

## Reliable data transfer with TCP

- pipeline errors
  - detection
  - handling

**Note:** Many of the lecture slides are based on presentations that accompany *Computer Networking: A Top Down Approach*, 6<sup>th</sup> edition, by Jim Kurose & Keith Ross, Addison-Wesley, 2013.

# TCP sender events

## Data comes down from application layer:

- Create segments with sequence numbers
  - sequence # is byte-stream number of first data byte in segment
- Start countdown timer (if not already running for a previous segment)
- Send segments

## If countdown timer expires:

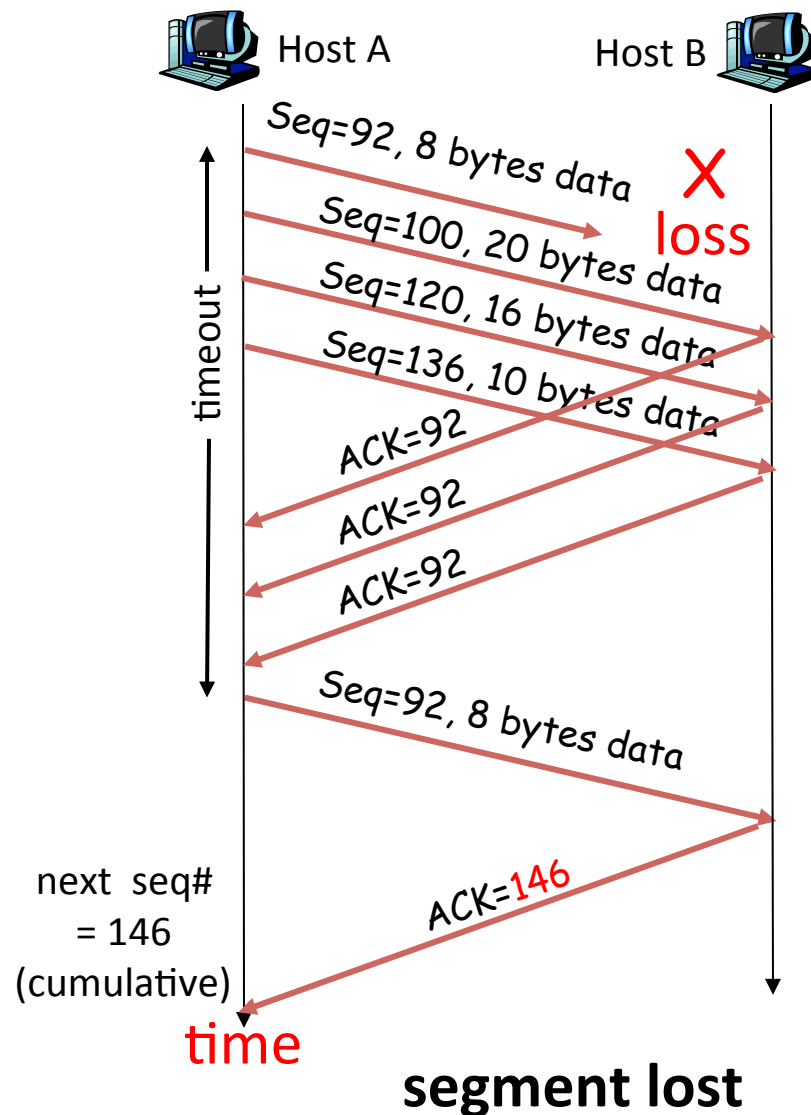
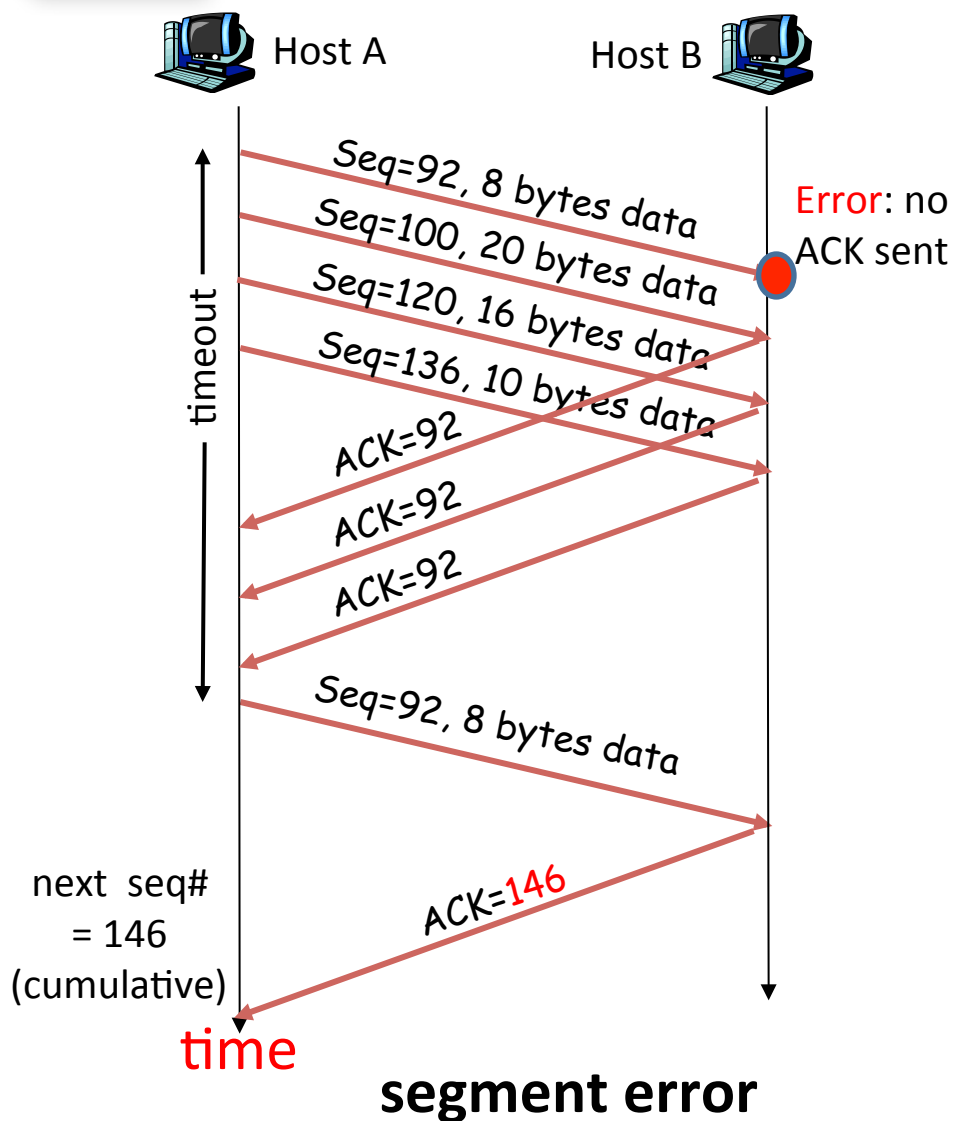
- Retransmit segment that caused timeout
- Restart timer

