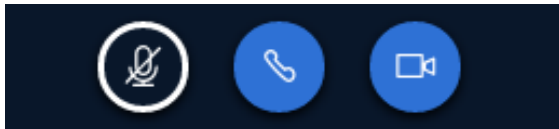


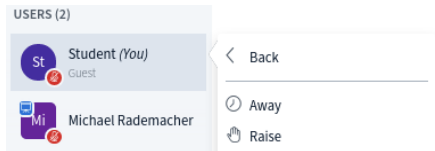
## Audio and Video

- **Find a suitable space:**
  - W/o. interruptions
  - W/o. background noise
  - Have good lighting on your face
- Use a **headset** if available
- **Turn ON your camera** unless you are experiencing connection problems
- **Turn OFF your microphone** and only unmute it if you want to participate



## Participation

- Click on your Name and “Set status” - “Raise”
- I will call your name
- Ask questions in the Shared Notes. Write “+1” if you have the same question.
- **Lectures (Audio, Video, all Chats) will be recorded.** If you do not like to appear in the recording, let me know



**Do NOT hesitate to participate!** You will get used to the system very quickly!

## Review of required background - Corona Quiz Edition!

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# Communication in Distributed Systems

2 Review of required background - **Corona Quiz Edition!**

 Dr.-Ing. Michael Rademacher

---

2020-04-28

## MANET Simulation Studies: The Incredibles \*

Stuart Kurkowski  
*skurkows@mines.edu*

Tracy Camp  
*tcamp@mines.edu*

Michael Colagrosso  
*mcolagro@mines.edu*

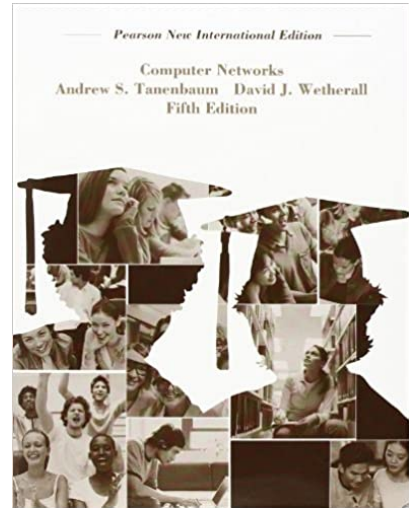
MCS Department, Colorado School of Mines, Golden, Colorado, USA

### Talking Points:

- What is the goal of this paper?
- What is the issue with “did not identify the simulator used in the research”
- What is an unbiased simulation study?
- What is a terminating and a steady-state simulation?
- What is the problem with Scenario Development? What is the main statement of Figure 2? How should we fix the problem?
- What is the main challenge when Analyzing the Output of a simulation?
- What did you learn for your paper?

# Today is Quiz Time!

- We will use the BigBlueButton poll feature for a little self-assessment
- The results of this self-assessment will **not** be relevant for your final grade
- The questions are **examples** for knowledge we expect you to have
- **Do Not Guess, do not Google**
- You need a pen and a paper
- I will add some explanations after each question when needed
- If you have problems answering the questions review some background by yourself



Andrew S Tanenbaum - Computer Networks

## Question

The \_\_\_\_\_ Layer handles transmission errors and regulates the flow of data on a point-to-point link.

A

Transport

B

Network

C

Data Link

D

Physical

## Answer

The \_\_\_\_\_ Layer handles transmission errors and regulates the flow of data on a point-to-point link.

