L5
SQL SQL SQL SQL SQL SQL SQL

Eugene Wu

Fall 2015

Didn't Lecture 3 Go Over SQL?

Two sublanguages

DDL Data Definition Language

define and modify schema (physical, logical, view)

CREATE TABLE, Integrity Constraints

DML Data Manipulation Language

get and modify data

simple SELECT, INSERT, DELETE

human-readable language

Gritty Details

DDL

NULL, Views

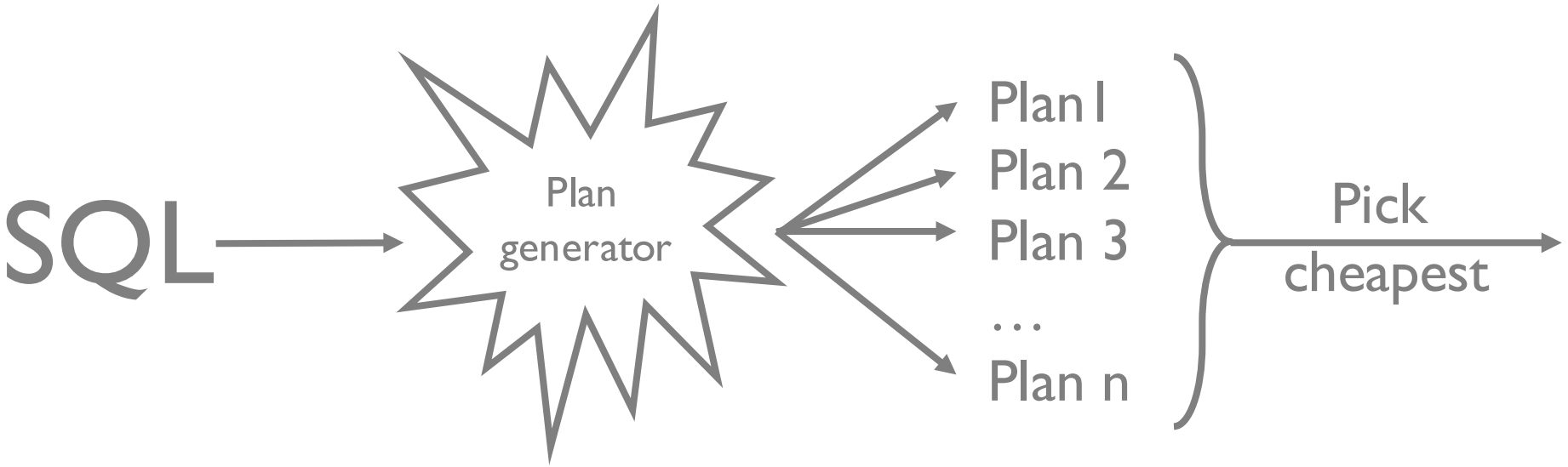
DML

Basics, SQL Clauses, Expressions, Joins, Nested Queries, Aggregation, With, Triggers

Didn't Lecture 3 Go Over SQL?

DBMS makes it run efficiently

- Key: precise query semantics
- Reorder/modify queries while answers stay same
- DBMS estimates costs for different evaluation plans



Didn't Lecture 3 Go Over SQL?

More expressive power than Rel Alg

- can be described by extensions of algebra
- One key difference: multisets rather than sets
- i.e. # duplicates in a table carefully accounted for

Most widely used *query language*, not just relational query language

Today's Database

Sailors

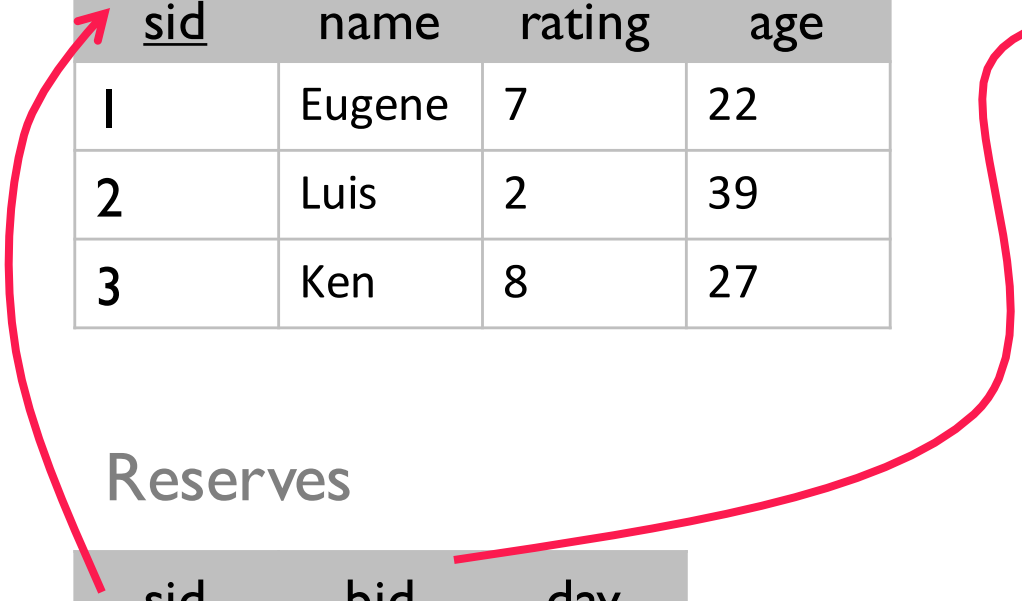
sid	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Boats

bid	name	color
101	Legacy	red
102	Melon	blue
103	Mars	red

Reserves

sid	bid	day
1	102	9/12
2	102	9/13
2	103	9/14



Is Reserves table correct?

