

L5

SQL SQL SQL SQL SQL SQL SQL

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# Didn't Lecture 3 Go Over SQL?

Two sublanguages

**DDL** Data Definition Language

define and modify schema (physical, logical, view)

CREATETABLE, Integrity Constraints

**DML** Data Manipulation Language

get and modify data

simple SELECT, INSERT, DELETE

*human-readable* language

# Gritty Details

## DDL

NULL, Views

## DML

Basics, SQL Clauses, Expressions, Joins, Nested Queries, Aggregation, With, Triggers

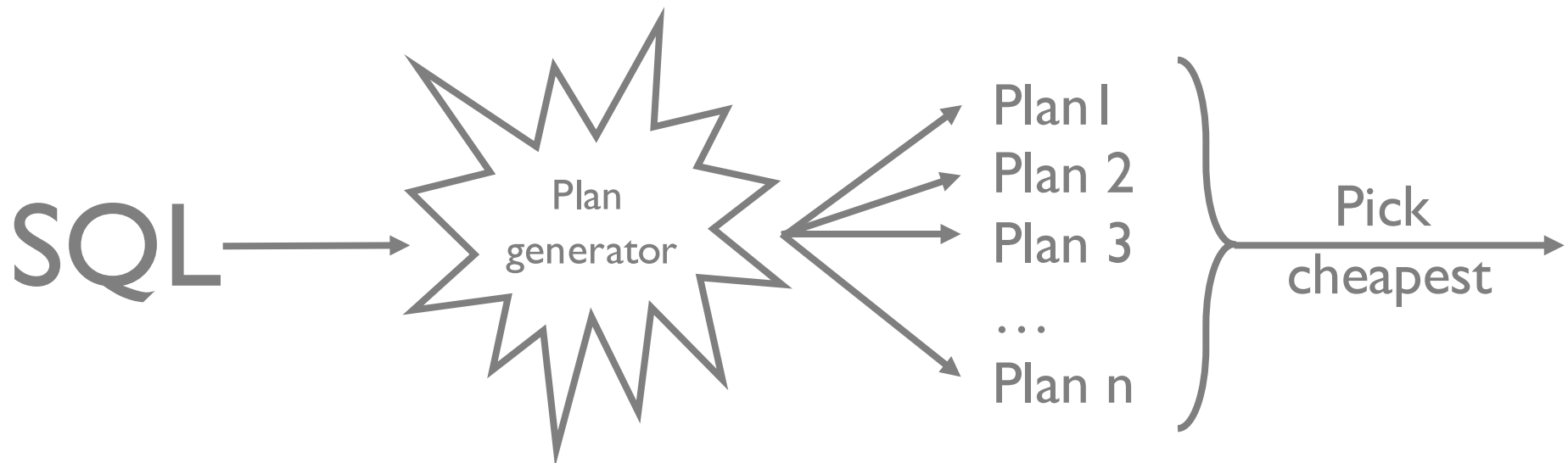
# Didn't Lecture 3 Go Over SQL?

DBMS makes it run efficiently

Key: precise query semantics

Reorder/modify queries while answers stay same

DBMS estimates costs for different evaluation plans



# Didn't Lecture 3 Go Over SQL?

More expressive power than Rel Alg

can be described by extensions of algebra

**One key difference:** multisets rather than sets

i.e. # duplicates in a table carefully accounted for

Most widely used *query language*, not just relational query language

# Today's Database

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Boats

<u>bid</u>	name	color
101	Legacy	red
102	Melon	blue
103	Mars	red

Reserves

<u>sid</u>	<u>bid</u>	day
1	102	9/12
2	102	9/13
2	103	9/14

Is Reserves table correct?

# Today's Database

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Boats

<u>bid</u>	name	color
101	Legacy	red
102	Melon	blue
103	Mars	red

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13
2	103	9/14

Is Reserves table correct?  
Day should be part of key

# Follow along at home!

`http://w4111db1.cloudapp.net:8000/`

or

```
psql -U demo -h w4111db1.cloudapp.net demo  
password: demo
```



# <30 year old sailors

```
SELECT *  
FROM Sailors  
WHERE age < 30
```

<u>sid</u>	name	rating	age
1	Eugene	7	22
3	Ken	8	27

```
SELECT name, age  
FROM Sailors  
WHERE age < 30
```

name	age
Eugene	22
Ken	27

# <30 year old sailors

```
SELECT *  
FROM Sailors  
WHERE age < 30
```

$\sigma_{\text{age} < 30} (\text{Sailors})$

```
SELECT name, age  
FROM Sailors  
WHERE age < 30
```

$\pi_{\text{name, age}} (\sigma_{\text{age} < 30} (\text{Sailors}))$

# Multiple Relations

```
SELECT S.name
FROM   Sailors AS S, Reserves AS R
WHERE  S.sid = R.sid AND R.bid = 102
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13
2	103	9/14



# Multiple Relations

```
SELECT S.name  
FROM   Sailors AS S, Reserves AS R  
WHERE  S.sid = R.sid AND R.bid = 102
```

$\pi_{\text{name}} (\sigma_{\text{bid}=2} (\text{Sailors} \bowtie_{\text{sid}} \text{Reserves}))$

# Structure of a SQL Query

## DISTINCT

Optional, answer should not have duplicates  
Default: duplicates not removed (multiset)

## target-list

List of expressions over attrs of tables in relation-list

SELECT    [DISTINCT] *target-list*  
FROM      *relation-list*  
WHERE     *qualification*

## relation-list

List of relation names  
Can define range-variable “AS X”

## qualification

Boolean expressions  
Combined w/ AND, OR, NOT  
attr *op* const  
attr<sub>1</sub> *op* attr<sub>2</sub>  
*op* is =, <, >, !=, etc

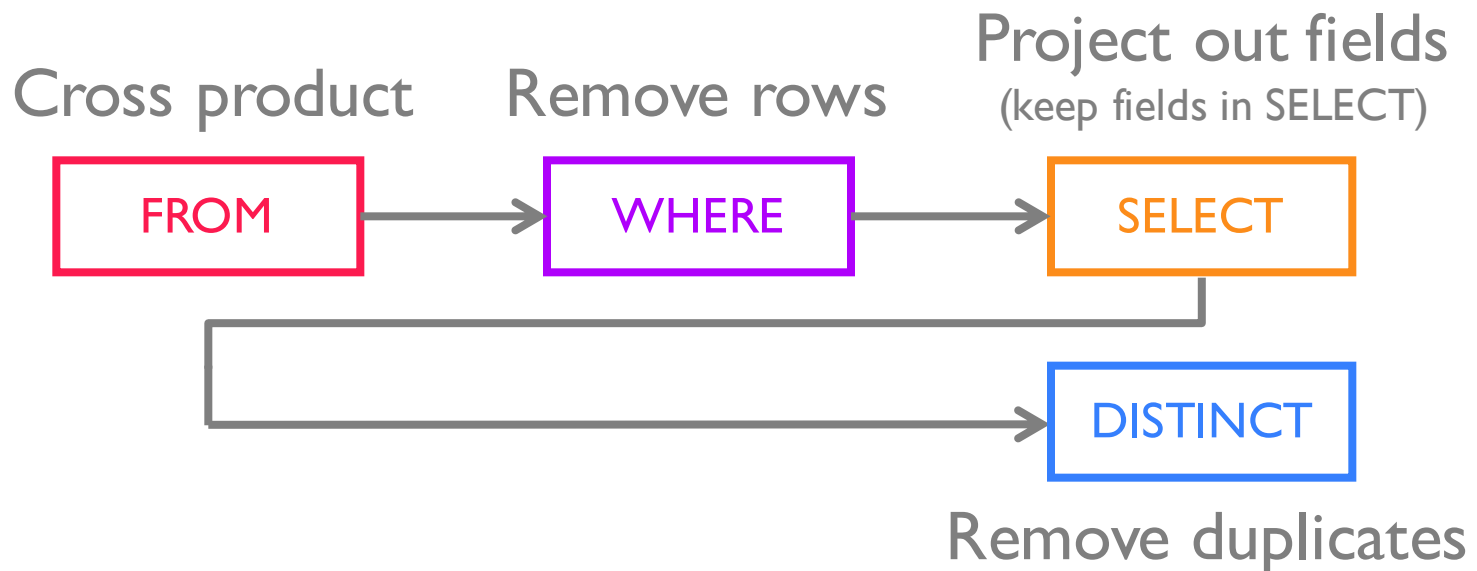
# Semantics

SELECT    [DISTINCT] *target-list*  
FROM      *relation-list*  
WHERE    *qualification*

FROM	compute cross product of relations
WHERE	remove tuples that fail qualifications
SELECT	remove fields not in target-list
DISTINCT	remove duplicate rows

# Conceptual Query Evaluation

SELECT	[DISTINCT] <i>target-list</i>
FROM	<i>relation-list</i>
WHERE	<i>qualification</i>
GROUP BY	<i>grouping-list</i>
HAVING	<i>group-qualification</i>



Not how actually executed! Above is likely very slow

# DISTINCT (vol.I)

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13
2	103	9/14

SELECT bid  
FROM Reserves

<u>bid</u>
102
102
103

SELECT DISTINCT bid  
FROM Reserves

<u>bid</u>
102
103



# Sailors that reserved 1+ boats