**A**

Transport

**B**

Network

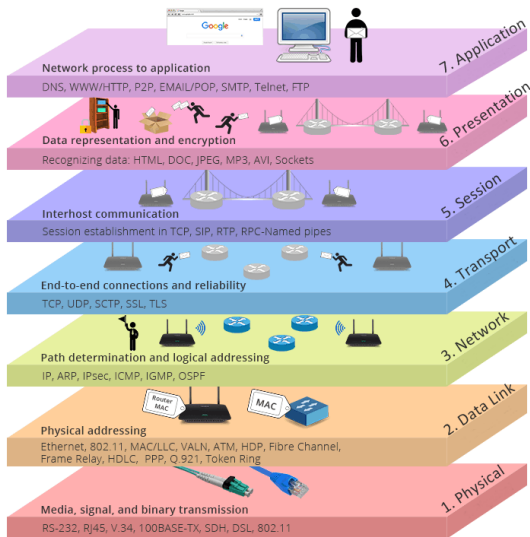
**C**

Data Link

**D**

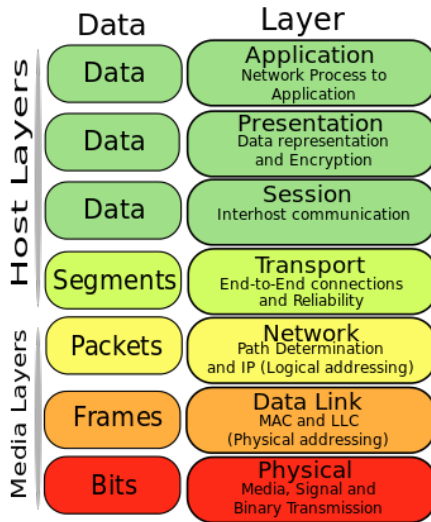
Physical

# Explanation



<https://securitystack.co/2019/02/08/security-stack-sheet-39/>

## OSI Model



<https://commons.wikimedia.org/wiki/File:Osi-model-jb.svg>



## Question

IPv4 is a \_\_\_\_ protocol.

**A**

reliable

**B**

connection-oriented

**C**

both A and B

**D**

none of the above

## Answer

IPv4 is a \_\_\_\_ protocol.

**A**

reliable

**B**

connection-oriented

**C**

both A and B

**D**

none of the above

## Reliability:

|         |   |   |   |   |   |   |   |   |   |                     |   |   |   |   |   |   |   |   |   |                 |   |   |   |   |   |   |   |   |   |                 |   |  |  |  |  |  |  |  |  |         |  |  |  |  |  |  |  |  |  |                 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|---------|---|---|---|---|---|---|---|---|---|---------------------|---|---|---|---|---|---|---|---|---|-----------------|---|---|---|---|---|---|---|---|---|-----------------|---|--|--|--|--|--|--|--|--|---------|--|--|--|--|--|--|--|--|--|-----------------|--|--|--|--|--|--|--|--|--|--|--|--|--|--|--|--|--|--|--|
| 0       |   |   |   |   |   |   |   |   |   | 1                   |   |   |   |   |   |   |   |   |   | 2               |   |   |   |   |   |   |   |   |   | 3               |   |  |  |  |  |  |  |  |  |         |  |  |  |  |  |  |  |  |  |                 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 0       | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0                   | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0               | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0               | 1 |  |  |  |  |  |  |  |  |         |  |  |  |  |  |  |  |  |  |                 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| Version |   |   |   |   |   |   |   |   |   | IHL                 |   |   |   |   |   |   |   |   |   | Type of Service |   |   |   |   |   |   |   |   |   | Total Length    |   |  |  |  |  |  |  |  |  |         |  |  |  |  |  |  |  |  |  |                 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|         |   |   |   |   |   |   |   |   |   | Identification      |   |   |   |   |   |   |   |   |   | Flags           |   |   |   |   |   |   |   |   |   | Fragment Offset |   |  |  |  |  |  |  |  |  |         |  |  |  |  |  |  |  |  |  |                 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|         |   |   |   |   |   |   |   |   |   | Time to Live        |   |   |   |   |   |   |   |   |   |                 |   |   |   |   |   |   |   |   |   | Protocol        |   |  |  |  |  |  |  |  |  |         |  |  |  |  |  |  |  |  |  | Header Checksum |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|         |   |   |   |   |   |   |   |   |   | Source Address      |   |   |   |   |   |   |   |   |   |                 |   |   |   |   |   |   |   |   |   |                 |   |  |  |  |  |  |  |  |  |         |  |  |  |  |  |  |  |  |  |                 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|         |   |   |   |   |   |   |   |   |   | Destination Address |   |   |   |   |   |   |   |   |   |                 |   |   |   |   |   |   |   |   |   |                 |   |  |  |  |  |  |  |  |  |         |  |  |  |  |  |  |  |  |  |                 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
|         |   |   |   |   |   |   |   |   |   | Options             |   |   |   |   |   |   |   |   |   |                 |   |   |   |   |   |   |   |   |   |                 |   |  |  |  |  |  |  |  |  | Padding |  |  |  |  |  |  |  |  |  |                 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

