

Compilation

0368-3133

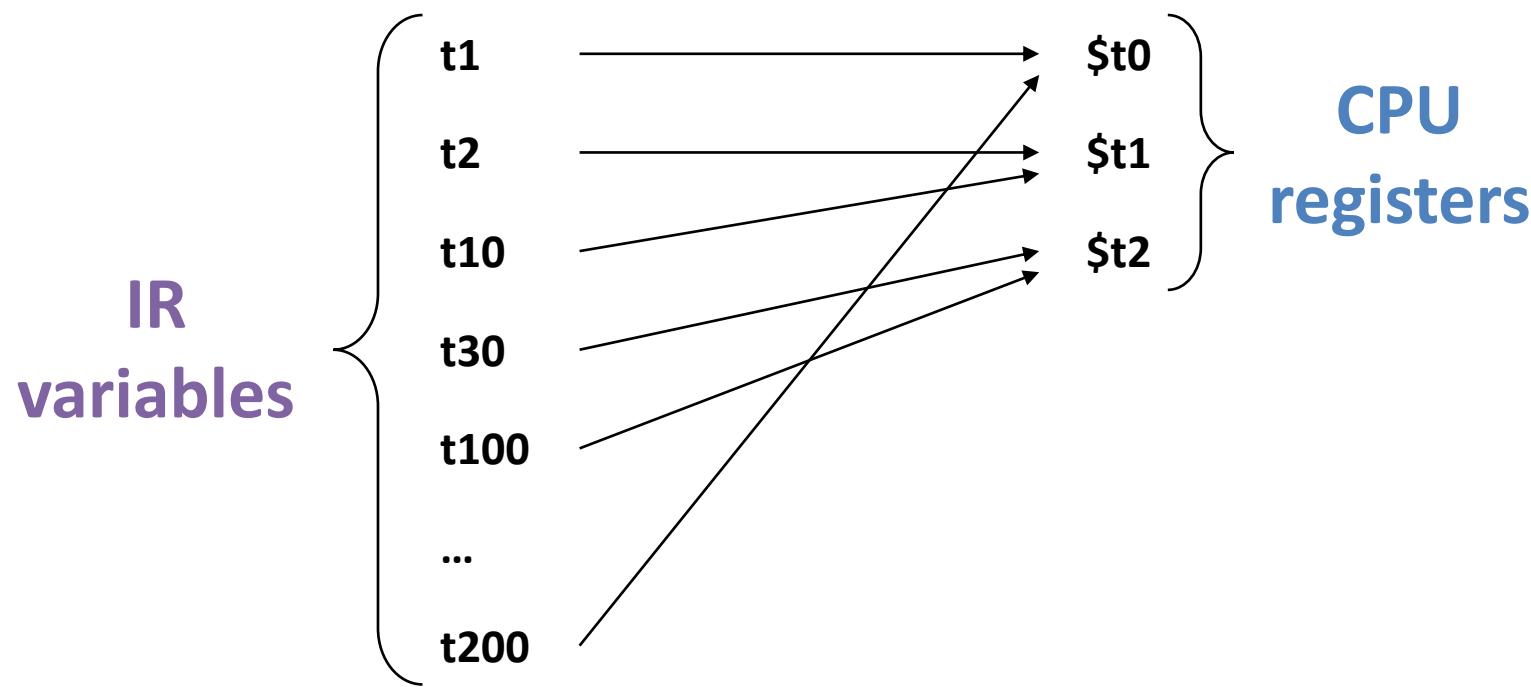
Tutorial 12:
Register Allocation

Today's Plan

- **Part I:** Register allocation for the project
 - We may give up and abort
- **Part II:** General register allocation
 - Spilling & coalescing
 - Never give up!

Register Allocation (in the project)

- Typically, more IR variables than CPU registers
- Reduce the number of IR variables:
 - Without affecting the behavior of the program



Register Allocation

For each function:

1. Construct the CFG (from the IR)
2. Run liveness analysis
3. Construct the interference graph
4. Compute a *k-coloring* of the graph
5. Use the coloring to build the required mapping

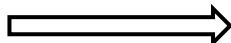
Register Allocation

For each function:

1. **Construct the CFG (from the IR)**
2. Run liveness analysis
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4. Compute a *k-coloring* of the graph
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Control Flow Graph

```
void f(int x, int y, int z) {  
    int a = x * (y - z);  
    if (a) {  
        a = a + 1;  
    }  
    int b = a;  
}
```



```
t1 = x  
t2 = y  
t3 = z  
t4 = sub t2, t3  
t5 = mul t1, t4  
a = t5  
t6 = a  
bne t6, 1, end  
t7 = a  
t8 = 1  
t9 = add t7, t8  
a = t9  
end:  
t10 = a  
b = t10
```

t1 = x

t2 = y

t3 = z

t4 = sub t2, t3

t5 = mul t1, t4

a = t5

t6 = a

bne t6, 1, end

t7 = a

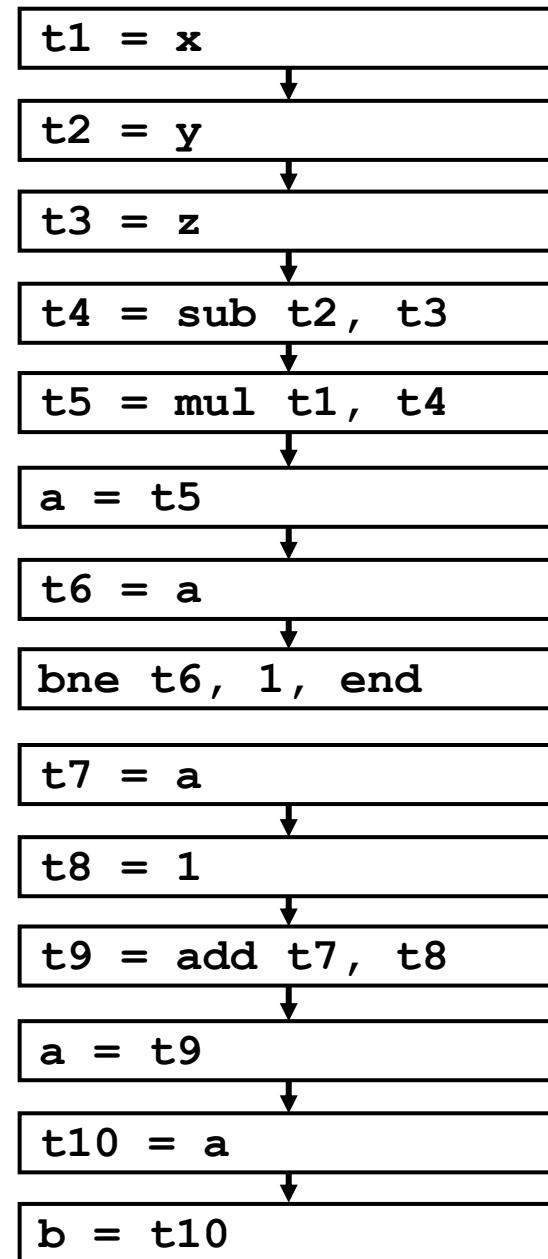
t8 = 1

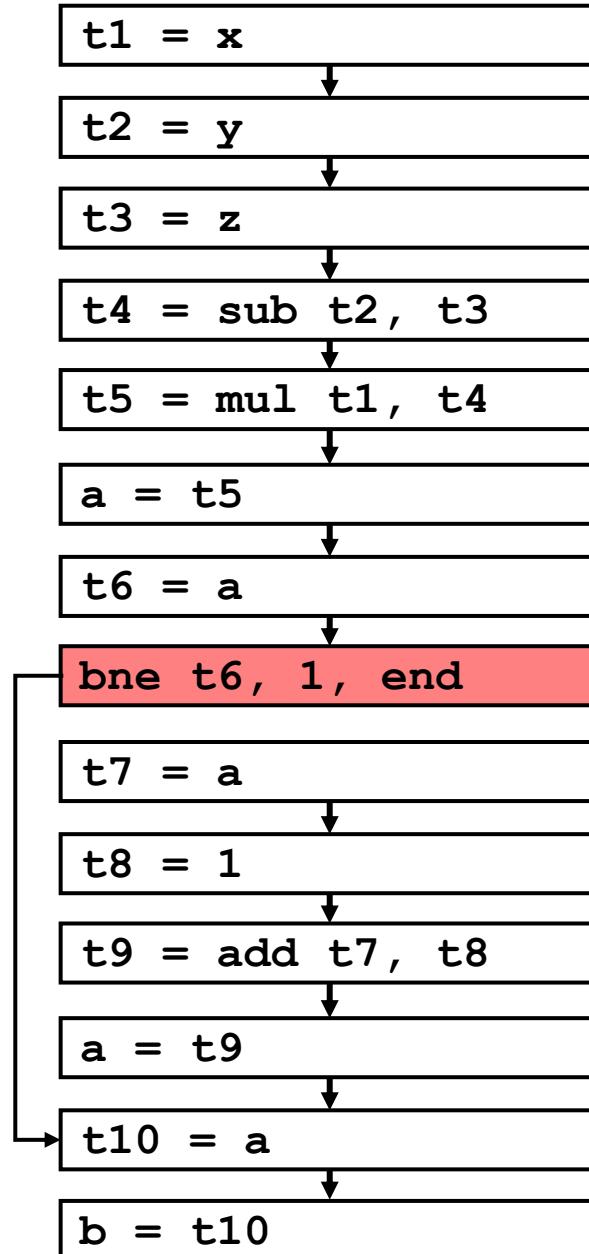
t9 = add t7, t8

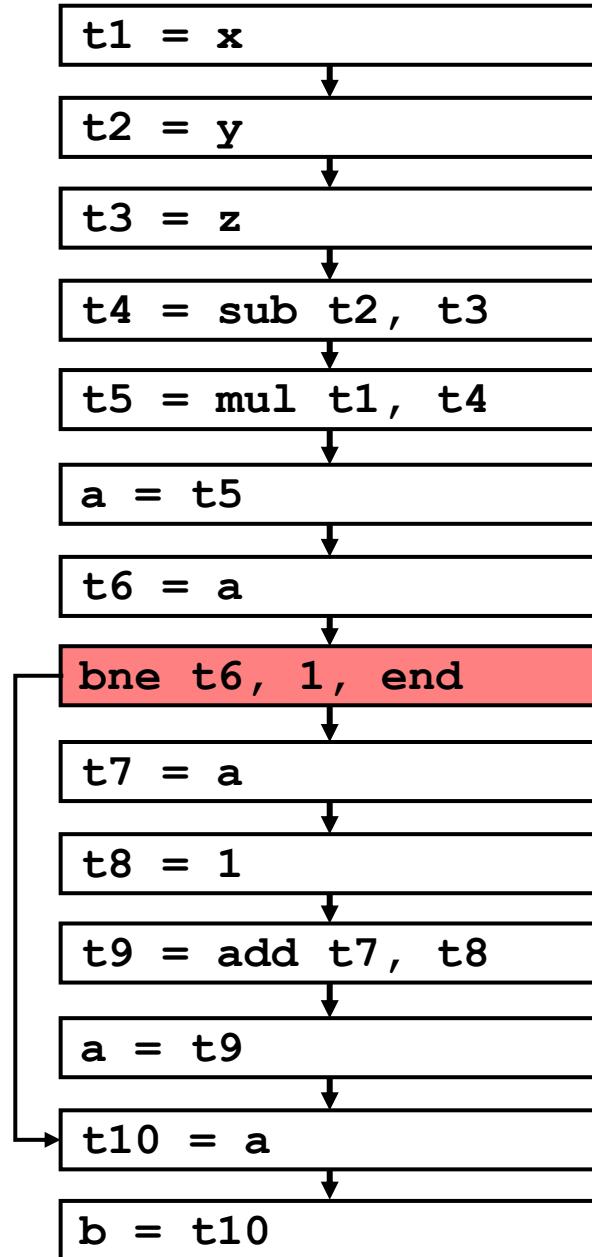
a = t9

t10 = a

b = t10







Register Allocation

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1. Construct the CFG (from the IR)
- 2. Run liveness analysis**
3. Construct the interference graph
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Liveness Analysis

- Determine *live variables* at each program location
- Variable x is *live* at location n if:
 - There is a path from n where x is read before it is overwritten

n: `t1 = 1`

`t2 = 9`

`t2 = add t2, t3`

t3 is **live**

n: `t1 = 1`

`t3 = 9`

`t4 = add t3, t4`

t3 is **dead**

Liveness Analysis

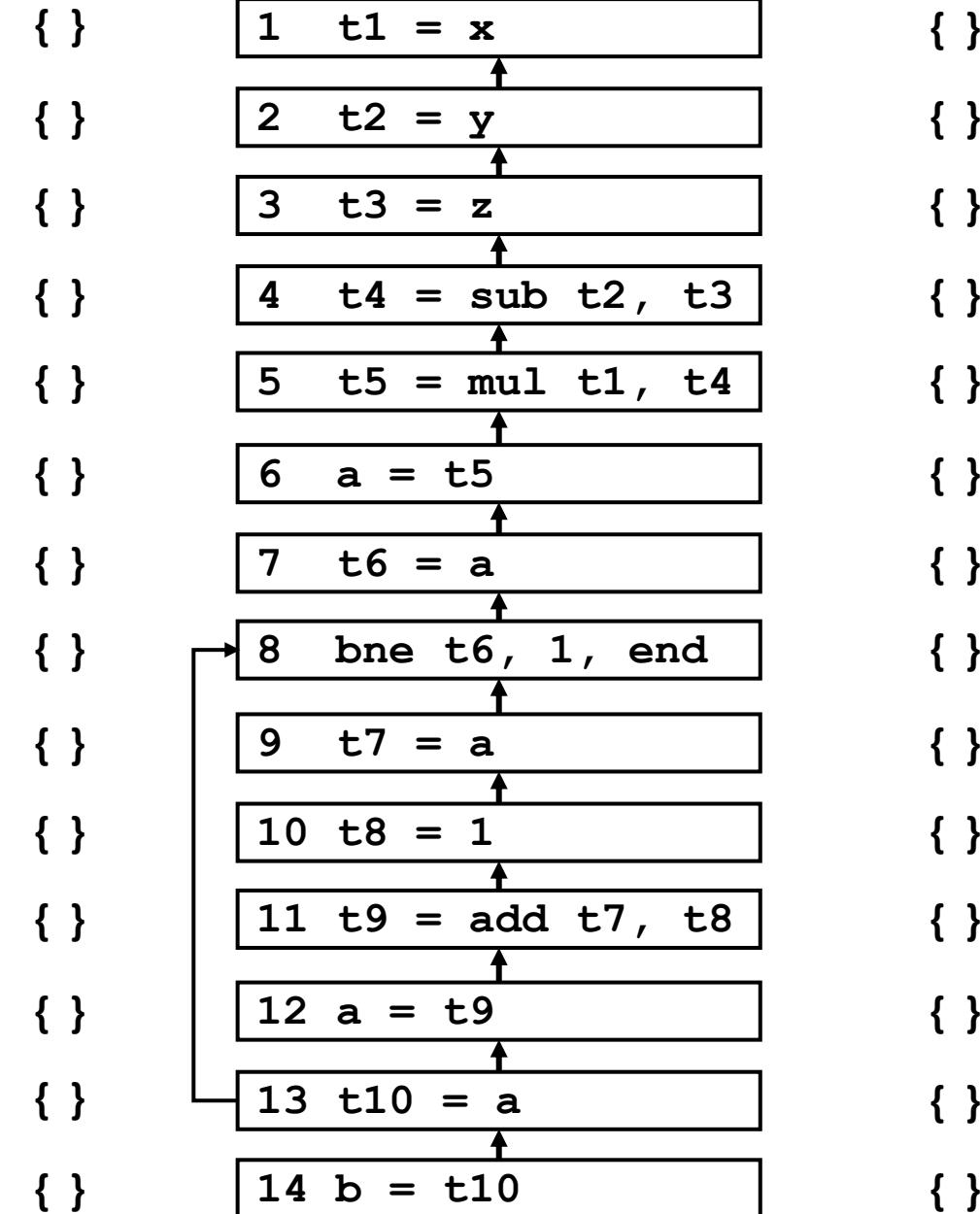
- $\text{Dom} = \wp(\text{Temps})$ (In our context, in general $\text{Dom} = \wp(\text{Vars})$)
- $\sqsubseteq = \subseteq, \sqcup = \cup$
- D = backward
- $I = \emptyset$

Statement	kill_i	gen_i
$x = \text{expr}$	$\{x\}$	Temporaries used in exp
stmt	\emptyset	Temporaries used in stmt