```
SELECT  S.sid  
FROM    Sailors AS S, Reserves AS R  
WHERE   S.sid = R.sid
```

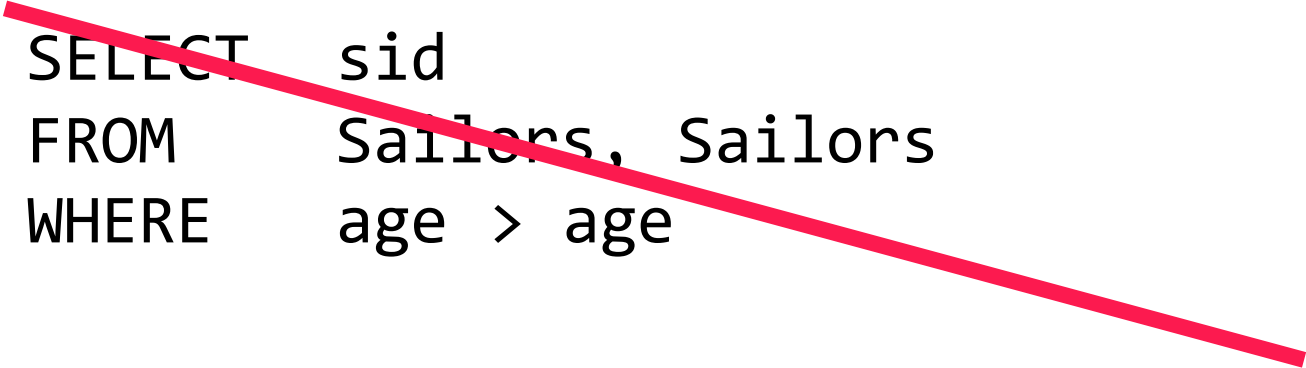
Would DISTINCT change anything in this query?

What if SELECT clause was SELECT S.name?

# Range Variables

Disambiguate relations

same table used multiple times (self join)



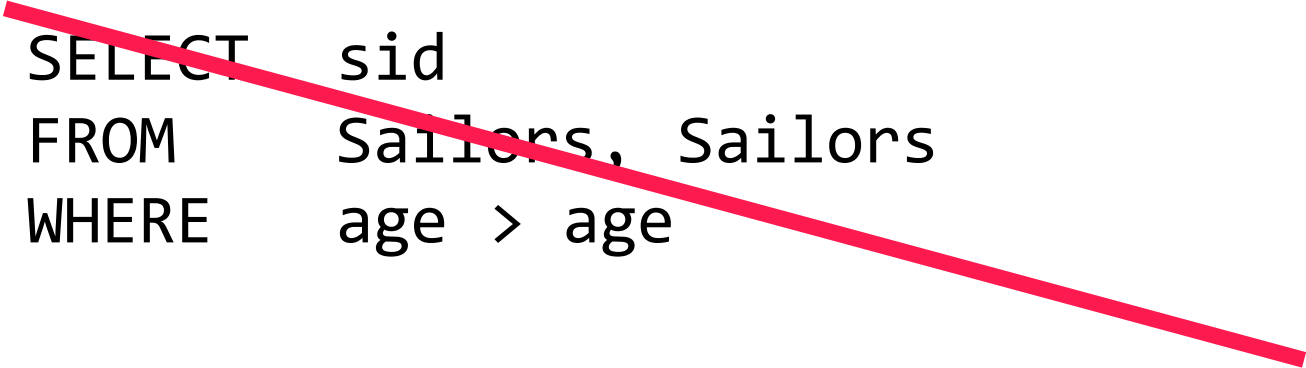
```
SELECT  sid
FROM    Sailors, Sailors
WHERE   age > age
```

```
SELECT  S1.sid
FROM    Sailors AS S1, Sailors AS S2
WHERE   S1.age > S2.age
```

# Range Variables

Disambiguate relations

same table used multiple times (self join)



```
SELECT  sid
FROM    Sailors, Sailors
WHERE   age > age
```

```
SELECT  S1.name, S1.age, S2.name, S2.age
FROM    Sailors AS S1, Sailors AS S2
WHERE   S1.age > S2.age
```

# Expressions (Math)

```
SELECT  S.age, S.age - 5 AS age2, 2*S.age AS age3
FROM    Sailors AS S
WHERE   S.name = 'eugene'
```

```
SELECT  S1.name AS name1, S2.name AS name2
FROM    Sailors AS S1, Sailors AS S2
WHERE   S1.rating*2 = S2.rating - 1
```

# Expressions (Strings)

```
SELECT  S.name  
FROM    Sailors AS S  
WHERE   S.name LIKE 'e_%'
```

‘\_’      any one character (• in regex)

‘%’      0 or more characters of any kind (•\* in regex)

Most DBMSes have rich string manipulation support e.g., regex

PostgreSQL documentation

<http://www.postgresql.org/docs/9.1/static/functions-string.html>

# Expressions (Date/Time)

```
SELECT    R.sid  
FROM      Reserves AS R  
WHERE     now() - R.date < interval '1 day'
```

TIMESTAMP, DATE, TIME types

now() returns timestamp at start of transaction

DBMSes provide rich time manipulation support

exact support may vary by vender

Postgresql Documentation

<http://www.postgresql.org/docs/9.1/static/functions-datetime.html>

# Expressions

Constant	
Col reference	Sailors.name
Arithmetic	Sailors.sid * 10
Unary operators	NOT, EXISTS
Binary operators	AND, OR, IN
Function calls	abs(), sqrt(), ...
Casting	1.7::int, '10-12-2015'::date

sid of Sailors that reserved red or blue boat

```
SELECT    R.sid
FROM      Boats B, Reserves R
WHERE     B.bid = R.bid AND
          (B.color = 'red' OR B.color = 'blue')
```

OR

```
SELECT    R.sid
FROM      Boats B, Reserves R
WHERE     B.bid = R.bid AND B.color = 'red'
UNION ALL
SELECT    R.sid
FROM      Boats B, Reserves R
WHERE     B.bid = R.bid AND B.color = 'blue'
```



sid of Sailors that reserved red or blue boat

```
SELECT    DISTINCT R.sid
FROM      Boats B, Reserves R
WHERE     B.bid = R.bid AND
          (B.color = 'red' OR B.color = 'blue')
```

OR

```
SELECT    R.sid
FROM      Boats B, Reserves R
WHERE     B.bid = R.bid AND B.color = 'red'
UNION
SELECT    R.sid
FROM      Boats B, Reserves R
WHERE     B.bid = R.bid AND B.color = 'blue'
```

sid of Sailors that reserved red and blue boat

```
SELECT    R.sid  
FROM      Boats B, Reserves R  
WHERE     B.bid = R.bid AND  
          (B.color = 'red' AND B.color = 'blue')
```

```
SELECT    R.sid  
FROM      Boats B, Reserves R  
WHERE     B.bid = R.bid AND B.color = 'red'  
INTERSECT ALL  
SELECT    R.sid  
FROM      Boats B, Reserves R  
WHERE     B.bid = R.bid AND B.color = 'blue'
```

sid of Sailors that reserved red and blue boat

Can use self-join instead

```
SELECT    R.sid
FROM      Boats B1, Reserves R1
WHERE
          B1.bid = R1.bid AND

          B1.color = 'red'
```

sid of Sailors that reserved red and blue boat

Can use self-join instead

```
SELECT    R.sid
FROM      Boats B1, Reserves R1, Boats B2, Reserves R2
WHERE
          B1.bid = R1.bid AND

          B1.color = 'red'
```

sid of Sailors that reserved red and blue boat

Can use self-join instead

```
SELECT    R.sid
FROM      Boats B1, Reserves R1, Boats B2, Reserves R2
WHERE
    B1.bid = R1.bid AND
    B2.bid = R2.bid AND
    B1.color = 'red' AND B2.color = 'blue'
```

sid of Sailors that reserved red and blue boat

Can use self-join instead

```
SELECT    R.sid
FROM      Boats B1, Reserves R1, Boats B2, Reserves R2
WHERE     R1.sid = R2.sid AND
          B1.bid = R1.bid AND
          B2.bid = R2.bid AND
          B1.color = 'red' AND B2.color = 'blue'
```

sids of sailors that haven't reserved a boat

```
SELECT    S.sid  
FROM      Sailors S
```

EXCEPT

```
SELECT    S.sid  
FROM      Sailors S, Reserves R  
WHERE     S.sid = R.sid
```

Can we write EXCEPT using more basic functionality?

# SET Comparison Operators

UNION, INTERSECT, EXCEPT

EXISTS, NOT EXISTS

IN, NOT IN

UNIQUE, NOT UNIQUE

*op* ANY, *op* ALL

$op \in \{ <, >, =, \leq, \geq, \neq, \dots \}$

Many of these rely on Nested Query Support



# Nested Queries

```
SELECT  S.sid
FROM    Sailors S
WHERE   S.sid IN (SELECT  R.sid
                  FROM    Reserves R
                  WHERE   R.bid = 101)
```

Many clauses can contain SQL queries  
WHERE, FROM, HAVING, SELECT

Conceptual model:

- for each Sailors tuple
- run the subquery and evaluate qualification

# Nested Correlated Queries

```
SELECT  S.sid
FROM    Sailors S
WHERE   EXISTS (SELECT  *
                  FROM    Reserves R
                  WHERE   R.bid = 101 AND
                          S.sid = R.sid)
```

Outer table referenced in nested query

Conceptual model:

for each Sailors tuple

run the subquery and evaluate qualification

# Nested Correlated Queries

```
SELECT  S.sid
FROM    Sailors S
WHERE   UNIQUE (SELECT  *
                  FROM    Reserves R
                  WHERE    R.bid = 101 AND
                          S.sid = R.sid)
```

UNIQUE checks that there are no duplicates

What does this do?

# Nested Correlated Queries

```
SELECT  S.sid
FROM    Sailors S
WHERE   UNIQUE (SELECT  R.sid
                  FROM    Reserves R
                  WHERE    R.bid = 101 AND
                          S.sid = R.sid)
```

UNIQUE checks that there are no duplicates

What does this do?

Sailors whose rating is greater than  
any sailor named “Bobby”

```
SELECT S1.name
FROM   Sailors S1
WHERE  S1.rating > ANY (SELECT   S2.rating
                        FROM     Sailors S2
                        WHERE      S2.name = 'Bobby')
```

# What about this?

