

# CRC (Cyclic Redundancy Check) - Study Notes

## Overview

CRC is a **powerful error detection method** widely used in real-life applications. It's the most commonly used error detection technique in practical environments.

## Why CRC is Powerful

CRC can detect:

- **All odd errors** (any odd number of bit errors)
- **Single bit errors**
- **Double bit errors**
- **Burst errors** of length equal to polynomial degree

## Key Concepts

### Basic Formula

- **Total bits sent =  $m + r$** 
  - $m$  = number of message bits
  - $r$  = number of redundant bits (CRC bits)

## Steps to Calculate CRC

### Step 1: Determine Number of Redundant Bits ( $r$ )

#### Case 1: Polynomial Form Given

- If divisor is given as polynomial (e.g.,  $x^4 + x^3 + 1$ )
- Number of redundant bits = **highest degree of polynomial**
- Example:  $x^4 + x^3 + 1 \rightarrow$  degree = 4  $\rightarrow$  append 4 zeros

#### Case 2: Binary Form Given

- If divisor is given directly in binary (e.g., 11001)
- Number of redundant bits = **number of bits in divisor - 1**
- Example: 11001 has 5 bits  $\rightarrow$  append  $5-1 = 4$  zeros

### Step 2: Convert Polynomial to Binary (if needed)

#### Method: Extract coefficients

- Example:  $x^4 + x^3 + 1$

- Coefficients:  $x^4(1)$ ,  $x^3(1)$ ,  $x^2(0)$ ,  $x^1(0)$ ,  $x^0(1)$
- Binary form: **11001**

### Step 3: Append Zeros to Message

- Append the required number of zeros to the original message
- Example: Message 1010011010 + 4 zeros = **10100110100000**

### Step 4: Perform Binary Division Using XOR

#### XOR Rules:

- Same values  $\rightarrow 0$  ( $1 \oplus 1 = 0$ ,  $0 \oplus 0 = 0$ )
- Different values  $\rightarrow 1$  ( $1 \oplus 0 = 1$ ,  $0 \oplus 1 = 1$ )

#### Division Process:

1. Always start from leading 1
2. Use XOR operation (not regular division)
3. Continue until no more bits can be processed
4. **Quotient is ignored** - only remainder matters

### Step 5: Replace Appended Zeros with Remainder

- Take the **last r bits** of the remainder (LSB side)
- Replace the appended zeros with these remainder bits
- This creates the **valid codeword** to be transmitted

## Worked Example

#### Given:

- Message: 1010011010
- Divisor polynomial:  $x^4 + x^3 + 1$

#### Solution:

1. Convert divisor:  $x^4 + x^3 + 1 \rightarrow 11001$
2. Append 4 zeros: 1010011010**0000**
3. Divide using XOR operations
4. Remainder: 000010  $\rightarrow$  Take last 4 bits: **0010**
5. Final codeword: 1010011010**0010**

## Error Detection at Receiver

1. Receiver divides the entire received codeword by same divisor
2. **If remainder = 0** → No error detected
3. **If remainder ≠ 0** → Error detected

## Efficiency Calculation

**Formula:**  $\text{Efficiency} = (m/(m+r)) \times 100\%$

Where:

- m = message bits
- r = redundant bits

**Example:** 10 message bits + 4 redundant bits =  $(10/14) \times 100\% = 71.43\%$

## Important Tips for Exams

### Common Mistakes to Avoid

- Don't confuse polynomial degree with number of binary digits
- Always start division from leading 1
- Take remainder from LSB side, not MSB
- Remember: quotient is not used in CRC

### Quick Method for Division

- Always ensure first bit is 1 before applying divisor
- Carry bits from above when needed to maintain divisor length
- Use XOR operations consistently

### Key Points to Remember

1. **Polynomial given** → append (highest degree) zeros
2. **Binary given** → append (number of bits - 1) zeros
3. **Division** → use XOR, start with leading 1
4. **Remainder** → take last r bits from LSB side
5. **Detection** → remainder = 0 means no error

## Applications

- Most widely used in real-life networking applications
- Standard method in computer networks
- Used in data storage systems

- Common in communication protocols
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*Note: CRC numerical problems are very common in GATE and UGC-NET exams. Master the calculation steps as theory questions are rare (99% numerical, 1% theory).*