Work list = {14}

out

in



Work list = {13}

out

{ }

1 t1 = x

{ }

2 t2 = y

{ }

3 t3 = z

{ }

4 t4 = sub t2, t3

{ }

5 t5 = mul t1, t4

{ }

6 a = t5

{ }

7 t6 = a

{ }

8 bne t6, 1, end

{ }

9 t7 = a

{ }

10 t8 = 1

{ }

11 t9 = add t7, t8

{ }

12 a = t9

{ t10 }

13 t10 = a

{ }

14 b = t10

in

{ }

{ }

{ }

{ }

{ }

{ }

{ }

{ }

{ }

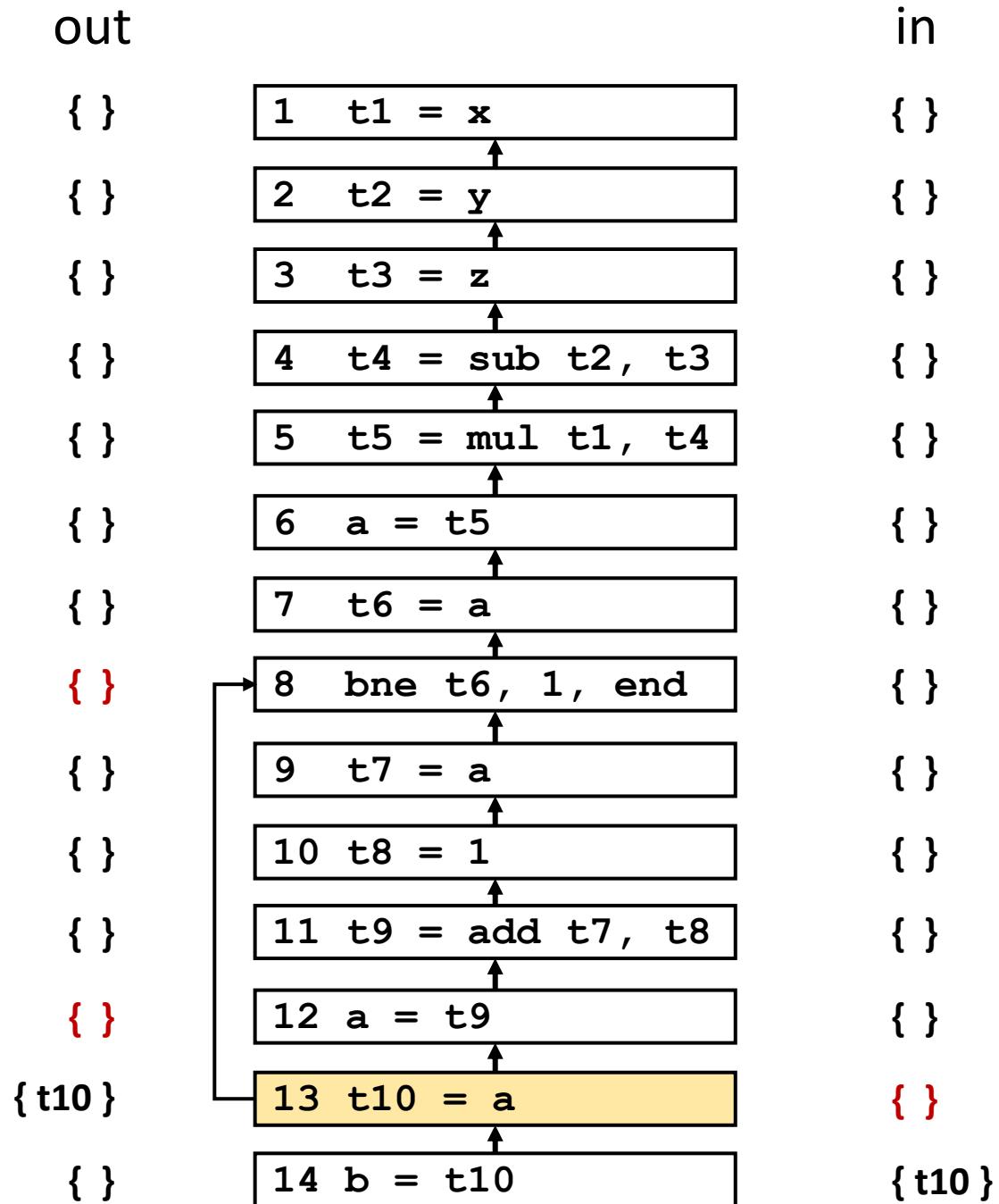
{ }

{ }

{ }

{ t10 }

Work list = {12,8}



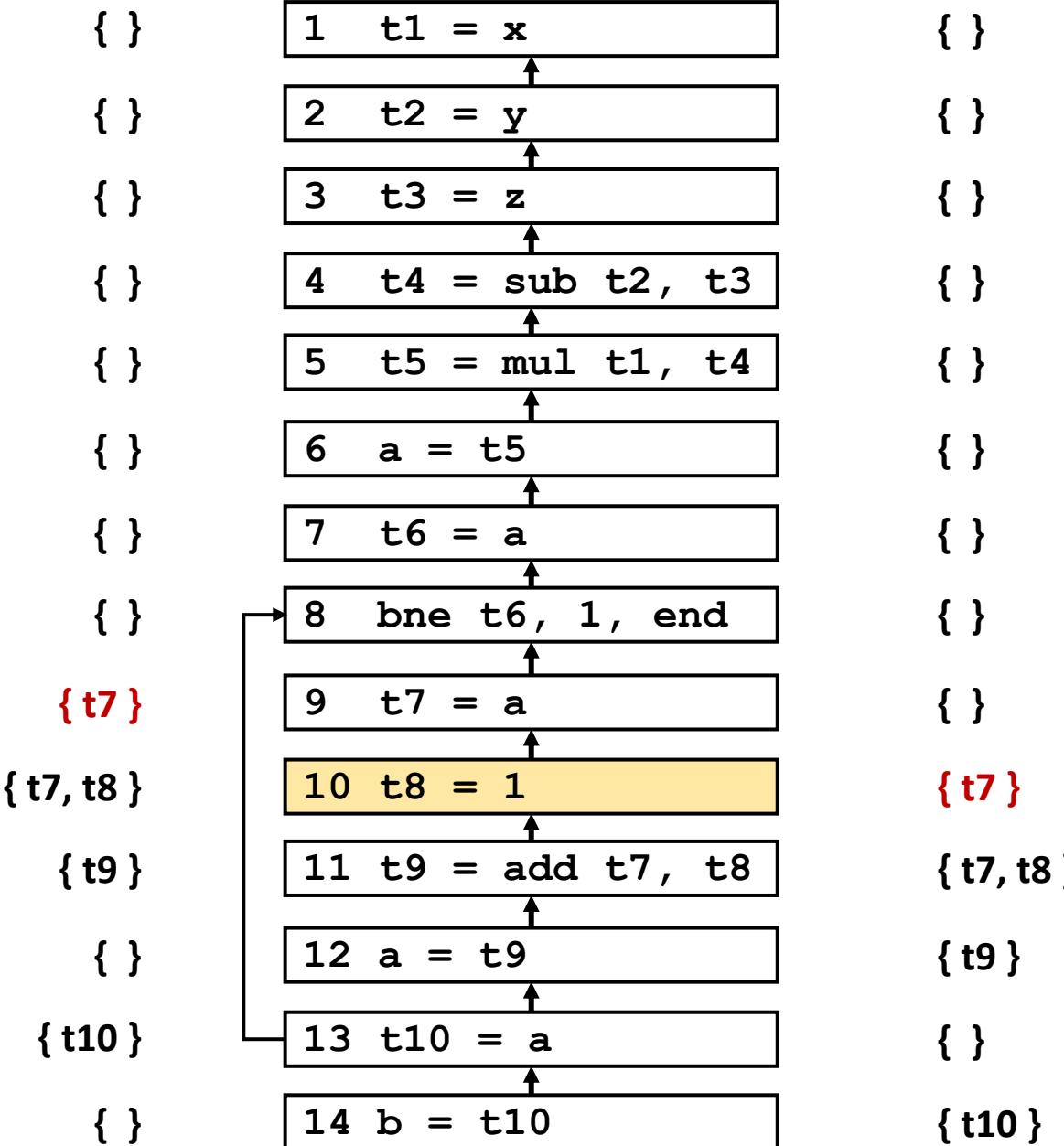
	out		in
Work list = {11,8}	{ }	1 t1 = x	{ }
	{ }	2 t2 = y	{ }
	{ }	3 t3 = z	{ }
	{ }	4 t4 = sub t2, t3	{ }
	{ }	5 t5 = mul t1, t4	{ }
	{ }	6 a = t5	{ }
	{ }	7 t6 = a	{ }
	{ }	8 bne t6, 1, end	{ }
	{ }	9 t7 = a	{ }
	{ }	10 t8 = 1	{ }
{t9}	{ }	11 t9 = add t7, t8	{ }
	{ }	12 a = t9	{t9}
{t10}	{ }	13 t10 = a	{ }
	{ }	14 b = t10	{t10}

	out		in
Work list = {10,8}	{ }	1 t1 = x	{ }
	{ }	2 t2 = y	{ }
	{ }	3 t3 = z	{ }
	{ }	4 t4 = sub t2, t3	{ }
	{ }	5 t5 = mul t1, t4	{ }
	{ }	6 a = t5	{ }
	{ }	7 t6 = a	{ }
	{ }	8 bne t6, 1, end	{ }
	{ t7, t8 }	9 t7 = a	{ }
	{ t9 }	10 t8 = 1	{ }
	{ t10 }	11 t9 = add t7, t8	{ t7, t8 }
	{ }	12 a = t9	{ t9 }
	{ }	13 t10 = a	{ }
	{ }	14 b = t10	{ t10 }

Work list = {9,8}

out

in



Work list = {8}

out

{ }

1 t1 = x

{ }

2 t2 = y

{ }

3 t3 = z

{ }

4 t4 = sub t2, t3

{ }

5 t5 = mul t1, t4

{ }

6 a = t5

{ }

7 t6 = a

{ }

8 bne t6, 1, end

{t7}

9 t7 = a

{t7, t8}

10 t8 = 1

{t9}

11 t9 = add t7, t8

{ }

12 a = t9

{t10}

13 t10 = a

{ }

14 b = t10

in

{ }

{ }

{ }

{ }

{ }

{ }

{ }

{ }

{ }

{t7}

{t7, t8}

{t9}

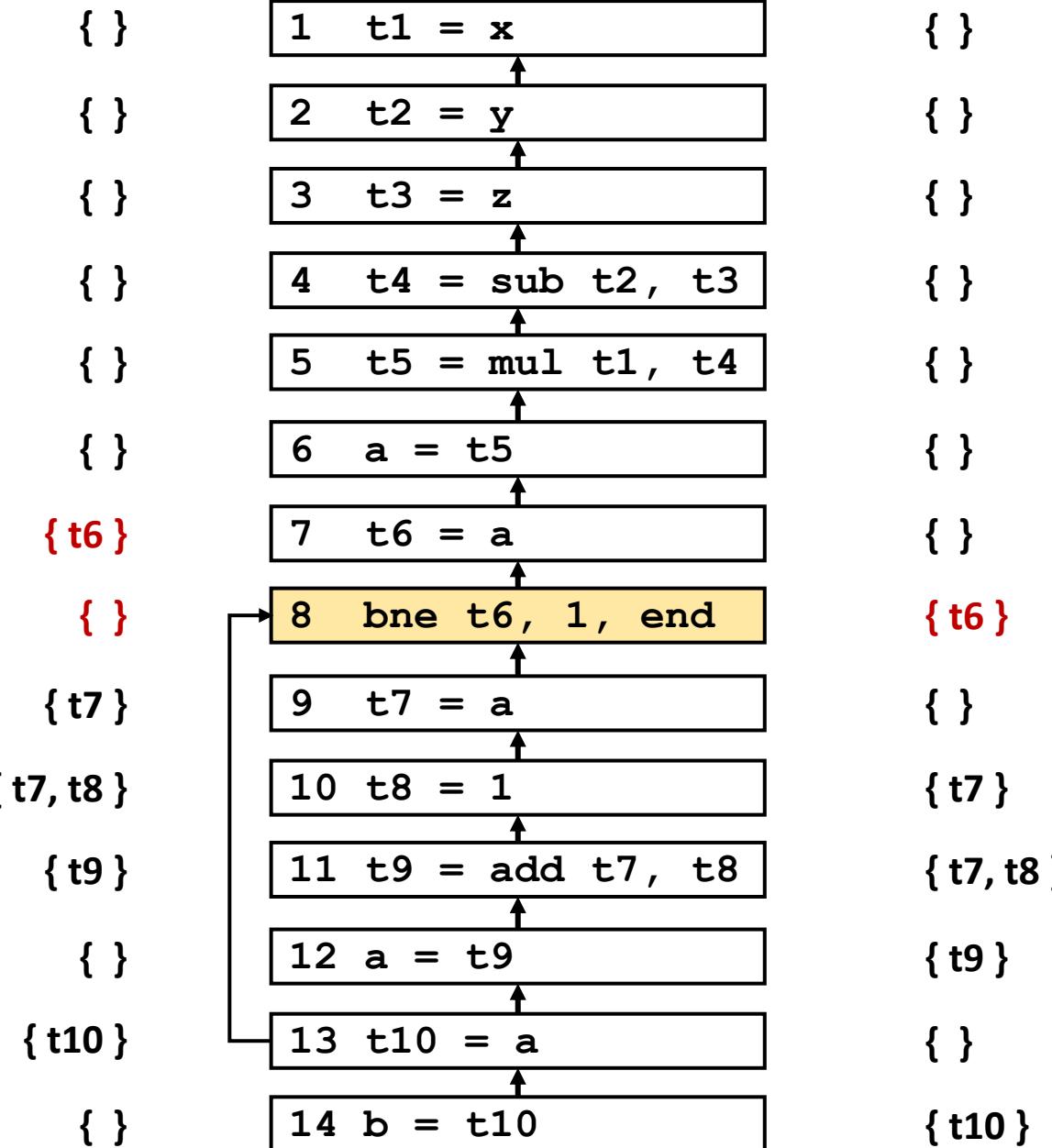
{ }

{t10}

Work list = {7}

out

in



Work list = {6}

out

{ }

1 t1 = x

{ }

2 t2 = y

{ }

3 t3 = z

{ }

4 t4 = sub t2, t3

{ }

5 t5 = mul t1, t4

{ }

6 a = t5

{ t6 }

7 t6 = a

{ }

8 bne t6, 1, end

{ t7 }

9 t7 = a

{ t7, t8 }

10 t8 = 1

{ t9 }

11 t9 = add t7, t8

{ }

12 a = t9

{ t10 }

13 t10 = a

{ }

14 b = t10

in

{ }

{ }

{ }

{ }

{ }

{ }

{ }

{ t6 }

{ }

{ t7 }

{ t7, t8 }

{ t9 }

{ }

{ t10 }

	out	in
Work list = {5}	{ }	{ }
	1 t1 = x	{ }
	2 t2 = y	{ }
	3 t3 = z	{ }
	4 t4 = sub t2, t3	{ }
{ t5 }	5 t5 = mul t1, t4	{ }
	6 a = t5	{ t5 }
{ t6 }	7 t6 = a	{ }
	8 bne t6, 1, end	{ t6 }
{ t7 }	9 t7 = a	{ }
{ t7, t8 }	10 t8 = 1	{ t7 }
{ t9 }	11 t9 = add t7, t8	{ t7, t8 }
	12 a = t9	{ t9 }
{ t10 }	13 t10 = a	{ }
	14 b = t10	{ t10 }

	out		in
Work list = {4}	{ }	1 t1 = x	{ }
	{ }	2 t2 = y	{ }
	{ }	3 t3 = z	{ }
{ t1, t4 }	4 t4 = sub t2, t3		{ }
	5 t5 = mul t1, t4		{ t1, t4 }
	{ t5 }	6 a = t5	{ t5 }
	{ t6 }	7 t6 = a	{ }
	{ }	8 bne t6, 1, end	{ t6 }
	{ t7 }	9 t7 = a	{ }
{ t7, t8 }	10 t8 = 1		{ t7 }
	{ t9 }	11 t9 = add t7, t8	{ t7, t8 }
	{ }	12 a = t9	{ t9 }
{ t10 }	13 t10 = a		{ }
	{ }	14 b = t10	{ t10 }

	out		in
Work list = {3}	{ }	1 t1 = x	{ }
	{ }	2 t2 = y	{ }
{ t2, t3, t1 }	3 t3 = z		{ }
{ t1, t4 }	4 t4 = sub t2, t3		{ t2, t3, t1 }
{ t5 }	5 t5 = mul t1, t4		{ t1, t4 }
{ }	6 a = t5		{ t5 }
{ t6 }	7 t6 = a		{ }
{ }	8 bne t6, 1, end		{ t6 }
{ t7 }	9 t7 = a		{ }
{ t7, t8 }	10 t8 = 1		{ t7 }
{ t9 }	11 t9 = add t7, t8		{ t7, t8 }
{ }	12 a = t9		{ t9 }
{ t10 }	13 t10 = a		{ }
{ }	14 b = t10		{ t10 }

	out		in
Work list = {2}	{ }	1 t1 = x	{ }
	{ t2, t1 }	2 t2 = y	{ }
	{ t2, t3, t1 }	3 t3 = z	{ t2, t1 }
	{ t1, t4 }	4 t4 = sub t2, t3	{ t2, t3, t1 }
	{ t5 }	5 t5 = mul t1, t4	{ t1, t4 }
	{ }	6 a = t5	{ t5 }
	{ t6 }	7 t6 = a	{ }
	{ }	8 bne t6, 1, end	{ t6 }
	{ t7 }	9 t7 = a	{ }
	{ t7, t8 }	10 t8 = 1	{ t7 }
	{ t9 }	11 t9 = add t7, t8	{ t7, t8 }
	{ }	12 a = t9	{ t9 }
	{ t10 }	13 t10 = a	{ }
	{ }	14 b = t10	{ t10 }

	out		in
Work list = {1}	{ t1 }	1 t1 = x	{ }
	{ t2, t1 }	2 t2 = y	{ t1 }
	{ t2, t3, t1 }	3 t3 = z	{ t2, t1 }
	{ t1, t4 }	4 t4 = sub t2, t3	{ t2, t3, t1 }
	{ t5 }	5 t5 = mul t1, t4	{ t1, t4 }
	{ }	6 a = t5	{ t5 }
	{ t6 }	7 t6 = a	{ }
	{ }	8 bne t6, 1, end	{ t6 }
	{ t7 }	9 t7 = a	{ }
	{ t7, t8 }	10 t8 = 1	{ t7 }
	{ t9 }	11 t9 = add t7, t8	{ t7, t8 }
	{ }	12 a = t9	{ t9 }
	{ t10 }	13 t10 = a	{ }
	{ }	14 b = t10	{ t10 }

	out		in
Work list = {}	{ t1 }	1 t1 = x	{ }
	{ t2, t1 }	2 t2 = y	{ t1 }
	{ t2, t3, t1 }	3 t3 = z	{ t2, t1 }
	{ t1, t4 }	4 t4 = sub t2, t3	{ t2, t3, t1 }
	{ t5 }	5 t5 = mul t1, t4	{ t1, t4 }
	{ }	6 a = t5	{ t5 }
	{ t6 }	7 t6 = a	{ }
	{ }	8 bne t6, 1, end	{ t6 }
	{ t7 }	9 t7 = a	{ }
	{ t7, t8 }	10 t8 = 1	{ t7 }
	{ t9 }	11 t9 = add t7, t8	{ t7, t8 }
	{ }	12 a = t9	{ t9 }
	{ t10 }	13 t10 = a	{ }
	{ }	14 b = t10	{ t10 }

out		in
{ t1 }	1 $t_1 = x$	{ }
{ t2, t1 }	2 $t_2 = y$	{ t1 }
{ t2, t3, t1 }	3 $t_3 = z$	{ t2, t1 }
{ t1, t4 }	4 $t_4 = \text{sub } t_2, t_3$	{ t2, t3, t1 }
{ t5 }	5 $t_5 = \text{mul } t_1, t_4$	{ t1, t4 }
{ }	6 $a = t_5$	{ t5 }
{ t6 }	7 $t_6 = a$	{ }
{ }	8 $\text{bne } t_6, 1, \text{end}$	{ t6 }
{ t7 }	9 $t_7 = a$	{ }
{ t7, t8 }	10 $t_8 = 1$	{ t7 }
{ t9 }	11 $t_9 = \text{add } t_7, t_8$	{ t7, t8 }
{ }	12 $a = t_9$	{ t9 }
{ t10 }	13 $t_{10} = a$	{ }
{ }	14 $b = t_{10}$	{ t10 }

$\text{in}_i =$
**live variables at
label i**

Register Allocation

For each function:

1. Construct the CFG (from the IR)
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Interference Graph

- Use the liveness analysis output to construct the **interference graph**
- Create a node for each IR variable (t_1, t_2, \dots)
- If t_1 and t_2 appear in the same liveness sets (in_i for some i):
 - Create an edge between t_1 and t_2