Today's Database

Sailors

sid	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Boats

bid	name	color
101	Legacy	red
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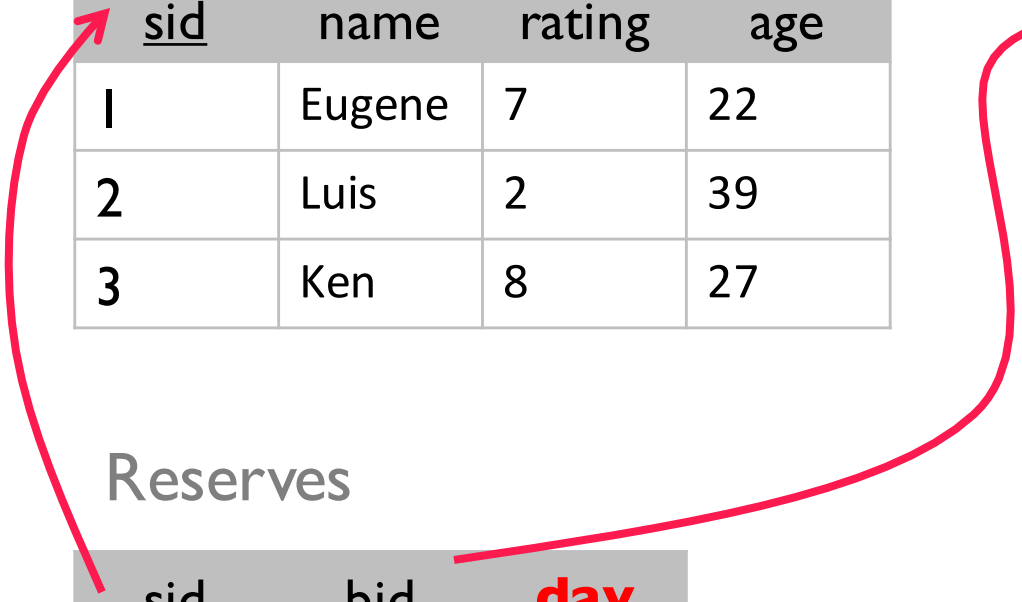
Reserves

sid	bid	day
1	102	9/12
2	102	9/13
2	103	9/14



Is Reserves table correct?
Day should be part of key

Follow along at home!

http://w4111db1.cloudapp.net:8000/

or

```
psql -U demo -h w4111db1.cloudapp.net demo
password: demo
```

<30 year old sailors

```
SELECT *
FROM Sailors
WHERE age < 30
```

sid	name	rating	age
1	Eugene	7	22
3	Ken	8	27

```
SELECT name, age
FROM Sailors
WHERE age < 30
```

name	age
Eugene	22
Ken	27

<30 year old sailors

```
SELECT *
FROM Sailors
WHERE age < 30
```

$\sigma_{age<30}$ (Sailors)

```
SELECT name, age
FROM Sailors
WHERE age < 30
```

$\pi_{name,\ age}$ ($\sigma_{age<30}$ (Sailors))

Multiple Relations

```
SELECT  S.name
FROM    Sailors AS S, Reserves AS R
WHERE   S.sid = R.sid AND R.bid = 102
```

Sailors

sid	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

sid	bid	day
1	102	9/12
2	102	9/13
2	103	9/14

Multiple Relations

```
SELECT S.name
FROM   Sailors AS S, Reserves AS R
WHERE  S.sid = R.sid AND R.bid = 102
```

$$\pi_{name}\,(\sigma_{bid=2}(\text{Sailors} \bowtie_{sid} \text{Reserves}))$$

Structure of a SQL Query

DISTINCT
Optional, answer should not have duplicates
Default: duplicates not removed (multiset)

target-list
List of expressions over attrs of tables in relation-list

```
SELECT   [DISTINCT] target-list
FROM     relation-list
WHERE    qualification
```

relation-list
List of relation names
Can define range-variable "AS X"

qualification
Boolean expressions
- Combined w/ AND, OR, NOT
- attr *op* const
- $attr_1$ *op* $attr_2$
- op is =, <, >, !=, etc

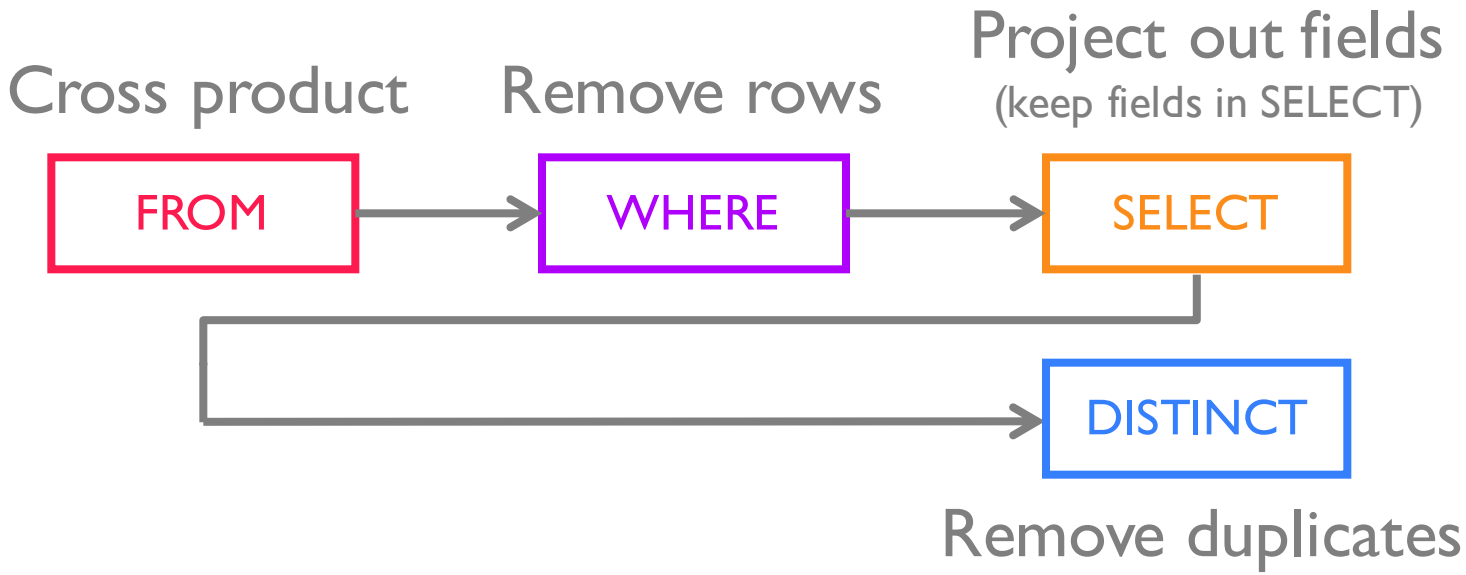
Semantics

SELECT [DISTINCT] *target-list*
FROM *relation-list*
WHERE *qualification*

FROM	compute cross product of relations
WHERE	remove tuples that fail qualifications
SELECT	remove fields not in target-list
DISTINCT	remove duplicate rows

Conceptual Query Evaluation

SELECT [DISTINCT] *target-list*
FROM *relation-list*
WHERE *qualification*
GROUP BY grouping-list
HAVING group-qualification



Not how actually executed! Above is likely very slow

DISTINCT (vol.1)

Reserves

sid	bid	day
1	102	9/12
2	102	9/13
2	103	9/14

```
SELECT   bid
FROM     Reserves
```

bid
102
102
103

```
SELECT   DISTINCT bid
FROM     Reserves
```

bid
102
103

Sailors that reserved 1+ boats

```
SELECT   S.sid
FROM     Sailors AS S, Reserves AS R
WHERE    S.sid = R.sid
```

Would DISTINCT change anything in this query?

