

COSC-6590/GSCS-6390

## Game Theory

Game Theoretic Potential Field for Autonomous  
Water Surface Vehicle Navigation Using  
Weather Forecasts

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## Problem Description

# Scenario: Autonomous Surface Vehicles' Navigation

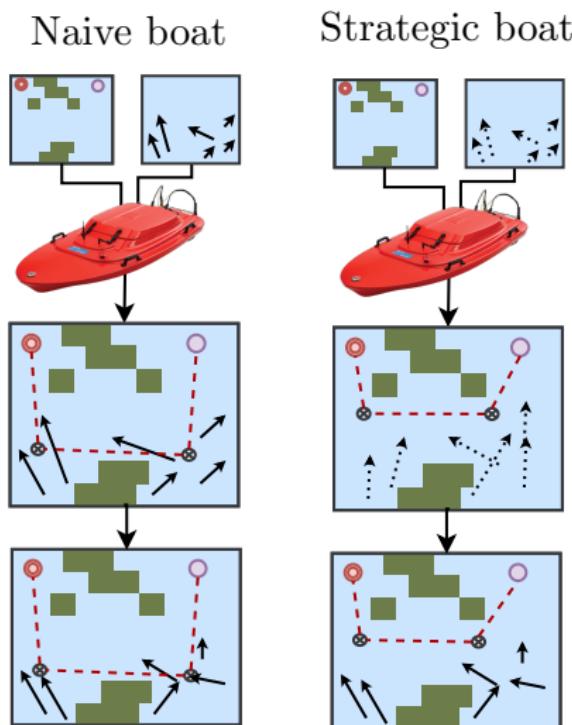


## Autonomous surface vehicles

- Complex coastal navigation
- Dynamic, uncertain weather
- Energy-efficient planning desired



# Addressing Environment Uncertainty



Naive boat

Strategic boat

## Naive boat

- Longer path to ride currents

What if weather is worse than expected?

- Boat regrets its plan

## Strategic boat

- Path avoids worst case

What if weather better than expected?

- Boat does not regret plan

# Contributions of this Research

## **Robust autonomous water surface vehicle navigation**

- Game theory: applied to dynamic programming motion planning for strategic planning in uncertain environments
- Dynamic programming: each iteration yields a feasible plan, with additional iterations optimizing the plan

## **Strategy makes use of real data**

- Large marine region and online water forecasts

## **Outcomes**

- Plan optimal paths that avoid worst-case weather
- High complexity is offset since effective plans can be generated using far fewer iterations than theoretical bound

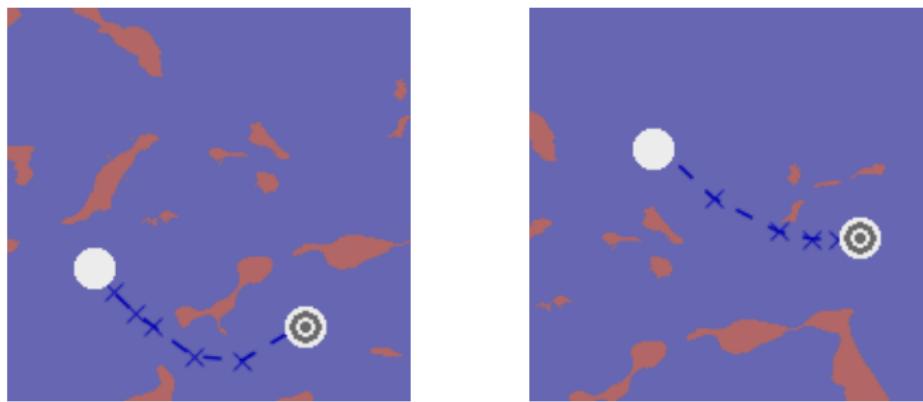
## Software Package Released for Robust Vehicle Navigation

