equals(...);
```

```
With Optional:
```

```
possible
         NoSuchElementException
Optional < St omName;
omName.get().equals(...);
if (omName.: Present()) {
                guals(...);
  omName.ge
             possible
```

Complex rules for using Optional correctly!



How <u>not</u> to use Optional

Stuart Marks's rules:

- Never, ever, use null for an Optional variable or return value.
- Never use Optional.get() unless you can prove that the Optional is present.
- Prefer alternative APIs over Optional.isPresent() and Optional.get().
- It's generally Let's enforce the
- If an Optiona rules with a tool. result of Opt

nal for the specific purpose of

chain, or has an intermediate

- Avoid using Optional in fields, method parameters, and collections.
- Don't use an Optional to wrap any collection type (List, Set, Map). Instead, use an empty collection to represent the absence of values.





Which rules to enforce with a tool

Stuart Marks's rules:

- 1. Never, ever, use null for an Optional variable or return value.
- 2. **Never** use Optional.get() unless you can prove that the Optional is present.
- 3. *Prefer* alternative APIs over Optional.isPresent() and Optional.get().
- 4. It's *generally a bad idea* to create an Optional for the specific purpose of chaining methods from it to get a value.
- 5. If an Optional chain has a nested Optional chain, or has an intermediate result of Optional, it's *probably too complex*.
- 6. Avoid using Optional in fields, method parameters, and collections.
- 7. **Don't** use an Optional to wrap any collection type (List, Set, Map). Instead, use an empty collection to represent the absence of values.



Which rules to enforce with a tool

Stuart Marks's rules:

- (1.) **Never**, ever, use null for an Optional variable or return value.
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- 5. If an Optiona type system properties. as an intermediate result of Opt
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specific purpose of