## If ACK received:

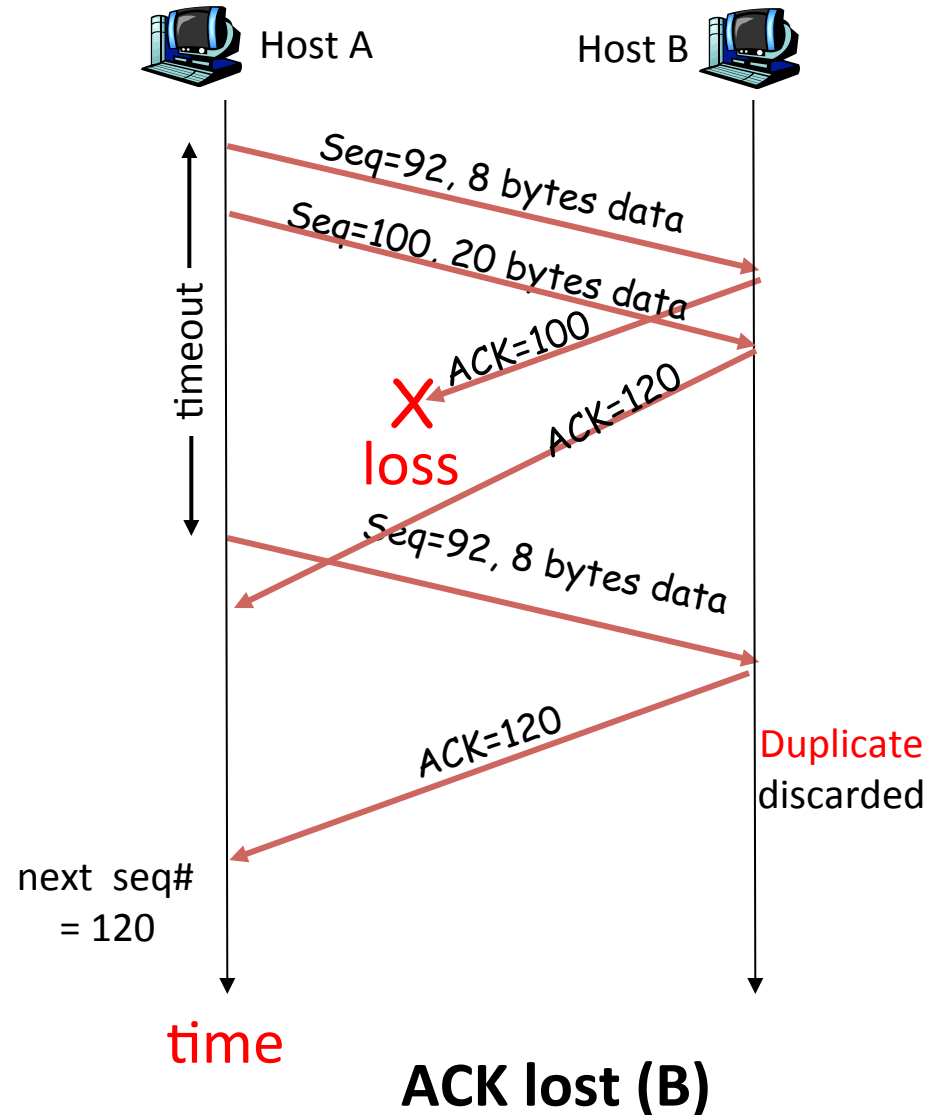
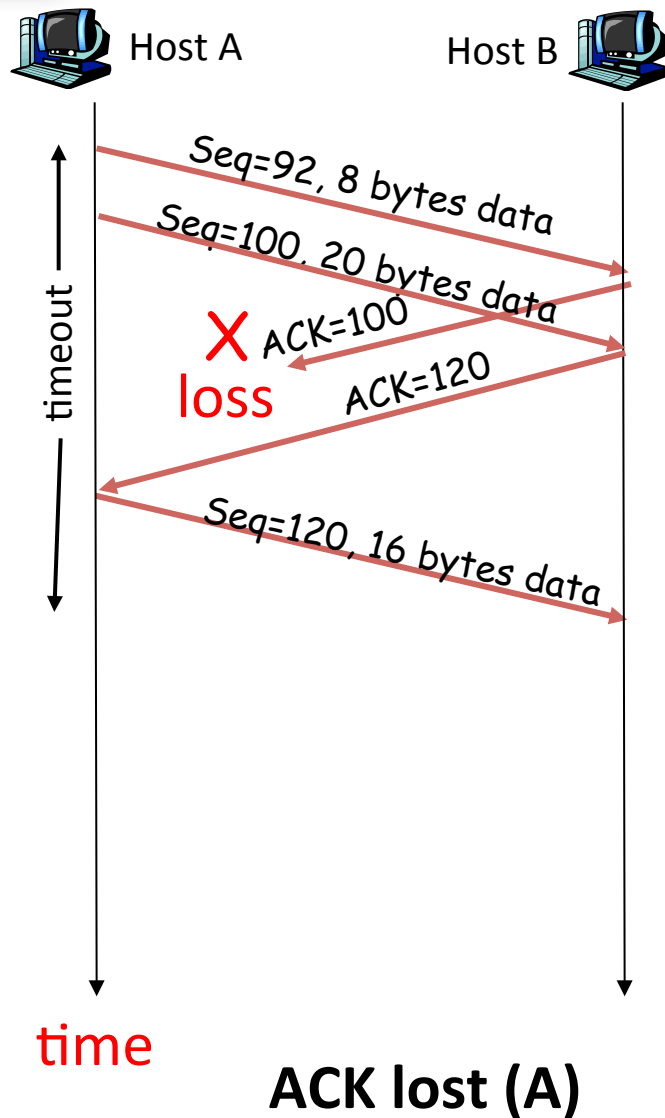
- Check to see if ACK includes previously unACK'ed segments
  - update what is known to be ACK'ed
  - restart timer if there are outstanding segments