Example Internet Datagram Header

<https://tools.ietf.org/html/rfc791>

A **reliable** protocol **notifies** the sender whether the delivery of data to intended recipients was **successful**.

## Connection-oriented / Connection-less:

### ■ Connection-oriented:

- Usually operates in three distinct phases
  1. Connection setup to initialize a state
  2. Data unit transfer
  3. Connection release

### ■ Connection-less:

- Immediate data unit transfer
- No connection setup
- Packets can arrive out-of-order through different paths

## Question

In \_\_\_\_\_ transmission, the channel capacity is shared by both communicating devices at all times.

**A**

simplex

**B**

half-duplex

**C**

full-duplex

**D**

half-simplex

In \_\_\_\_\_ transmission, the channel capacity is shared by both communicating devices at all times.

**A**

simplex

**B**

half-duplex

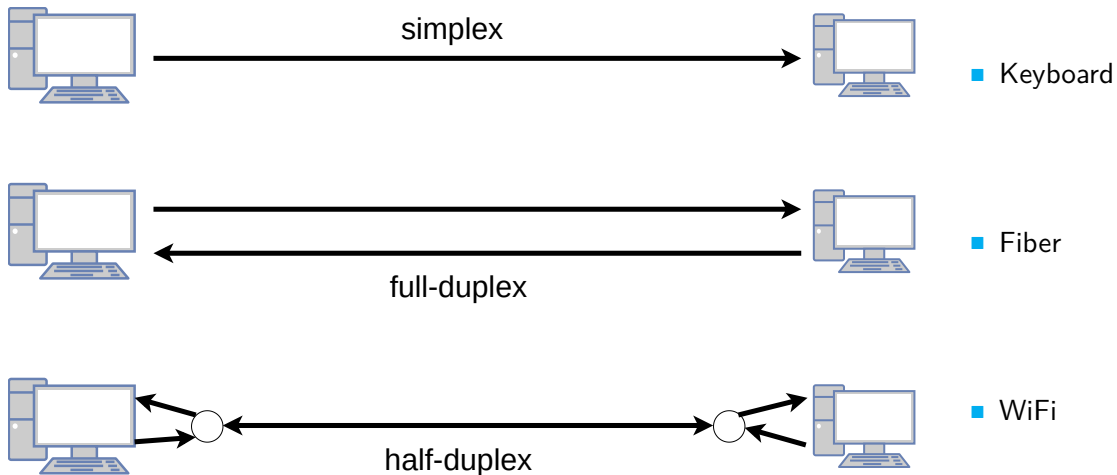
**C**

full-duplex

**D**

half-simplex

# Explanation



## Question

A client-server system uses a satellite network, with the satellite at a height of 30,000 km. What is the best-case round-trip delay in response to a request. Provide a reasonable estimation.

**A**

0.4 ms

**B**

4 ms

**C**

400 ms

**D**

4 s

## Answer

A client-server system uses a satellite network, with the satellite at a height of 30,000 km. What is the best-case round-trip delay in response to a request. Provide a reasonable estimation.