What if SELECT clause was SELECT S.name?

Range Variables

Disambiguate relations

same table used multiple times (self join)

```
SELECT   sid
FROM     Sailors, Sailors
WHERE    age > age
```

```
SELECT   S1.sid
FROM     Sailors AS S1, Sailors AS S2
WHERE    S1.age > S2.age
```

Range Variables

Disambiguate relations

same table used multiple times (self join)

```
SELECT   sid
FROM     Sailors, Sailors
WHERE    age > age
```

```
SELECT   S1.name, S1.age, S2.name, S2.age
FROM     Sailors AS S1, Sailors AS S2
WHERE    S1.age > S2.age
```

Expressions (Math)

```
SELECT   S.age, S.age – 5 AS age2, 2*S.age AS age3
FROM     Sailors AS S
WHERE    S.name = ‘eugene’
```

```
SELECT   S1.name AS name1, S2.name AS name2
FROM     Sailors AS S1, Sailors AS S2
WHERE    S1.rating*2 = S2.rating - 1
```

Expressions (Strings)

```
SELECT   S.name
FROM     Sailors AS S
WHERE    S.name LIKE 'e_%'
```

'_' any one character (. in regex)

'%' 0 or more characters of any kind (.* in regex)

Most DBMSes have rich string manipulation support e.g., regex

PostgreSQL documentation

http://www.postgresql.org/docs/9.1/static/functions-string.html

Expressions (Date/Time)

```
SELECT   R.sid
FROM     Reserves AS R
WHERE    now() - R.date < interval ‘1 day’
```

TIMESTAMP, DATE, TIME types

now() returns timestamp at start of transaction

DBMSes provide rich time manipulation support

exact support may vary by vender

Postgresql Documentation

http://www.postgresql.org/docs/9.1/static/functions-datetime.html

Expressions

Constant	1
Col reference	Sailors.name
Arithmetic	Sailors.sid * 10
Unary operators	NOT, EXISTS
Binary operators	AND, OR, IN
Function calls	abs(), sqrt(), ...
Casting	1.7::int, '10-12-2015'::date

sid of Sailors that reserved red or blue boat

```
SELECT   R.sid
FROM     Boats B, Reserves R
WHERE    B.bid = R.bid AND
         (B.color = ‘red’ OR B.color = ‘blue’)
```

OR

```
SELECT   R.sid
FROM     Boats B, Reserves R
WHERE    B.bid = R.bid AND B.color = ‘red’
UNION ALL
SELECT   R.sid
FROM     Boats B, Reserves R
WHERE    B.bid = R.bid AND B.color = ‘blue’
```

sid of Sailors that reserved red or blue boat

```
SELECT   DISTINCT R.sid
FROM     Boats B, Reserves R
WHERE    B.bid = R.bid AND
         (B.color = ‘red’ OR B.color = ‘blue’)
```

OR

```
SELECT   R.sid
FROM     Boats B, Reserves R
WHERE    B.bid = R.bid AND B.color = ‘red’
UNION
SELECT   R.sid
FROM     Boats B, Reserves R
WHERE    B.bid = R.bid AND B.color = ‘blue’
```

sid of Sailors that reserved red and blue boat

```
SELECT   R.sid
FROM     Boats B, Reserves R
WHERE    B.bid = R.bid AND
         (B.color = ‘red’ AND B.color = ‘blue’)
```

```
SELECT   R.sid
FROM     Boats B, Reserves R
WHERE    B.bid = R.bid AND B.color = ‘red’
INTERSECT ALL
SELECT   R.sid
FROM     Boats B, Reserves R
WHERE    B.bid = R.bid AND B.color = ‘blue’
```

sid of Sailors that reserved red and blue boat

Can use self-join instead

```
SELECT   R.sid
FROM     Boats B1, Reserves R1
WHERE
         B1.bid = R1.bid AND

         B1.color = ‘red’
```

sid of Sailors that reserved red and blue boat

Can use self-join instead

```
SELECT   R.sid
FROM     Boats B1, Reserves R1, Boats B2,Reserves R2
WHERE
         B1.bid = R1.bid AND

         B1.color = ‘red’
```

sid of Sailors that reserved red and blue boat

Can use self-join instead

```
SELECT   R.sid
FROM     Boats B1, Reserves R1, Boats B2,Reserves R2
WHERE
         B1.bid = R1.bid AND
         B2.bid = R2.bid AND
         B1.color = ‘red’ AND B2.color = ‘blue’
```

sid of Sailors that reserved red and blue boat

Can use self-join instead

```
SELECT   R.sid
FROM     Boats B1, Reserves R1, Boats B2,Reserves R2
WHERE    R1.sid = R2.sid AND
         B1.bid = R1.bid AND
         B2.bid = R2.bid AND
         B1.color = ‘red’ AND B2.color = ‘blue’
```

sids of sailors that haven't reserved a boat

```
SELECT   S.sid
FROM     Sailors S

EXCEPT

SELECT   S.sid
FROM     Sailors S, Reserves R
WHERE    S.sid = R.sid
```

Can we write EXCEPT using more basic functionality?