Interference Graph

{ t1 }

{ t2, t1 }

{ t2, t3, t1 }

{ t1, t4 }

{ t5 }

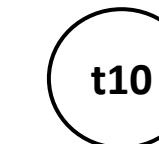
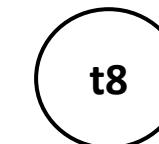
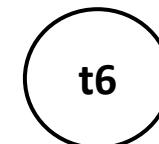
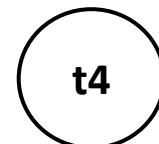
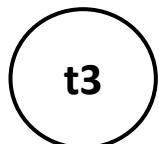
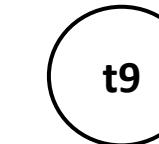
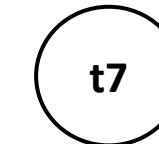
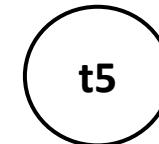
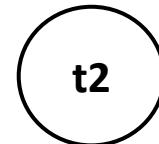
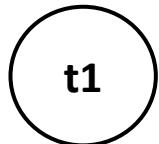
{ t6 }

{ t7 }

{ t7, t8 }

{ t9 }

{ t10 }



Interference Graph

{ t1 }

{ t2, t1 }

{ t2, t3, t1 }

{ t1, t4 }

{ t5 }

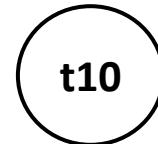
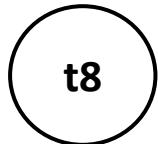
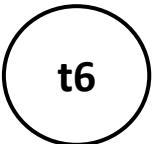
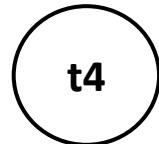
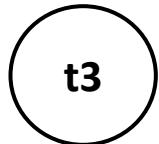
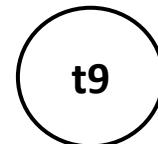
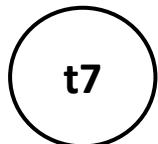
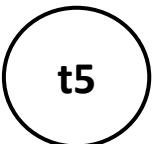
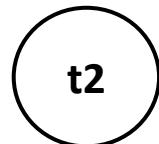
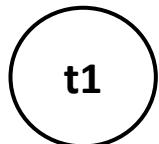
{ t6 }

{ t7 }

{ t7, t8 }

{ t9 }

{ t10 }



Interference Graph

{ t1 }

{ t2, t1 }

{ t2, t3, t1 }

{ t1, t4 }

{ t5 }

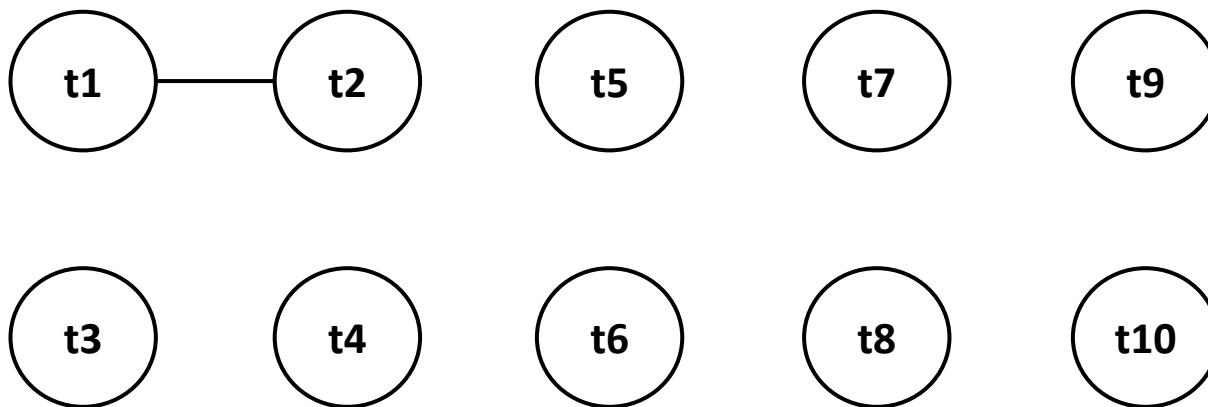
{ t6 }

{ t7 }

{ t7, t8 }

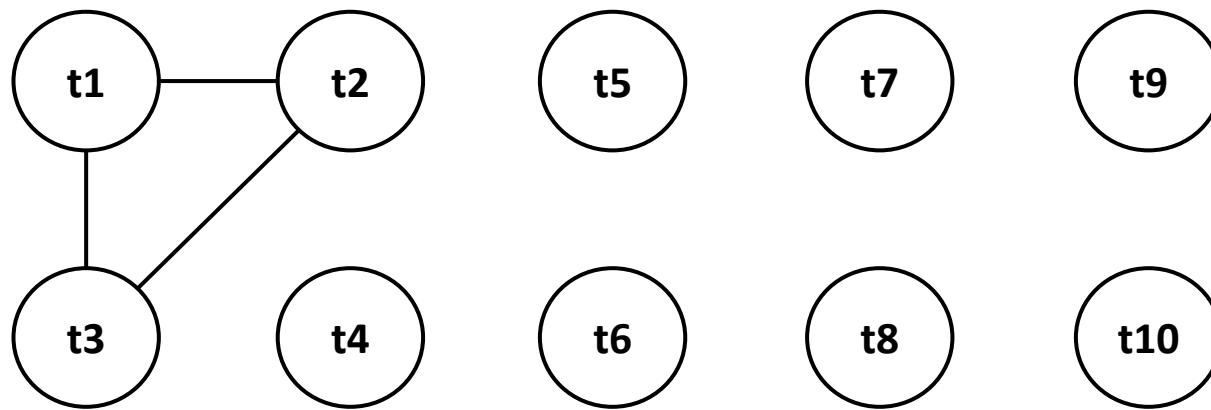
{ t9 }

{ t10 }



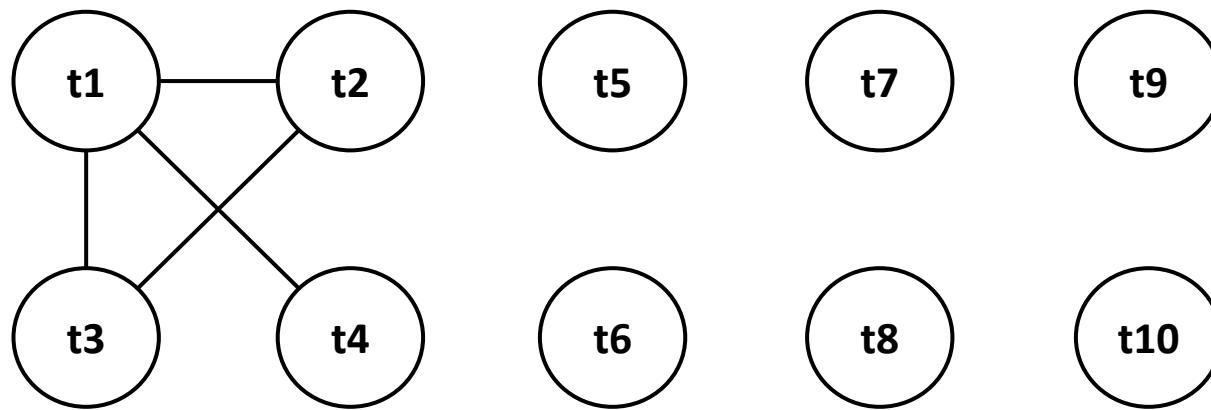
Interference Graph

{ t1 }
{ t2, t1 }
{ t2, t3, t1 }
{ t1, t4 }
{ t5 }
{ t6 }
{ t7 }
{ t7, t8 }
{ t9 }
{ t10 }



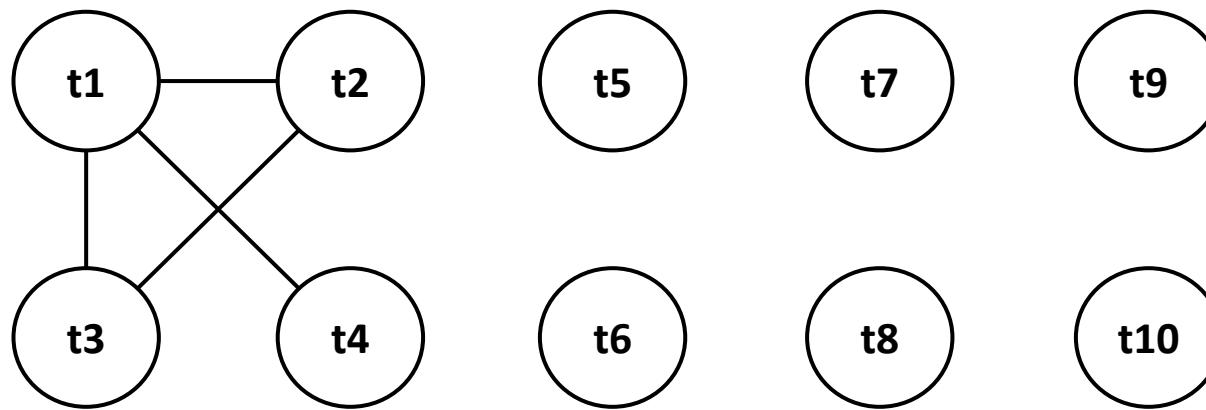
Interference Graph

{ t1 }
{ t2, t1 }
{ t2, t3, t1 }
{ t1, t4 }
{ t5 }
{ t6 }
{ t7 }
{ t7, t8 }
{ t9 }
{ t10 }



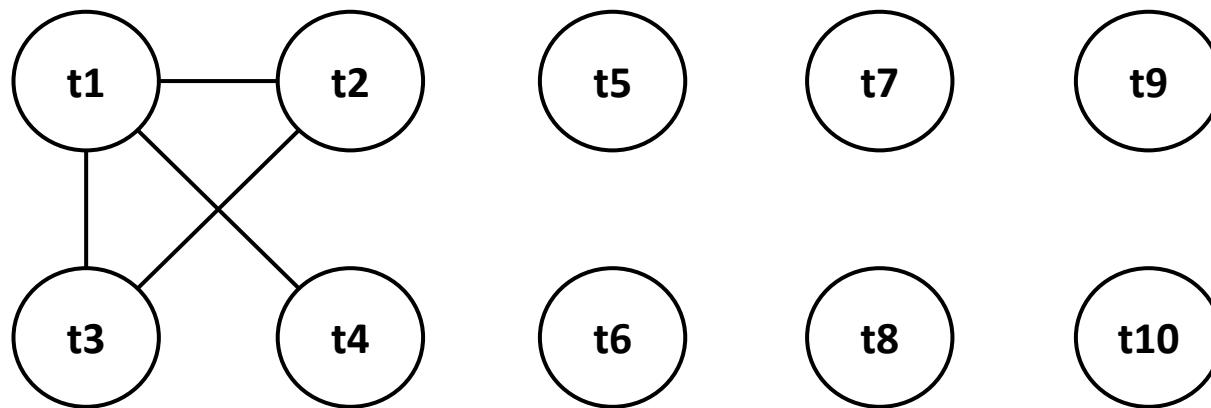
Interference Graph

{ t1 }
{ t2, t1 }
{ t2, t3, t1 }
{ t1, t4 }
{ t5 }
{ t6 }
{ t7 }
{ t7, t8 }
{ t9 }
{ t10 }



Interference Graph

{ t1 }
{ t2, t1 }
{ t2, t3, t1 }
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Interference Graph

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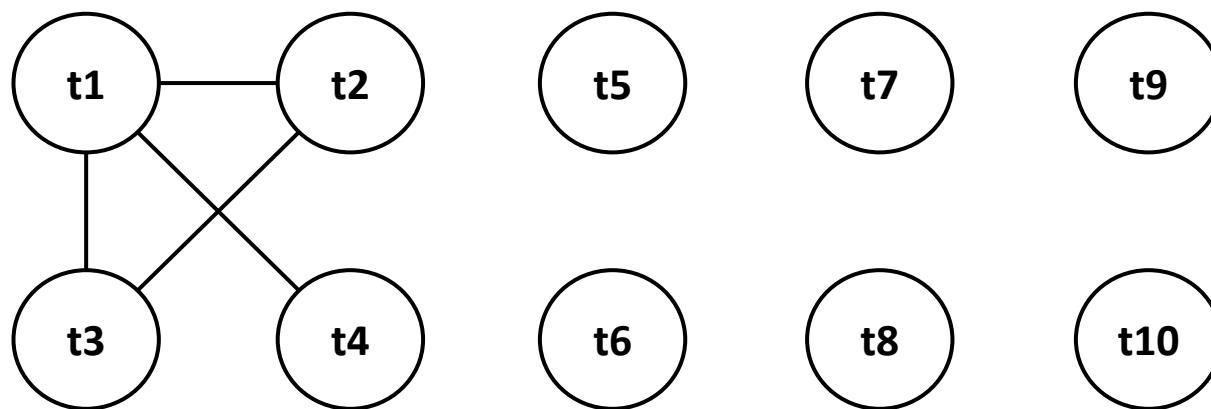
{ t6 }

{ t7 }

{ t7, t8 }

{ t9 }

{ t10 }



Interference Graph

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{ t2, t3, t1 }

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{ t5 }

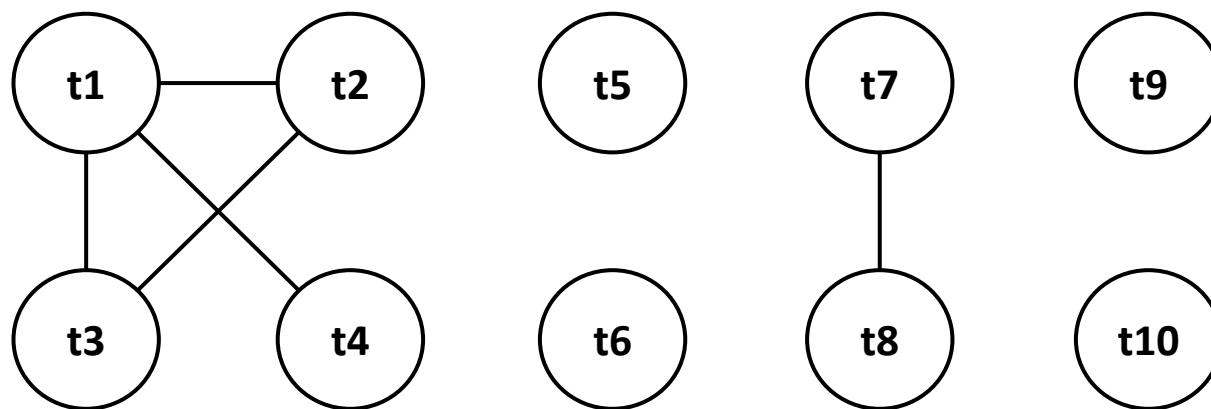
{ t6 }

{ t7 }

{ t7, t8 }

{ t9 }

{ t10 }



Interference Graph

{ t1 }

{ t2, t1 }

{ t2, t3, t1 }

{ t1, t4 }

{ t5 }

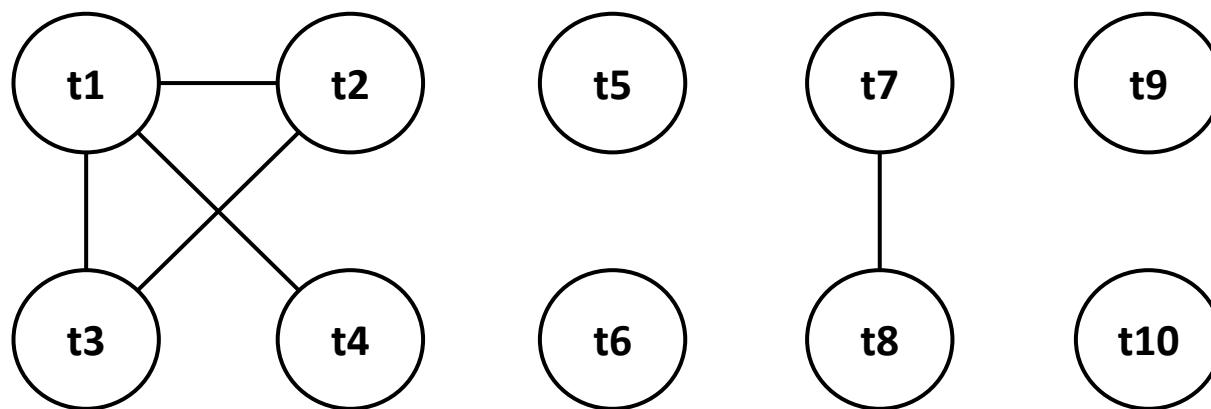
{ t6 }

{ t7 }

{ t7, t8 }

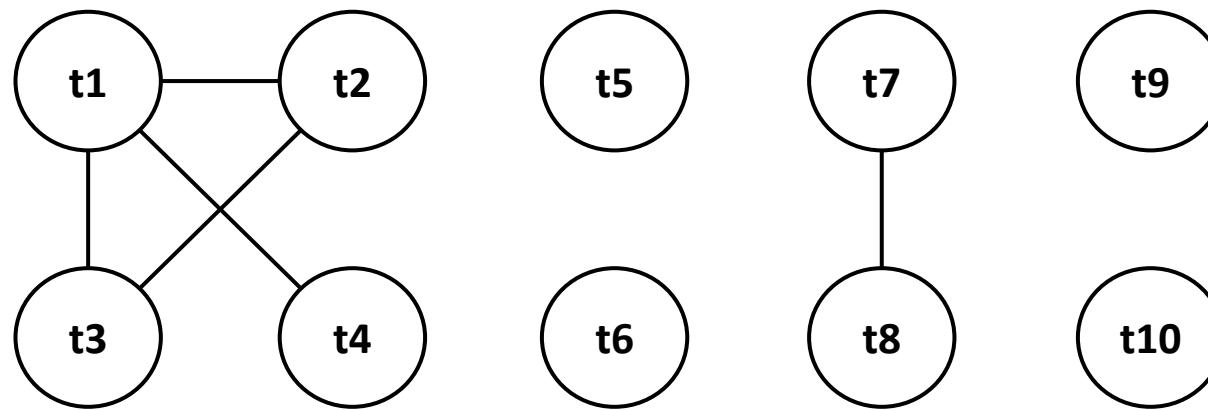
{ t9 }

{ t10 }



Interference Graph

{ t1 }
{ t2, t1 }
{ t2, t3, t1 }
{ t1, t4 }
{ t5 }
{ t6 }
{ t7 }
{ t7, t8 }
{ t9 }
{ t10 }



Register Allocation

1. For each function:
 1. Construct the CFG (from the IR)
 2. Run liveness analysis
 3. Construct the interference graph
 - 4. Compute a *k-coloring* of the graph**
 5. Use the coloring to build the required mapping

Graph Coloring

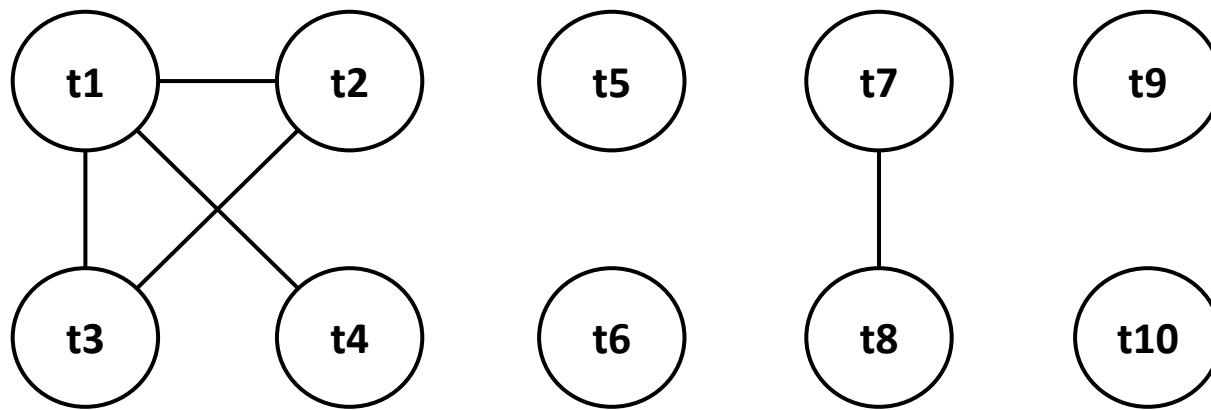
- Graph coloring is **hard** (NP-complete problem)
- We need a **heuristic**...