**A**

0.4 ms

**B**

4 ms

**C**

400 ms

**D**

4 s



## Explanation

$$2 * \frac{30\,000 \text{ km}}{3 \times 10^5 \text{ km/s}} = 2 * \frac{3 \text{ km}}{30 \text{ km/s}} = 2 * 0.1 \text{ s} = 0.2 \text{ s} = 200 \text{ ms}$$

| Route                        | Distance  | Time in vacuum | Time in fiber | RTT in fiber |
|------------------------------|-----------|----------------|---------------|--------------|
| New York<br>to San Francisco | 4,148 km  | 14 ms          | 21 ms         | 42 ms        |
| New York<br>to London        | 5,585 km  | 19 ms          | 28 ms         | 56 ms        |
| New York<br>to Sydney        | 15,993 km | 53 ms          | 80 ms         | 160 ms       |
| Equatorial<br>circumference  | 40,075 km | 133.7 ms       | 200 ms        | 200 ms       |

## Question

How long does it take to transmit a 1 Gigabyte<sup>a</sup> file using a link with a goodput of 50 Mbps.

<sup>a</sup>Assume the decimal definition, i.e. 1000 bytes = 1 kilobyte.

A

16 s

B

160 s

C

320 s

D

640 s

How long does it take to transmit a 1 Gigabyte<sup>a</sup> file using a link with a goodput of 50 Mbps.

<sup>a</sup>Assume the decimal definition, i.e. 1000 bytes = 1 kilobyte.

**A**

16 s

**B**

160 s

**C**

320 s

**D**

640 s

# Explanation

## Wired communication technologies

|             |            |                   |              |
|-------------|------------|-------------------|--------------|
| Fiber optic | PON        | 155 Mbps–2.5 Gbps | Up to 60 km  |
|             | WDM        | 40 Gbps           | Up to 100 km |
|             | SONET/SDH  | 10 Gbps           | Up to 100 km |
| DSL         | ADSL       | 1–8 Mbps          | Up to 5 km   |
|             | HDSL       | 2 Mbps            | Up to 3.6 km |
|             | VDSL       | 15–100 Mbps       | Up to 1.5 km |
| Coaxial     | DOCSIS     | 172 Mbps          | Up to 28 km  |
| Cable       |            |                   |              |
| PLC         | HomePlug   | 14–200 Mbps       | Up to 200 m  |
|             | Narrowband | 10–500 kbps       | Up to 3 km   |
| Ethernet    | 802.3x     | 10 Mbps–10 Gbps   | Up to 100 m  |

## Wireless communication technologies

|           |   |                       |              |
|-----------|---|-----------------------|--------------|
| Z-Wave    | Z-Wave                                  | 40 kbps               | Up to 30 m   |
| Bluetooth | 802.15.1                                | 721 kbps              | Up to 100 m  |
| ZigBee    | ZigBee                                  | 250 kbps              | Up to 100 m  |
|           | ZigBee Pro                              | 250 kbps              | Up to 1600 m |
|           | 802.11x                                 | 2–600 Mbps            | Up to 100 m  |
| WiMAX     | 802.16                                  | 75 Mbps               | Up to 50 km  |
| Wireless  | Various (e.g., RF mesh, 802.11, 802.15, | Depending on selected | Depending on |
| Mesh      | 802.16)                                 | protocols             | deployment   |
| Cellular  | 2G                                      | 14.4 kbps             | Up to 50 km  |
|           | 2.5G                                    | 144 kbps              |              |
|           | 3G                                      | 2 Mbps                |              |
|           | 3.5G                                    | 14 Mbps               |              |
|           | 4G                                      | 100 Mbps              |              |
| Satellite | Satellite Internet                      | 1 Mbps                | 100–6000 km  |

$$\begin{aligned}
 \text{time} &= \frac{\text{size}}{\text{goodput}} \\
 &= \frac{8 \times 10^9 \text{ Bit}}{50 \times 10^6 \text{ Bit/s}} \\
 &= 160 \text{ s}
 \end{aligned}$$

<https://doi.org/10.1016/j.comnet.2014.03.029>

## Question

The \_\_\_\_\_ identifies a process on a host.

**A**

physical address

**B**

IP address

**C**

port

**D**

none of the above

The \_\_\_\_\_ identifies a process on a host.

