Deep Dive on Amazon Redshift

Storage Subsystem and Query Life Cycle

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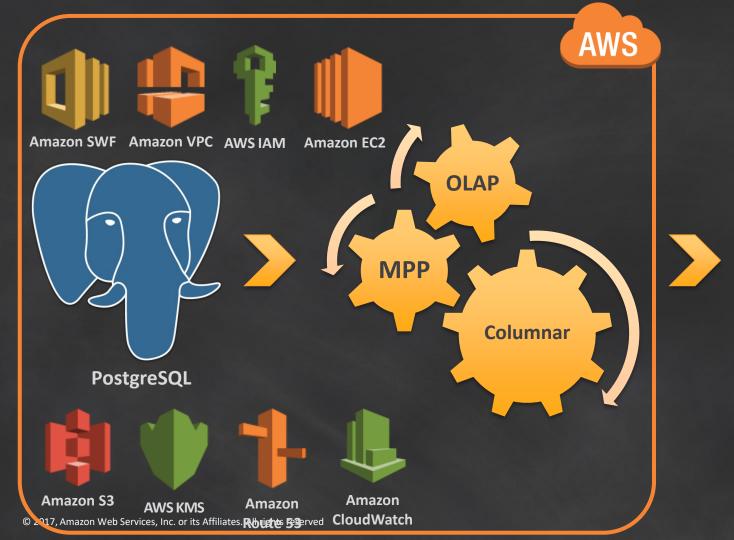


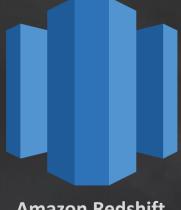
October 2017

Deep Dive Overview

- Amazon Redshift History and Development
- Cluster Architecture
- Concepts and Terminology
- Storage Deep Dive
- Query Life Cycle
- Data Ingestion Best Practices
- Recently Released Features
- Additional Resources
- Open Q&A









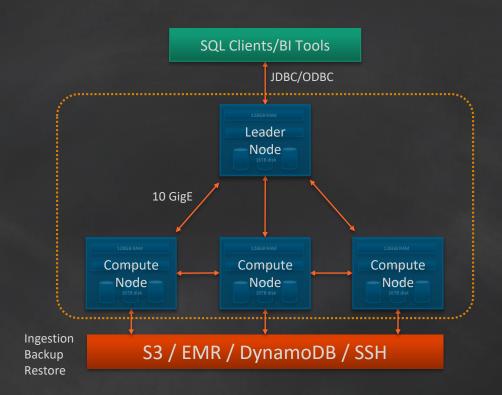


October 2017 February 2013 > 100 Significant Patches > 150 Significant Features

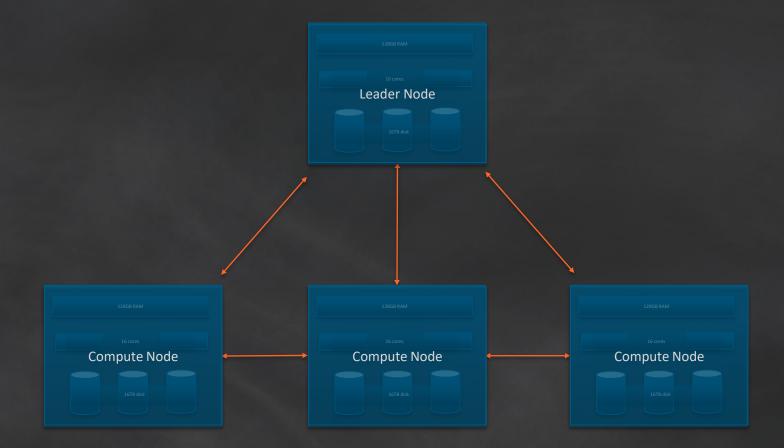


Redshift Cluster Architecture

- Massively parallel, shared nothing
- Leader node
 - SQL endpoint
 - Stores metadata
 - Coordinates parallel SQL processing
- Compute nodes
 - Local, columnar storage
 - Executes queries in parallel
 - Load, backup, restore









- Parser & Rewriter
- Planner & Optimizer
- Code Generator
 - <u>Input</u>: Optimized plan
 - Output: >=1 C++ functions
- Compiler