Note: ACK's are cumulative. The ACK # is always the next expected byte number. This implies that all previous bytes have been accounted for.

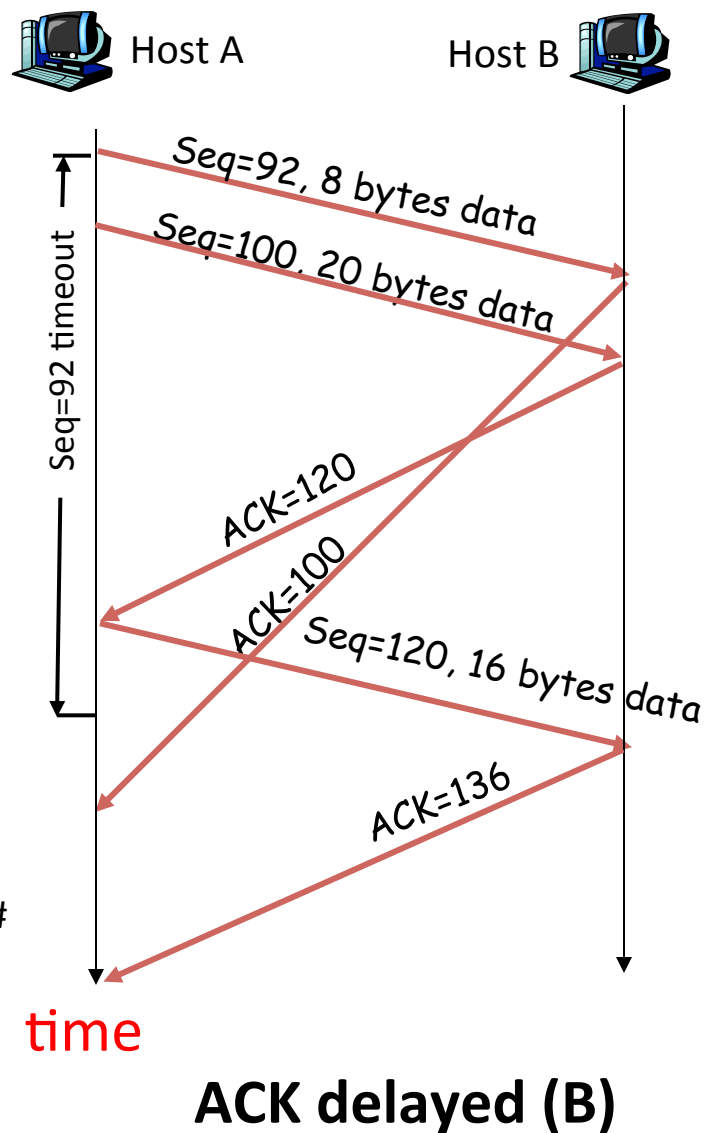
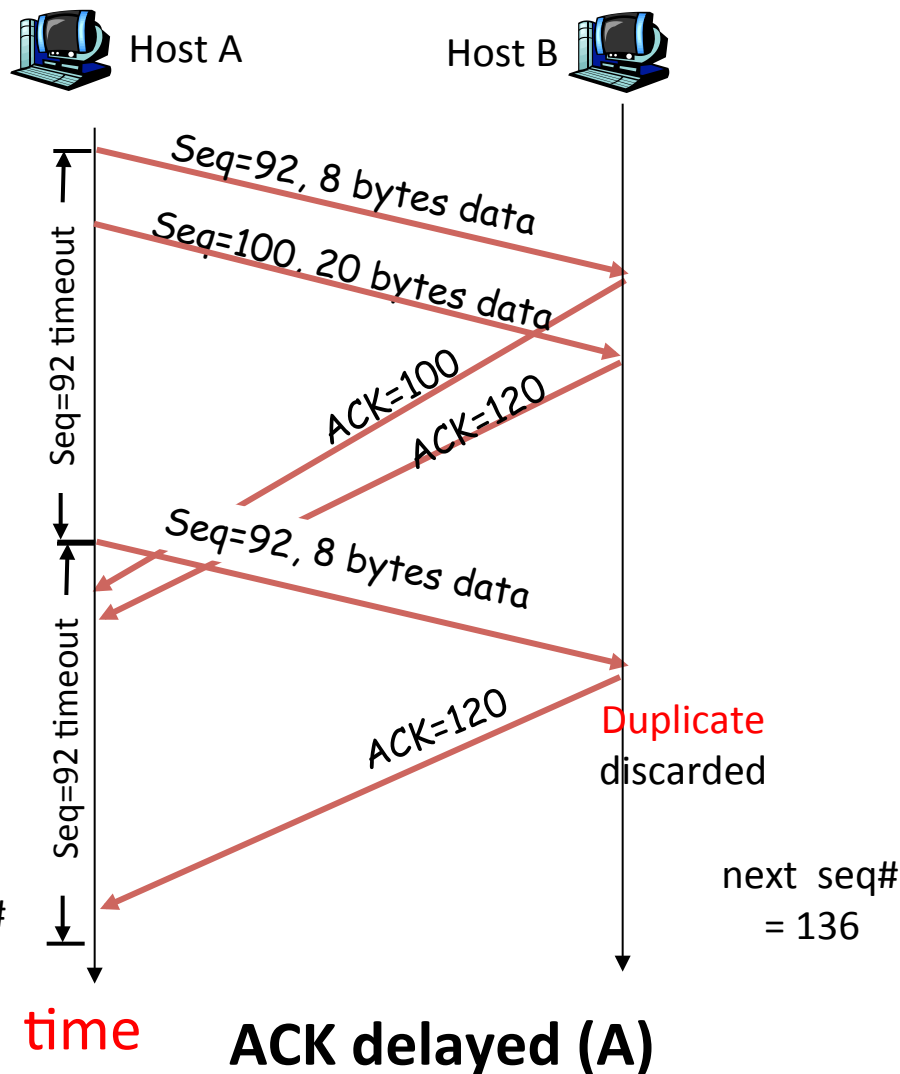
# Pipeline retransmission error scenarios (cumulative ACK)