**A**

physical address

**B**

IP address

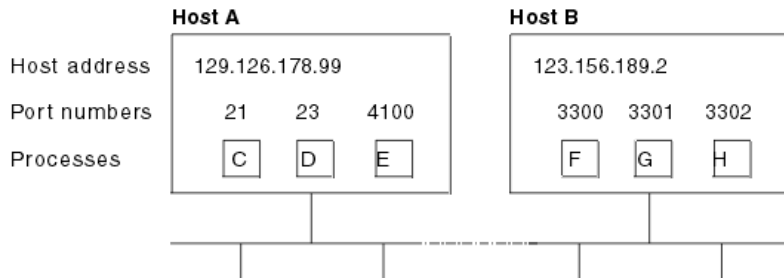
**C**

port

**D**

none of the above

# Explanation



[https://www.ibm.com/support/knowledgecenter/SSGMCP\\_5.6.0/fundamentals/graphics/DFHTM13.gif](https://www.ibm.com/support/knowledgecenter/SSGMCP_5.6.0/fundamentals/graphics/DFHTM13.gif)

| name   | port | protocol |
|--------|------|----------|
| http   | 80   | tcp      |
| ipp    | 631  | udp      |
| snmp   | 161  | udp      |
| ntp    | 123  | udp      |
| telnet | 23   | tcp      |
| https  | 443  | tcp      |
| ftp    | 21   | tcp      |
| smtp   | 25   | tcp      |
| pop3   | 110  | tcp      |

## Question

A periodic signal completes one cycle in 0.001 s. What is the frequency?

**A**

1 Hz

**B**

100 Hz

**C**

1 kHz

**D**

1 MHz



## Answer

A periodic signal completes one cycle in 0.001 s. What is the frequency?

**A**

1 Hz

**B**

100 Hz

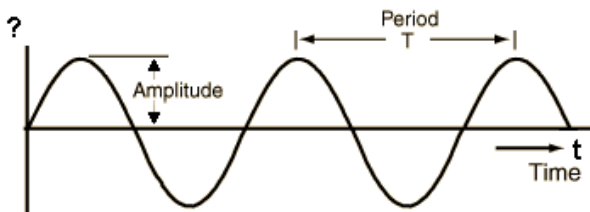
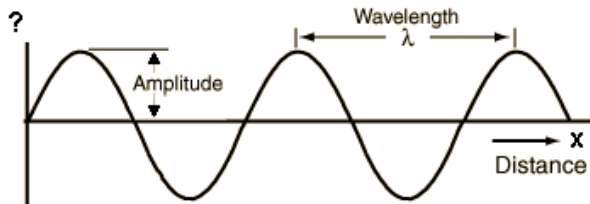
**C**

1 kHz

**D**

1 MHz

# Explanation



- Period:  $T = 1/f$
- Frequency  $f = 1/T = c/\lambda$
- Wavelength  $\lambda = c/f$
- Wave Speed  $c = \lambda f$

## Question

For a \_\_\_\_\_ channel, we need to use the Shannon–Hartley theorem to find the maximum bit rate.

**A**

noisy

**B**

noiseless

**C**

bandpass

**D**

low-pass

For a \_\_\_\_\_ channel, we need to use the Shannon–Hartley theorem to find the maximum bit rate.