- <https://github.com/ekrell/fujin>

# Game Theory Motion Planning

# Conventional Path Planning

Typically, a single path is generated from a pre-defined start



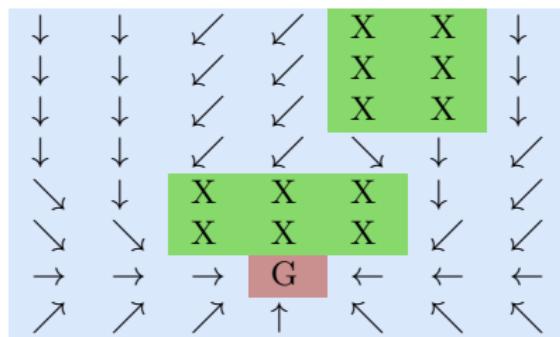
Example of paths found using Particle Swarm Optimization

## Common Problem:

- if the robot gets off-course, it must generate new solution

# Our Approach: Vector Field Motion Plan

## Motion plan to reach goal G



- Here, plan for entire region
- Path from any reachable cell
- Compare: artificial potential field

## Dynamic Programming

- generate optimal motion plan in a bottom-up fashion, starting from the goal

# Dynamic Programming

Cost Map

4	4	4	X	X	X
4	3	3	X	X	3
4	3	2	2	X	2
4	X	2	1	1	1
5	X	X	1	G	1
4	3	2	1	1	1

Path via Cost Lookup

4	4	4	X	X	X
4	3	3	X	X	3
4	3	2	2	X	2
4	X	2	1	1	1
5	X	X	1	G	1
4	3	2	1	1	1

- Each cell has cost to reach **G**
- Each cost depends on other, previously solved, cells
- Resulting map can generate paths

# Environment Uncertainty

## Sources of uncertainty for the scenario at hand

- Obtained from model error (*w.r.t* monitoring stations)
- Learned by robot by several missions

Predicted water force



Error margin of water force



## Forecast uncertainty within estimated error range

- Otherwise, weather unlimited → predictions useless

# Game Theory-inspired robust actions

Each cell's cost the solution of a 0-sum, 2-player game

- Uncertainty: weather has discrete range of actions
- The range based on the prediction's error margin
- Higher action resolution → higher game complexity

Halt	INF	INF	INF	INF	INF
↑	146	146	147	148	149
↓	91	90	89	89	88
←	149	148	147	147	146
→	132	132	133	134	135

Game between 4-way robot and weather, 5 error choices

- Row-player minimizes: **blue**
- Column-player maximizes: **red**
- Mini-max solution: **purple**

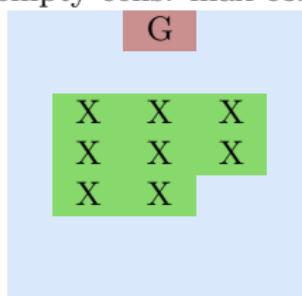
## Algorithms & Problem Formulation

# Algorithm Overview: Single Iteration

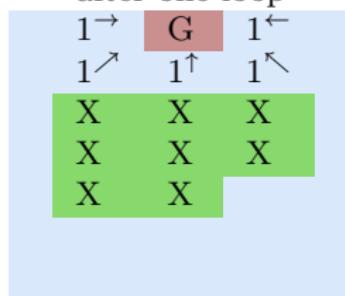
Bottom-up DP: solves cells to move, starting from goal

- Cell cost: work of action at cell + all other cells to goal

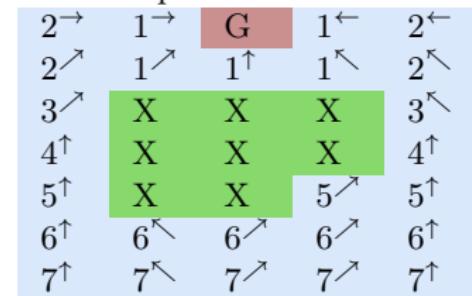
Initial map  
empty cells: max cost



Bottom-up DP  
after one loop



Bottom-up DP  
completed iteration



Problem: each solution based on incomplete information!

