COMP3121 Assignment 2 A17S1N2

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Question 1

a)

let
$$A^n = \begin{pmatrix} F(n+1) & F(n) \\ F(n) & F(n-1) \end{pmatrix}$$
 (1)

When n=1

$$A^1 = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \tag{2}$$

Assume n = k

$$A^{k} = \begin{pmatrix} F(k+1) & F(k) \\ F(k) & F(k-1) \end{pmatrix}$$

$$\tag{3}$$

let n = k + 1

$$A^{k+1} = AA^k \tag{4}$$

$$= \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} F(k+1) & F(k) \\ F(k) & F(k-1) \end{pmatrix}$$
 (5)

$$= \begin{pmatrix} F(k+1) + F(k) & F(k) + F(k-1) \\ F(k+1) & F(k) \end{pmatrix}$$
 (6)

$$= \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} F(k+1) & F(k) \\ F(k) & F(k-1) \end{pmatrix}$$

$$= \begin{pmatrix} F(k+1) + F(k) & F(k) + F(k-1) \\ F(k+1) & F(k) \end{pmatrix}$$

$$= \begin{pmatrix} F(k+2) & F(k+1) \\ F(k+1) & F(k) \end{pmatrix}$$
(5)
$$= \begin{pmatrix} F(k+2) & F(k+1) \\ F(k+1) & F(k) \end{pmatrix}$$
(7)

Therefore the forumla is true by induction for all n > 0.

b) F(n) can be found in $log_2(n)$ matrix multiplications using a recursive algorithm

```
matrix = ((1, 1), (1,0))
func(n):
    if n == 1:
        return matrix
    if n is even:
        return func(n/2)^2
    if n is odd:
        return func(n-1) * matrix
```

Question 2

a) First we calculate the Karastuba trick.

$$(a+b)(c+d) = ac + ad + bd + bc \tag{1}$$

$$ad + bd = (a+b)(c+d) - ac - bc$$
(2)

Then we substitute (2) into (4)

$$(a+ib)(c+id) = ac + adi + bdi + bc$$
(3)

$$= ac + i(ad + bd) + bc \tag{4}$$

$$= ac + i((a+b)(c+d) - ac - bc) + bc$$
 (5)

b) First we calculate the Karastuba trick.

$$(a+b)^2 = a^2 + b^2 + 2ab (1)$$

$$2ab = (a+b)^2 - a^2 - b^2 (2)$$

Then we substitute (2) into (3)

$$(a+ib)^2 = a^2 + 2abi - b^2 (3)$$

$$= a^{2} + i((a+b)^{2} - a^{2} - b^{2}) - b^{2}$$
(4)

c) By re-arranging by the laws of exponents:

$$(a+ib)^2(c+id)^2 = ((a+ib)(c+id))^2$$

thus from above, we then calculate the middle multiplication using 3 real number multiplications, and then we find the square as above using 2 more real number multiplications.

Question 3

Expand P(x) and Q(x) as follows.

$$P(x) = a_0 + x^{17}(a_{17} + a_{19}x^2 + a_{21}x^4 + a_{23}x^6)$$
$$Q(x) = b_0 + x^{17}(b_{17} + b_{19}x^2 + b_{21}x^4 + b_{23}x^6)$$

let $y = x^2$ so that

$$R_a(y) = a_{17} + a_{19}y + a_{21}y^2 + a_{23}y^3$$

$$R_b(y) = b_{17} + b_{19}y + b_{21}y^2 + b_{23}y^3$$

then

$$P(x)Q(x) = a_0b_0 + x^{17}(a_0R_b(x^2) + b_0R_a(x^2)) + x^{34}R_a(x^2)R_b(x^2)$$

through brute force, we then calculate

 a_0b_0 to require 1 multiplication

$$a_0R_b(x^2)$$
 and $b_0R_a(x^2)$ to require 4 multiplications each

then to multiply $R_a(x^2)R_b(x^2)$ (of degree 3), we require 2(3) + 1 = 7 multiplications using the generalised Karatsuba method.

and so we get 1 + 4 + 4 + 7 = 16 multiplications of large numbers.

Question 4

As P(x) has all 15 roots of unity, and $x^{15} - 1$ and P(x) are both of the same degree and are both monic then it follows that

$$P(x) = x^{15} - 1$$

Question 5

For any input $(a_0, a_1, a_2, ..., a_{2^n-1})$, we can describe a_i 's new position by converting i to n binary places and finding the reversed sequence. e.g. $6 \to 110 \to 011 \to 3$

Question 6

let

$$f_m = \sum_{i+j=m} (j+1)q_j q_i \tag{1}$$

$$p_j = (j+1)q_j \tag{2}$$

then substitute (2) into (1)

$$\sum_{i+j=m} p_j q_i = \vec{p} * \vec{q} \tag{3}$$

as (3) is a linear convolution, f_m can be computed in O(nlog n)

Question 7

- a) Assuming the furtest distance from 1 corner to it's opposite corner, an inefficient algorithm would require (2+3)*2+(n-2)*4 such queries
- b) text b goes here

Question 8

text goes here