```
Addr
Obj
                                                                                                                           = Set of Addresses \cup \{\text{null}_a\}
                                                                                                                           = rType, Fields
                                                                                           \in <sup>r</sup>Type
                                                                                                                           = OwnerAddr ClassId<\(\bar{r}\)Type>
   P \in Program ::= \overline{Class}, ClassId, \underline{Expr} Fs
                                                                                           \in
                                                                                                  Fields
                                                                                                                           = FieldId \rightarrow Addr
                Class
                                            class ClassId<TVarId
Cls ∈
                                                                                                                           = Addr \cup {any<sub>a</sub>}
                                                                                          \in
                                                                                                  OwnerAddr
                                            extends ClassId< Typ
                                                                                                                           = TVarId rType; ParId Addr
                                                                                                   <sup>r</sup>Env
                                            { FieldId SType; Met
                <sup>s</sup>Type
                                            SNType | TVarId
                                                                                                                                           h, {}^{r}\Gamma, e_0 \rightsquigarrow h', \iota_0
                                            OM ClassId < Type >
                 <sup>s</sup>NType
                                                                                                                                                \iota_0 \neq \mathtt{null}_a
                 OM
                                                                    h, {}^{r}\Gamma, e_0 \rightsquigarrow h_0, \iota_0
                                                                                                                    OS-Read \frac{\iota = h'(\iota_0) \downarrow_2 (f)}{h, {}^{\mathbf{r}}\Gamma, e_0.f \leadsto h', \iota}
                Meth
                                                                         \iota_0 \neq \mathtt{null}_a
                 MethSig
                                                                   h_0, {}^{\mathbf{r}}\Gamma, e_2 \rightsquigarrow h_2, \iota
                                           OS-Upd \frac{h' = h_2[\iota_0.f := \iota]}{h, {}^{r}\Gamma, e_0.f = e_2 \leadsto h',}
         ∈ Purity
        ∈ Expr
                                                                                                                           \Gamma \vdash e_0 : N_0 \qquad N_0 = u_0 C_0 < >
                                            Expr.MethId<sType>(Expr)
                                                                                                                                    \mathtt{T}_1 = fType(\mathtt{C}_0,\mathtt{f})
                                            new SType (SType) Expr
                                                                                                                                    \Gamma \vdash e_2 : N_0 \triangleright T_1
                                   \text{GT-Read} \frac{\Gamma \vdash \mathsf{e}_0 : \mathsf{N}_0 \quad \mathsf{N}_0 = \_(\text{GT-Upd} \frac{\mathsf{u}_0 \neq \mathsf{any} \quad rp(\mathsf{u}_0, \mathsf{T}_1)}{\Gamma \vdash \mathsf{e}_0.\mathsf{f} : \mathsf{N}_0 \triangleright fType(\mathsf{C}_0, \mathsf{f})} 
 εГ
                 Env
                                   ::= TVarId sNType; ParId sType
 h \vdash {}^{\mathbf{r}}\Gamma : {}^{\mathbf{s}}\Gamma
 h \vdash \iota_1 : dyn({}^{\mathfrak s}N, h, {}^{\mathfrak l}_{1,1})
 h \vdash \iota_2 : dyn(^{s}T, \iota_1, h(\iota_1)\downarrow_1)
                                                               \implies h \vdash \iota_2 : dun({}^{\mathtt{s}}\mathtt{N} \triangleright^{\mathtt{s}}\mathtt{T}.h.{}^{\mathtt{r}}\Gamma)
 {}^{\mathtt{s}}\mathtt{N} = \mathtt{u}_N \; \mathtt{C}_N < \!\!\!>
                                                        u_N = this_u \Rightarrow {}^{\mathbf{r}}\Gamma(this)
                                                                         dom(C) = \overline{X}
                                                                                                                                   free(^{s}T) \subseteq \overline{X} \circ \overline{X'}
 free(^{s}T) \subseteq dom(C_N)
                                            DYN-
                                                        dyn(^{s}T, \iota, ^{r}T, (\overline{X'} ^{r}T'; \_)) = {^{s}T[\iota'/this}, \iota'/peer, \iota/rep, any, /any, /^{r}T/X, /rT'/X']
```

Heap

= Addr \rightarrow Obj



- 1. **Type hierarchy** (subtyping)
- 2. **Type rules** (what operations are illegal)
- 3. **Type introduction** (what types for literals, ...)
- 4. **Dataflow** (run-time tests)

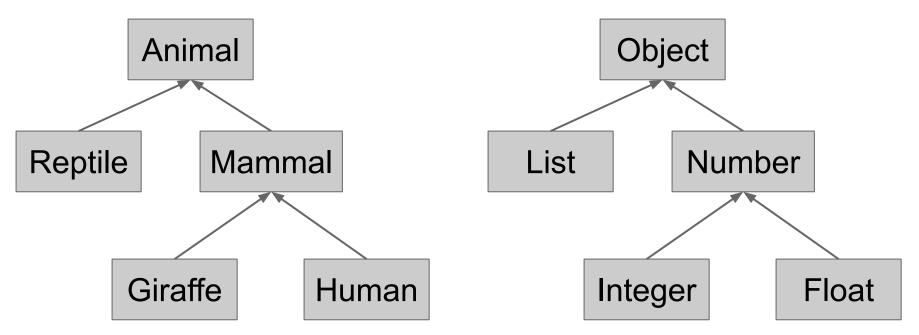
We will define two type systems: Nullness and Optional



- 1. Type hierarchy (subtyping)
- 2. Type rules (what operations are illegal)
- 3. Type introduction (what types for literals, ...)
- 4. Dataflow (run-time tests)



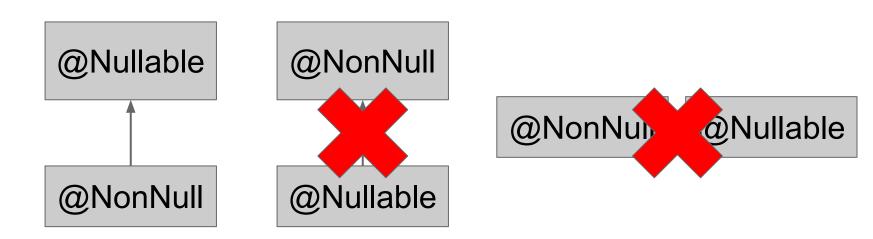
1. Type hierarchy



- the types
- their relationships (lower = fewer values, more properties)



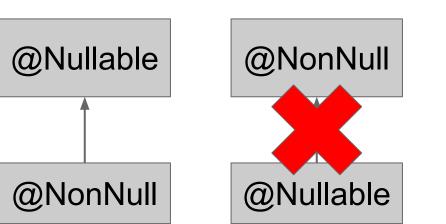
Type hierarchy for nullness

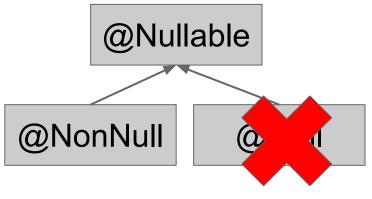


- the types
- their relationships



Type hierarchy for nullness



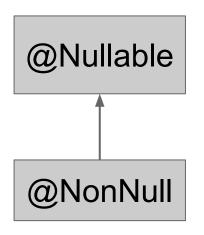


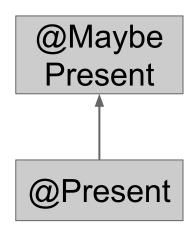
- the types
- their relationships



Type hierarchy for Optional

"Never use Optional.get() unless you can prove that the Optional is present."





- the types
- their relationships



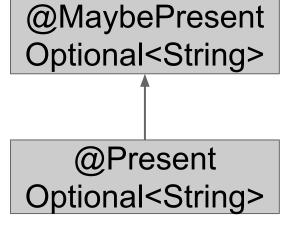
Type = type qualifier + Java basetype



Default qualifier = @MaybePresent

Type

- @MaybePresent Optional<String>
- Optional<String>



equivalent



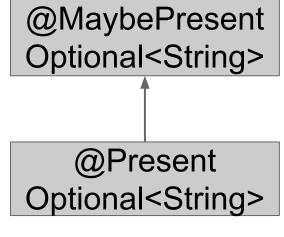
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Type

- @MaybePresent Optional<String>
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equivalent



- 1. Type hierarchy (subtyping)
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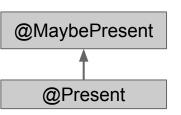
2. Type rules

To prevent <u>null pointer exceptions</u>:

```
    expr.field
    expr.getValue()
    receiver must be non-null
```

- synchronized (expr) { ... }
 monitor must be non-null
- ...





"Never use Optional.get() unless you can prove that the Optional is present."

Only call Optional.get() on a receiver of type

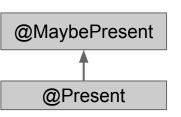
@Present Optional.

```
class Optional<T> {
   T get() { ... }
}
```

```
example call:
myOptional.get()
```

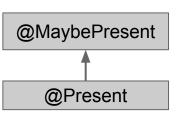
```
example call:
a.equals(b)
```





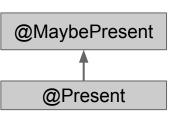
```
Only call Optional.get() on a receiver of type
@Present Optional.
                           example call:
                           myOptional.get()
class Optional<T> {
  T get(Optional<T> this) { ... }
```





```
Only call Optional.get() on a receiver of type
@Present Optional.
                          example call:
                          myOptional.get()
class Optional<T> {
  T get(@Present Optional<T> this) {...}
```





```
Only call Optional.get() on a receiver of type
@Present Optional.
                           example call:
                           myOptional.get()
class Optional<T> {
  T get(@Present Optional<T> this) {...}
  T orElseThrow(@Present O... this, ...) {...}
```



- 1. Type hierarchy (subtyping)
- 2. Type rules (what operations are illegal)
- 3. Type introduction (what types for literals...)
- 4. Dataflow (run-time tests)



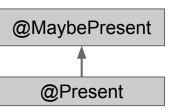
Type introduction rules

For Nullness type system:

- null : @Nullable
- "Hello World" : @NonNull



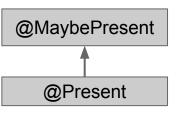
Type introduction for Optional



```
Optional<T> of(T value) {...}
Optional<T> ofNullable(T value){...}
```



Type introduction for Optional