- Solution: iterative dynamic programming

# Algorithm Overview: Multiple Iterations

Initial map strong currents			First iteration high costs near current					Second iteration three actions changed				
	G		2 $\rightarrow$	1 $\rightarrow$	G	1 $\leftarrow$	2 $\leftarrow$	2 $\rightarrow$	1 $\rightarrow$	G	1 $\leftarrow$	2 $\leftarrow$
			2 $\nearrow$	1 $\nearrow$	1 $\uparrow$	1 $\nwarrow$	2 $\nwarrow$	2 $\nearrow$	1 $\nearrow$	1 $\uparrow$	1 $\nwarrow$	2 $\nwarrow$
↓	X	X	X	10 $\nearrow$	X	X	X	3 $\nwarrow$	10 $\nearrow$	X	X	X
↓	X	X	X	18 $\uparrow$	X	X	X	4 $\uparrow$	18 $\uparrow$	X	X	X
↓	X	X		26 $\uparrow$	X	X	5 $\nearrow$	5 $\uparrow$	26 $\uparrow$	X	X	5 $\nearrow$
				27 $\uparrow$	28 $\nwarrow$	6 $\nearrow$	6 $\nearrow$	6 $\uparrow$	8 $\nwarrow$	7 $\rightarrow$	6 $\nearrow$	6 $\nearrow$
				28 $\uparrow$	7 $\nearrow$	7 $\nearrow$	7 $\nearrow$	7 $\uparrow$	8 $\nearrow$	7 $\nearrow$	7 $\nearrow$	7 $\nearrow$

Each iteration's choices dependent on neighborhood

- Center: chooses high-cost since neighbors not solved
- Right: cells choose lower-cost neighbors next time
- Several iterations for lower-cost strategies to propagate

# Problem Formulation: Environment

Occupancy Grid  $\mathcal{R} := M \times N$  matrix where

- Value of 1 indicates occupied
- Value of 0 indicates free

$J :=$  number of force vectors affecting  $R$

- $u$  components: force grids  $F^u := \{F_1^u, F_2^u, \dots, F_J^u\}$ 
  - each component is an  $M \times N$  matrix of weather force
- $v$  components: force grids  $F^v := \{F_1^v, F_2^v, \dots, F_J^v\}$ 
  - each component is an  $M \times N$  matrix of weather force

Error Grids  $E := \{E_1, E_2, \dots, E_J\}$ :  $M \times N$  errors of each force

- $E_i \in E$  is the error range of  $F_i^u$  and  $F_i^v$ .

$x_{\text{goal}} := (\text{row} \in M, \text{col} \in N)$ : coordinates of goal in  $\mathcal{R}$ .

# Problem Formulation: Players & State Transition

Players  $P := \{P_1 := \text{robot} ; P_2 := \text{environment}\}$

Stages  $\mathbf{K} := \{1 \dots K\}$

Action space for  $P_1$ :  $U := \{U_1^1 \dots U_K^B\}$

- $B$ : number of actions - discrete heading angles, and *halt*
- Control action for  $P_1$  is  $u_k^b \in U$ .

Action space for  $P_2$ :  $\Theta := \{\Theta_1^1 \dots \Theta_K^A\}$

- $A$ : number of actions - a tuple of selected latitudinal and longitudinal components
- Control action for  $P_2$  is  $\theta_k^\alpha$

State space  $X$

- At each stage  $k$ , the game state is  $x_k$
- Each state has an associated robot location

$$x_{\text{loc}} := (\text{row} \in M, \text{col} \in N)$$

State transition function  $f_k : x_{k+1} = f_k(x_k, u_k, \theta_k^\alpha)$

# Problem Formulation: 2-Player, Zero-Sum Games

Game  $G := B \times A$

- Rows  $\rightarrow P_1$  action choices
- Columns  $\rightarrow P_2$  action choices

$$g_{b,a} := \begin{cases} 0 & \text{if } x_{\text{loc}} = x_{\text{goal}} \\ t_{b,a} & \text{if } R(x_{\text{loc}}) = 0 \\ D_{\max} & \text{if } R(x_{\text{loc}}) = 1 \end{cases}$$

Cost of joint strategy  $t_{b,a} := work_{b,a} + cost2go(f_k(x_k, u_k, \theta_k^\alpha))$

- $work_{b,a}$ : Applied work done by  $P_1$
- $cost2go(x_k)$ : cost to go to goal after reaching state  $x_k$ .

$actgrid(x_k)$ :  $P_1$  action after reaching state  $x_k$   
 ↑ the motion plan itself.