SET Comparison Operators

UNION, INTERSECT, EXCEPT

EXISTS, NOT EXISTS

IN, NOT IN

UNIQUE, NOT UNIQUE

op ANY, *op* ALL

$op \in \{ <, >, =, \leq, \geq, \neq, \ldots \}$

Many of these rely on Nested Query Support

Nested Queries

```
SELECT   S.sid
FROM     Sailors S
WHERE    S.sid IN (SELECT    R.sid
                   FROM      Reserves R
                   WHERE     R.bid = 101)
```

Many clauses can contain SQL queries

WHERE, FROM, HAVING, SELECT

Conceptual model:

```
for each Sailors tuple
   run the subquery and evaluate qualification
```

Nested Correlated Queries

```
SELECT   S.sid
FROM     Sailors S
WHERE    EXISTS (SELECT   *
                 FROM     Reserves R
                 WHERE    R.bid = 101 AND
                          S.sid = R.sid)
```

Outer table referenced in nested query

Conceptual model:

```
for each Sailors tuple
   run the subquery and evaluate qualification
```

Nested Correlated Queries

```
SELECT   S.sid
FROM     Sailors S
WHERE    UNIQUE (SELECT   *
                 FROM     Reserves R
                 WHERE    R.bid = 101 AND
                          S.sid = R.sid)
```

UNIQUE checks that there are no duplicates

What does this do?

Nested Correlated Queries

```
SELECT   S.sid
FROM     Sailors S
WHERE    UNIQUE (SELECT   s.sid
                 FROM     Reserves R
                 WHERE    R.bid = 101 AND
                          S.sid = R.sid)
```

UNIQUE checks that there are no duplicates

What does this do?

Sailors whose rating is greater than any sailor named “Bobby”

```
SELECT S1.name
FROM   Sailors S1
WHERE  S1.rating > ANY (SELECT   S2.rating
                        FROM     Sailors S2
                        WHERE    S2.name = ‘Bobby’)
```

What about this?

```
SELECT S1.name
FROM   Sailors S1
WHERE  S1.rating > ALL (SELECT   S2.rating
                        FROM     Sailors S2
                        WHERE    S2.name = ‘Bobby’)
```

Rewrite INTERSECT using IN

```
SELECT  S.sid
FROM    Sailors  S
WHERE   S.rating  > 2

INTERSECT

SELECT  R.sid
FROM    Reserves  R
```

```
SELECT  S.sid
FROM    Sailors  S
WHERE   S.rating  > 2 AND
        S.sid  IN (
    SELECT  R.sid
    FROM    Reserves  R
)
```

Similar trick for EXCEPT → NOT IN

What if want *names* instead of sids?

Sailors that reserved all boats (Division)

Hint: double negation

reserved all boats == no boat w/out reservation

```
SELECT    S.name
FROM      Sailors S
WHERE     NOT EXISTS (

    (SELECT   B.bid FROM   Boats B)

    EXCEPT

    (SELECT   R.bid
     FROM Reserves R
     WHERE R.sid = S.sid)
)
```

Sailors that reserved all boats (Division)

Hint: double negation

reserved all boats == no boat w/out reservation

```
SELECT  S.name
FROM    Sailors  S
WHERE   NOT EXISTS (
```

Sailors S such that

There's no boat without

A reservation by S

Sailors that reserved all boats (Division)

Hint: double negation

reserved all boats == no boat w/out reservation

```
SELECT  S.name
FROM    Sailors  S
WHERE   NOT EXISTS (SELECT B.bid
                    FROM   Boats  B
                    WHERE  NOT EXISTS (
```

Sailors S such that

There's no boat without

A reservation by S

Sailors that reserved all boats (Division)

Hint: double negation

reserved all boats == no boat w/out reservation

```
SELECT  S.name
FROM    Sailors  S
WHERE   NOT  EXISTS  (SELECT B.bid
                      FROM   Boats  B
                      WHERE  NOT  EXISTS  (SELECT R.bid
                                           FROM Reserves  R
                                           WHERE  R.sid  =  S.sid))
```

Sailors S such that

There's no boat without

A reservation by S

Serious people can count: Aggregation

```
SELECT  COUNT(*)
FROM    Sailors S

SELECT  AVG(S.age)
FROM    Sailors S
WHERE   S.rating = 10

SELECT COUNT(DISTINCT S.name)
FROM    Sailors S
WHERE   S.name LIKE 'D%'

SELECT  S.name
FROM    Sailors
WHERE   S.rating = (SELECT MAX(S2.rating)
                    FROM   Sailors S2)
```

COUNT([DISTINCT] A)

SUM([DISTINCT] A)

AVG([DISTINCT] A)

MAX/MIN(A)

STDDEV(A)

CORR(A,B)

PostgreSQL documentation

http://www.postgresql.org/docs/9.4/static/functions-aggregate.html

Name and age of oldest sailor(s)

```
SELECT  S.name, MAX(S.age)
FROM    Sailors S
```

```
SELECT  S.name, S.age
FROM    Sailors S
WHERE   S.age >= ALL (SELECT   S2.age
                      FROM     Sailors S2)
```

```
SELECT  S.name, S.age
FROM    Sailors S
WHERE   S.age = (SELECT    MAX(S2.age)
                 FROM      Sailors S2)
```

```
SELECT  S.name, S.age
FROM    Sailors S
ORDER BY   S.age DESC
LIMIT 1
```

← When does this not work?

NULL

Field values sometimes unknown or inapplicable

SQL provides a special value *null* for such situations.

The presence of null complicates many issues e.g.,

Is age = null true or false?

Is null = null true or false?

Is null = 8 OR 1 = 1 true or false?

3 Valued Logic (true, false, unknown)

Special syntax "IS NULL" and "IS NOT NULL"

How does WHERE remove rows?

if qualification doesn't evaluate to true

New operators (in particular, outer joins) possible/needed.