**A**

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**C**

bandpass

**D**

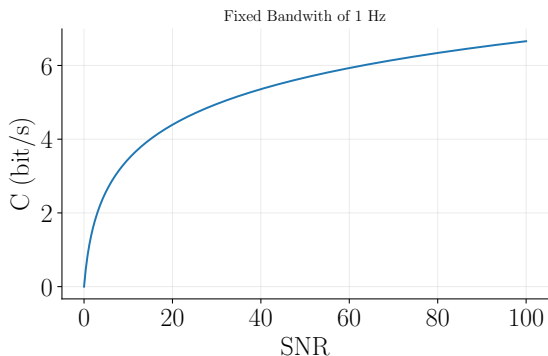
low-pass

$$C = B \log_2 \left( 1 + \frac{S_P}{N_0 * B} \right)$$

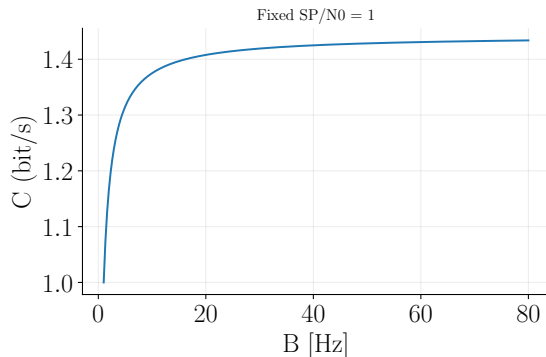
- C: Capacity (bits/s)
- B: Bandwidth (Hz)
- $N_0$ : Noise (W/Hz)
- $S_P$ : Signal power (W)
- $\frac{S_P}{N_0 * B}$ : SNR

- Increasing the **SNR** increases the channel capacity
- Increasing the **channel bandwidth** increases the channel capacity
- Depending on whether B or  $S_P$  is more precious, we can increase one and reduce the other, and yet maintain the same capacity
- **But, there are limits**

## Explanation 2



Limit of increasing the SNR



Limit of increasing the Bandwidth

- Observation: If  $B \rightarrow \infty$ ,  $C$  approaches  $\frac{SP}{N_0} \log_2(e)$ .
- Can you proof this?
- **Write me a mail if you can!**

## Question

The \_\_\_\_\_ address, also known as the link address or physical address, is the address of a node as defined by its LAN or WAN.

A

MAC

B

IP

C

port

D

specific

The \_\_\_\_\_ address, also known as the link address or physical address, is the address of a node as defined by its LAN or WAN.

**A**

MAC

**B**

IP

**C**

port

**D**

specific



## Question

The decibel is a way of describing a \_\_\_\_\_

**A**

logarithmic / ratio

**B**

linear / ratio

**C**

linear / sum

**D**

logarithmic / sum

## Answer

The decibel is a way of describing a \_\_\_\_\_

**A**

logarithmic / ratio

**B**

linear / ratio

**C**

linear / sum

**D**

logarithmic / sum

# Explanation

- **dB is a logarithmic ratio of powers**

$$\frac{P_1}{P_2} [dB] = 10 \log_{10} \left( \frac{P_1 [W]}{P_2 [W]} \right)$$

- $P_1 = 2 * P_2$  is equivalent to saying  $\frac{P_1}{P_2} \hat{=} 3dB$ , because

$$10 \log_{10} \left( \frac{P_1}{P_2} \right) [dB] = 10 \log_{10} \left( \frac{2 * P_2}{P_2} \right) = 10 \log_{10}(2) \approx 3 [dB]$$

- dBm and dBW are (logarithmic) units to measure powers
- The ratio of the power you want to express relative to a reference power
- 1W for dBW and 1mW for dBm

A table to remember:

| Power Ratio | in dB  |
|-------------|--------|
| 1/10        | -10 dB |
| 1/2         | -3 dB  |
| 1           | 0 dB   |
| 2           | 3 dB   |
| 10          | 10 dB  |
| 100         | 20 dB  |

- What is 36 dBm in Watt?
- $20 + 10 + 3 + 3$
- $100 * 10 * 2 * 2$
- $4000mW = 4W$

## Question

Using \_\_\_\_\_, the chance of collision can be reduced since a station senses the medium before trying to use it.

**A**

TDMA

**B**

CSMA

**C**

FDMA

**D**

CDMA

## Answer

Using \_\_\_\_\_, the chance of collision can be reduced since a station senses the medium before trying to use it.