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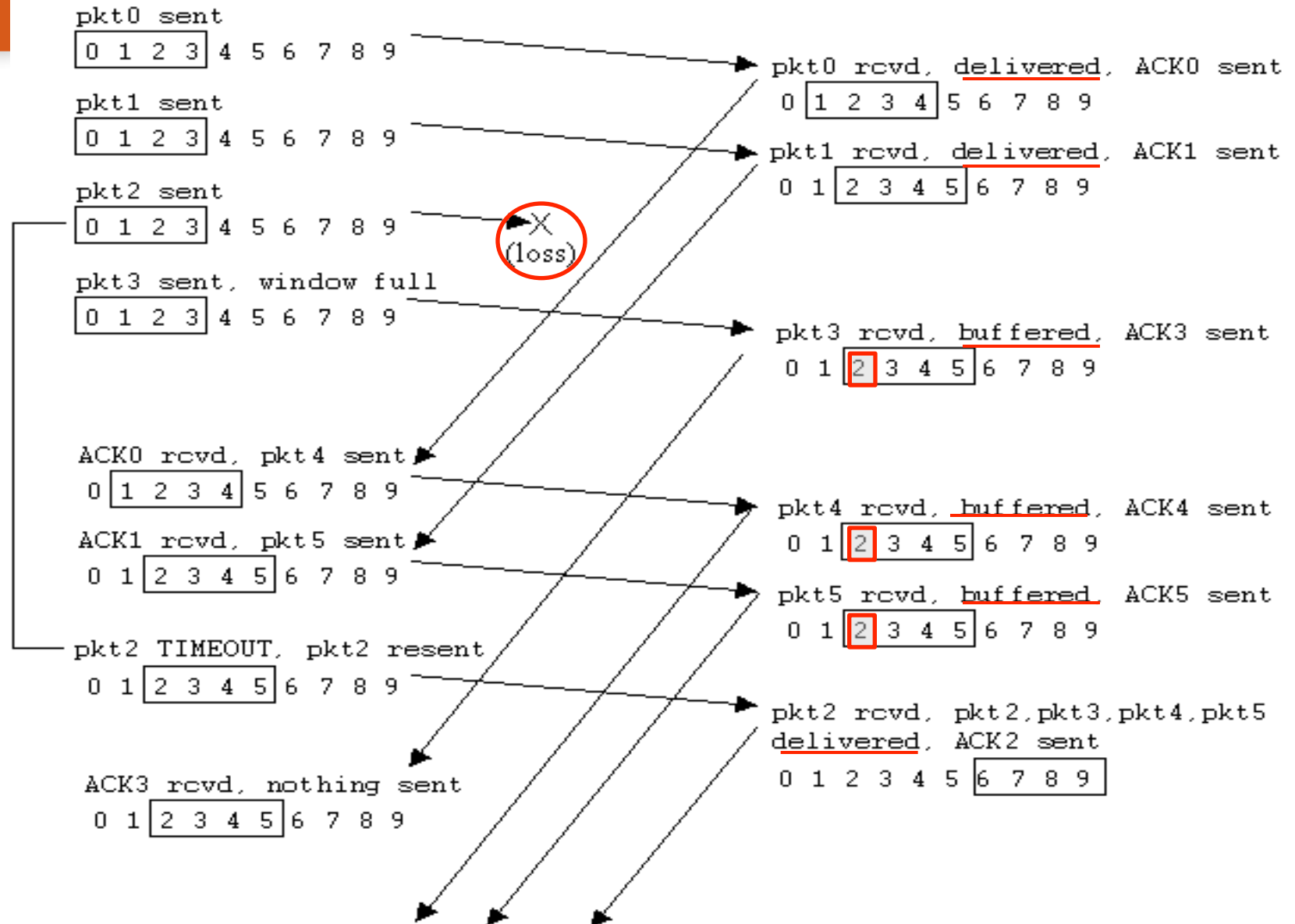
# Pipeline retransmission error scenarios (cumulative ACK)



# Selective Repeat protocol

- Sender can have up to “window size” un-ACK’ ed packets in pipeline
- Receiver ACKs individual packets
- Sender maintains timer for each un-ACK’ ed packet
  - When timer expires, retransmit only the un-ACK’ ed packet