Graph Coloring

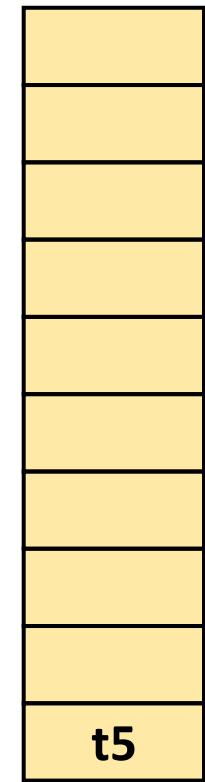
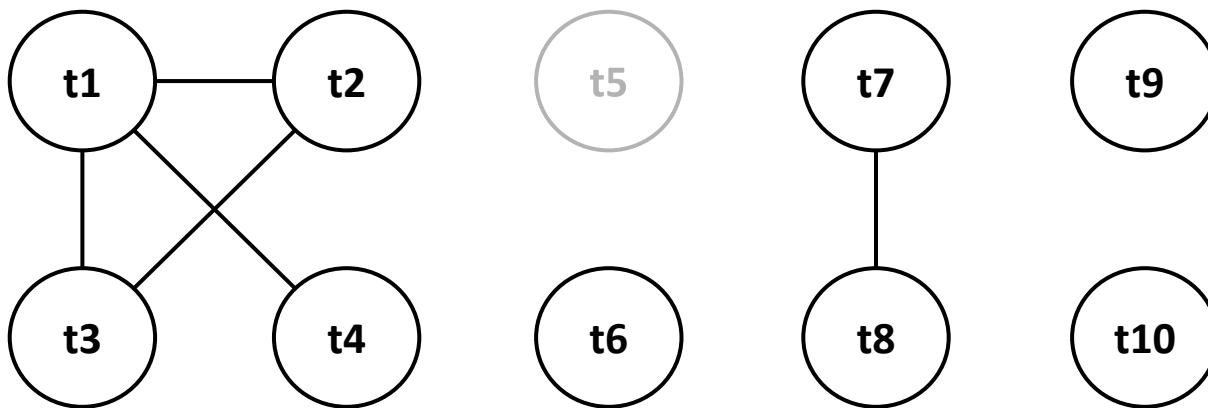
Chaitin's Algorithm for k -coloring (simplified):

- While there is a node with less than k neighbors:
 - Remove it and its edges from the graph
 - Push it on the stack
- If the entire graph was removed, then it is k -colorable
- While the stack is not empty:
 - Pop a node from the stack and assign it an available color

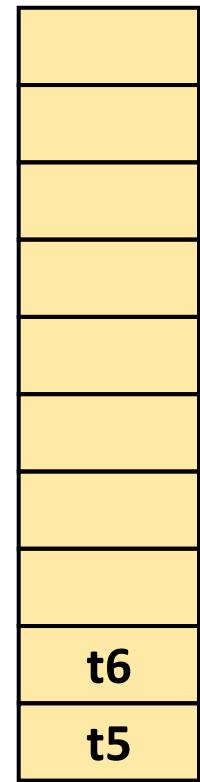
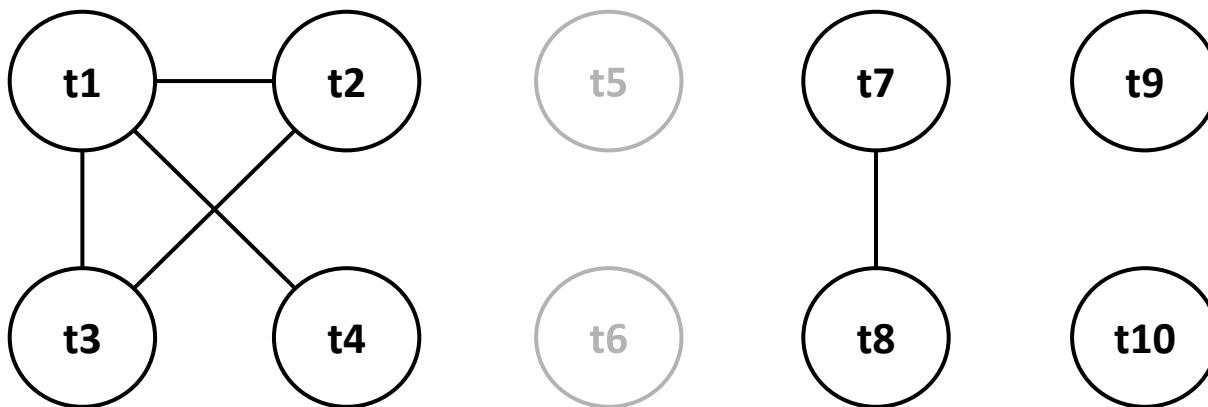
Graph Coloring ($k = 3$)



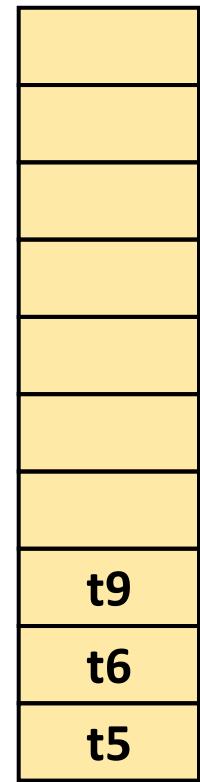
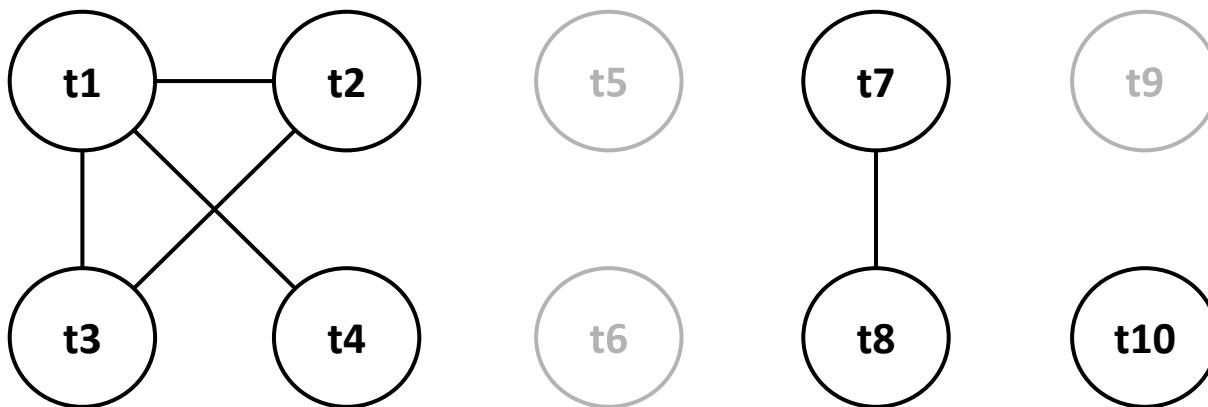
Graph Coloring ($k = 3$)



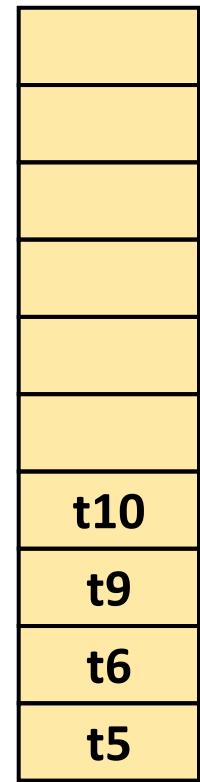
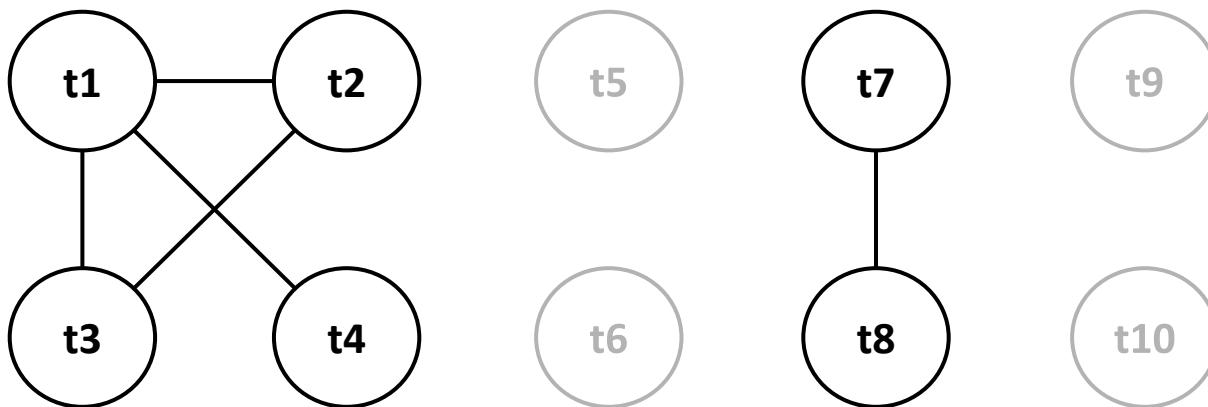
Graph Coloring ($k = 3$)



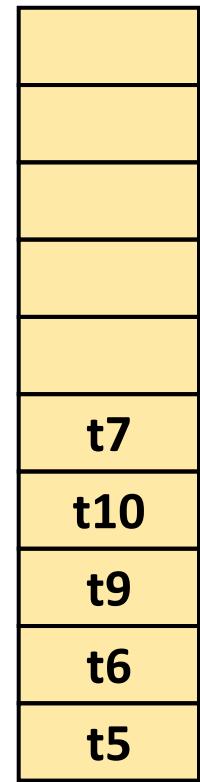
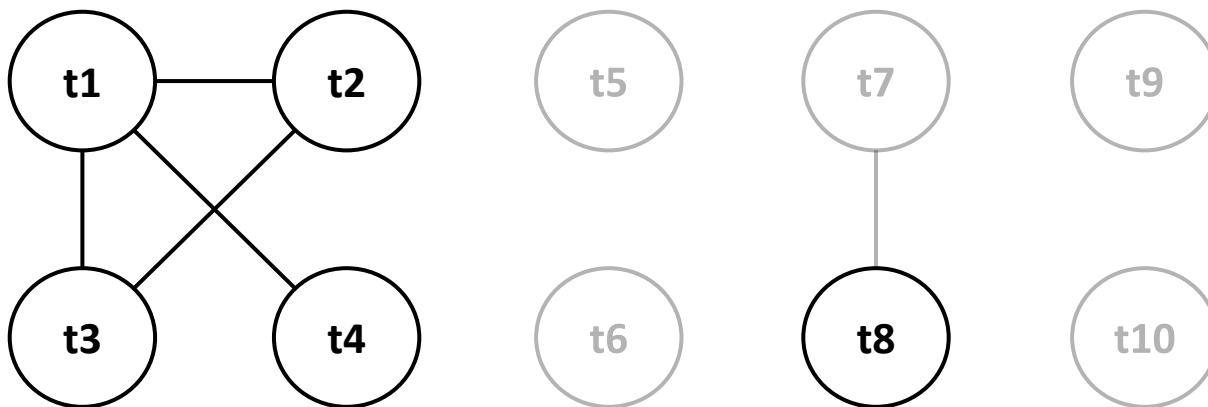
Graph Coloring ($k = 3$)



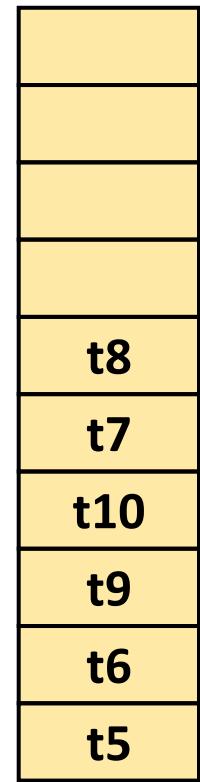
Graph Coloring ($k = 3$)



Graph Coloring ($k = 3$)



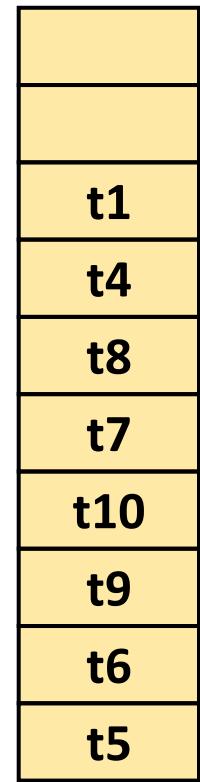
Graph Coloring ($k = 3$)



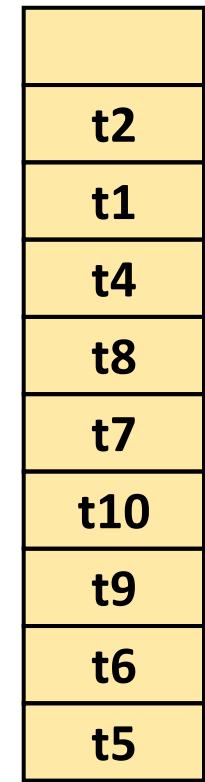
Graph Coloring ($k = 3$)



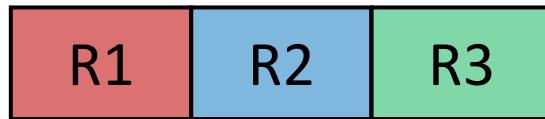
Graph Coloring ($k = 3$)



Graph Coloring ($k = 3$)



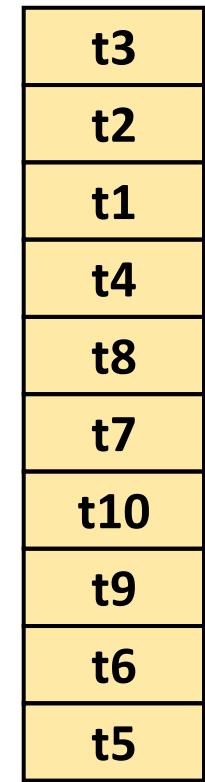
Graph Coloring ($k = 3$)



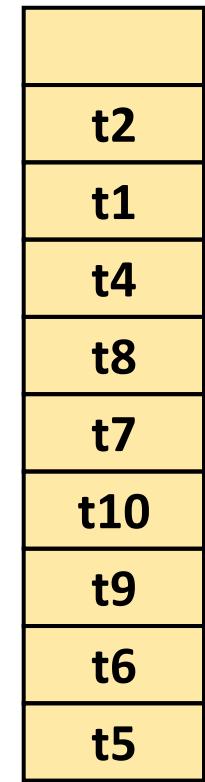
Graph Coloring ($k = 3$)

Pop nodes and assign colors...

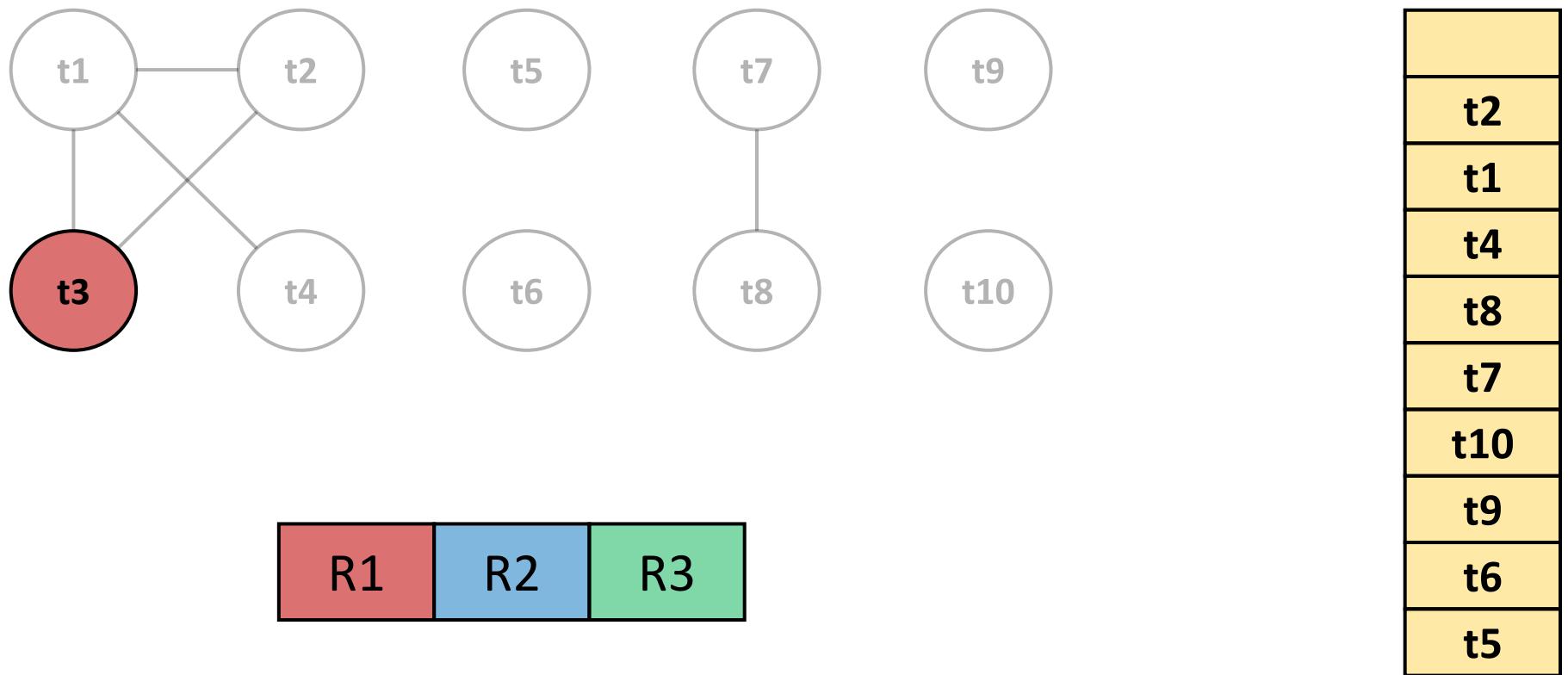
Graph Coloring ($k = 3$)