- Task Scheduler
- WLM
 - Admission
 - Scheduling
- PostgreSQL Catalog Tables



Leader Node



- Query execution processes
- Backup & restore processes
- Replication processes
- Local Storage
 - Disks
 - Slices
 - Tables
 - Columns
 - Blocks



Leader Node



- Query execution processes
- Backup & restore processes
- Replication processes
- Local Storage
 - Disks
 - Slices
 - Tables
 - Columns
 - Blocks









- Columnar storage
- Data compression
- Zone maps



```
CREATE TABLE deep_dive (
    aid INT --audience_id
    ,loc CHAR(3) --location
    ,dt DATE --date
);
```

aid	loc	dt
1	SFO	2016-09-01
2	JFK	2016-09-14
3	SFO	2017-04-01
4	JFK	2017-05-14

- Accessing dt with row storage:
 - Need to read everything
 - Unnecessary I/O



- Columnar storage
- Data compression
- Zone maps



```
CREATE TABLE deep_dive (
    aid INT --audience_id
    ,loc CHAR(3) --location
    ,dt DATE --date
);
```

aid	loc	dt
1	SFO	2016-09-01
2	JFK	2016-09-14
3	SFO	2017-04-01
4	JFK	2017-05-14

- Accessing dt with columnar storage:
 - Only scan blocks for relevant column



- Columnar storage
- Data compression
- Zone maps



```
CREATE TABLE deep_dive (
aid INT ENCODE LZO
,loc CHAR(3) ENCODE BYTEDICT
,dt DATE ENCODE RUNLENGTH
);
```

aid	loc	dt
1	SFO	2016-09-01
2	JFK	2016-09-14
3	SFO	2017-04-01
4	JFK	2017-05-14

- Columns grow and shrink independently
- Effective compression ratios due to like data
- Reduces storage requirements
- Reduces I/O



- Columnar storage
- Data compression
- Zone maps



```
CREATE TABLE deep_dive (
    aid INT --audience_id
    ,loc CHAR(3) --location
    ,dt DATE --date
);
```

aid	loc	dt
1	SFO	2016-09-01
2	JFK	2016-09-14
3	SFO	2017-04-01
4	JFK	2017-05-14

- In-memory block metadata
- Contains per-block MIN and MAX value
- Effectively prunes blocks which cannot contain data for a given query
- Eliminates unnecessary I/O



Zone Maps and Sorting

SELECT COUNT(*) FROM deep dive WHERE dt = '06-09-2013'

Unsorted Table



MIN: 01-JUNE-2013

MAX: 20-JUNE-2013



MIN: 08-JUNE-2013

MAX: 30-JUNE-2013



MIN: 12-JUNE-2013

MAX: 20-JUNE-2013



MIN: 02-JUNE-2013

MAX: 25-JUNE-2013

Sorted By Date



MIN: 01-JUNE-2013

MAX: 06-JUNE-2013



MIN: 07-JUNE-2013

MAX: 12-JUNE-2013



MIN: 13-JUNE-2013

MAX: 18-JUNE-2013



MIN: 19-JUNE-2013

MAX: 24-JUNE-2013

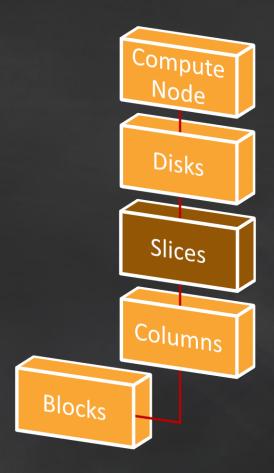
Terminology and Concepts: Data Sorting

- Goal:
 - Make queries run faster by optimizing the effectiveness of zone maps.
 - Typically on the columns that are filtered on (where clause predicates)
- Impact:
 - Enables range restricted scans to prune blocks by leveraging zone maps
 - Overall reduction in block I/O
- Achieved with the table property SORTKEY defined on one or more columns
- Optimal SORTKEY is dependent on:
 - Query patterns
 - Data profile
 - Business requirements



Terminology and Concepts: Slices

- A slice can be thought of like a "virtual compute node"
 - Unit of data partitioning
 - Parallel query processing
- Facts about slices:
 - Each compute node has either 2, 16, or 32 slices
 - Table rows are distributed to slices
 - A slice processes only its own data