```
SELECT S1.name
FROM   Sailors S1
WHERE  S1.rating > ALL (SELECT   S2.rating
                        FROM     Sailors S2
                        WHERE      S2.name = 'Bobby')
```

# Rewrite INTERSECT using IN

```
SELECT S.sid  
FROM   Sailors S  
WHERE  S.rating > 2
```

INTERSECT

```
SELECT R.sid  
FROM   Reserves R
```

```
SELECT S.sid  
FROM   Sailors S  
WHERE  S.rating > 2 AND  
       S.sid IN (  
           SELECT R.sid  
           FROM   Reserves R  
       )
```

Similar trick for EXCEPT → NOT IN

What if want *names* instead of sids?

# Sailors that reserved all boats (Division)

Hint: double negation

reserved all boats == no boat w/out reservation

```
SELECT    S.name
FROM      Sailors S
WHERE     NOT EXISTS (

    (SELECT  B.bid FROM    Boats B)

    EXCEPT

    (SELECT  R.bid
     FROM    Reserves R
     WHERE   R.sid = S.sid)
)
```



# HWI bugs

## Conflicting CHECK constraints

```
Prof(  
    type text,  
    check(text in ('junior', 'senior')),  
    check(text = 'junior' and hired is not null),  
    check(text = 'senior' and tenure_year is not null)
```



conflicting

# HWI bugs

At most once *per semester* translated as at most once

```
CREATE TABLE Offers (  
    deptid text,  
    courseid text,  
    semester text,  
    year int,  
    . . .  
    PRIMARY KEY(deptid, courseid)  
);
```

Wrong

# HWI bugs

At most once *per semester* translated as at most once

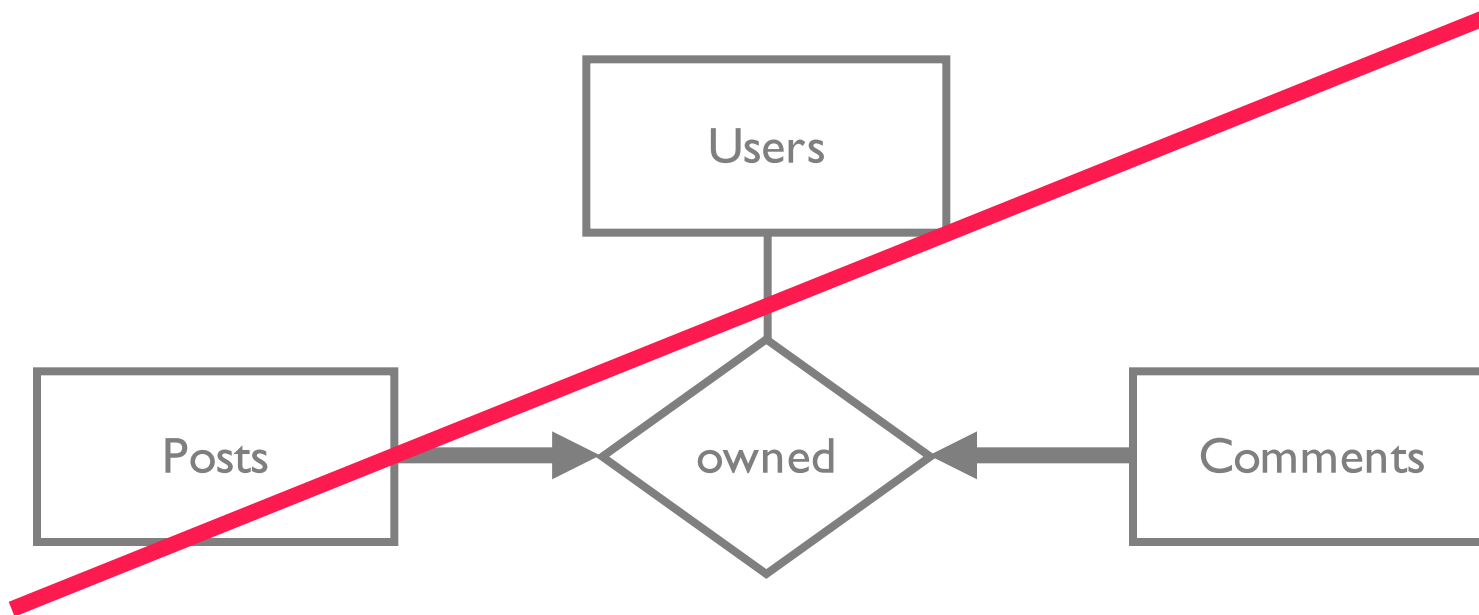
```
CREATE TABLE Offers (  
    deptid text,  
    courseid text,  
    semester text,  
    year int,  
    . . .  
    PRIMARY KEY(deptid, courseid, semester, year)  
);
```

# HWI bugs

Reddit:

Comments owned by one user

Posts owned by one user

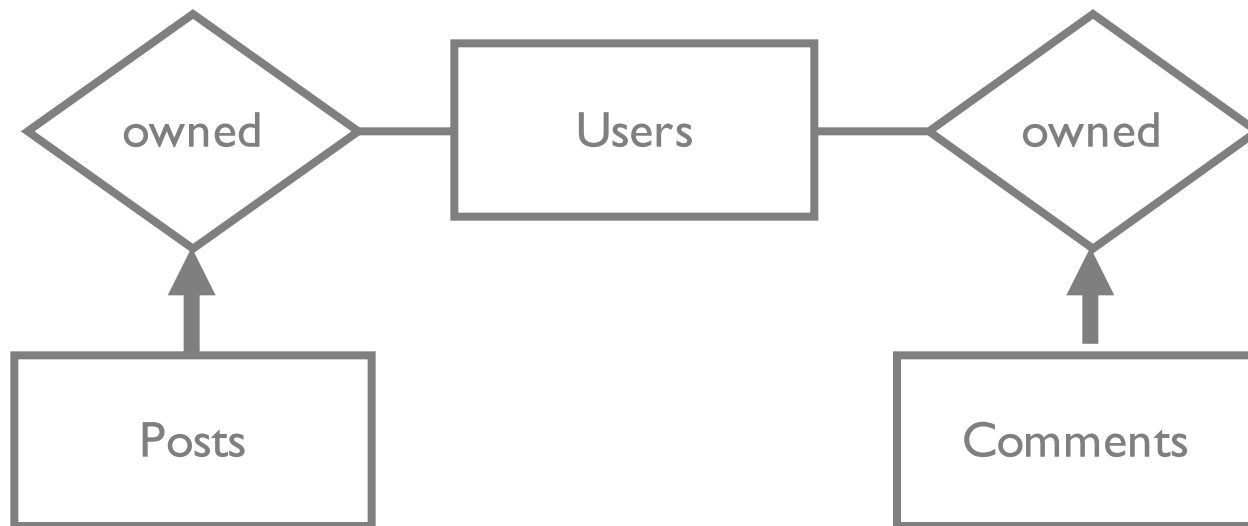


# HWI bugs

Reddit:

Comments owned by one user

Posts owned by one user



# Sailors that reserved all boats (Division)

Hint: double negation

reserved all boats == no boat w/out reservation

```
SELECT S.name  
FROM   Sailors S  
WHERE  NOT EXISTS (
```

Sailors S such that

There's no boat without

A reservation by S

# Sailors that reserved all boats (Division)

Hint: double negation

reserved all boats == no boat w/out reservation

```
SELECT S.name
FROM   Sailors S
WHERE  NOT EXISTS (SELECT B.bid
                   FROM   Boats B
                   WHERE  NOT EXISTS (
```

Sailors S such that

There's no boat without

A reservation by S

# Sailors that reserved all boats (Division)

Hint: double negation

reserved all boats == no boat w/out reservation