**A**

TDMA

**B**

CSMA

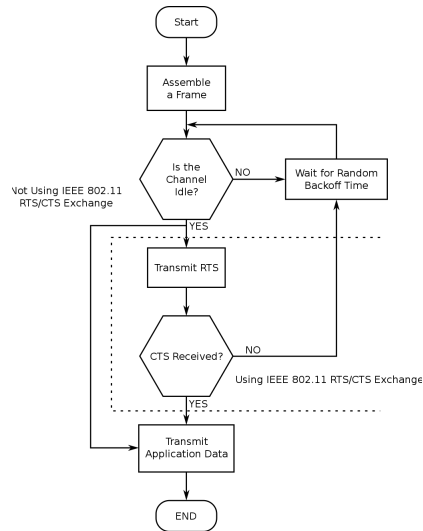
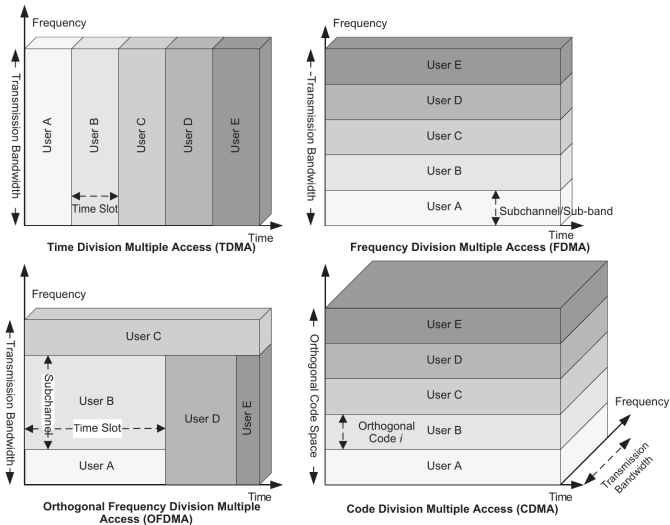
**C**

FDMA

**D**

CDMA

# Explanation



<https://doi.org/10.1016/B978-0-12-374964-2.10009-8>

[https://en.wikipedia.org/wiki/Carrier-sense\\_multiple\\_access\\_with\\_collision\\_avoidance](https://en.wikipedia.org/wiki/Carrier-sense_multiple_access_with_collision_avoidance)

## Question

IPv4 as best-effort delivery service includes \_\_\_\_\_.

**A**

error checking

**B**

error correction

**C**

datagram acknowledgment

**D**

none of the above

IPv4 as best-effort delivery service includes \_\_\_\_\_.

**A**

error checking

**B**

error correction

**C**

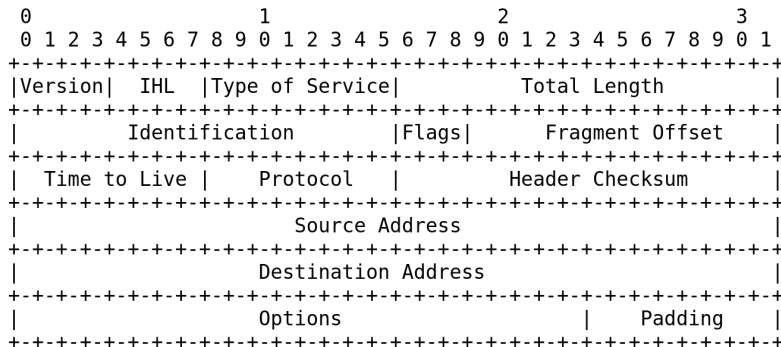
datagram acknowledgment

**D**

none of the above



# Explanation



Example Internet Datagram Header

<https://tools.ietf.org/html/rfc791>

“The type of service is used to specify the treatment of the datagram during its transmission through the internet system.”

## Question

UDP is called a \_\_\_\_\_ transport protocol.

**A**

connectionless, reliable

**B**

connection-oriented, unreliable

**C**

connectionless, unreliable

**D**

none of the above

UDP is called a \_\_\_\_\_ transport protocol.

**A**

connectionless, reliable

**B**

connection-oriented, unreliable

**C**

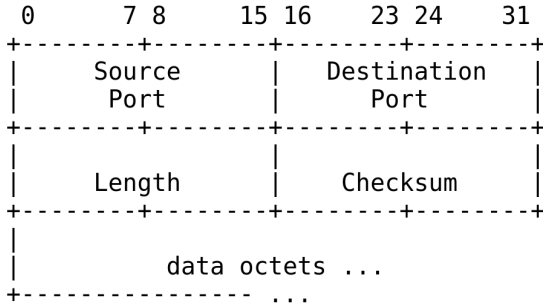
connectionless, unreliable

**D**

none of the above

# Explanation

- UDP is equal to “fire and forget”
- Ordering of packets is not guaranteed
- Lightweight with very few overhead



<https://tools.ietf.org/html/rfc768>

**To avoid Handshakes  
all TCP connections  
are converted to UDP**

**Stay safe - prevent  
COVID-19/Corona Virus**

## Question

TCP uses \_\_\_\_\_ to check the safe and sound arrival of data.