# Algorithms, Quick Look

## Algorithm 1: DynamicPlanner

- Initialize  $cost2go$  values to maximum,  $actgrid$  values to NULL
- For each dynamic programming iteration  $I_{DP}$ :
  - $cost2go, actgrid \leftarrow \text{DynamicPlannerIter}(cost2go, actgrid)$
- return  $cost2go, actgrid$

## Algorithm 2: DynamicPlannerIter

- $Q \leftarrow$  empty FIFO queue
- Assign  $cost2go_{goal} := 0$ ,  $actgrid_{goal} := \text{"halt"}$
- Add neighbors of  $x_{goal}$  to  $Q$
- While  $Q$  is not empty:
  - $c \leftarrow \text{dequeue } Q$
  - Enqueue neighbors of  $c$  that have never been enqueued
  - $cost2go_c, actgrid_c \leftarrow \text{NashSolver}(c)$
- return  $cost2go, actgrid$

# Algorithms, Quick Look

## Algorithm 3: NashSolver

- Init  $P_1$  mixed policy as  $B$ -length 0-vector,     $P_2$  as  $A$ -length 0-vector
- $P_1$ 's selected action (row)  $\leftarrow 0$ ,     $P_2$ 's selected action (column)  $\leftarrow 0$
- For each iteration  $I_{nash}$ :
  - $P_1$  selects counter row to minimize game value on  $P_2$ 's column
  - $P_2$  selects counter column to minimize game value on  $P_1$ 's row
  - $P_1$  increments mixed policy at index specified by row
  - $P_2$  increments mixed policy at index specified by column
- $P_1$  action  $\leftarrow$  most frequently chosen row
- $P_2$  action  $\leftarrow$  most frequently chosen column
- Cost  $\leftarrow$  game value at  $P_1$  action and  $P_2$  action
- return cost (*cost2go*),  $P_1$  action (*actgrid*)

Based on 2-player, 0-sum game solver code by Raymond Hettinger:

[code.activestate.com/recipes/496825-game-theory-payoff-matrix-solver](http://code.activestate.com/recipes/496825-game-theory-payoff-matrix-solver)

# Algorithm 1: DynamicPlanner

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**Algorithm 1: DynamicPlanner:** Motion planner

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**Input:**  $x_{goal}$ ,  $\mathcal{R}$ ,  $D_{max}$ ,  $I_{DP}$

**Output:**  $cost2go$ ,  $actgrid$

```
1 Set M  $\leftarrow$  number of rows in  $R$ ;  
2 Set N  $\leftarrow$  number of cols in  $R$ ;  
3 Initialize grid  $cost2go \leftarrow M \times N$ , all cells having value  $D_{max}$ ;  
4 Initialize grid  $actgrid \leftarrow M \times N$  NULL matrix;  
    /* Iteratively update  $cost2go$ ,  $actgrid$  */  
5 for  $i$  in range  $0 \dots I_{DP}$  do  
    /* Call Algorithm 2 */  
6    $cost2go, actgrid \leftarrow \textbf{DynamicPlannerIter}(x_{goal},$   
      $cost2go, actgrid, D_{max})$ ;  
7 return  $cost2go$ ,  $actgrid$ 
```

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# Algorithm 2: DynamicPlannerIter

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**Algorithm 2: DynamicPlannerIter:**


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Assigns costs, actions to each cell in occupancy grid  $\mathcal{R}$

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**Input:**  $x_{\text{goal}}$ ,  $\text{cost2go}$ ,  $\text{actgrid}$ ,  $D_{\max}$

**Output:**  $\text{cost2go}$ ,  $\text{actgrid}$

*/\* Q: cells that need assignment \*/*

1 Initialize FIFO queue  $Q \leftarrow \emptyset$ ;

*/\* V: remembered all added cells \*/*

2 Initialize set  $V \leftarrow \emptyset$ ;

3 Start at  $x_{\text{goal}}$ ;

4 Set  $\text{cost2go}_{\text{goal}} = 0$ ,  $\text{actgrid}_{\text{goal}} = \text{"halt"}$ ;

5 Add neighborhood grid cells to  $Q$  and to  $V$ ;

6 **while**  $Q \neq \emptyset$  **do**

7   Cell  $c \leftarrow$  Dequeue  $Q$ ;

8   Add neighborhood cells to  $Q$  if  $c$  not in  $V$ ;

9   Add neighborhood cells to  $V$  if  $c$  not in  $V$ ;

*/\* Call Algorithm 3 \*/*

10    $g_{b,a}, P_1^{\text{policy}}, P_2^{\text{policy}} \leftarrow \text{NashSolver}(c)$ ;