NULL

(null > 0) = null

(null + 1) = null

(null = 0) = null

(null AND true) = null

null is null = true

Some truth tables

AND	T	F	NULL
T	T	F	NULL
F	F	F	F
NULL	NULL	F	NULL

OR	T	F	NULL
T	T	T	T
F	T	F	NULL
NULL	T	NULL	NULL

JOINS

```
SELECT (column_list)
FROM table_name
    [INNER | {LEFT |RIGHT | FULL } {OUTER}] JOIN table_name
    ON qualification_list
WHERE ...
```

INNER is default

Difference in how to deal with NULL values

PostgreSQL documentation:

http://www.postgresql.org/docs/9.4/static/tutorial-join.html

Inner/Natural Join

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors S, Reserves r
WHERE   s.sid = r.sid

SELECT s.sid, s.name, r.bid
FROM    Sailors s INNER JOIN Reserves r
ON      s.sid = r.sid

SELECT  s.sid, s.name, r.bid
FROM    Sailors s NATURAL JOIN Reserves r
```

All Equivalent!

Natural Join means equi-join for each pair of attrs with same name

Sailor names and their reserved boat ids

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors s INNER JOIN Reserves r
ON      s.sid = r.sid
```

Sailors

sid	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

sid	bid	day
1	102	9/12
2	102	9/13

Result

sid	name	bid
1	Eugene	102
2	Luis	102

Sailor names and their reserved boat ids

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors s INNER JOIN Reserves r
ON      s.sid = r.sid
```

Sailors

sid	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

sid	bid	day
1	102	9/12
2	102	9/13

Result

sid	name	bid
1	Eugene	102
2	Luis	102

Prefer INNER JOIN over NATURAL JOIN. Why?

Sailor names and their reserved boat ids

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors s INNER JOIN Reserves r
ON      s.sid = r.sid
```

Sailors

sid	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

sid	bid	day
1	102	9/12
2	102	9/13

Result

sid	name	bid
1	Eugene	102
2	Luis	102

Notice: No result for Ken!

Left Outer Join (or No Results for Ken)

Returns all matched rows *and all unmatched rows from table on left of join clause*

(*at least one* row for each row in left table)

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors s LEFT OUTER JOIN Reserves r
ON      s.sid = r.sid
```

All sailors & bid for boat in their reservations

Bid set to NULL if no reservation

Left Outer Join

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors s LEFT OUTER JOIN Reserves r
ON      s.sid = r.sid
```

Sailors

sid	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

sid	bid	day
1	102	9/12
2	102	9/13

Result

sid	name	bid
1	Eugene	102
2	Luis	102
3	Ken	NULL

Right Outer Join

Same as LEFT OUTER JOIN, but guarantees result for rows in table on right side of JOIN

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors s RIGHT OUTER JOIN Reserves r
ON      s.sid = r.sid
```

FULL OUTER JOIN

Returns all matched *or* unmatched rows from both sides of JOIN

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors s FULL OUTER JOIN Reserves r
ON      s.sid = r.sid
```