```
SELECT S.name
FROM   Sailors S
WHERE  NOT EXISTS (SELECT B.bid
                   FROM   Boats B
                   WHERE  NOT EXISTS (SELECT R.bid
                                     FROM   Reserves R
                                     WHERE  R.sid = S.sid))
```

Sailors S such that

There's no boat without

A reservation by S



# NULL

Field values sometimes unknown or inapplicable

SQL provides a special value *null* for such situations.

The presence of null complicates many issues e.g.,

Is age = null true or false?

Is null = null true or false?

Is null = 8 OR 1 = 1 true or false?

Special syntax “IS NULL” and “IS NOT NULL”

3 Valued Logic (true, false, unknown)

How does WHERE remove rows?

if qualification doesn't evaluate to true

New operators (in particular, outer joins) possible/needed.

# NULL

(null > 0) = null  
(null + 1) = null  
(null = 0) = null  
(null AND true) = null  
null is null = true

## Some truth tables

<b>AND</b>	T	F	NULL
T	T	F	NULL
F	F	F	F
NULL	NULL	F	NULL

<b>OR</b>	T	F	NULL
T	T	T	T
F	T	F	NULL
NULL	T	NULL	NULL

# JOINS

```
SELECT [DISTINCT] target_list
FROM table_name
    [INNER | {LEFT | RIGHT | FULL } {OUTER}] JOIN table_name
    ON qualification_list
WHERE ...
```

INNER is default

Difference in how to deal with NULL values

PostgreSQL documentation:

<http://www.postgresql.org/docs/9.4/static/tutorial-join.html>

# Inner/Natural Join

```
SELECT s.sid, s.name, r.bid  
FROM   Sailors S, Reserves r  
WHERE  s.sid = r.sid
```

```
SELECT s.sid, s.name, r.bid  
FROM   Sailors s INNER JOIN Reserves r  
ON     s.sid = r.sid
```

All  
Equivalent!

```
SELECT s.sid, s.name, r.bid  
FROM   Sailors s NATURAL JOIN Reserves r
```

**Natural Join** means equi-join for each pair of  
attrs with same name

# Sailor names and their reserved boat ids

```
SELECT s.sid, s.name, r.bid
FROM   Sailors s INNER JOIN Reserves r
ON     s.sid = r.sid
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13

Result

sid	name	bid
1	Eugene	102
2	Luis	102

# Sailor names and their reserved boat ids

```
SELECT s.sid, s.name, r.bid
FROM   Sailors s INNER JOIN Reserves r
ON     s.sid = r.sid
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13

Result

sid	name	bid
1	Eugene	102
2	Luis	102

Prefer INNER JOIN over NATURAL JOIN. Why?

# Sailor names and their reserved boat ids

```
SELECT s.sid, s.name, r.bid
FROM   Sailors s INNER JOIN Reserves r
ON     s.sid = r.sid
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13

Result

sid	name	bid
1	Eugene	102
2	Luis	102

Notice: No result for Ken!

# Left Outer Join (or No Results for Ken)

Returns all matched rows *and all unmatched rows from table on left of join clause*

*(at least one row for each row in left table)*

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors s LEFT OUTER JOIN Reserves r
ON      s.sid = r.sid
```

All sailors & bid for boat in their reservations

Bid set to NULL if no reservation



# Left Outer Join

```
SELECT s.sid, s.name, r.bid
FROM   Sailors s LEFT OUTER JOIN Reserves r
ON     s.sid = r.sid
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13

Result

sid	name	bid
1	Eugene	102
2	Luis	102
3	Ken	NULL

# Can Left Outer Join be expressed with Cross-Product?

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
------------	------------	------------

Sailors x Reserves

Sailors s **LEFT OUTER JOIN** Reserves r  
ON s.sid = r.sid

Result

sid	name	bid
-----	------	-----

Result

sid	name	bid
1	Eugene	NULL
2	Luis	NULL
3	Ken	NULL

# Can Left Outer Join be expressed with Cross-Product?

Sailors			
<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves		
<u>sid</u>	<u>bid</u>	<u>day</u>

Sailors ⋈ Reserves  
U

$(\text{Sailors} - (\text{Sailors} \bowtie \text{Reserves})) \times \{(\text{null}, \dots)\}$

How to compute this with a query?

# Right Outer Join

Same as LEFT OUTER JOIN, but guarantees result for rows in table on **right side of JOIN**

```
SELECT s.sid, s.name, r.bid
FROM   Reserves r RIGHT OUTER JOIN Sailors S
ON     s.sid = r.sid
```

# FULL OUTER JOIN

Returns all matched *or* unmatched rows from both sides of JOIN

```
SELECT  s.sid, s.name, r.bid
FROM    Sailors s FULL OUTER JOIN Reserves r
ON      s.sid = r.sid
```

# FULL OUTER JOIN

```
SELECT s.sid, s.name, r.bid
FROM   Sailors s Full OUTER JOIN Reserves r
ON     s.sid = r.sid
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
1	102	9/12
2	102	9/13
4	109	9/20

Result

sid	name	bid
1	Eugene	102
2	Luis	102
3	Ken	NULL
NULL	NULL	109

Why is sid NULL?

# Serious people can count: Aggregation

```
SELECT COUNT(*)  
FROM   Sailors S
```

COUNT([DISTINCT] A)

```
SELECT AVG(S.age)  
FROM   Sailors S  
WHERE  S.rating = 10
```

SUM([DISTINCT] A)

AVG([DISTINCT] A)

MAX/MIN(A)

STDDEV(A)

```
SELECT COUNT(DISTINCT S.name)  
FROM   Sailors S  
WHERE  S.name LIKE 'D%'
```

CORR(A,B)

```
SELECT S.name  
FROM   Sailors  
WHERE  S.rating = (SELECT MAX(S2.rating)  
                  FROM   Sailors S2)
```

PostgreSQL documentation

<http://www.postgresql.org/docs/9.4/static/functions-aggregate.html>

# Name and age of oldest sailor(s)

```
SELECT S.name, MAX(S.age)  
FROM   Sailors S
```

```
SELECT S.name, S.age  
FROM   Sailors S  
WHERE  S.age >= ALL (SELECT S2.age  
                     FROM   Sailors S2)
```

```
SELECT S.name, S.age  
FROM   Sailors S  
WHERE  S.age = (SELECT MAX(S2.age)  
               FROM   Sailors S2)
```

```
SELECT S.name, S.age  
FROM   Sailors S  
ORDER BY S.age DESC  
LIMIT 1
```

← When does this not work?



# GROUP BY

```
SELECT min(s.age)
FROM   Sailors s
```

Minimum age among all sailors

What if want min age *per rating level*?

We don't even know how many rating levels exist!

If we did, could write (awkward):

```
for rating in [0..10]
  SELECT min(s.age)
  FROM   Sailors s
  WHERE  s.rating = <rating>
```

# GROUP BY

```
SELECT count(*)  
FROM   Reserves R
```

Total number of reservations

What if want reservations per boat?

May not even know all our boats (depends on data)!

If we did, could write (awkward):

```
for boat in [0..10]  
    SELECT count(*)  
    FROM   Reserves R  
    WHERE  R.bid = <boat>
```

# GROUP BY