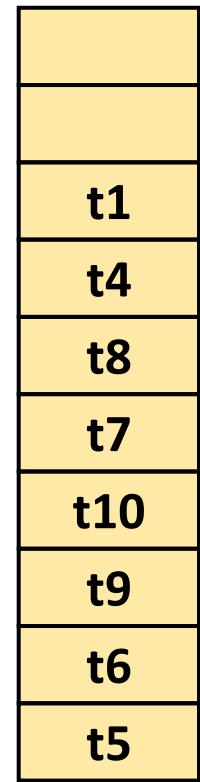
Graph Coloring ($k = 3$)



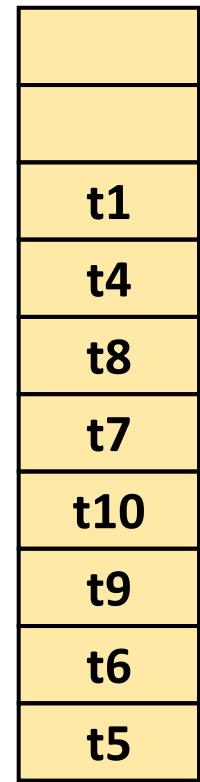
Graph Coloring ($k = 3$)



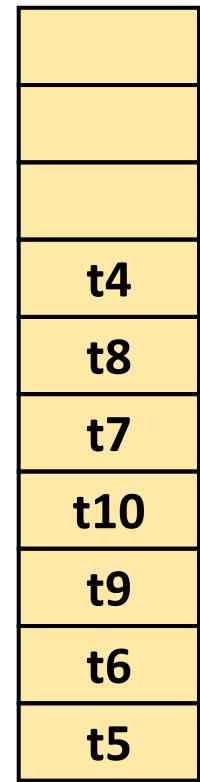
Graph Coloring ($k = 3$)



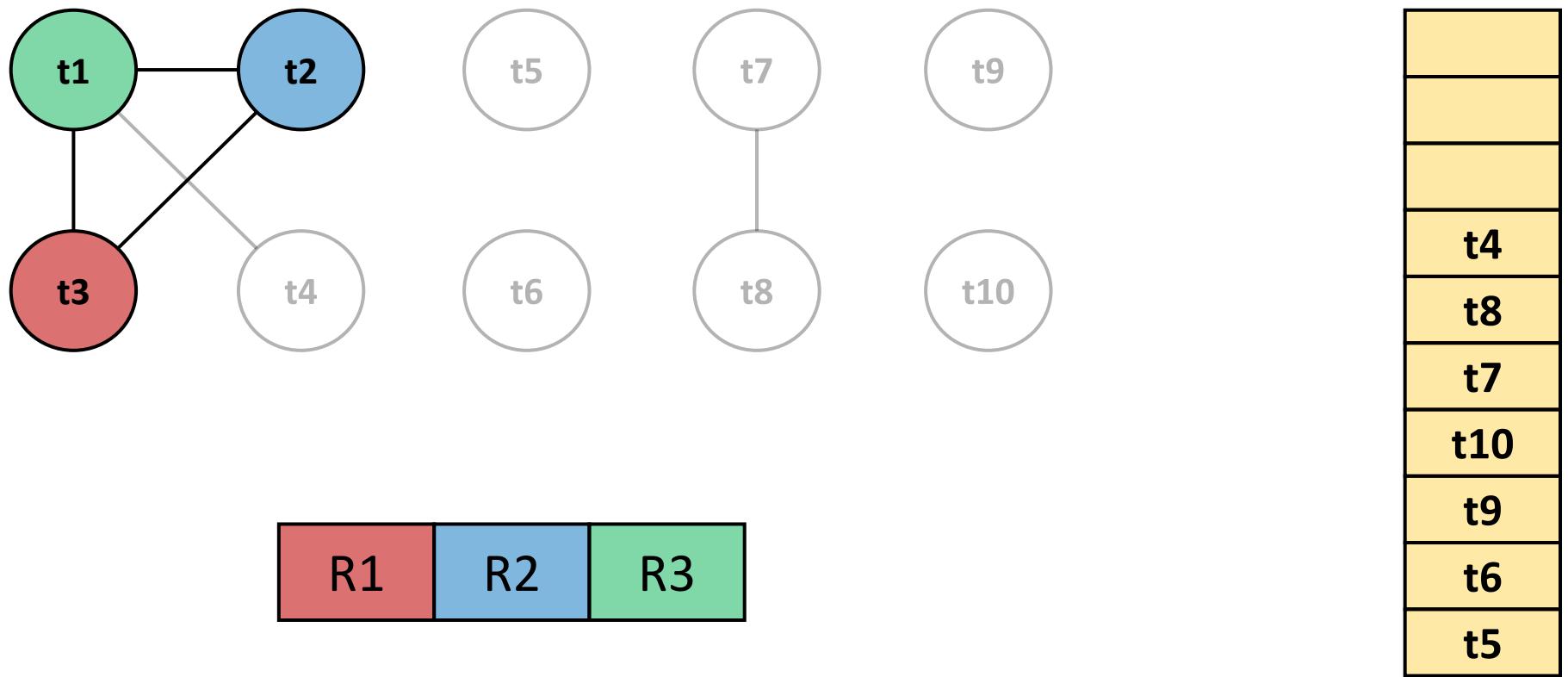
Graph Coloring ($k = 3$)



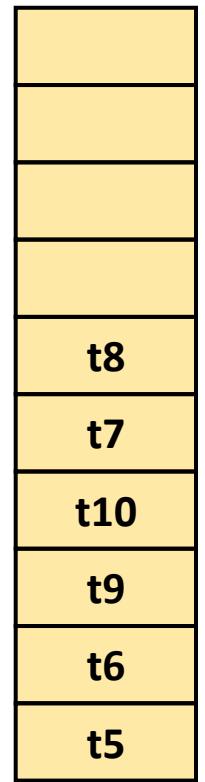
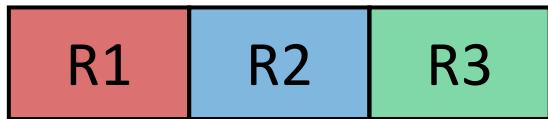
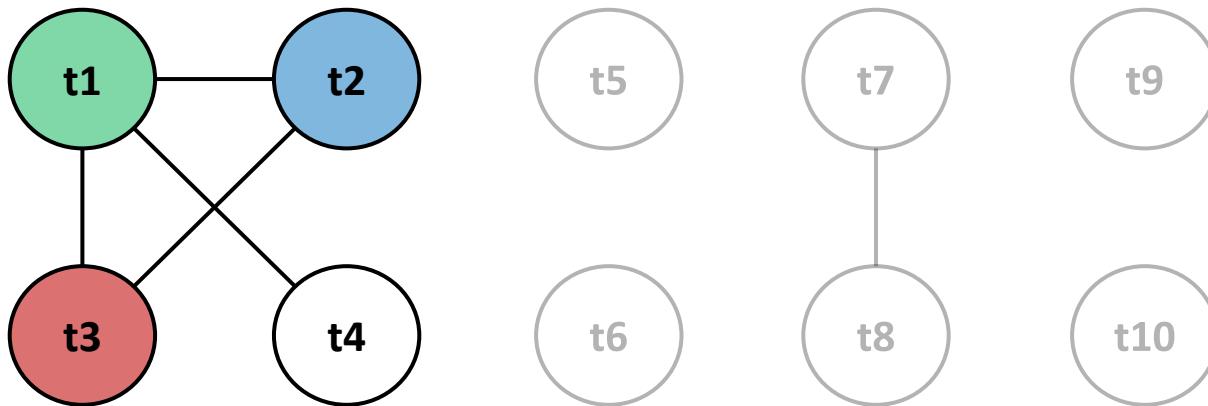
Graph Coloring ($k = 3$)



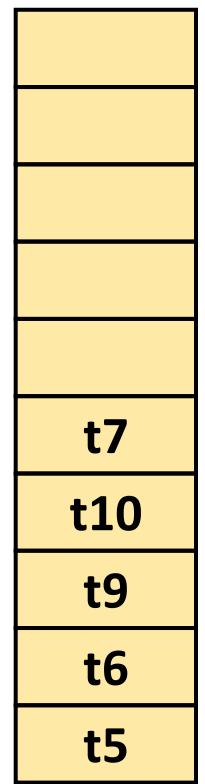
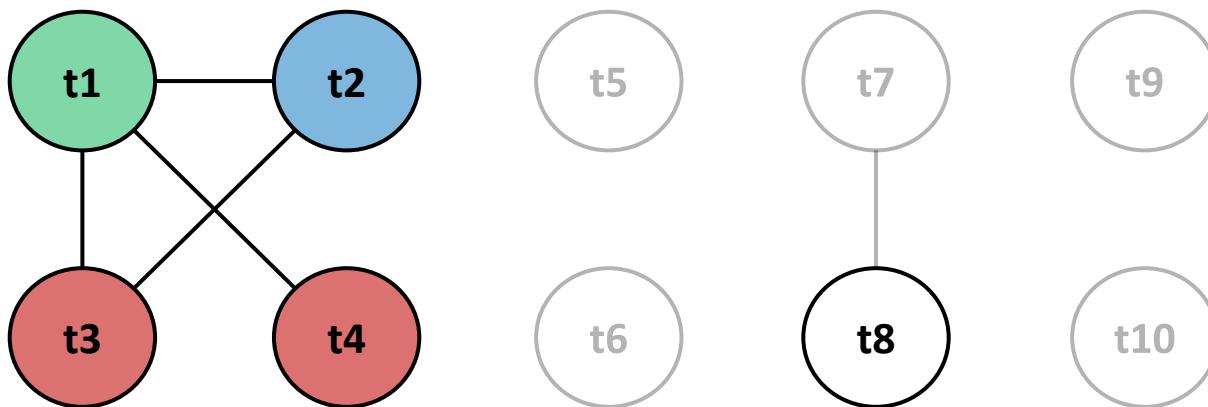
Graph Coloring ($k = 3$)



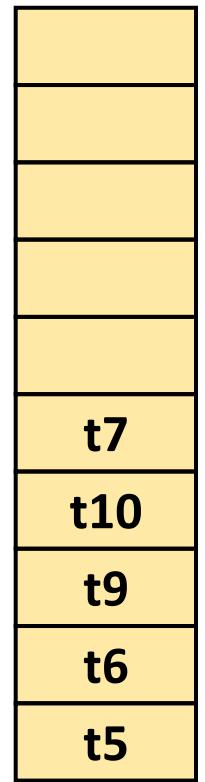
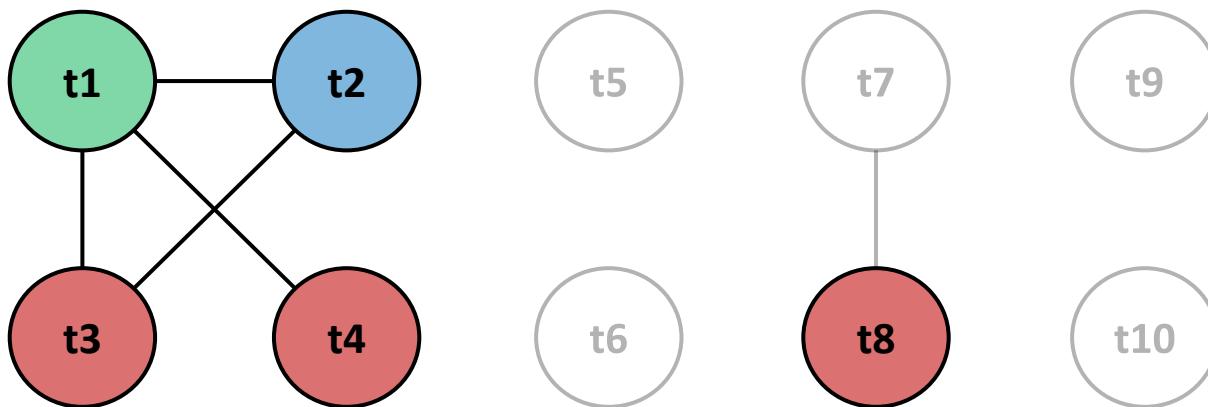
Graph Coloring ($k = 3$)



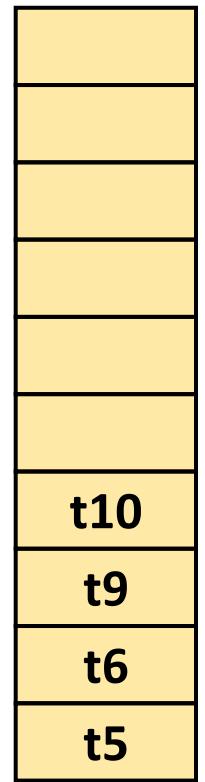
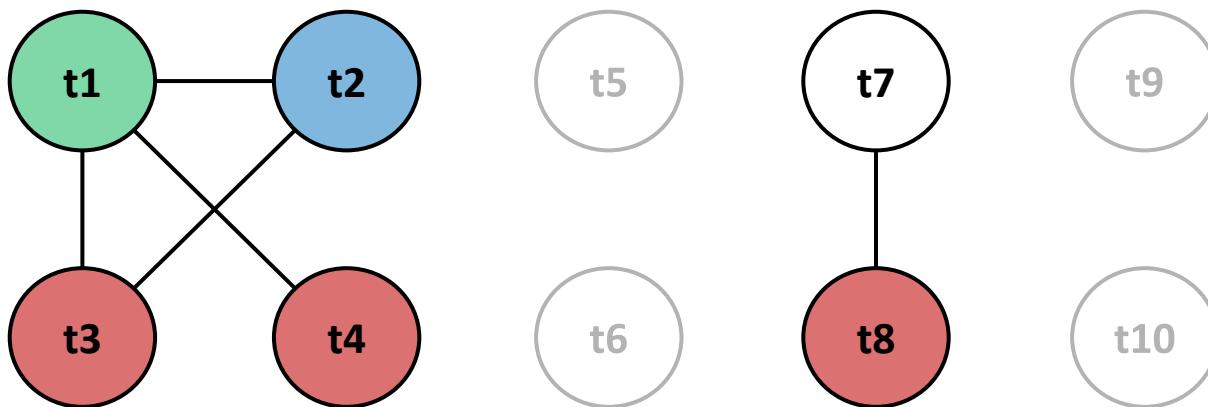
Graph Coloring ($k = 3$)



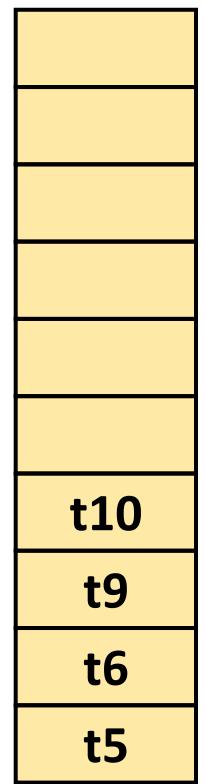
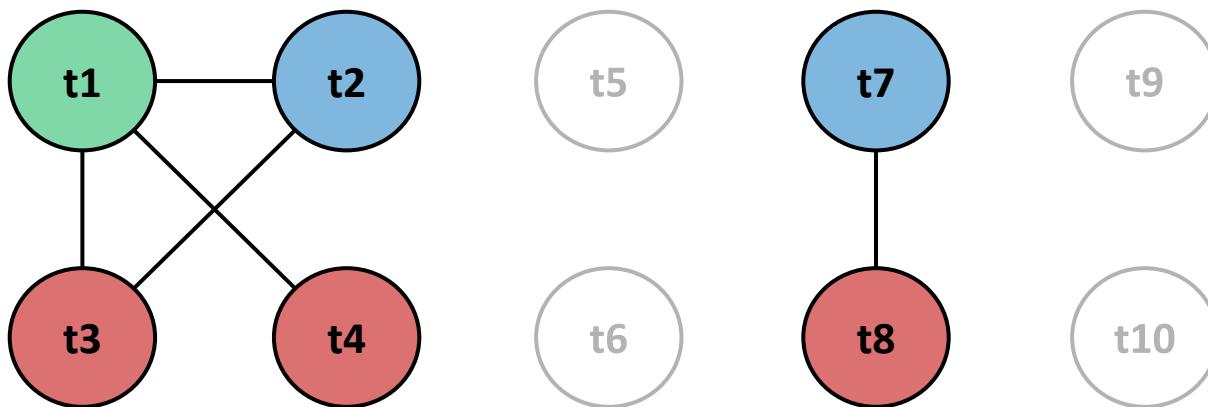
Graph Coloring ($k = 3$)



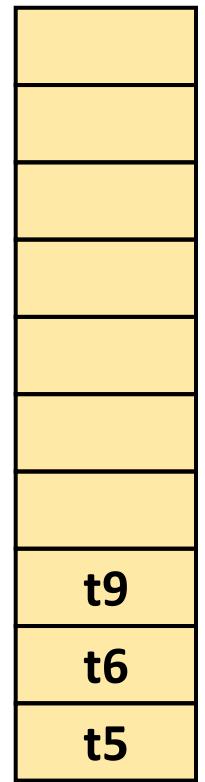
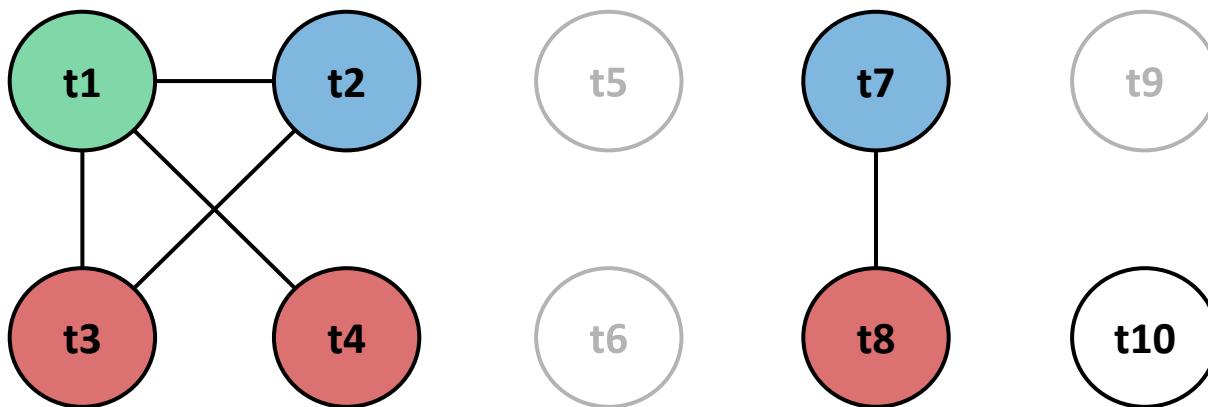
Graph Coloring ($k = 3$)