```
@Present Optional<T> of(T value) {...}
Optional<T> ofNullable(@Nullable T value){...}
```



- 1. Type hierarchy (subtyping)
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Flow-sensitive type refinement

After an operation, give an expression a more specific type

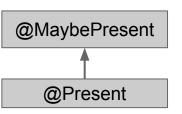
```
@Nullable Object x;
if (x != null) {
    ... x is @NonNull here
}

y = new SomeType();
    ... y is @NonNull here

y = unknownValue;
    x is @Nullable again

y is @Nullable again
```

Type refinement for Optional



"Never use Optional.get() unless you can prove that the Optional is present."

After receiver.isPresent() returns true, the receiver's type is @Present

```
@MaybePresent Optional<String> x;
if (x.isPresent()) {
    ...     x is @Present here
}
...     x is @MaybePresent again
```



Now, let's implement it

Follow the instructions in the Checker Framework Manual

https://checkerframework.org/manual/#creating-a-checker



Design the type system first

Before you start coding, first write the user manual.

What problem are you solving?

What qualifiers will you need?

What rules do you need to enforce?



Implement type qualifiers and hierarchy

```
@Documented
@Retention(RetentionPolicy.RUNTIME)
@Target({ElementType.TYPE USE,
     ElementType.TYPE PARAMETER})
@SubtypeOf({MaybePresent.class})
public @interface Present {}
```



Use the Subtyping Checker

Experiment with the type qualifiers using the Subtyping Checker:

```
javac -processor ...SubtypingChecker \
  -Aquals=...Present,...MaybePresent \
  SubtypeCheck.java
```



Name your checker

Simplify usage by adding a main class:

```
public class OptionalChecker
    extends BaseTypeChecker {}
```



Annotate parts of the JDK

Add jdk.astub JDK specifications

```
class Optional<T extends Object> {
    static <T extends Object> @Present Optional<T> of(T value) ...
    T get(@Present Optional<T> this) ...
    ...
}
```

https://github.com/typetools/jdk/blob/master/src/java.base/share/classes/java/util/Optional.java



Type rules and type introductions

Not needed for the Optional Checker



Implement dataflow refinement

Declarative specification possible:

```
@EnsuresQualifierIf(result = true,
  expression = "this",
  qualifier = Present.class)
public boolean isPresent() {
    return value != null;
```



You can use the Optional Checker

Distributed with the Checker Framework Checks 6 of the 7 rules for using Optional

https://checkerframework.org/manual/#optional-checker



Solution: Pluggable Type Checking

- 1. Design a type system to solve a specific problem
- 2. Write type qualifiers in code (or, use type inference)

```
void foo (@Nullable Date date) {
  date.setSeconds(0); // compile-time error
```

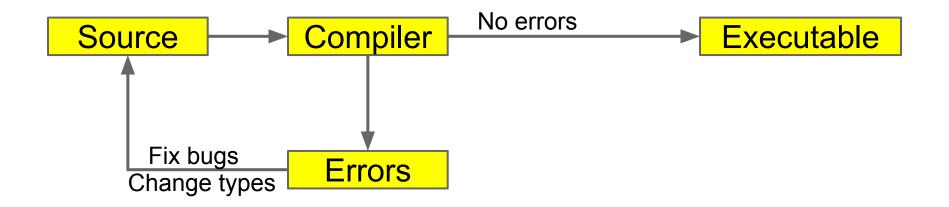
3. Type checker warns about violations (bugs)