```
SELECT      [DISTINCT] target-list
FROM        relation-list
WHERE       qualification
GROUP BY    grouping-list
HAVING      group-qualification
```

*grouping-list* is a list of expressions that defines groups  
set of tuples w/ same value for all attributes in *grouping-list*

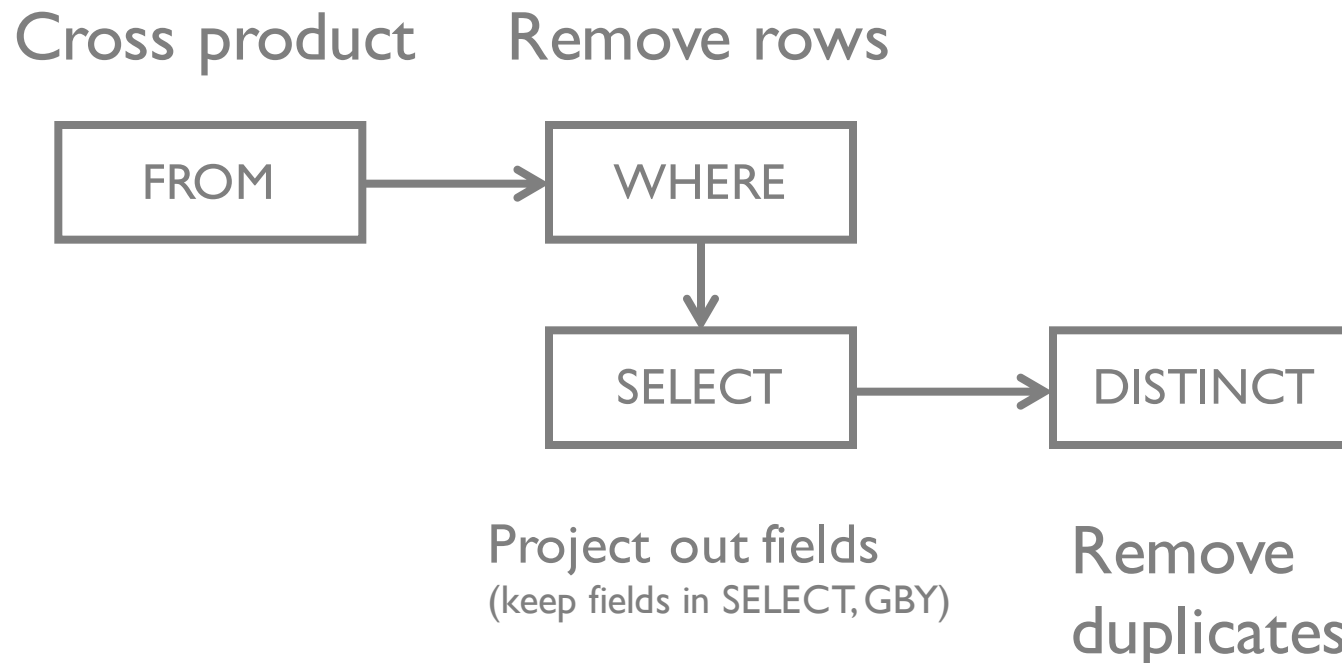
*target-list* contains

*attribute-names*  $\subseteq$  *grouping-list*

*aggregation expressions*

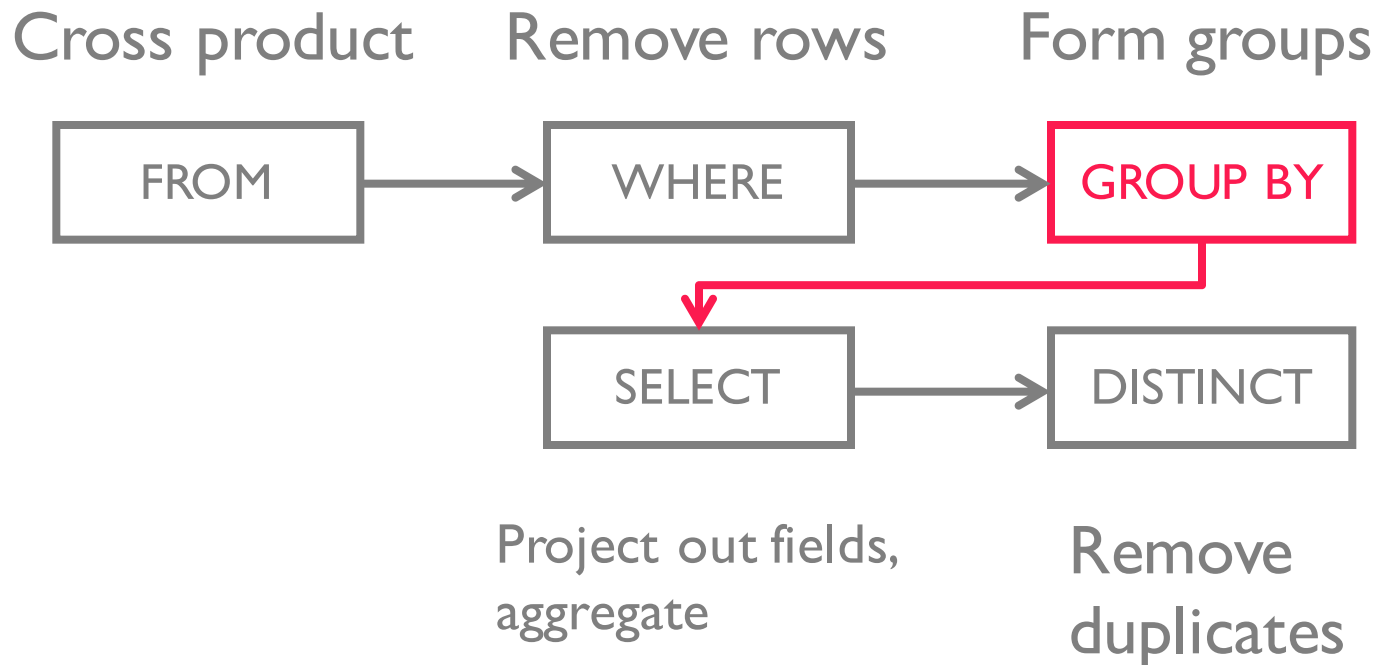
# Conceptual Query Evaluation

SELECT	[DISTINCT] <i>target-list</i>
FROM	<i>relation-list</i>
WHERE	<i>qualification</i>
GROUP BY	<i>grouping-list</i>
HAVING	<i>group-qualification</i>



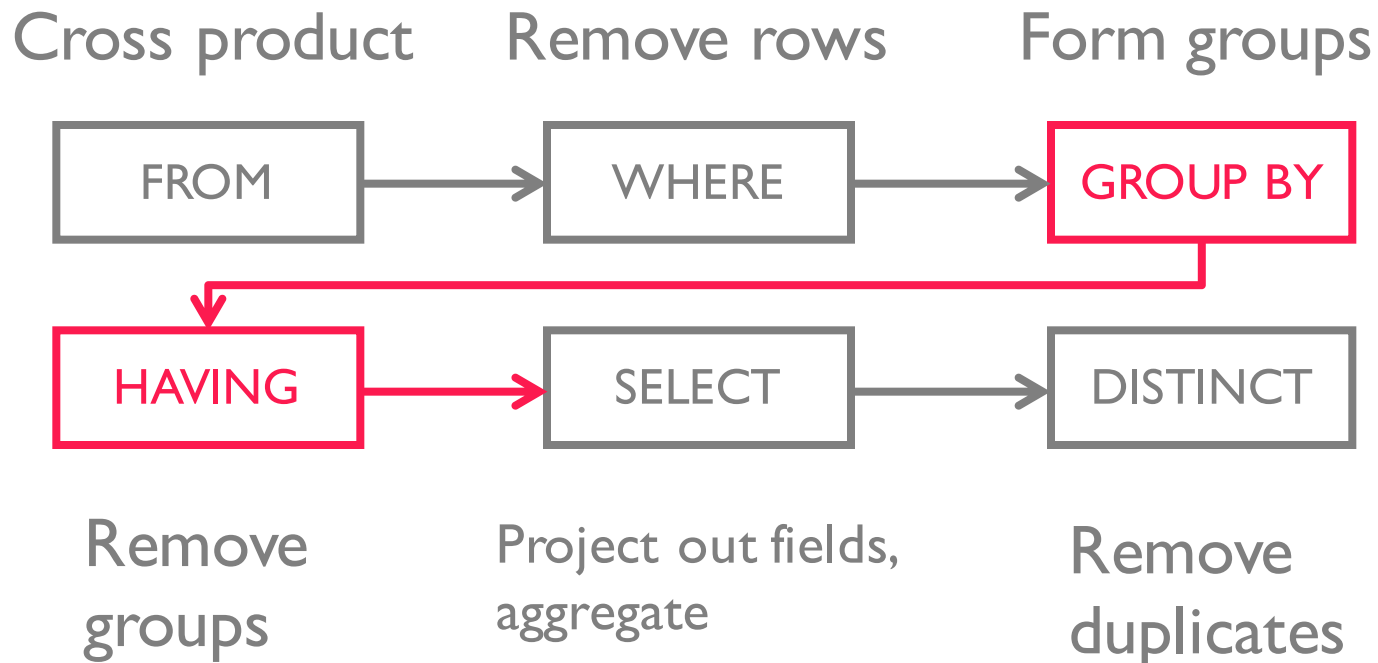
# Conceptual Query Evaluation

SELECT	[DISTINCT] <i>target-list</i>
FROM	<i>relation-list</i>
WHERE	<i>qualification</i>
<b>GROUP BY</b>	<b><i>grouping-list</i></b>
HAVING	<i>group-qualification</i>

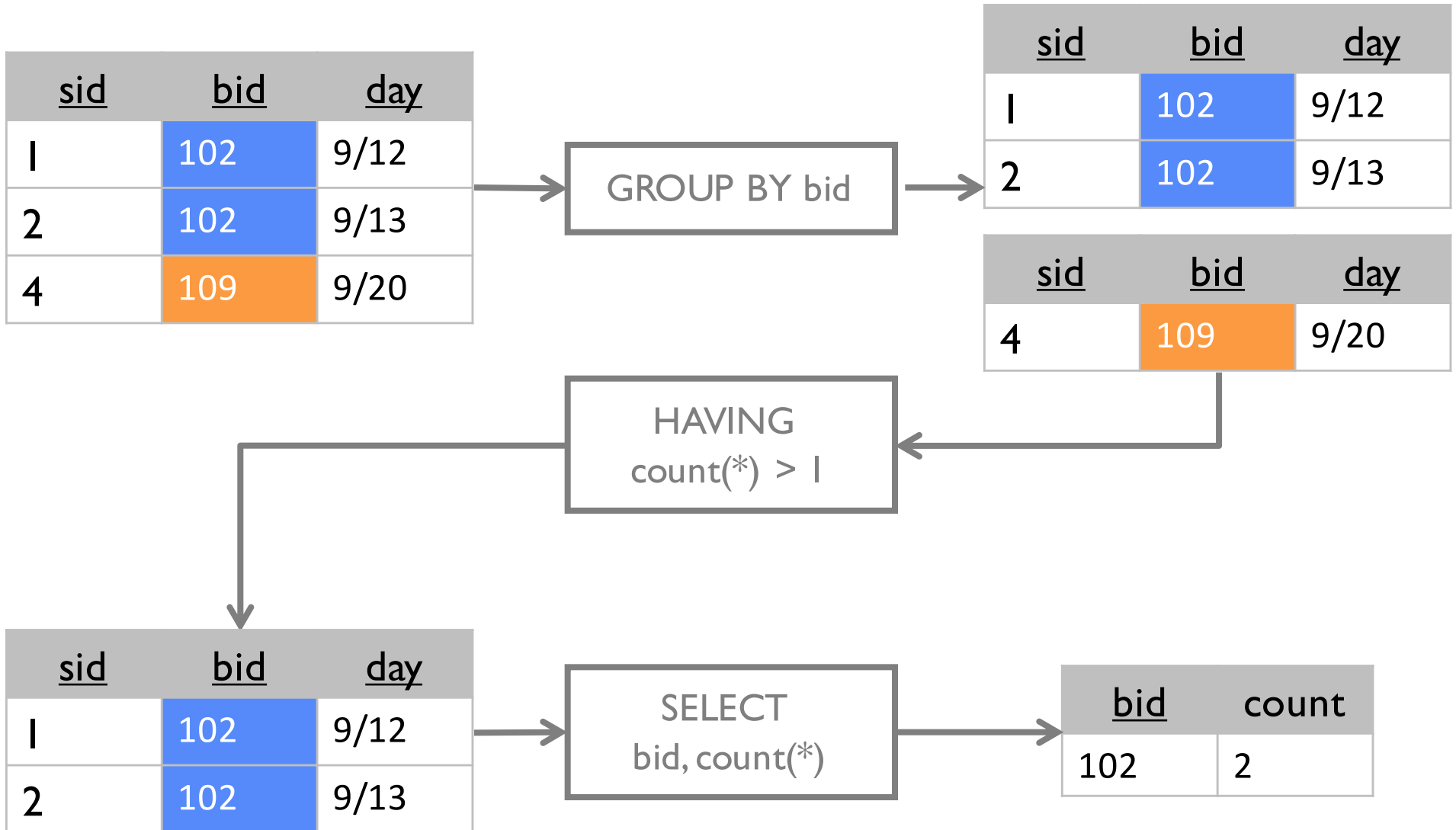


# Conceptual Query Evaluation

SELECT	[DISTINCT] <i>target-list</i>
FROM	<i>relation-list</i>
WHERE	<i>qualification</i>
<b>GROUP BY</b>	<i>grouping-list</i>
<b>HAVING</b>	<i>group-qualification</i>



# Conceptual Evaluation



# GROUP BY