FULL OUTER JOIN

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors s Full OUTER JOIN Reserves r
ON      s.sid = r.sid
```

Sailors

sid	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

sid	bid	day
1	102	9/12
2	102	9/13
4	109	9/20

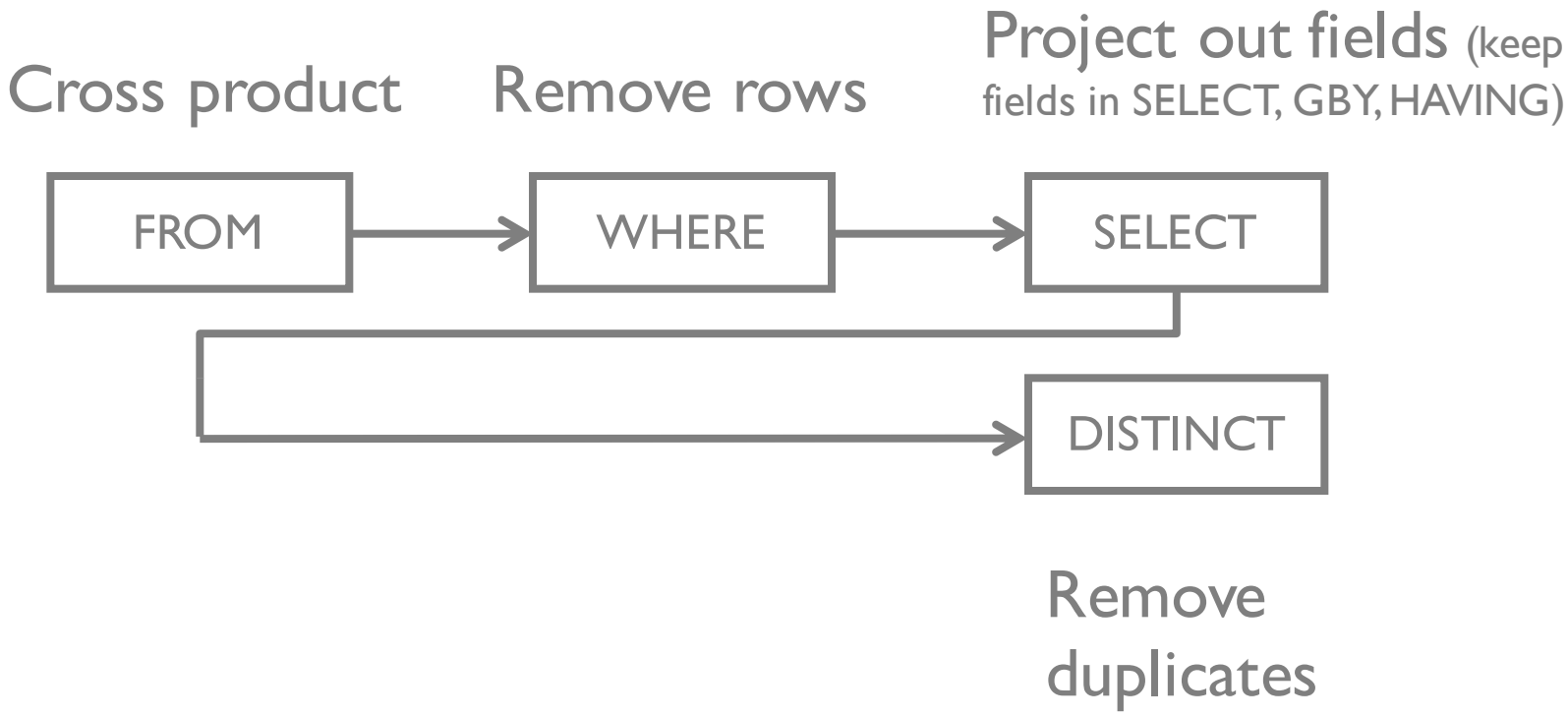
Result

sid	name	bid
1	Eugene	102
2	Luis	102
3	Ken	NULL
NULL	NULL	4

Why is sid NULL?

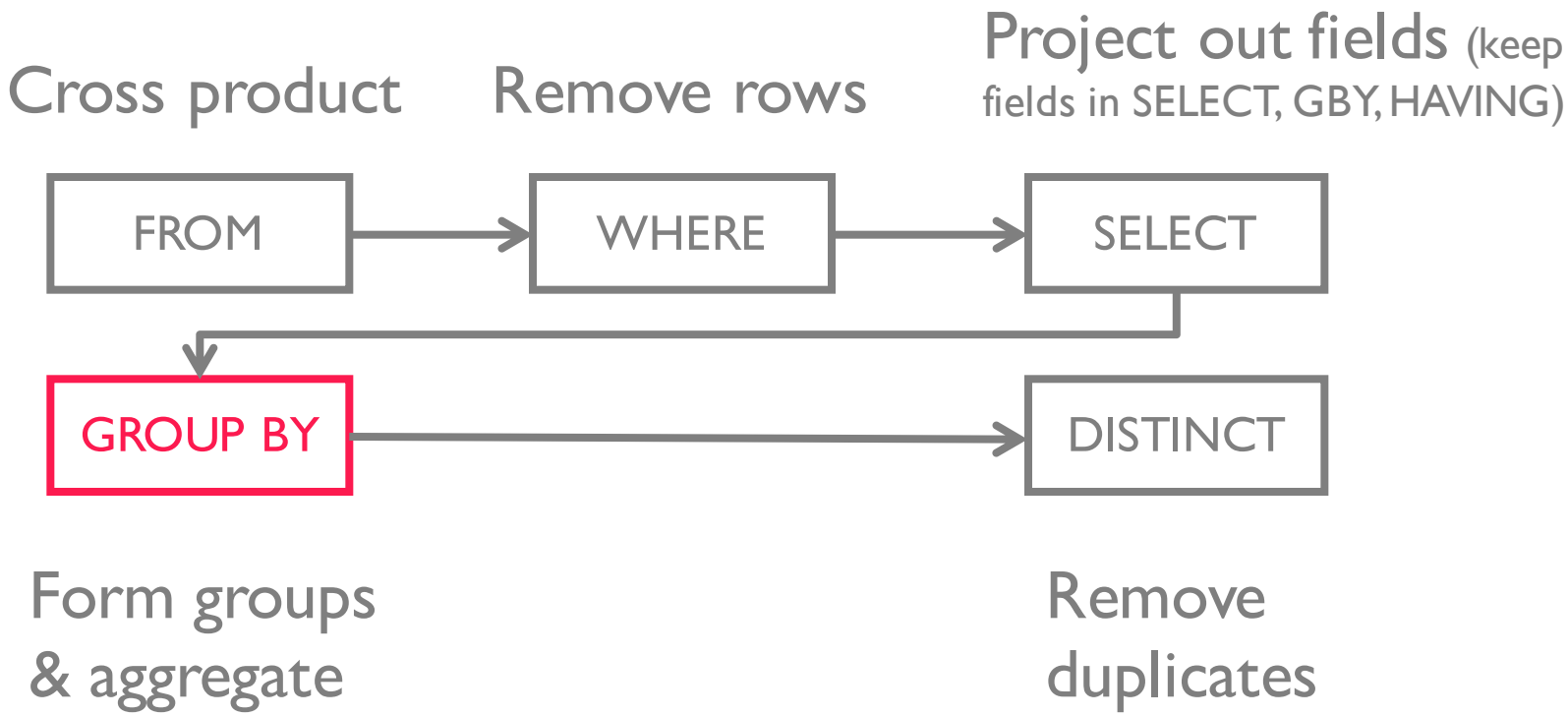
Conceptual Query Evaluation

SELECT [DISTINCT] *target-list*
FROM *relation-list*
WHERE *qualification*
GROUP BY grouping-list
HAVING group-qualification



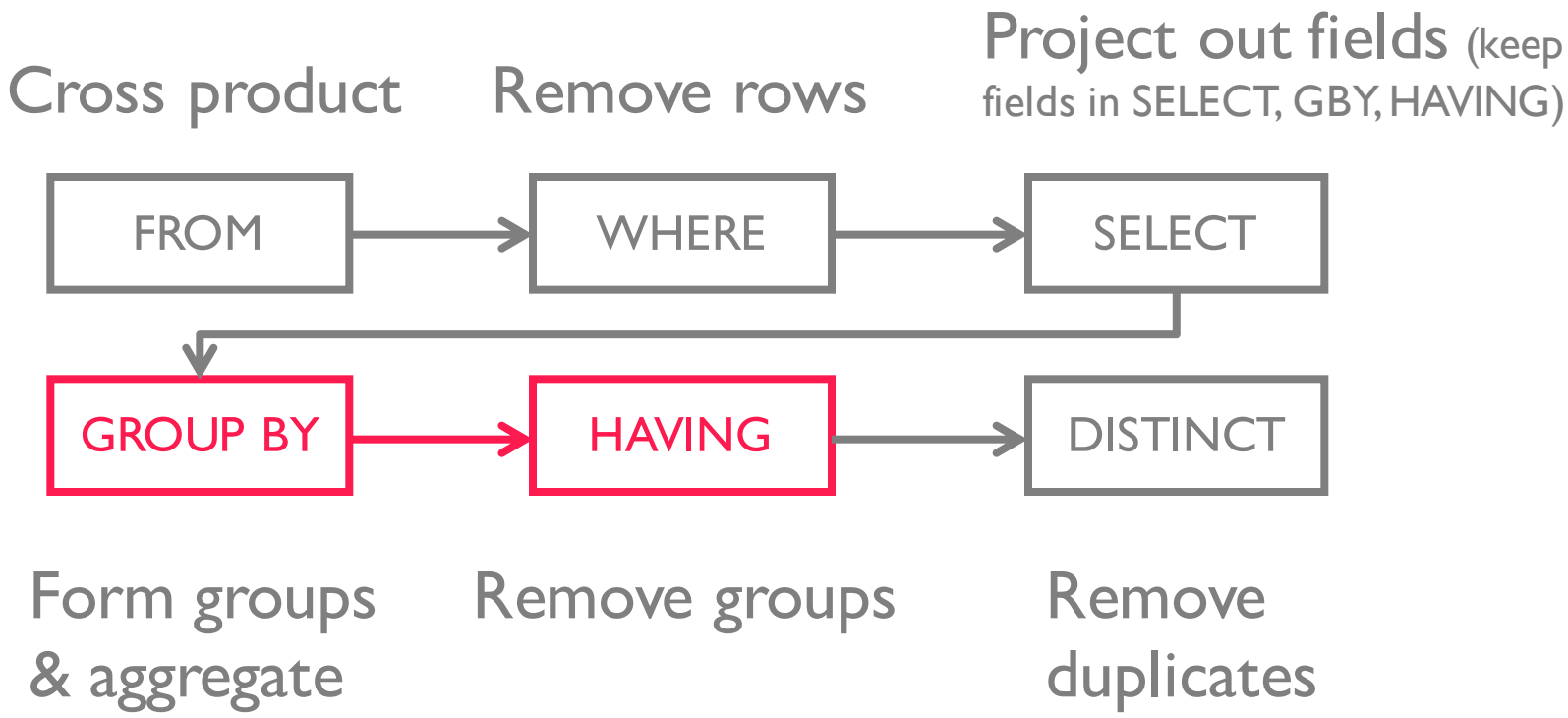
Conceptual Query Evaluation

SELECT [DISTINCT] *target-list*
FROM *relation-list*
WHERE *qualification*
GROUP BY *grouping-list*
HAVING *group-qualification*



Conceptual Query Evaluation

SELECT [DISTINCT] *target-list*
FROM *relation-list*
WHERE *qualification*
GROUP BY *grouping-list*
HAVING *group-qualification*



GROUP BY