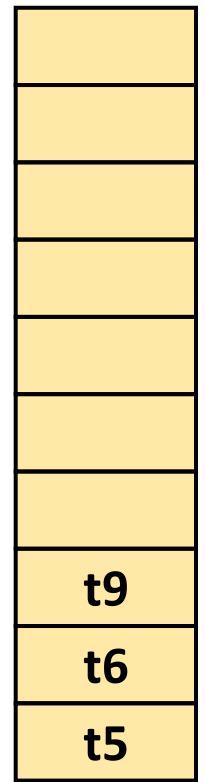
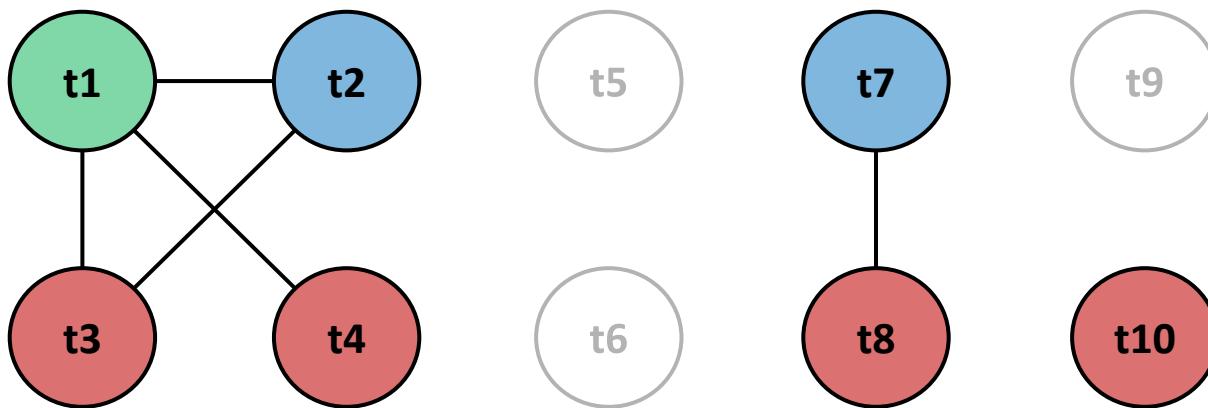
Graph Coloring ($k = 3$)



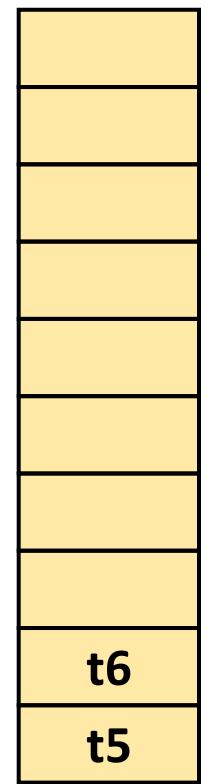
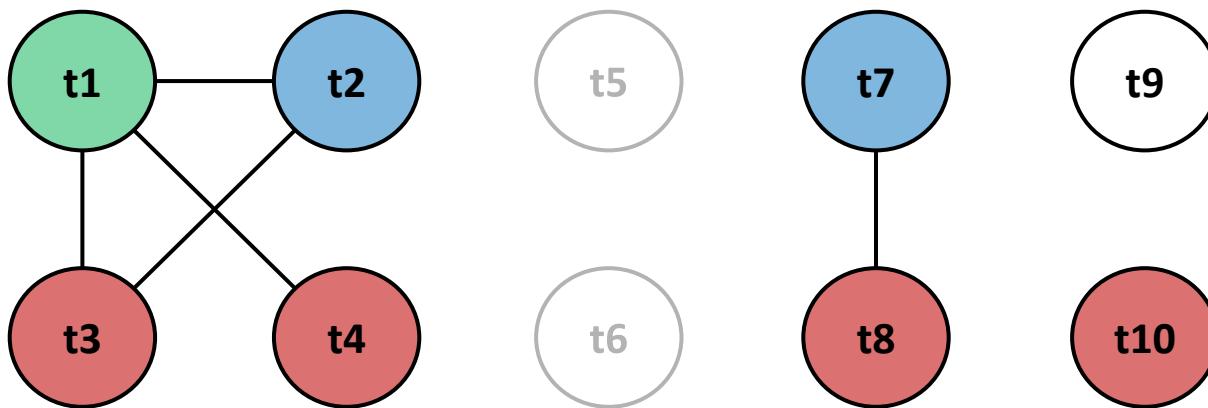
Graph Coloring ($k = 3$)



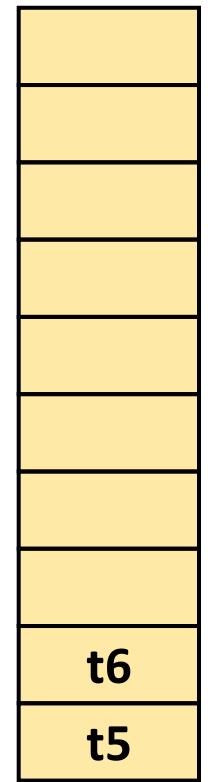
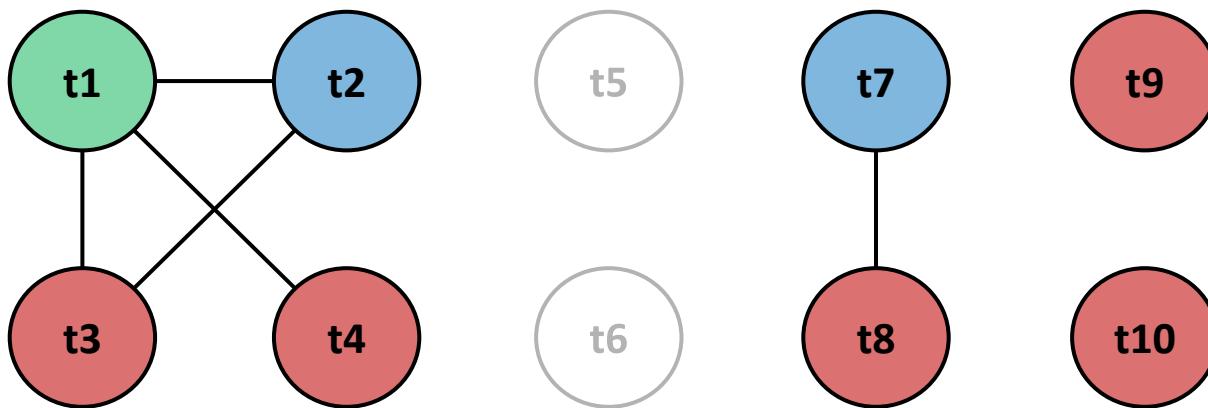
Graph Coloring ($k = 3$)



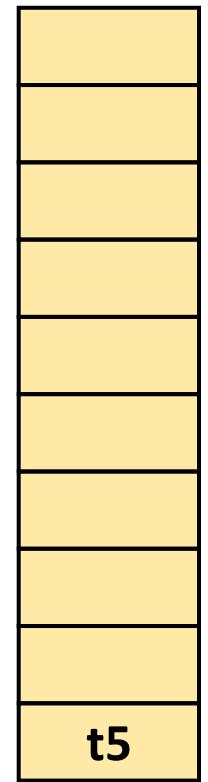
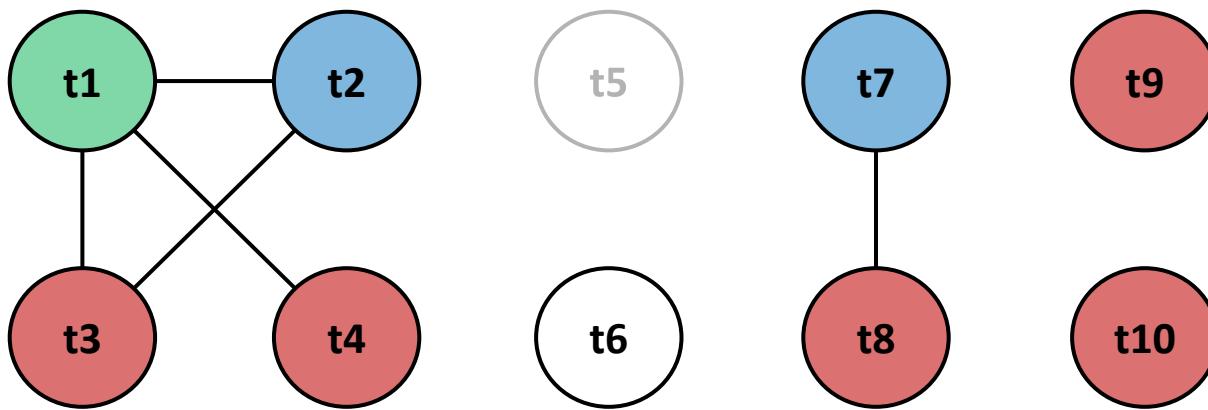
Graph Coloring ($k = 3$)



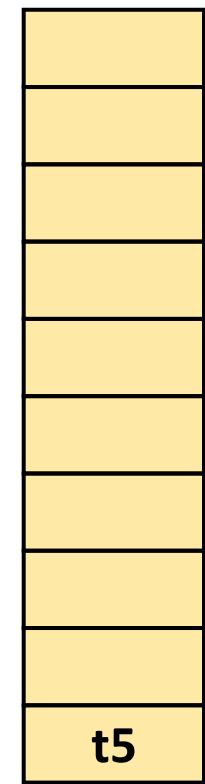
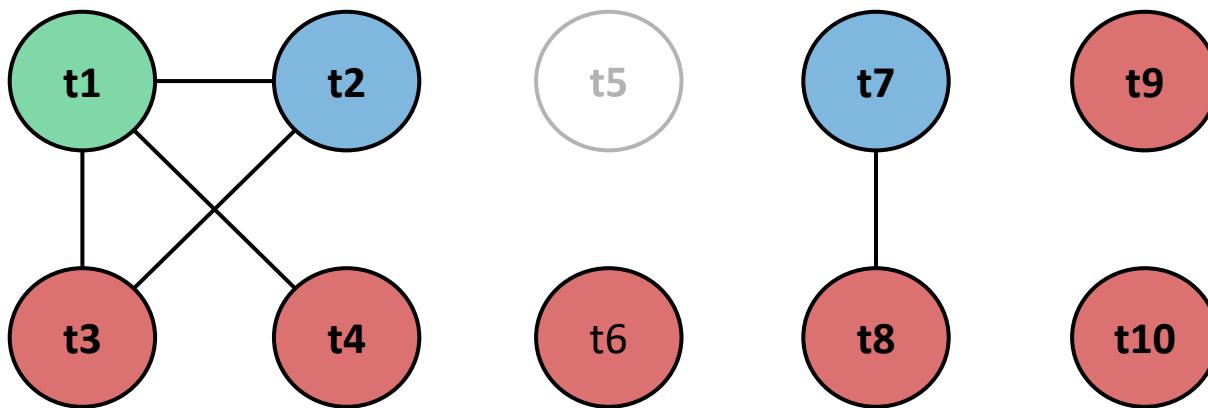
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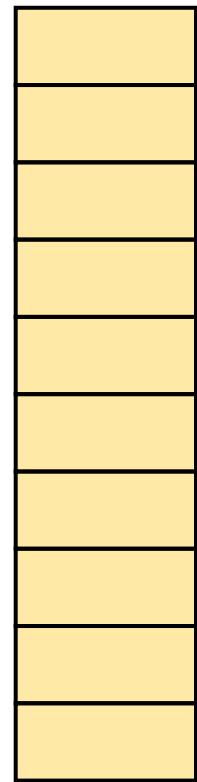
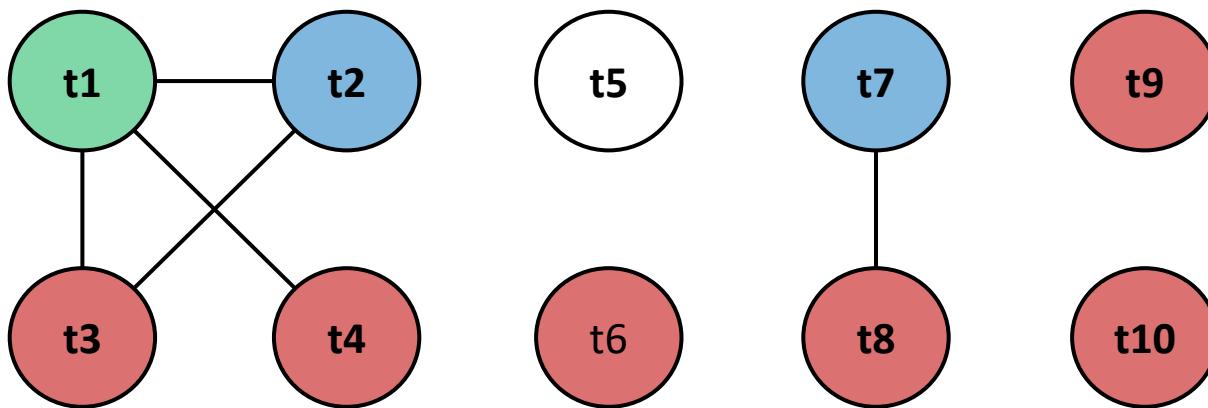
Graph Coloring ($k = 3$)



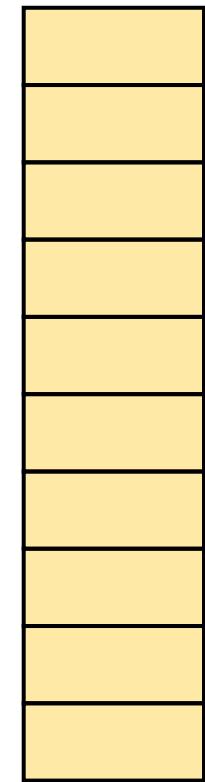
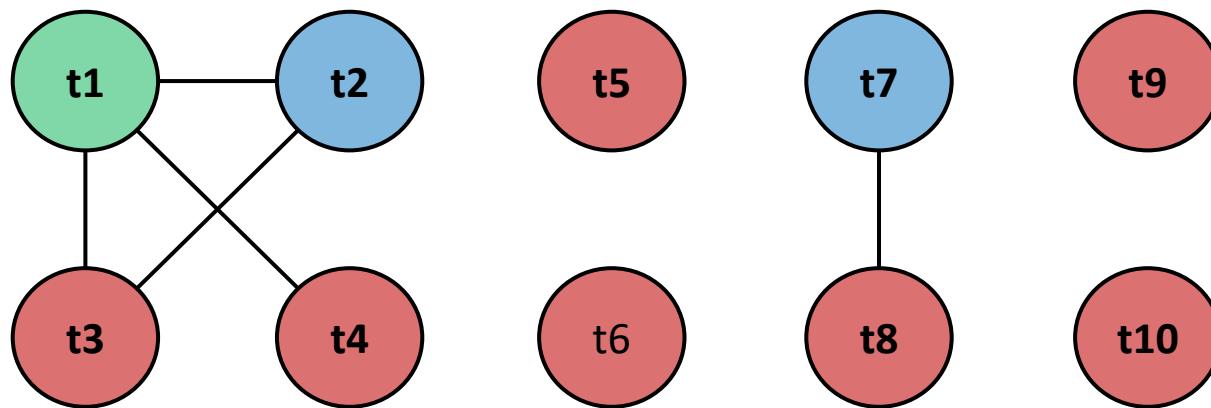
Graph Coloring ($k = 3$)



Graph Coloring ($k = 3$)



Graph Coloring ($k = 3$)



Register Allocation

For each function:

1. Construct the CFG (from the IR)
2. Run liveness analysis
3. Construct the interference graph
4. Compute a *k-coloring* of the graph
5. **Use the coloring to build the required mapping**

Register Allocation

- According to the coloring, our register allocation is:

IR Register	Color	MIPS Register
t1	R3	\$t2
t2	R2	\$t1
t3	R1	\$t0
t4	R1	\$t0
t5	R1	\$t0
t6	R1	\$t0
t7	R2	\$t1
t8	R1	\$t0
t9	R1	\$t0
t10	R1	\$t0

Register Allocation

```
t1 = x
t2 = y
t3 = z
t4 = sub t2, t3
t5 = mul t1, t4
a = t5
t6 = a
bne t6, 1, end
t7 = a
t8 = 1
t9 = add t7, t8
a = t9
end:
t10 = a
b = t10
```

IR Register	MIPS Register
t1	\$t2
t2	\$t1
t3	\$t0
t4	\$t0
t5	\$t0
t6	\$t0
t7	\$t1
t8	\$t0
t9	\$t0
t10	\$t0

```
lw $t1, 8($fp)
lw $t2, 12($fp)
lw $t3, 16($fp)
sub $t4, $t2, $t3
mul $t5, $t1, $t4
sw $t5, -44($fp)
lw $t6, -44($fp)
bne $t6, 1, end
lw $t7, -44($fp)
li $t8, 1
add $t9, $t7, $t8
sw $t9, -44($fp)
end:
lw $t10, -44($fp)
sw $t10, -48($fp)
```

Register Allocation

```
t1 = x
t2 = y
t3 = z
t4 = sub t2, t3
t5 = mul t1, t4
a = t5
t6 = a
bne t6, 1, end
t7 = a
t8 = 1
t9 = add t7, t8
a = t9
end:
t10 = a
b = t10
```

IR Register	MIPS Register
t1	\$t2
t2	\$t1
t3	\$t0
t4	\$t0
t5	\$t0
t6	\$t0
t7	\$t1
t8	\$t0
t9	\$t0
t10	\$t0

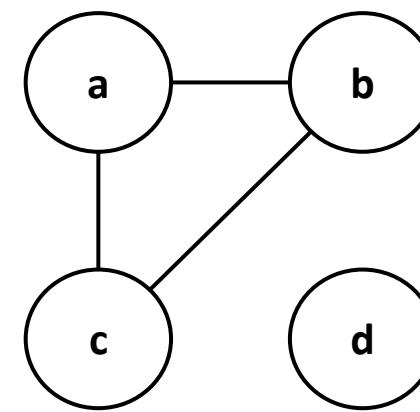
```
lw $t2, 8($fp)
lw $t1, 12($fp)
lw $t0, 16($fp)
sub $t0, $t1, $t0
mul $t0, $t2, $t0
sw $t0, -44($fp)
lw $t0, -44($fp)
bne $t0, 1, end
lw $t1, -44($fp)
li $t0, 1
add $t0, $t1, $t0
sw $t0, -44($fp)
end:
lw $t0, -44($fp)
sw $t0, -48($fp)
```

General Algorithm

- Spilling
 - We choose a “victim” variable to be stored in memory
 - Remove nodes from the graph
- Coalescing
 - Get rid of redundant MOV instructions ($x = y$)
 - Merge nodes in the graph

Example: Interference Graph

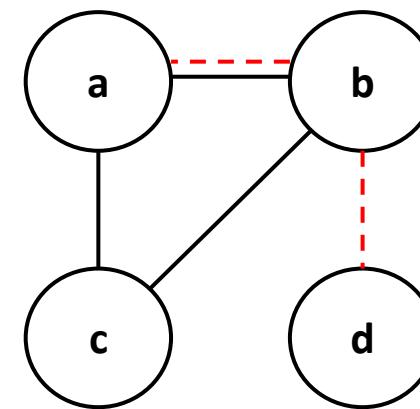
```
a = b  
b = a + b  
d = b  
c = c + a
```