# Selective Repeat protocol example

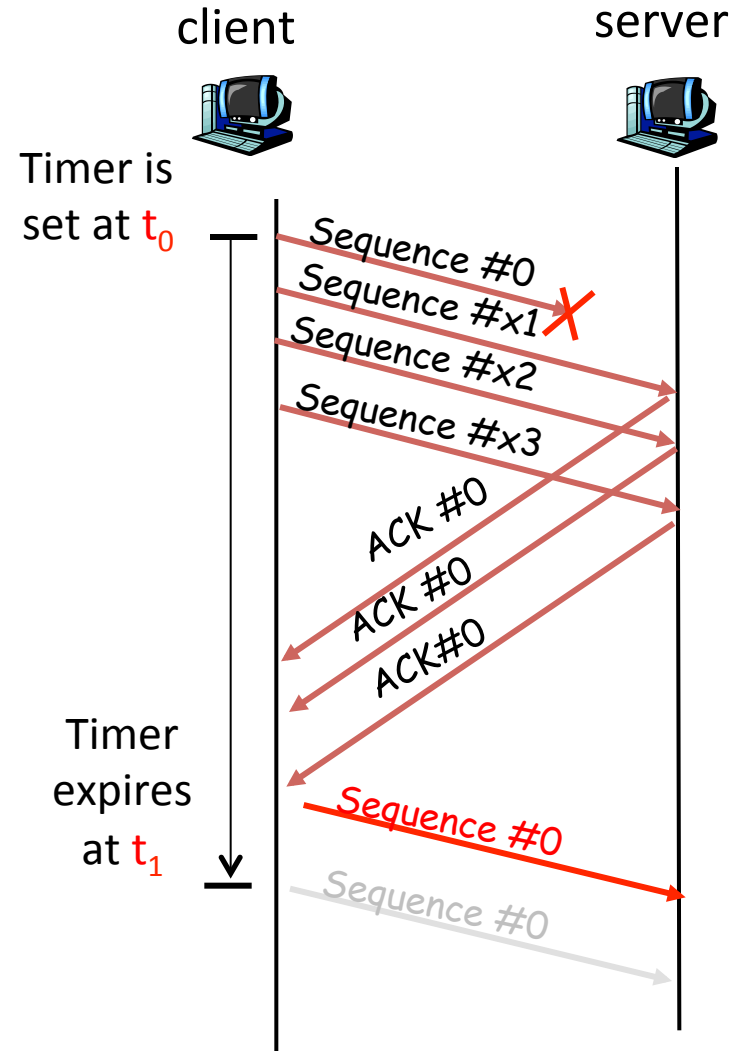


# TCP Fast Retransmit

- Suppose that the packet with sequence #0 gets lost
  - **Q:** When will the packet with sequence #0 get retransmitted?
  - **A:** typically at  $t_1$ . We think it is lost when the timer expires
- Can we do better??
  - Why wait till timeout?
    - We already know the packet is lost.

Remember:

- **Sequence#** is the number (in the data stream) of the first byte of the sent segment
- **ACK#** is the number of the next byte expected by the receiver.