```
SELECT min(s.age)
FROM   Sailors s
```

Minimum age among all sailors

What if want min age *per rating level*?

We don't even know how many rating levels exist!

If we did, could write (awkward):

```
for rating in [0…10]
    SELECT min(s.age)
    FROM   Sailors s
    WHERE  s.rating = <rating>
```

GROUP BY

```
SELECT count(*)
FROM   Reserves R
```

Total number of reservations

What if want reservations per boat?

May not even know all our boats (depends on data)!

If we did, could write (awkward):

```
for boat in [0…10]
    SELECT count(*)
    FROM   Reserves R
    WHERE  R.bid = <boat>
```

GROUP BY

```
SELECT    [DISTINCT] target-list
FROM      relation-list
WHERE     qualification
GROUP BY  grouping-list
HAVING    group-qualification
```

Target-list contains

attribute-names $\subseteq$ *grouping-list*

aggregation expressions

grouping-list is a list of expressions that defines groups

set of tuples w/ same value for all attributes in grouping-list

GROUP BY

```
SELECT    bid, count(*)
FROM      Reserves R
GROUP BY  bid
```

Minimum age for each rating

```
SELECT    bid, count(*)
FROM      Reserves R
GROUP BY  bid
HAVING    count(*) > 1
```

Minimum age for each boat with more than 1 reservation

HAVING

group-qualification used to remove groups

similar to WHERE clause

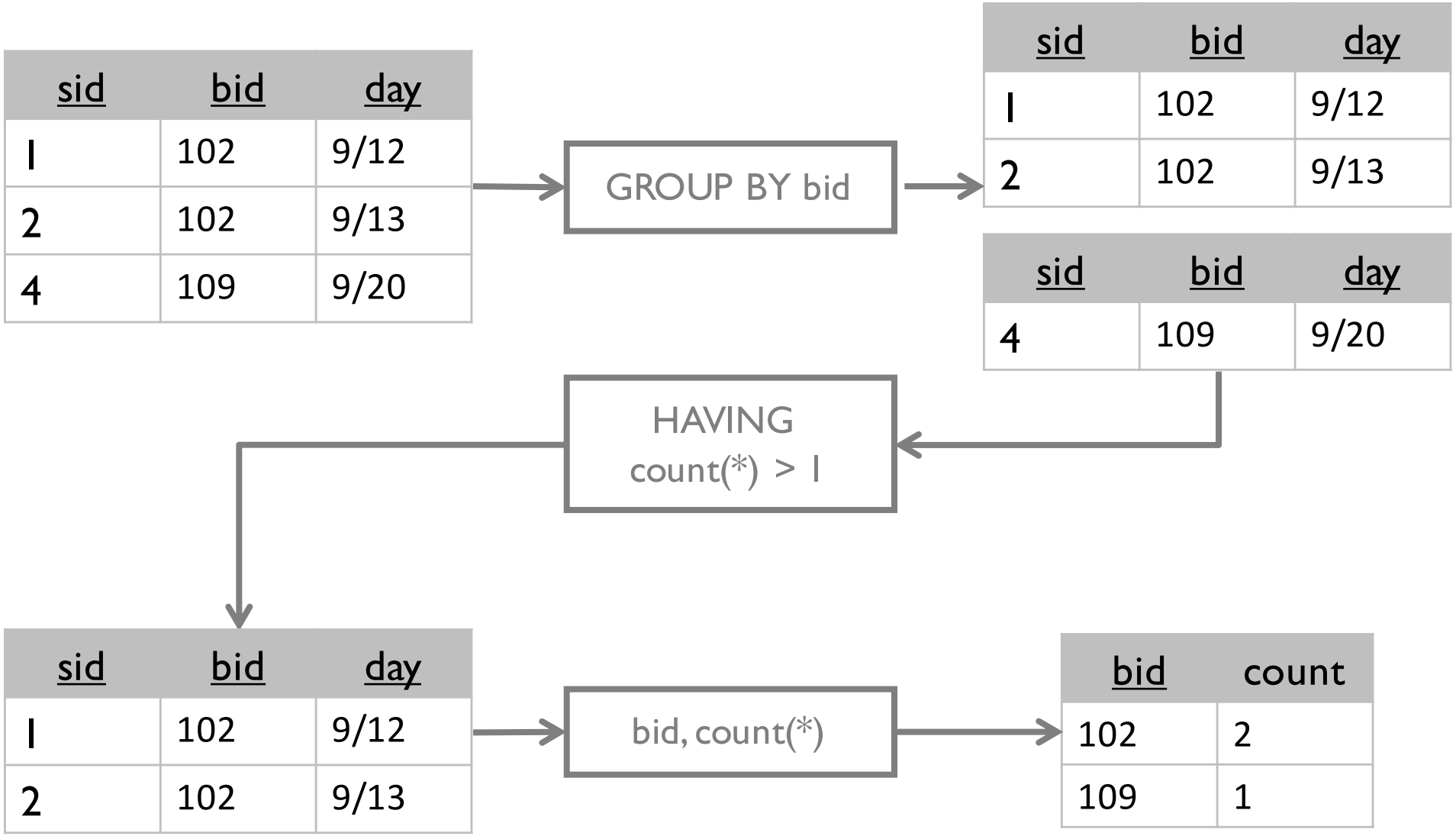
Expressions must have *one value per group*. Either

An aggregation function or

In *grouping-list*