```
% javac -processor NullnessChecker MyFile.java
```

MyFile.java:149: dereference of possibly-null reference bb2
 allVars = bb2.vars;

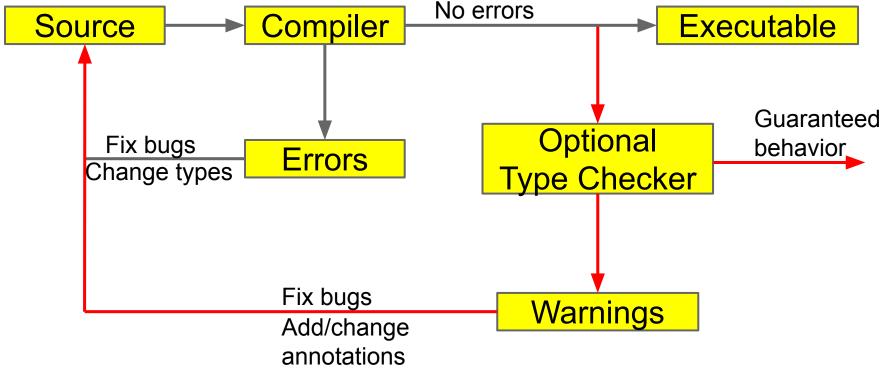


Type Checking



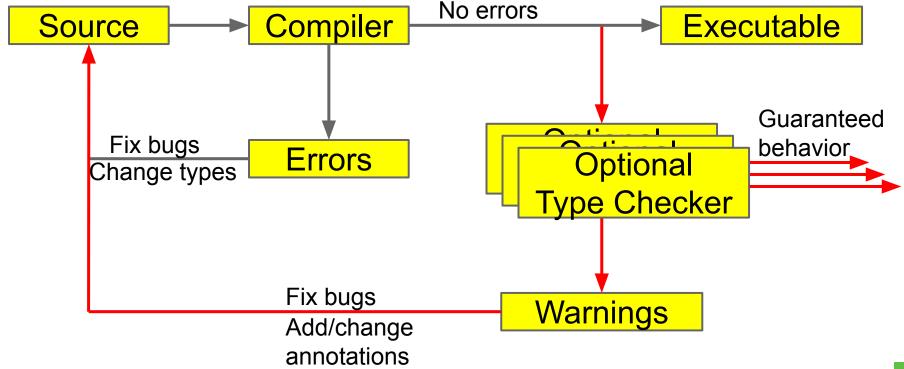


Optional Type Checking





Optional Type Checking





Benefits of type systems

- Find bugs in programs
 - Guarantee the absence of errors
- Improve documentation
 - Improve code structure & maintainability
- Aid compilers, optimizers, and analysis tools
 - E.g., could reduce number of run-time checks
- Possible negatives:
 - Must write the types (or use type inference)
 - False positives are possible (can be suppressed)



The Checker Framework

A framework for pluggable type checkers "Plugs" into the OpenJDK or OracleJDK compiler

javac -processor MyChecker ...

Standard error format allows tool integration



Ant, Maven, Gradle integration

```
cpresetdef name="jsr308.javac">
  <javac fork="yes"</pre>
   executable="${checkerframework}/checker/bin/${cfJavac}" >
    <!-- JSR-308-related compiler arguments -->
    <compilerarg value="-version"/>
    <compilerarg value="-implicit:class"/>
 </javac>
                                       <dependencies>
</presetdef>
                                         ... existing <dependency> items ...
                                         <!-- annotations from the Checker Framework:
                                                nullness, interning, locking, ... -->
                                           <dependency>
                                             <groupId>org.checkerframework
                                             <artifactId>checker-qual</artifactId>
                                             <version>1.9.7
                                           </dependency>
```

</dependencies>

Live demo: http://CheckerFramework.org/live/

Checker Framework Live Demo

Write Java code here:

```
import org.checkerframework.checker.nullness.qual.Nullable;
class YourClassNameHere {
    void foo(Object nn, @Nullable Object nbl) {
        nn.toString(); // OK
        nbl.toString(); // Error
    }
}
```

Choose a type system: Nullness Checker ▼

Check

Examples:

Nullness: <u>NullnessExample</u> | <u>NullnessExampleWithWarnings</u>

MapKey: MapKeyExampleWithWarnings

Interning: InterningExample | InterningExampleWithWarnings



Pluggable type-checking improves code

Checker Framework for creating type checkers

• Featureful, effective, easy to use, scalable

Prevent bugs at compile time

Create custom type-checkers

Improve your code!

http://CheckerFramework.org/





Comparison: other nullness tools

	Null pointer errors		False	Annotations
	Found	Missed	warnings	written
Checker Framework	9	0	4	35
FindBugs	0	9	1	0
Jlint	0	9	8	0
PMD	0	9	0	0
Eclipse, in 2017	0	9	8	0
Intellij (@NotNull	0	9	1	0
default), in 2017	3	6	1	925 + 8

51 Checking the Lookup program for file system searching (4kLOC)



Example type systems