```
SELECT    bid, count(*)  
FROM      Reserves R  
GROUP BY  bid
```

Minimum age for each rating

```
SELECT    bid, count(*)  
FROM      Reserves R  
GROUP BY  bid  
HAVING    count(*) > 1
```

Minimum age for each boat with  
more than 1 reservation



# HAVING

*group-qualification* used to remove groups  
similar to WHERE clause

Expressions must have *one value per group*. Either

An aggregation function or

In *grouping-list*

```
SELECT    bid, count(*)  
FROM      Reserves R  
GROUP BY  bid  
HAVING    color = 'red'
```

# AVG age of sailors reserving red boats, by rating

```
SELECT  
FROM      Sailors S, Boats B, Reserves R  
WHERE     S.sid = R.sid AND  
          R.bid = B.bid AND  
          B.color = 'red'
```

## AVG age of sailors reserving red boats, by rating

```
SELECT    S.rating, avg(S.age) AS age
FROM      Sailors S, Boats B, Reserves R
WHERE     S.sid = R.sid AND
          R.bid = B.bid AND
          B.color = 'red'
GROUP BY  S.rating
```

What if move B.color='red' to HAVING clause?

# Ratings where the avg age is min over all ratings



```
SELECT S.rating
FROM   Sailors S
WHERE  S.age = (
        SELECT MIN(AVG(S2.age))
        FROM   Sailors S2
      )
```



```
SELECT S.rating
FROM   (SELECT S.rating, AVG(S.age) as avgage
        FROM   Sailors S
        GROUP BY S.rating) AS tmp
WHERE  tmp.avgage = (
        SELECT MIN(tmp2.avgage) FROM (
            SELECT S.rating, AVG(S.age) as avgage
            FROM   Sailors S
            GROUP BY S.rating
        ) AS tmp2
      )
```

# Ratings where the avg age is min over all ratings



```
SELECT S.rating
FROM   Sailors S
WHERE  S.age = (
        SELECT MIN(AVG(S2.age))
        FROM   Sailors S2
      )
```



```
SELECT S.rating
FROM   (SELECT S.rating, AVG(S.age) as avgage
        FROM   Sailors S
        GROUP BY S.rating) AS tmp
WHERE  tmp.avgage <= ALL (
        SELECT tmp2.avgage FROM (
        SELECT S.rating, AVG(S.age) as avgage
        FROM   Sailors S
        GROUP BY S.rating
        ) AS tmp2
      )
```

# Setting up Proj I Part 2

Users assigned to schemas (namespaces).

Noticed user didn't have an assigned schema

User created their tables under Public schema.

Uh oh! Did I miss anyone else

Students without a schema

username
vx3948
et1827
etu4938

```
SELECT username  
FROM pg_user;
```

schemaname
et2039
sa2037
kt6765

```
SELECT schemaname  
FROM pg_tables;
```

# Setting up Proj I Part 2

```
FROM    (SELECT username FROM pg_user) AS U  
        (SELECT schemaname FROM pg_tables) AS S
```

username
vx3948
et1827
etu4938

```
SELECT username  
FROM pg_user;
```

schemaname
et2039
sa2037
kt6765

```
SELECT schemaname  
FROM pg_tables;
```

# Setting up Proj I Part 2

```
FROM (SELECT username FROM pg_user) AS U
      RIGHT OUTER JOIN
      (SELECT schemaname FROM pg_tables) AS S
      ON U.username = S.schemaname
```

username
vx3948
et1827
etu4938

```
SELECT username
FROM pg_user;
```

schemaname
et2039
sa2037
kt6765

```
SELECT schemaname
FROM pg_tables;
```



# Setting up Proj I Part 2

```
SELECT U.username, S.schemaname
FROM   (SELECT username FROM pg_user) AS U
       RIGHT OUTER JOIN
       (SELECT schemaname FROM pg_tables) AS S
       ON U.username = S.schemaname
WHERE  S.schemaname is null;
```

username
vx3948
et1827
etu4938

```
SELECT username
FROM pg_user;
```

schemaname
et2039
sa2037
kt6765

```
SELECT schemaname
FROM pg_tables;
```

# Setting up Proj I Part 2

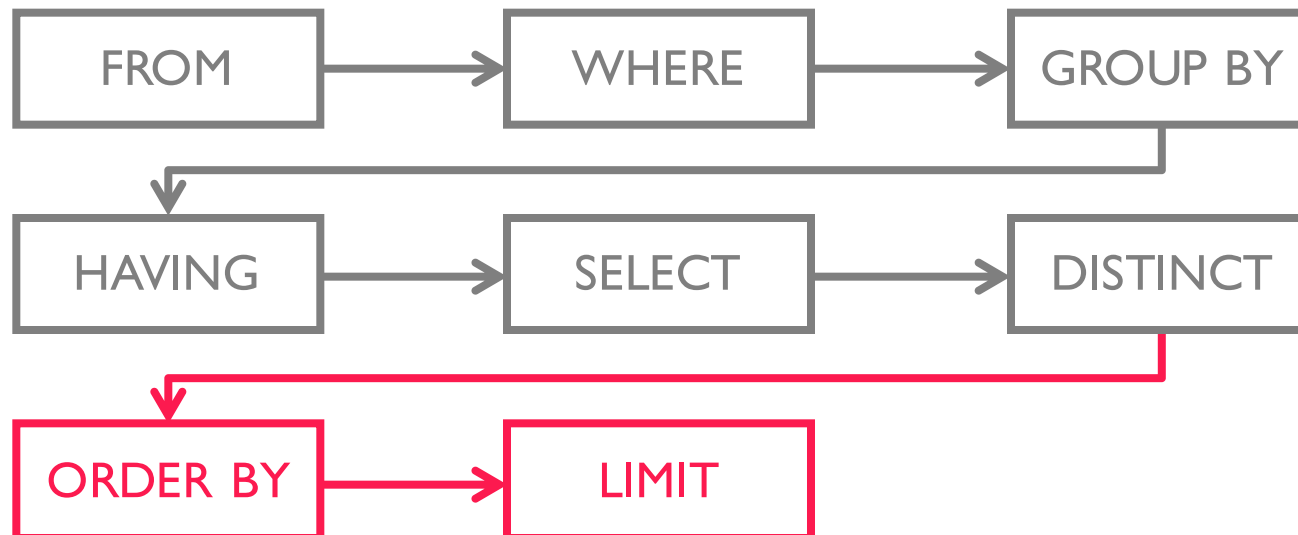
```
SELECT count(*)  
FROM (  
    SELECT U.username, S.schemaname  
    FROM    (SELECT username FROM pg_user) AS U  
            RIGHT OUTER JOIN  
            (SELECT schemaname FROM pg_tables) AS S  
            ON U.username = S.schemaname  
    WHERE   S.schemaname is null;  
) AS nestedquery
```

124

but 190 users!

# ORDER BY, LIMIT

SELECT      [DISTINCT] *target-list*  
FROM        *relation-list*  
WHERE       *qualification*  
GROUP BY   *grouping-list*  
HAVING      *group-qualification*  
*ORDER BY*   *order-list*  
*LIMIT*      *limit-expr* [*OFFSET offset-expr*]



# ORDER BY