```
SELECT    bid, count(*)
FROM      Reserves R
GROUP BY  bid
HAVING    color = 'red'
```

Conceptual Evaluation



Number of reservations for each red boat

```
SELECT    S.bid, COUNT(*) AS count
FROM      Sailors S, Boats B, Reserves R
WHERE     S.sid = R.sid AND
          R.bid = B.bid AND
          B.color = ‘red’
GROUP BY  B.bid
```

What if we move `B.color = ‘red’` from WHERE clause to the HAVING clause?

What if we drop Sailors from query?

Ratings where the average age is minimum over all ratings

```
SELECT S.rating
FROM   Sailors S
WHERE  S.age = (
           SELECT MIN(AVG(S2.age)
           FROM Sailors S2
       )
```

```
SELECT S.rating
FROM   (SELECT S.rating, AVG(S.age) as avgage
        FROM Sailors S
        GROUP BY S.rating) AS tmp
WHERE  tmp.avgage = (
           SELECT MIN(tmp.avgage) FROM tmp
       )
```

Integrity Constraints

Conditions that every legal instance must satisfy

- Inserts/Deletes/Updates that violate ICs rejected
- Helps ensure app semantics or prevent inconsistencies

We've discussed

- domain/type constraints, primary/foreign key
- general constraints

Beyond Keys: General Constraints

```
CREATE TABLE Sailors(
    sid int,
    …
    PRIMARY KEY (sid),
    CHECK (rating >= 1 AND rating <= 10)
```

```
CREATE TABLE Reserves(
    sid int,
    bid int,
    day date,
    PRIMARY KEY (bid, day),
    CONSTRAINT no_red_reservations
    CHECK ('red' <> (SELECT B.color
                     FROM Boats B
                     WHERE B.bid = bid))
```

Nested subqueries

Named constraints

Multi-Relation Constraints

of boats + # of sailors should be less than 100

```
CREATE TABLE Sailors (
    sid int,
    bid int,
    day date,
    PRIMARY KEY (bid, day),
    CHECK (
        (SELECT COUNT(S.sid) FROM Sailors S)
        +
        (SELECT COUNT(B.bid) FROM Boats B)
        < 100
    )
```

What if Sailors is empty?

ASSERTIONS: Multi-Relation Constraints

```
CREATE ASSERTION small_club
CHECK (
    (SELECT COUNT(*) FROM Sailors S)
    +
    (SELECT COUNT(*) FROM Boats B)
    < 100
)
```

ASSERTIONs are not associated with any table

Triggers

def: procedure that runs automatically if specified changes in DBMS happen

```
CREATE TRIGGER name
    [BEFORE | AFTER | INSTEAD OF] event_list
    ON table
    FOR EACH (ROW | STATEMENT)
    WHEN trigger_qualifications
    procedure
```

Event activates the trigger

Condition tests if triggers should run

Action what to do

Copy new young sailors into special table

```
CREATE TRIGGER youngSailorUpdate
    AFTER INSERT ON SAILORS
REFERENCING NEW TABLE NewInserts
FOR EACH STATEMENT
    INSERT
        INTO YoungSailors(sid, name, age, rating)
        SELECT sid, name, age, rating
        FROM NewInserts N
        WHERE N.age <= 18
```

Triggers

Can be complicated to reason about

- Triggers may (e.g., insert) cause other triggers to run
- If >1 trigger match an action, which is run first? ¯_(ツ)_/¯

```
CREATE TRIGGER recursiveTrigger
    AFTER INSERT ON SAILORS
FOR EACH ROW
    INSERT
        INTO Sailors(sid, name, age, rating)
        SELECT sid, name, age, rating
        FROM Sailors S
```

WITH

```
WITH RedBoats(bid, count) AS
    (SELECT   B.bid, count(*)
     FROM     Boats B, Reserves R
     WHERE    R.bid = B.bid AND B.color = ‘red’
    GROUP BY  B.bid)
SELECT    name, count
FROM      Boats B,  RedBoats RB
WHERE     B.bid = RB.bid AND count < 10
```

Names of unpopular boats

Views

```
CREATE VIEW   view_name
AS select_statement
```

Instead of table of inserted records, "tables" defined as query results

Makes development simpler

Used for security

Not *materialized*

Views

```
CREATE VIEW boat_counts
AS  SELECT  bid, count(*)
    FROM        Reserves  R
    GROUP BY    bid
    HAVING      count(*)  > 10
```

Used like a normal query

```
SELECT  bname
FROM    boat_counts  bc, Boats  B
WHERE   bc.bid  = B.bid
```

Names of popular boats

```
SELECT  bname
FROM
    (SELECT bid, count(*)
     FROM Reserves  R
     GROUP BY bid
     HAVING count(*)  > 10) bc,
    Boats  B
WHERE   bc.bid  = B.bid
```

CREATE TABLE

```
CREATE TABLE <table_name> AS
    <SELECT STATEMENT>
```

Guess the schema:

```
CREATE TABLE used_boats1 AS
    SELECT r.bid
    FROM   Sailors s,
           Reservations r
    WHERE  s.sid = r.sid
```

used_boats1(bid int)

```
CREATE TABLE used_boats2 AS
    SELECT r.bid as foo
    FROM   Sailors s,
           Reservations r
    WHERE  s.sid = r.sid
```

used_boats2(foo int)

How is this different than views?

- What if we insert a new record into Reservations?

Summary

SQL is pretty complex

- Superset of Relational Algebra SQL99 turing complete!
- Human readable

More than one way to skin a horse

- Many alternatives to write a query
- Optimizer (theoretically) finds most efficient plan