11   Set  $\text{cost2go}(c) = g_{b,a}$ ;

12   Set  $\text{actgrid}(c) = P_1^{\text{policy}}$ ;

13 **return**  $\text{cost2go}$ ,  $\text{actgrid}$

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# Algorithm 3: NashSolver

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**Algorithm 3: NashSolver:**


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Approximate, iterative 0-sum 2-player game solver

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**Input:**  $G, I_{\text{nash}}$

**Output:** cost  $g_{b,a}$ , policies  $P_1^{\text{policy}} \in U$  and  $P_2^{\text{policy}} \in \Theta$

- 1  $B \leftarrow$  number of rows in  $G$ ;
- 2  $A \leftarrow$  number of columns in  $G$ ;
- 3 Init  $P_1^{\text{policy}} \leftarrow B$ -length 0-vector;
- 4 Init  $P_2^{\text{policy}} \leftarrow A$ -length 0-vector;
- 5 /\* Record whenever action selected \*/
- 6 Set  $P_1^a, P_2^a$  action  $\leftarrow 0$ ;
- 7 for  $i$  in range  $0 \dots I_{\text{nash}}$  do
- 8    $P_1^a \leftarrow {}_b(G(b, P_2^a)), b \in B$  ;
- 9    $P_2^a \leftarrow {}_a(G(P_1^a, a)), a \in A$  ;
- 10   /\* Increment action count \*/
- 11    $P_1^{\text{policy}}[P_1^a] \leftarrow 1$ ;
- 12    $P_2^{\text{policy}}[P_2^a] \leftarrow 1$ ;
- 13 return  $P_1^{\text{policy}}, P_2^{\text{policy}}$

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# Theorems

## Theorem 1.- Global Optimum Convergence

After a finite number of dynamic programming iterations, the motion plan converges to a global optimum. The maximum number of executions is proportional to the region dimensions.

Large regions will be a large planning search space

## Theorem 2.- Single-Iteration Feasibility

A single iteration of dynamic programming creates an unobstructed plan to the goal from any reachable cell. Each terminates in a deterministic, finite number of iterations.

Each iteration a feasible motion plan: intermediate solution

# Computational Complexity: Big-O Analysis

$$\mathcal{O}^{NASH} = \mathcal{O}((B + A) \times I_{nash})$$

$$\mathcal{O}^{DPI} = \mathcal{O}(M \times N \times \mathcal{O}^{NASH})$$

Substituting  $\mathcal{O}^{NASH}$

$$\mathcal{O}^{DPI} = \mathcal{O}(M \times N \times ((B + A) \times I_{nash}))$$

Also

$$\mathcal{O}^{DP} = \mathcal{O}(I_{DP} \times \mathcal{O}^{DPI})$$

Substituting  $\mathcal{O}^{DPI}$

$$\mathcal{O}^{DP} = \mathcal{O}(I_{DP} \times (M \times N \times ((B + A) \times I_{nash})))$$

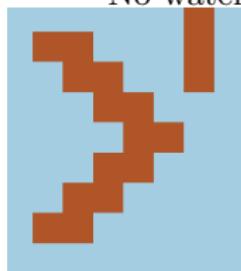
**Very high complexity!**

- Can intermediate solutions give usable motion plans?