Example: MOV Edges

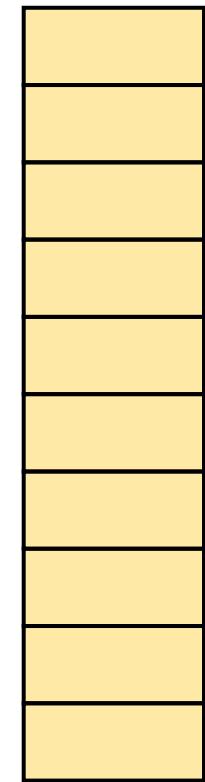
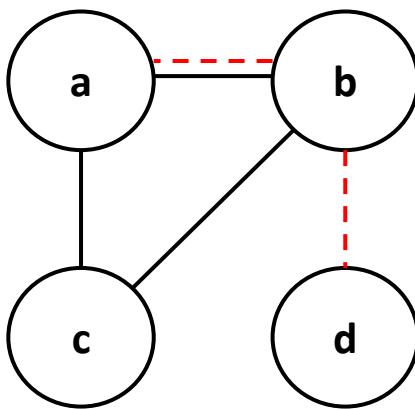
- Add dashed edges to MOV instructions

```
a = b  
b = a + b  
d = b  
c = c + a
```



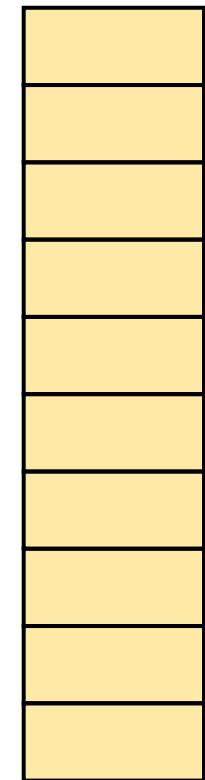
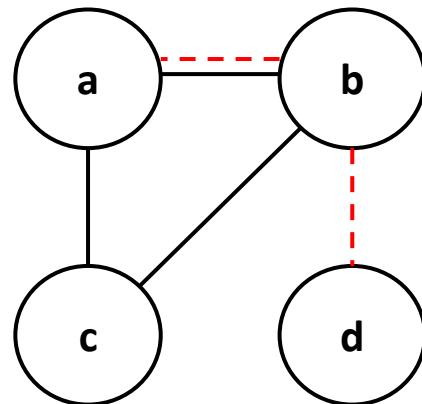
Example

- We want to compute $k = 2$ coloring



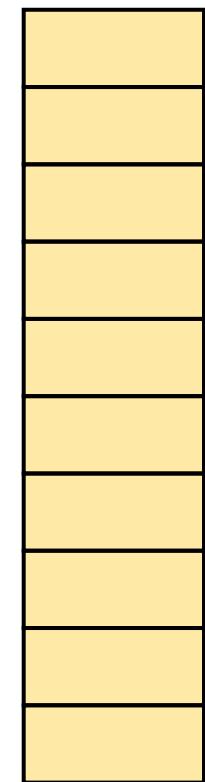
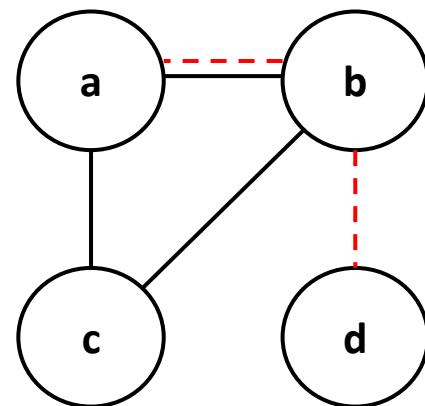
Example: Simplify

- We want to compute $k = 2$ coloring
- **No** node to simplify
 - All nodes have degree ≥ 2 or connected with a MOV edge



Example: Coalesce

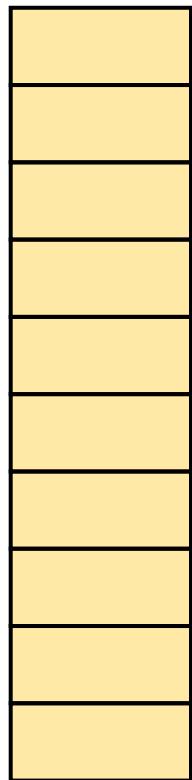
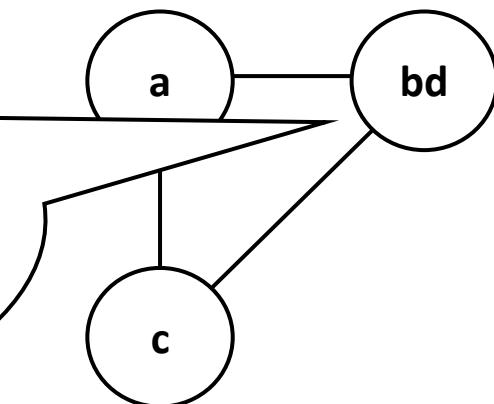
- We want to compute $k = 2$ coloring
- Cannot coalesce a & b – **constrained MOV**
- Can we coalesce b & d?



Example: Coalesce

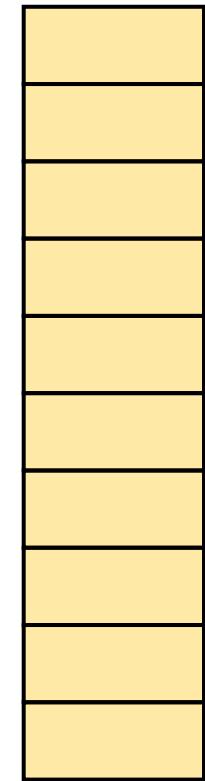
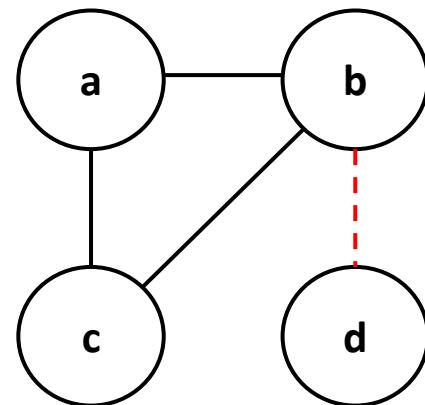
- We want to compute $k = 2$ coloring
- Cannot coalesce a & b – **constrained MOV**
- Can we coalesce b & d?

Briggs heuristic:
coalesce nodes x and y if
the resulting node xy has
less than k neighbors
with degree $\geq k$



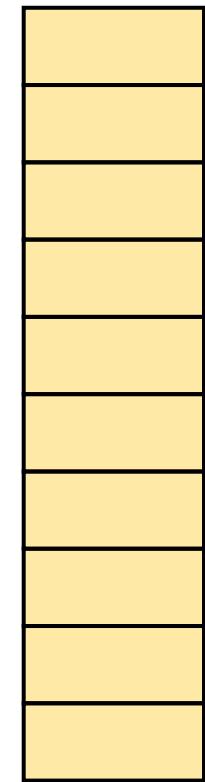
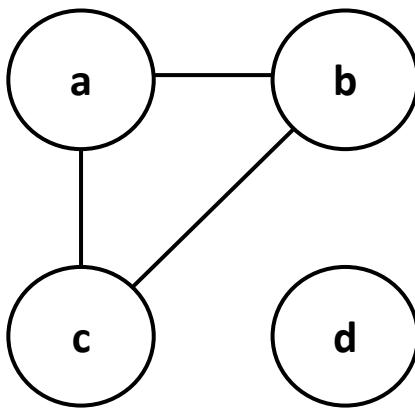
Example: Conservative Coalescing

- We want to compute $k = 2$ coloring
- Following Briggs heuristic, we don't Coalesce b & d



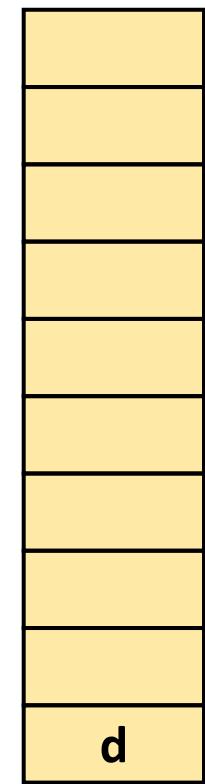
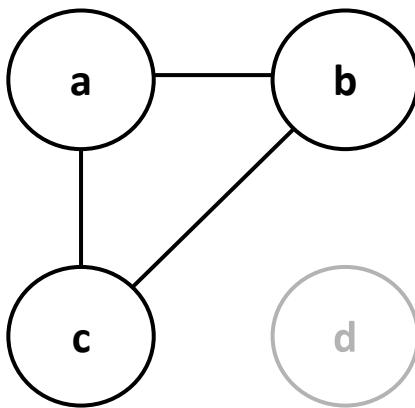
Example: Freeze

- We want to compute $k = 2$ coloring
- Remove MOV edge between b & d



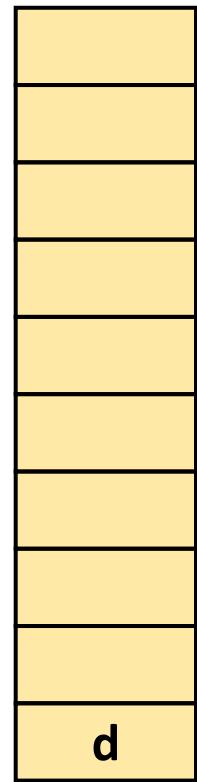
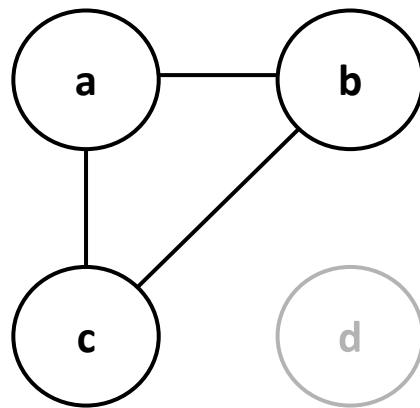
Example: Simplify

- We want to compute $k = 2$ coloring
- Simplify
- No more nodes to simplify...



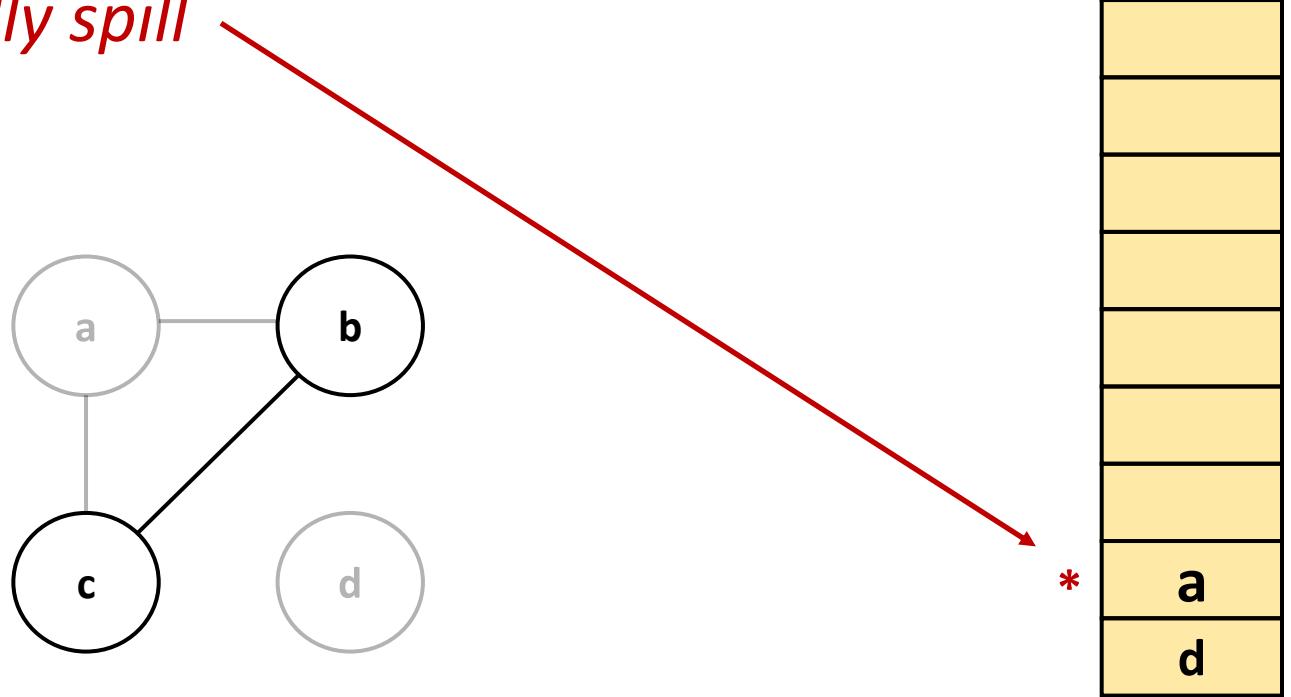
Example: Spilling

- We want to compute $k = 2$ coloring
- Choose a node to *potentially spill*



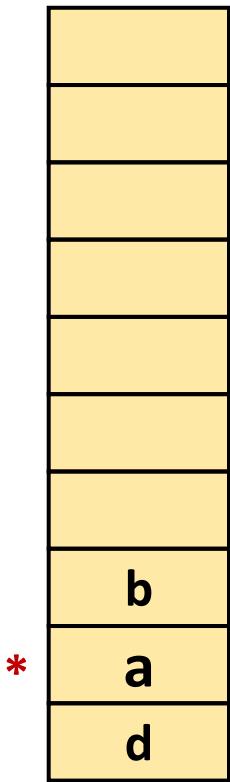
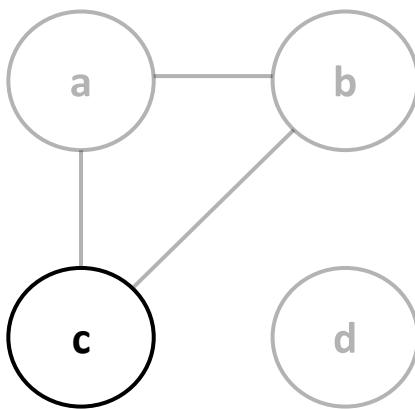
Example: Spilling

- We want to compute $k = 2$ coloring
- Choose a node to *potentially spill*



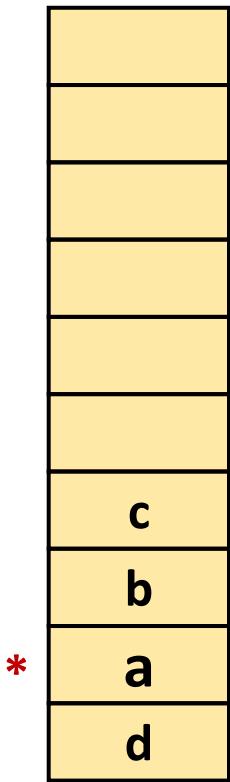
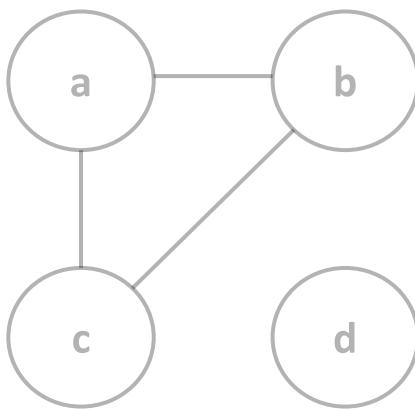
Example: Simplify

- We want to compute $k = 2$ coloring



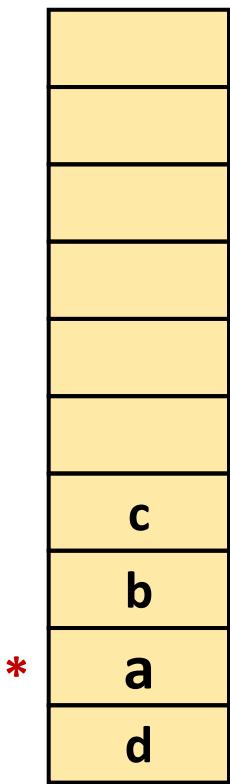
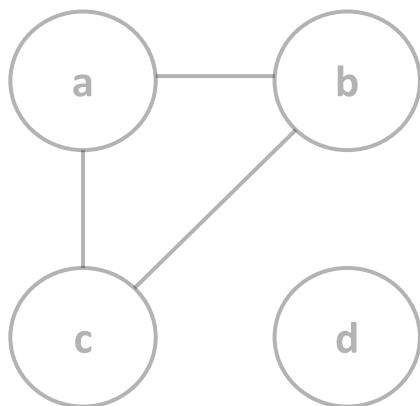
Example: Simplify

- We want to compute $k = 2$ coloring



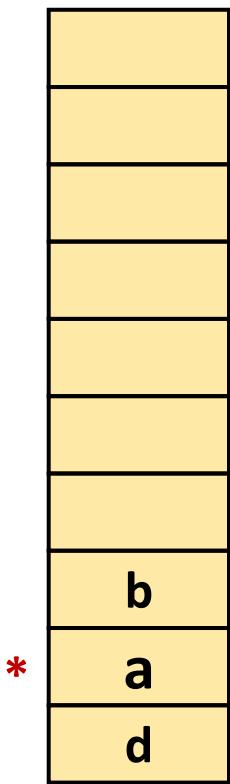
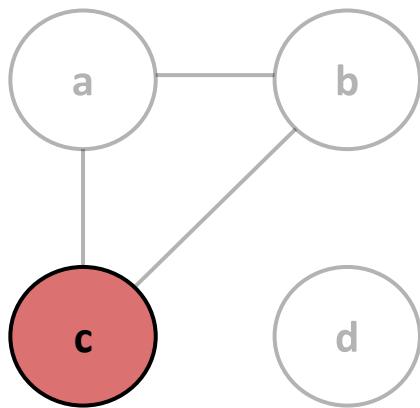
Example: Select

- We want to compute $k = 2$ coloring
- Pop and (try to) assign colors



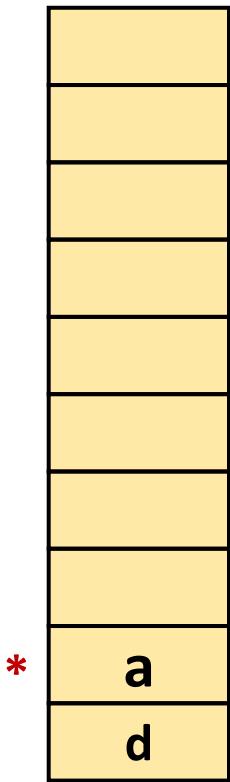
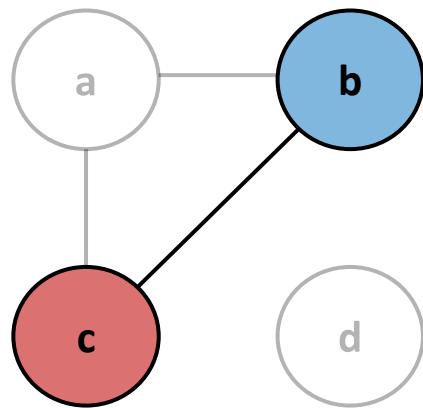
Example: Select