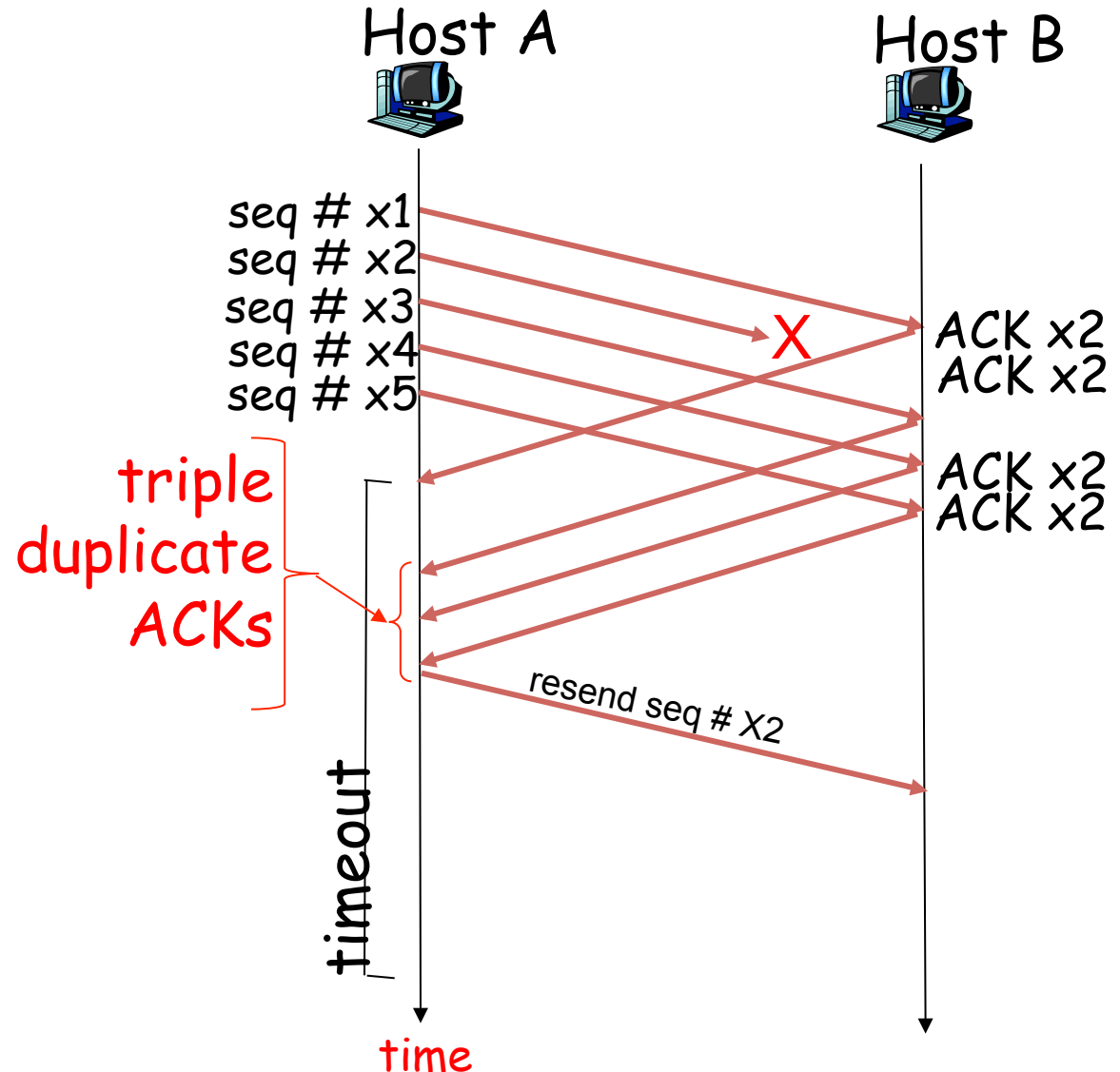


# TCP Fast Retransmit

If we receive many duplicate ACKs for sequence #x2 ...

... it means packet with sequence #x2 is lost.

Fast retransmit => better performance



# TCP Fast Retransmit recap

- Receipt of duplicate ACKs indicate loss of segments
  - Sender usually pipelines segments
  - If segment is lost, there will likely be many duplicate ACKs.

## This is how TCP works:

- If sender receives 3 ACKs for the same data, it supposes that segment after ACK'ed data was lost
- fast retransmit:
  - resend segment before timer expires
  - better performance

- TCP
  - cumulative ACKS
  - count-down timer
  - handling segment errors, segment loss
  - handling lost ACK, delayed ACK
- Retransmission protocols
  - Selective-Repeat
  - fast retransmission
- Next lecture: review for midterm exam