## Experimental Results

# Simple Region Results – Synthetic Data

No water currents

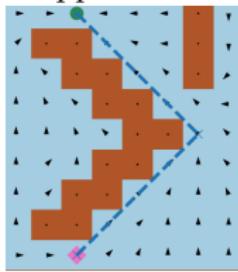
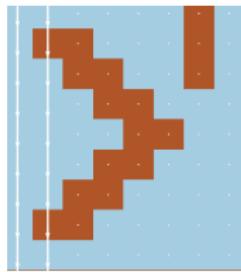


**Goal:** green circle

**Start:** pink diamond

- No currents → shortest path

Certain currents that oppose boat

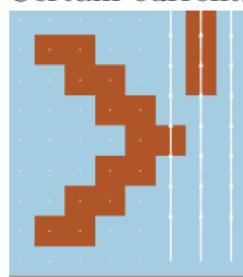


Current direction is directly  
against the boat

- Boat chooses a longer path  
to save energy

# Simple Region Results – Synthetic Data

Certain currents that help boat



Currents directed towards goal

- Boat chooses a longer path to ride the currents

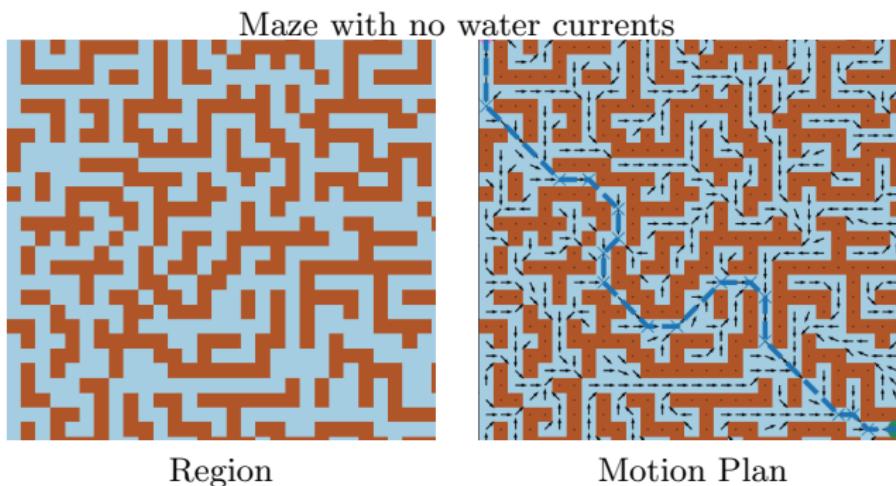
Uncertain, antagonistic currents



Antagonistic weather modifies currents to oppose boat

- Error range: bounds weather
- Ride currents when it can, shortest-path when it cannot