```
SELECT    S.name  
FROM      Sailors S  
ORDER BY  (S.rating/2)::int ASC,  
          S.age DESC
```

List of *order-list* expressions dictates ordering precedence  
Sorted in ascending by age/rating ratio  
If ties, sorted high to low rating

# ORDER BY

```
SELECT    S.name, (S.rating/2)::int, S.age
FROM      Sailors S
ORDER BY  (S.rating/2)::int ASC,
          S.age DESC
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Result

name	int4	age
Luis	1	39
Ken	4	27
Eugene	4	22

# ORDER BY

```
SELECT    S.name, (S.rating/2)::int, S.age
FROM      Sailors S
ORDER BY  (S.rating/2)::int ASC,
          S.age ASC
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Result

name	int4	age
Luis	1	39
<b>Eugene</b>	<b>4</b>	<b>22</b>
<b>Ken</b>	<b>4</b>	<b>27</b>

# LIMIT

```
SELECT    S.name, (S.rating/2)::int, S.age
FROM      Sailors S
ORDER BY  (S.rating/2)::int ASC,
          S.age DESC
LIMIT    2
```

Only the first 2 results

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Result

name	int4	age
Luis	1	39
Ken	4	27

# LIMIT

```
SELECT    S.name, (S.rating/2)::int, S.age
FROM      Sailors S
ORDER BY  (S.rating/2)::int ASC,
          S.age DESC
LIMIT    2 OFFSET 1
```

Only the first 2 results

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Result

name	int4	age
Ken	4	27
Eugene	4	22



# LIMIT

```
SELECT    S.name, (S.rating/2)::int, S.age
FROM      Sailors S
ORDER BY  (S.rating/2)::int ASC,
          S.age DESC
LIMIT     (SELECT count(S2.*) / 2
          FROM Sailors AS S2)
```

Can have expressions instead of constants

Result

name	int4	age
Luis	1	39

# Integrity Constraints

Conditions that every legal instance must satisfy

Inserts/Deletes/Updates that violate ICs rejected

Helps ensure app semantics or prevent inconsistencies

We've discussed

domain/type constraints, primary/foreign key

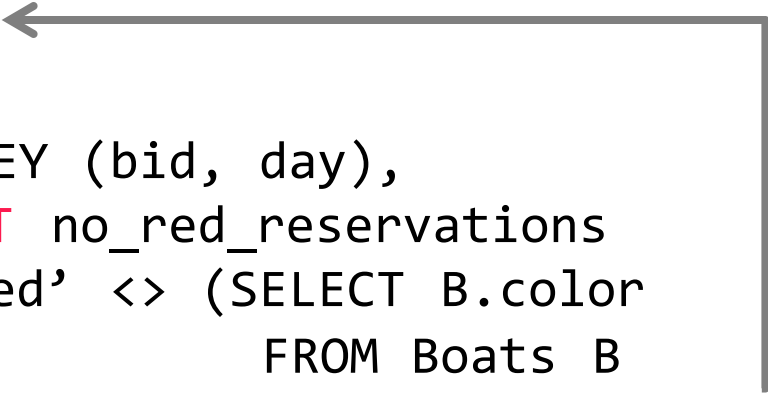
general constraints ←

# Beyond Keys: Table Constraints

Runs when table is not empty

```
CREATE TABLE Sailors(  
    sid int,  
    ...  
    PRIMARY KEY (sid),  
    CHECK (rating >= 1 AND rating <= 10)
```

```
CREATE TABLE Reserves(  
    sid int,  
    bid int, ←  
    day date,  
    PRIMARY KEY (bid, day),  
    CONSTRAINT no_red_reservations  
    CHECK ('red' <> (SELECT B.color  
                     FROM Boats B  
                     WHERE B.bid = bid))
```



Nested subqueries  
Named constraints

# Multi-Relation Constraints

# of sailors + # of boats should be less than 100

```
CREATE TABLE Sailors (  
    sid int,  
    bid int,  
    day date,  
    PRIMARY KEY (bid, day),  
    CHECK (  
        (SELECT COUNT(S.sid) FROM Sailors S)  
        +  
        (SELECT COUNT(B.bid) FROM Boats B)  
        < 100  
    )  
)
```

What if Sailors is empty?

# ASSERTIONS: Multi-Relation Constraints

```
CREATE ASSERTION small_club
CHECK (
    (SELECT COUNT(*) FROM Sailors S)
    +
    (SELECT COUNT(*) FROM Boats B)
    < 100
)
```

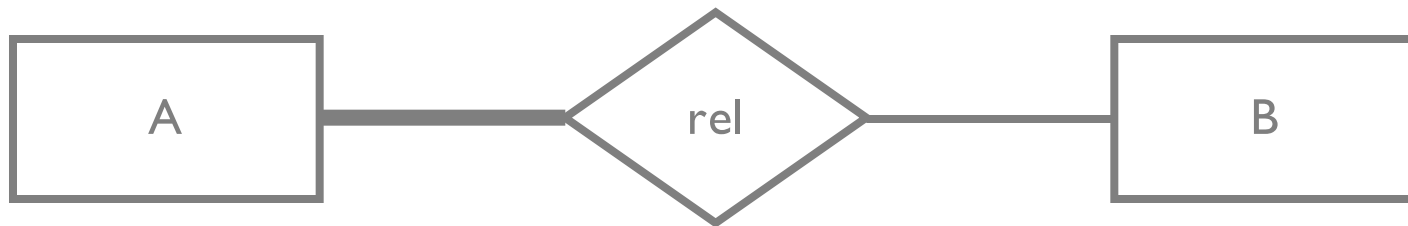
**ASSERTION**s are not associated with any table

# WHAT!

So many things we can't express or don't work!

Assertions

Nested queries in CHECK constraints



# Advanced Stuff

User defined functions

Triggers

WITH

Views

# User Defined Functions (UDFs)

Custom functions that can be called in database

Many languages: SQL, python, C, perl, etc

```
CREATE FUNCTION function_name(p1 type, p2 type, ...)  
RETURNS type
```



# User Defined Functions (UDFs)

Custom functions that can be called in database

Many languages: SQL, python, C, perl, etc

```
CREATE FUNCTION function_name(p1 type, p2 type, ...)  
RETURNS type  
AS $$
```

```
-- logic
```

```
$$ LANGUAGE language_name;
```

# User Defined Functions (UDFs)

Custom functions that can be called in database

Many languages: SQL, python, C, perl, etc

```
CREATE FUNCTION function_name(p1 type, p2 type, ...)
```

```
RETURNS type
```

```
AS $$
```

```
-- logic
```

```
$$ LANGUAGE language_name;
```

# A simple UDF (lang = SQL)

```
CREATE FUNCTION mult1(v int) RETURNS int
AS $$
SELECT v * 100;
$$ LANGUAGE SQL;
```

← Last statement  
is returned

```
CREATE FUNCTION function_name(p1 type, p2 type, ...)
RETURNS type
AS $$
```

```
-- logic
```

```
$$ LANGUAGE language_name;
```

# A simple UDF (lang = SQL)

```
CREATE FUNCTION mult1(v int) RETURNS int
AS $$
SELECT v * 100;
$$ LANGUAGE SQL;
```

```
SELECT mult1(S.age)
FROM   sailors AS S
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Result

int4
220
390
270

# A simple UDF (lang = SQL)

```
CREATE FUNCTION mult1(v int) RETURNS int
AS $$
SELECT $1 * 100;
$$ LANGUAGE SQL;
```

```
SELECT mult1(S.age)
FROM   sailors AS S
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Result

int4
220
390
270

# Process a Record (lang = SQL)

```
CREATE FUNCTION mult2(x sailors) RETURNS int  
AS $$  
SELECT (x.sid + x.age) / x.rating;  
$$ LANGUAGE SQL;
```

```
SELECT mult2(S.*)  
FROM   sailors AS S
```

Sailors

<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Result

int4
3.285
20.5
3.75

# Process a Record (lang = SQL)

```
CREATE FUNCTION mult2(sailors) RETURNS int  
AS $$  
SELECT ($1.sid + $1.age) / $1.rating;  
$$ LANGUAGE SQL;
```

```
SELECT mult2(S.*)  
FROM   sailors AS S
```

Sailors


<u>sid</u>	name	rating	age
1	Eugene	7	22
2	Luis	2	39
3	Ken	8	27

Result

int4
3.285
20.5
3.75

# Procedural Language/SQL(lang = plsql)

```
CREATE FUNCTION proc(v int) RETURNS int
AS $$
DECLARE
    -- define variables
BEGIN
    -- PL/SQL code
END;
$$ LANGUAGE plpgsql;
```



Boilerplate



# Procedural Language/SQL(lang = plsql)

```
CREATE FUNCTION proc(v int) RETURNS int
AS $$
DECLARE
    -- define variables.  VAR TYPE [= value]
    qty int = 10;
BEGIN
    qty = qty * v;
    INSERT INTO blah VALUES(qty);
    RETURN qty + 2;
END;
$$ LANGUAGE plpgsql;
```

# Procedural Code (lang = plpython2u)