- We want to compute $k = 2$ coloring
- Pop and (try to) assign colors



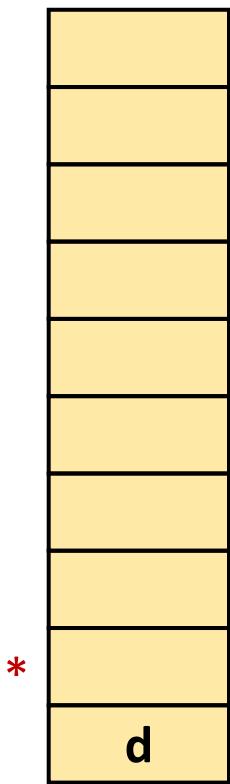
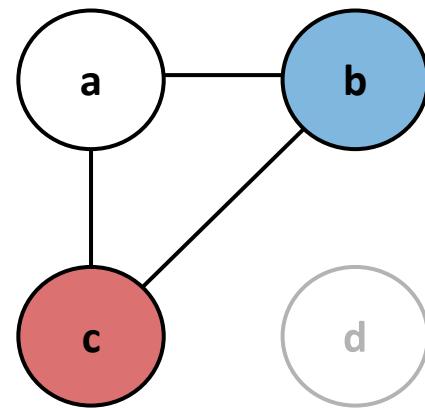
Example: Select

- We want to compute $k = 2$ coloring
 - Pop and (try to) assign colors



Example: Actually Spill

- We want to compute $k = 2$ coloring
- Pop and (try to) assign colors
- No color to assign a, need to *actually spill*



Example: Actually Spill

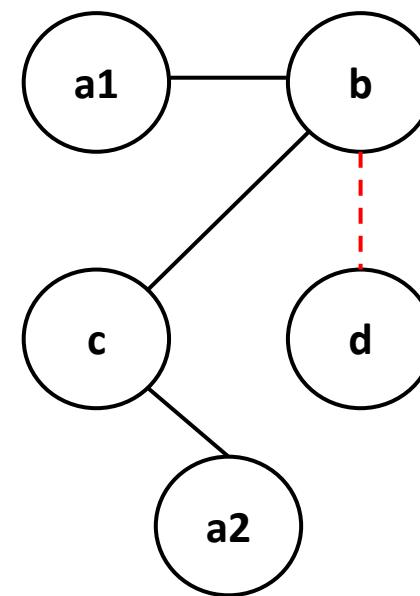
- Rewrite the IR & repeat

```
a = b  
b = a + b  
d = b  
c = c + a
```

Example: Actually Spill

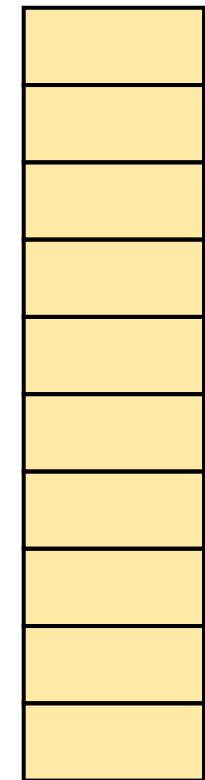
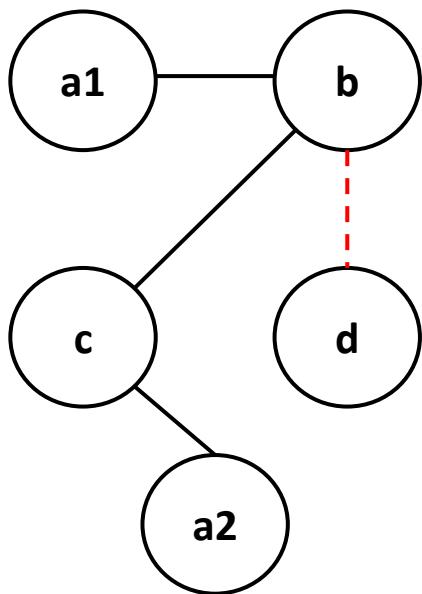
- Rewrite the IR & repeat

```
a1 = b
b = a1 + b
*(@a) = a1
d = b
a2 = *(@a)
c = c + a2
```



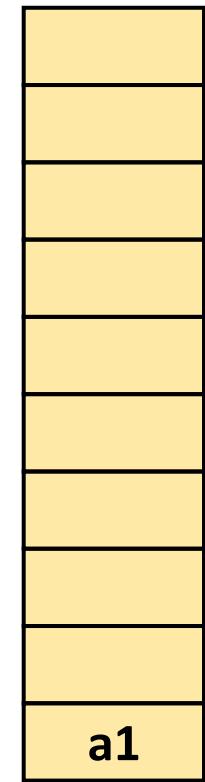
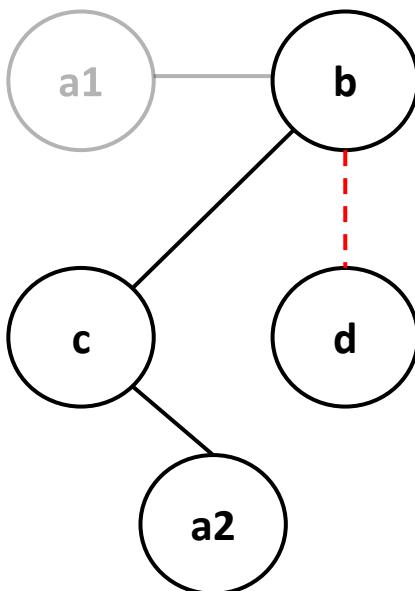
Example: Try Again

- We want to compute $k = 2$ coloring



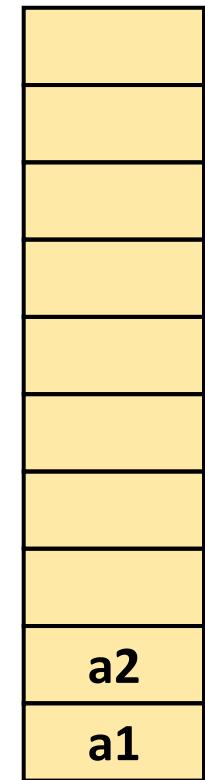
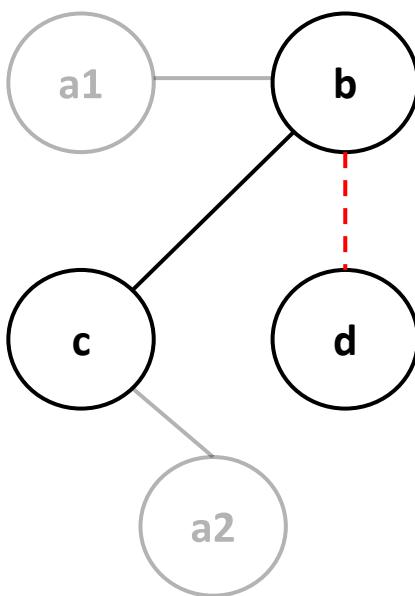
Example: Simplify

- We want to compute $k = 2$ coloring
- Simplify



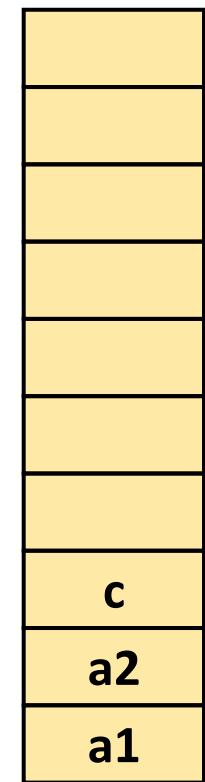
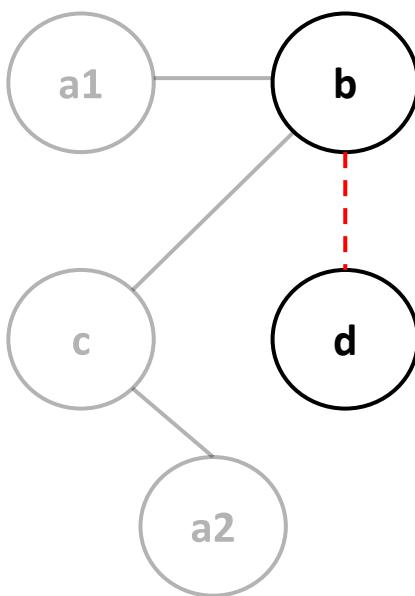
Example: Simplify

- We want to compute $k = 2$ coloring
- Simplify



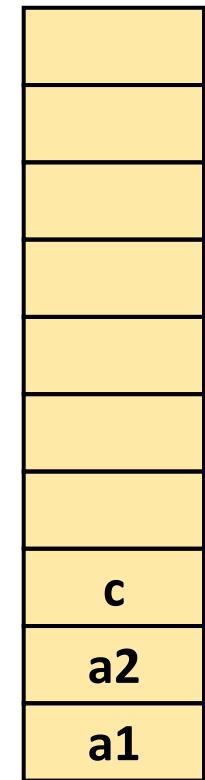
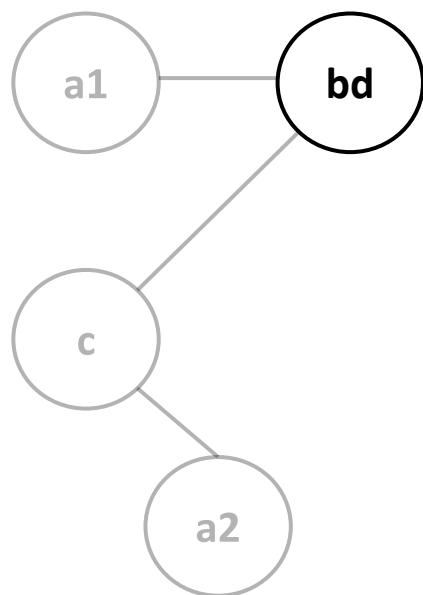
Example: Simplify

- We want to compute $k = 2$ coloring
- Simplify



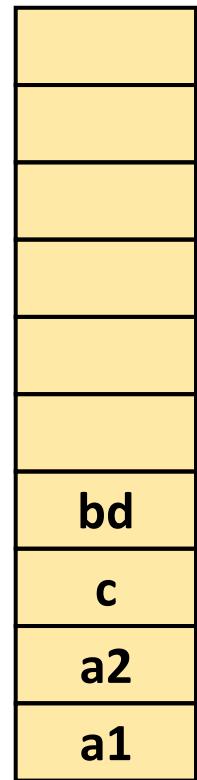
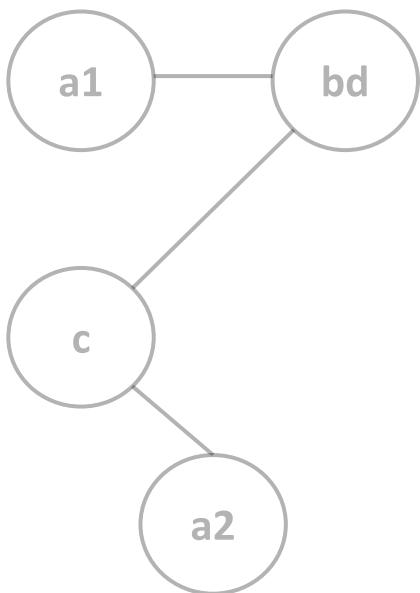
Example: Coalesce

- We want to compute $k = 2$ coloring
- Coalesce b & d



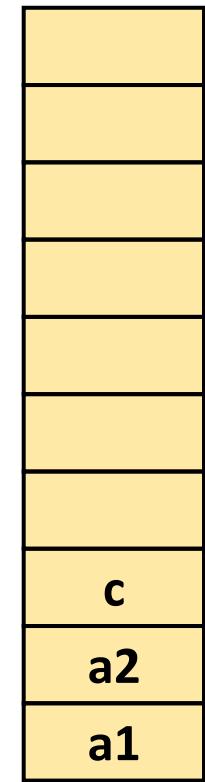
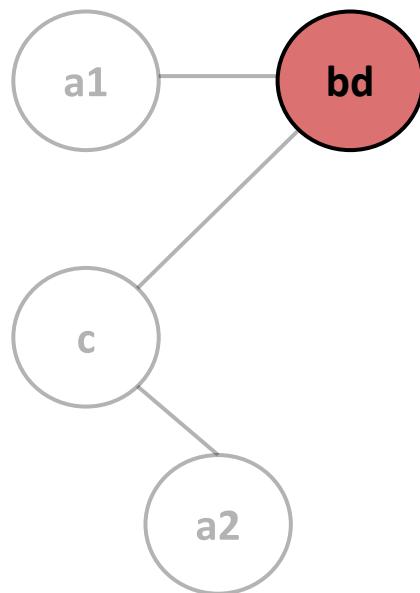
Example: Simplify

- We want to compute $k = 2$ coloring
- Simplify



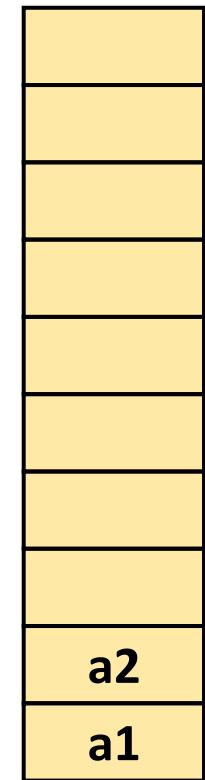
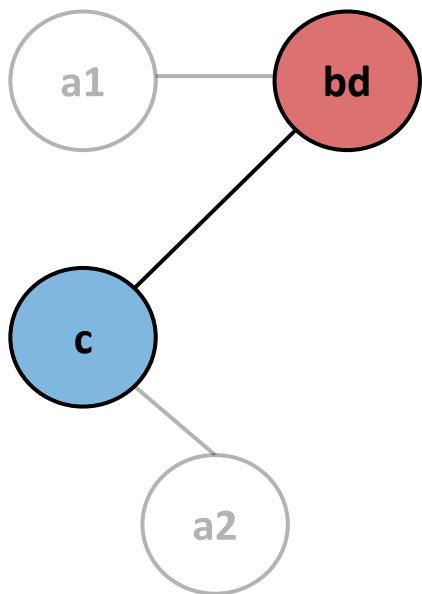
Example: Select

- We want to compute $k = 2$ coloring
- Pop and assign colors



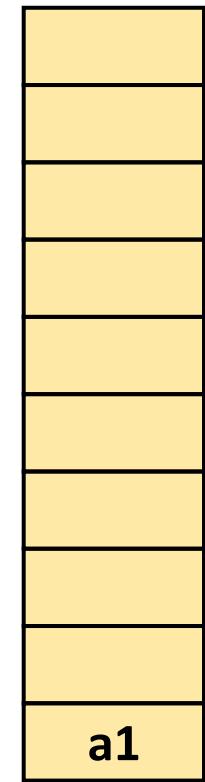
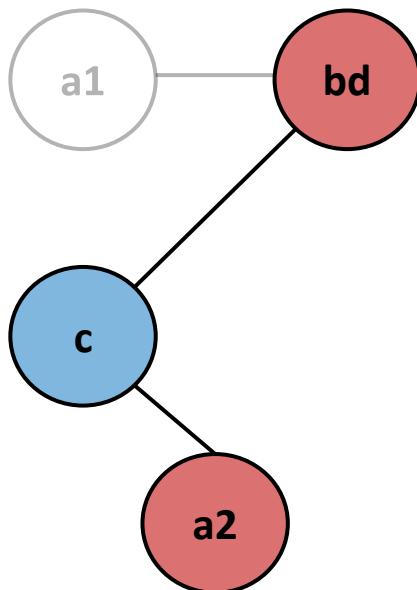
Example: Select

- We want to compute $k = 2$ coloring
 - Pop and assign colors



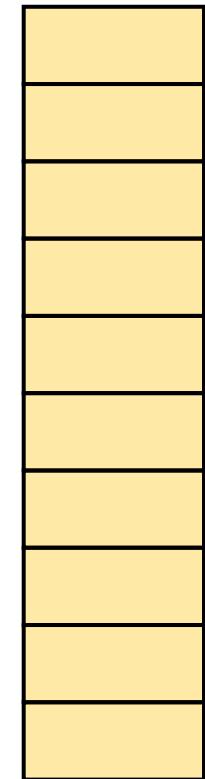
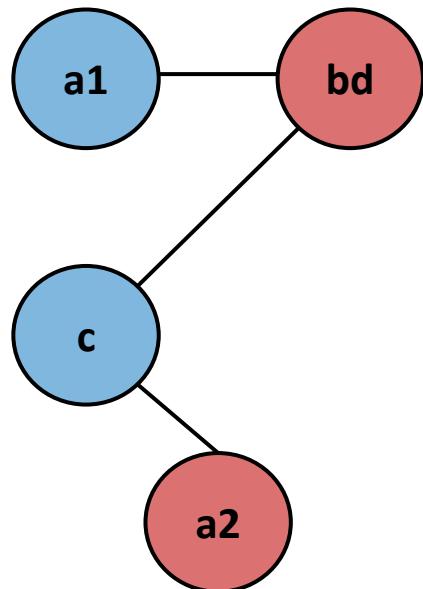
Example: Select

- We want to compute $k = 2$ coloring
- Pop and assign colors

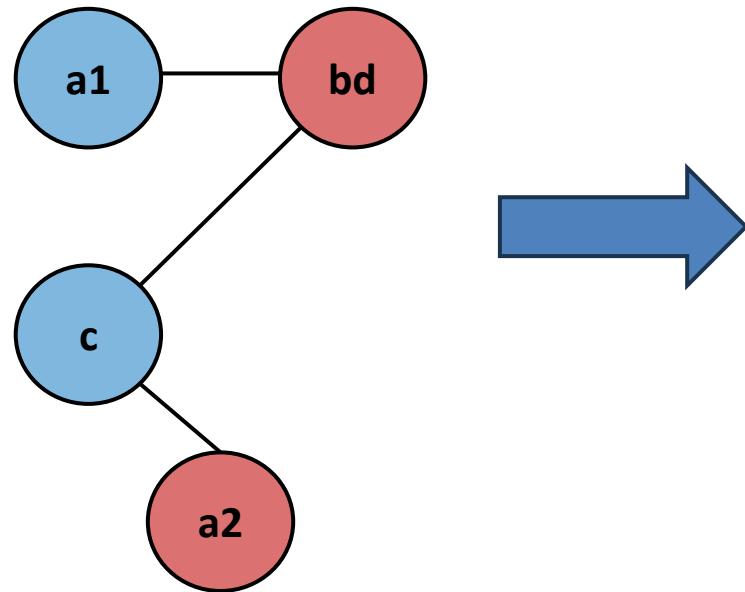


Example: Select

- We want to compute $k = 2$ coloring
- Pop and assign colors



Example: Final Register Allocation



```
R2 = R1  
R1 = R2 + R1  
*(@a) = R2  
R1 = *(@a)  
R2 = R2 + R1
```

In the Project

- We do register allocation on the IR
- Allocate physical registers **only** to IR registers
- Implement only the simplified register allocation algorithm
without spilling or coalescing
- If need to actually spill (coloring failed) output an **error message**