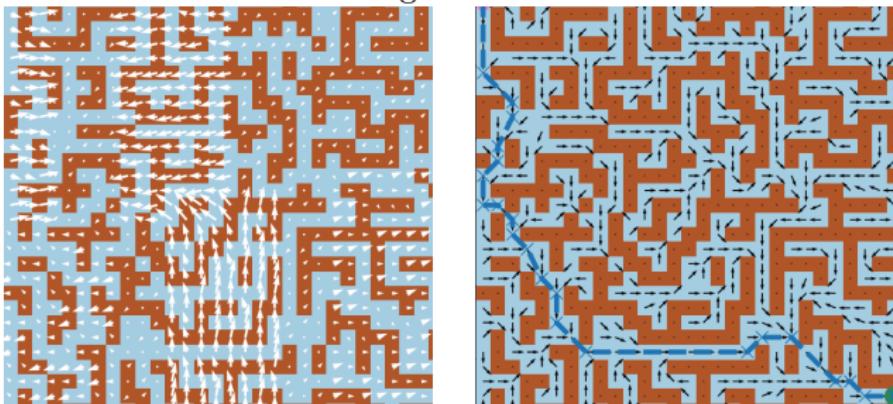
# Maze Region Results – Synthetic Data



Dynamic programming gives optimal maze solution

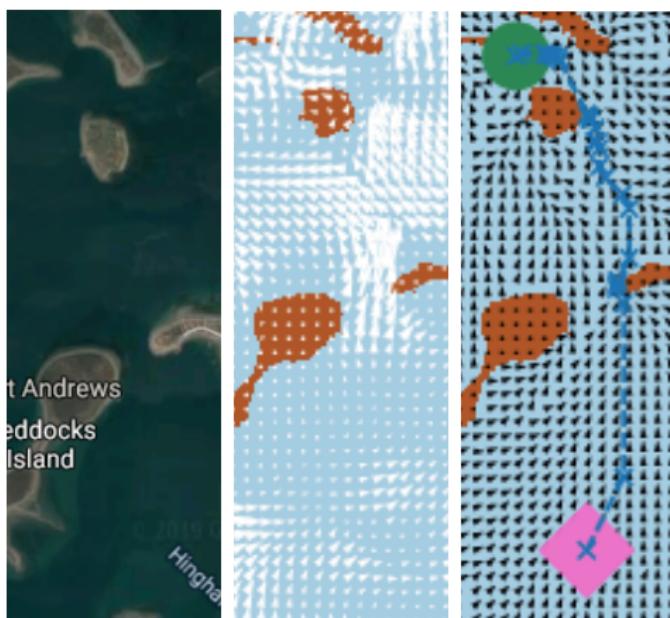
# Maze Region Results – Synthetic Data

Maze with antagonistic water currents



Currents can have a dramatic impact on best solution

# Marine Region – Real Data Results



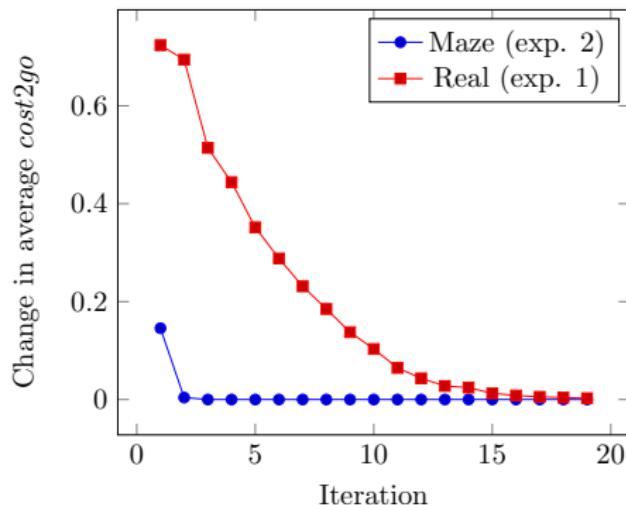
Massachusetts Bay region

- Northeast Coastal Ocean Forecast System (NECOFS)
- Antagonistic currents

Plan Shown: 20 iterations

- Strategic planning evident

# Dynamic Programming Convergence

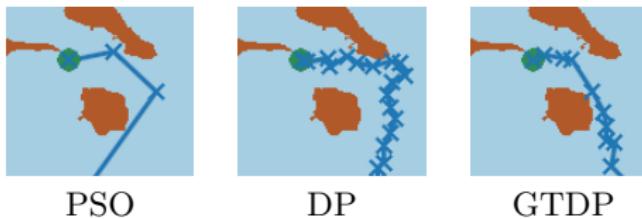


Real marine region

- Convergence shows very little improvement in average cost after only 20 iterations
- Feasibility of onboard planning despite high complexity

# Comparison with Particle Swarm Optimization

- PSO: 500 iterations (no uncertainty)
- DP: dynamic programming, 20 iterations (no uncertainty)
- GTDP: game theory dynamic programming, 20 iterations



Each path applied to both the certain and worst-case scenarios

Solution	Work, static forces	Work, antagonistic forces
<i>PSO</i>	294349	320234
<i>DP</i>	345085	368574
<i>GTDP</i>	297142	319969

Blue cells indicate the scenario used for generating that solution

## Conclusions

# Conclusions & Future Work

## Conclusions

- Game theory allows robot to handle worst-case weather
- Real data suggests boat can benefit from strategic planning
- High complexity offset by ability to use early iterations
- Online forecasts such as NECOFs enable better autonomous navigation

## Future Work

- Incorporate dynamic water currents
- Incorporate dynamic vehicle model
- Dynamically consider currents too strong for boat
- Extend to 3D (underwater and aerial applications)

# Conclusions & Future Work

## Towards a Real Robotic Boat

- Building airboat for shallow-water applications
- Modified Zelos ProBoat and EMILY ERS



End of Presentation

## **Game Theoretic Potential Field for Autonomous Water Surface Vehicle Navigation Using Weather Forecasts**

Questions?