```
CREATE FUNCTION proc(v int) RETURNS int
AS $$
import random
return random.randint(0, 100) * v
$$ LANGUAGE plpython2u;
```

Very powerful – can do anything so must be careful

run in a python interpreter with no security protection

plpy module provides database access

```
plpy.execute("select 1")
```

# Procedural Code (lang = plpython2u)

```
CREATE FUNCTION proc(word text) RETURNS text
AS $$
import requests
resp = requests.get('http://google.com/search?q=%s' % v)
return resp.content
$$ LANGUAGE plpython2u;
```

Very powerful – can do anything so must be careful

run in a python interpreter with no security protection

plpy module provides database access

```
plpy.execute("select 1")
```

# Triggers (logical)

def: procedure that runs automatically if specified changes in DBMS happen

```
CREATE TRIGGER name
```

**Event** activates the trigger

**Condition** tests if triggers should run

**Action** what to do

# Triggers (logical)

def: procedure that runs automatically if specified changes in DBMS happen

```
CREATE TRIGGER name  
    [BEFORE | AFTER | INSTEAD OF] event_list  
    ON table
```

**Event** activates the trigger

**Condition** tests if triggers should run

**Action** what to do

# Triggers (logical)

def: procedure that runs automatically if specified changes in DBMS happen

```
CREATE TRIGGER name  
    [BEFORE | AFTER | INSTEAD OF] event_list  
    ON table  
  
    WHEN trigger_qualifications
```

**Event** activates the trigger

**Condition** tests if triggers should run

**Action** what to do

# Triggers (logical)

def: procedure that runs automatically if specified changes in DBMS happen

```
CREATE TRIGGER name  
    [BEFORE | AFTER | INSTEAD OF] event_list  
    ON table  
    [FOR EACH ROW]  
    WHEN trigger_qualifications  
    procedure
```

**Event** activates the trigger

**Condition** tests if triggers should run

**Action** what to do

# Copy new young sailors into special table (logical)

```
CREATE TRIGGER youngSailorUpdate
  AFTER INSERT ON SAILORS
  REFERENCING NEW TABLE NewInserts
  FOR EACH STATEMENT
  INSERT
    INTO YoungSailors(sid, name, age, rating)
    SELECT sid, name, age, rating
    FROM NewInserts N
    WHERE N.age <= 18
```

**Event** activates the trigger

**Condition** tests if triggers should run

**Action** what to do



# Copy new young sailors into special table (logical)

```
CREATE TRIGGER youngSailorUpdate
  AFTER INSERT ON SAILORS
  FOR EACH ROW
  WHEN NEW.age <= 18
  INSERT
    INTO YoungSailors (sid, name, age, rating)
    VALUES (NEW.sid, NEW.name, NEW.age, NEW.rating)
```

**Event** activates the trigger

**Condition** tests if triggers should run

**Action** what to do

# Triggers (logical)

Can be complicated to reason about

Triggers may (e.g., insert) cause other triggers to run

If >1 trigger match an action, which is run first?

¬\_(ツ)\_/

```
CREATE TRIGGER recursiveTrigger
  AFTER INSERT ON SAILORS
FOR EACH ROW
  INSERT INTO Sailors(sid, name, age, rating)
    SELECT sid, name, age, rating
    FROM Sailors S
```

# Triggers (postgres)

```
CREATE TRIGGER name  
    [BEFORE | AFTER | INSTEAD OF] event_list  
    ON table  
    FOR EACH (ROW | STATEMENT)  
    WHEN trigger_qualifications  
    EXECUTE PROCEDURE user_defined_function();
```

PostgreSQL only runs *trigger* UDFs

<http://www.postgresql.org/docs/9.1/static/sql-createtrigger.html>

<http://www.postgresql.org/docs/9.1/static/plpgsql-trigger.html>

# Trigger Example

```
CREATE FUNCTION copyrecord() RETURNS trigger
AS $$
BEGIN
    INSERT INTO blah VALUES(NEW.a);
    RETURN NEW;
END;
$$ LANGUAGE plpgsql;
```

Signature: no args, return type is trigger

Returns NULL or same record structure as modified row

Special variables: OLD, NEW

```
CREATE TRIGGER t_copyinserts BEFORE INSERT ON a
    FOR EACH ROW
    EXECUTE PROCEDURE copyrecord();
```

<http://www.postgresql.org/docs/9.1/static/sql-createtrigger.html>

<http://www.postgresql.org/docs/9.1/static/plpgsql-trigger.html>

# Total boats and sailors < 100

```
CREATE FUNCTION checktotal() RETURNS trigger
AS $$
BEGIN
    IF ((SELECT COUNT(*) FROM sailors) +
        (SELECT COUNT(*) FROM boats) < 100) THEN
        RETURN NEW
    ELSE
        RETURN null;
    END IF;
END;
$$ LANGUAGE plpgsql;
```

```
CREATE TRIGGER t_checktotal BEFORE INSERT ON sailors
FOR EACH ROW
EXECUTE PROCEDURE checktotal();
```

<http://www.postgresql.org/docs/9.1/static/sql-createtrigger.html>

<http://www.postgresql.org/docs/9.1/static/plpgsql-trigger.html>

# You can get into trouble...

```
CREATE FUNCTION addme_bad() RETURNS trigger
AS $$
BEGIN
    INSERT INTO a VALUES (NEW.*);
    RETURN NEW;
END;
$$ LANGUAGE plpgsql;
```

```
CREATE TRIGGER t_addme_bad BEFORE INSERT ON a
    FOR EACH ROW
        EXECUTE PROCEDURE addme_bad();
```

<http://www.postgresql.org/docs/9.1/static/sql-createtrigger.html>

<http://www.postgresql.org/docs/9.1/static/plpgsql-trigger.html>

# You can get into trouble...

```
CREATE FUNCTION addme_ok() RETURNS trigger
AS $$
BEGIN
    IF (SELECT COUNT(*) FROM a) < 100 THEN
        INSERT INTO a VALUES (NEW.a + 1);
    END IF;
    RETURN NEW;
END;
$$ LANGUAGE plpgsql;
```

```
CREATE TRIGGER t_addme_ok BEFORE INSERT ON a
FOR EACH ROW
EXECUTE PROCEDURE addme_ok();
```

<http://www.postgresql.org/docs/9.1/static/sql-createtrigger.html>

<http://www.postgresql.org/docs/9.1/static/plpgsql-trigger.html>

# You can get into trouble...

```
CREATE FUNCTION addme_works() RETURNS trigger
AS $$
BEGIN
    IF (SELECT COUNT(*) FROM a) < 100 THEN
        INSERT INTO a VALUES (NEW.a + 1);
    END IF;
    RETURN NEW;
END;
$$ LANGUAGE plpgsql;
```

```
CREATE TRIGGER t_addme_works AFTER INSERT ON a
FOR EACH ROW
EXECUTE PROCEDURE addme_works();
```

<http://www.postgresql.org/docs/9.1/static/sql-createtrigger.html>

<http://www.postgresql.org/docs/9.1/static/plpgsql-trigger.html>



# WITH

```
WITH RedBoats(bid, count) AS
    (SELECT    B.bid, count(*)
     FROM      Boats B, Reserves R
     WHERE     R.bid = B.bid AND B.color = 'red'
     GROUP BY  B.bid)
SELECT    name, count
FROM      Boats AS B, RedBoats AS RB
WHERE     B.bid = RB.bid AND count < 2
```

Names of unpopular boats

# Views

```
CREATE VIEW view_name  
AS select_statement
```

“tables” defined as query results rather than inserted base data

Makes development simpler

Used for security

Not *materialized*

References to *view\_name* replaced with *select\_statement*

Similar to WITH, lasts longer than one query

# Views

```
CREATE VIEW boat_counts
AS SELECT      bid, count(*)
   FROM        Reserves R
   GROUP BY    bid
   HAVING      count(*) > 10
```

## Used like a normal table

```
SELECT bname
FROM   boat_counts bc, Boats B
WHERE  bc.bid = B.bid
```

Names of popular boats

```
SELECT bname
FROM
    (SELECT bid, count(*)
     FROM Reserves R
     GROUP BY bid
     HAVING count(*) > 10) bc,
    Boats B
WHERE  bc.bid = B.bid
```

Rewritten expanded query

# CREATE TABLE

```
CREATE TABLE <table_name> AS  
  <SELECT STATEMENT>
```

Guess the schema:

```
CREATE TABLE used_boats1 AS  
  SELECT r.bid  
  FROM   Sailors s,  
         Reservations r  
  WHERE  s.sid = r.sid
```

used\_boats1(bid int)

```
CREATE TABLE used_boats2 AS  
  SELECT r.bid as foo  
  FROM   Sailors s,  
         Reservations r  
  WHERE  s.sid = r.sid
```

used\_boats2(foo int)

How is this different than views?

What if we insert a new record into Reservations?

# Summary

SQL is pretty complex

Superset of Relational Algebra SQL99 turing complete!

Human readable

More than one way to skin a horse

Many alternatives to write a query

Optimizer (theoretically) finds most efficient plan