**A**

an acknowledgment mechanism

**B**

out-of-band signalling

**C**

the services of another protocol

**D**

none of the above

TCP uses \_\_\_\_\_ to check the safe and sound arrival of data.

**A**

an acknowledgment mechanism

**B**

out-of-band signalling

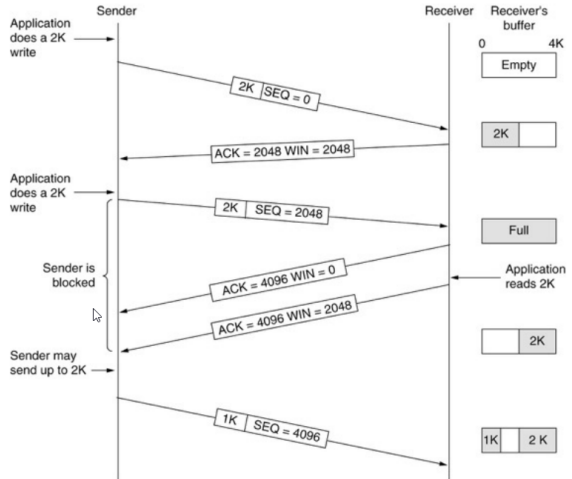
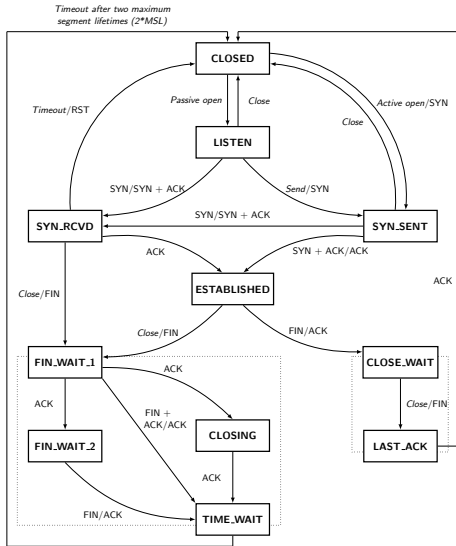
**C**

the services of another protocol

**D**

none of the above

# Explanation



# HTTP over UDP: an Experimental Investigation of QUIC

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Quavlive, Italy

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Quavlive, Italy

mascolo@poliba.it

**Abstract:** It is often said that the Internet is ubiquitous in our daily lives, but this holds true only for those who can easily access it. In fact, billions of people are still digitally disconnected, as bringing connectivity to certain zones does not make a good business case. The only solution for these unsatisfied potential users is to directly undertake the building of the infrastructure required to obtaining access to the Internet, typically forming groups in order to share the corresponding cost. This article presents a global classification and a summary of the main characteristics of different Alternative Network deployments that have arisen in recent years with an aim to provide Internet services in places where mainstream network deployments do not exist or are not adequate solutions. The "Global Access to the Internet for All" Research Group of the Internet Research Task Force, where all authors actively participate, is interested in documenting these emerging deployments. As an outcome of this work, a classification has converged by consensus, where five criteria have been identified and, based on them, four different types of Alternative Networks have been identified and described with real-world examples. Such a classification is useful for a deeper understanding of the common characteristics behind existing and emerging Alternative Networks.[1]





Thank you for your attention.  
Are there any questions left?



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Rathausallee 10  
Technopark  
Sankt Augustin



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www.mc-lab.de  
<https://michael-rademacher.net>

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*IEEE Commun. Mag.* (2017), 2–8.