```
Null dereferences (@Nullable)
   >200 errors in Google Collections, javac, ...
Equality tests (@Interned)
   >200 problems in Xerces, Lucene, ...
Concurrency / locking (@GuardedBy)
   >500 errors in BitcoinJ, Derby, Guava, Tomcat, ...
Fake enumerations / typedefs (@Fenum)
   problems in Swing, JabRef
```



String type systems

```
Regular expression syntax (@Regex)
   56 errors in Apache, etc.; 200 annos required
printf format strings (@Format)
   104 errors, only 107 annotations required
Method signature format (@FullyQualified)
   28 errors in OpenJDK, ASM, AFU
Compiler messages (@CompilerMessageKey)
   8 wrong keys in Checker Framework
```



Security type systems

Command injection vulnerabilities (@OsTrusted)

5 missing validations in Hadoop

Information flow privacy (@Source)

SPARTA detected malware in Android apps



It's easy to write your own type system!



Checkers are usable

- Type-checking is familiar to programmers
- Modular: fast, incremental, partial programs
- Annotations are not too verbose
 - @Nullable: 1 per 75 lines
 - @Interned: 124 annotations in 220 KLOC revealed 11 bugs
 - **@Format**: 107 annotations in 2.8 MLOC revealed 104 bugs
 - Possible to annotate part of program
 - Fewer annotations in new code
- Few false positives
- First-year CS majors preferred using checkers to not
- Practical: in use in Silicon Valley, on Wall Street, etc.



What a checker guarantees

The program satisfies the type property. There are:

- no bugs (of particular varieties)
- no wrong annotations
- Caveat 1: only for code that is checked
 - Native methods (handles reflection!)
 - Code compiled without the pluggable type checker
 - Suppressed warnings
 - Indicates what code a human should analyze
 - Checking part of a program is still useful
- Caveat 2: The checker itself might contain an error



Since Java 5: declaration annotations

Only for declaration locations:

```
@Deprecated
class Foo {
  @Getter @Setter private String query;
  @SuppressWarnings("unchecked")
  void foo() { ... }
```



But we couldn't express

A <u>non-null</u> reference to my data

An <u>interned</u> string

A <u>non-null</u> List of <u>English</u> strings

A <u>non-empty</u> array of <u>English</u> strings



With Java 8 Type Annotations we can!

```
A non-null reference to my data
   @NonNull Data mydata;
An interned String
   @Interned String query;
A non-null List of English Strings
   @NonNull List<@English String> msgs;
A non-empty array of English strings
   @English String @NonEmpty [] a;
```



Java 8 extends annotation syntax

Annotations on all occurrences of types:

```
@Untainted String query;
List<@NonNull String> strings;
myGraph = (@Immutable Graph) tmp;
class UnmodifiableList<T>
  implements @Readonly List<T> {}
```

Stored in classfile

Handled by javac, javap, javadoc, ...



Annotating external libraries

When type-checking clients, need library spec. Can write manually or automatically infer Two syntaxes:

- As separate text file (stub file)
- Within its .jar file (from annotated partial source code)



Checker Framework facilities

- Full type systems: inheritance, overriding, ...
- Generics (type polymorphism)
 - Also qualifier polymorphism
- Qualifier defaults
- Pre-/post-conditions
- Warning suppression



Verification

- Goal: prove that no bug exists
- Specifications: user provides
- False negatives: none
- False positives: user suppresses warnings
- Downside: user burden

Bug-finding

- **Goal**: find some bugs at low cost
- Specifications: infer likely specs
- False negatives: acceptable
- False positives: heuristics focus on most important bugs
- Downside: missed bugs

Meither is "better"; each is appropriate in certain circumstances.



Checker Framework Community

Open source project:

https://github.com/typetools/checker-framework

- Monthly release cycle
- >15,700 commits, 105 authors
- Welcoming & responsive community



Pluggable type-checking improves code

Checker Framework for creating type checkers

• Featureful, effective, easy to use, scalable

Prevent bugs at compile time

Create custom type-checkers

Improve your code!

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