Data Distribution

Distribution style is a table property which dictates how that table's data is distributed throughout the cluster:

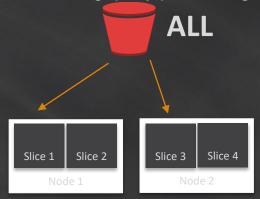
• KEY: Value is hashed, same value goes to same location (slice)

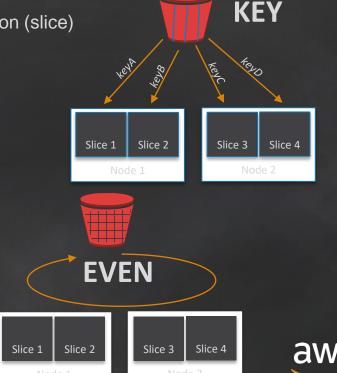
ALL: Full table data goes to first slice of every node

EVEN: Round robin

Goals:

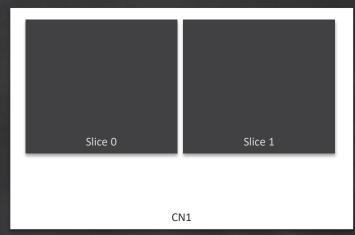
- Distribute data evenly for parallel processing
- Minimize data movement during query processing

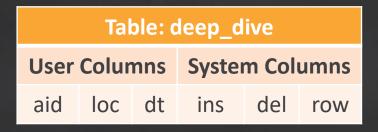




Data Distribution: Example

```
CREATE TABLE deep_dive (
    aid INT --audience_id
    ,loc CHAR(3) --location
    ,dt DATE --date
) DISTSTYLE (EVEN|KEY|ALL);
```



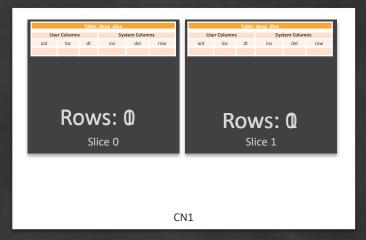






Data Distribution: **EVEN** Example

```
CREATE TABLE deep_dive (
    aid INT --audience_id
    ,loc CHAR(3) --location
    ,dt DATE --date
) DISTSTYLE EVEN;
```





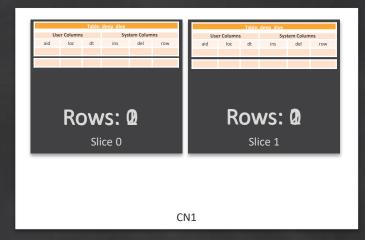


(3 User Columns + 3 System Columns) x (4 slices) = 24 Blocks (24MB)

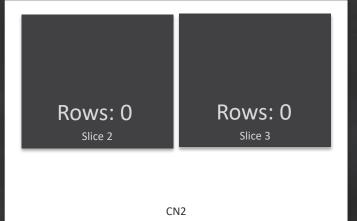


Data Distribution: **KEY** Example #1

```
CREATE TABLE deep_dive (
    aid INT --audience_id
    ,loc CHAR(3) --location
    ,dt DATE --date
) DISTSTYLE KEY DISTKEY (loc);
```





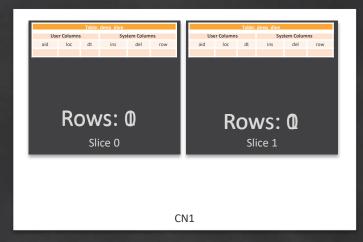


(3 User Columns + 3 System Columns) x (2 slices) = 12 Blocks (12 MB)

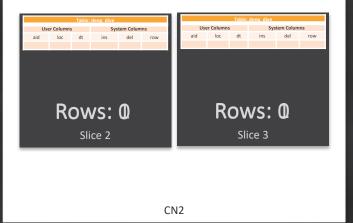


Data Distribution: **KEY** Example #2

```
CREATE TABLE deep_dive (
    aid INT --audience_id
    ,loc CHAR(3) --location
    ,dt DATE --date
) DISTSTYLE KEY DISTKEY (aid);
```





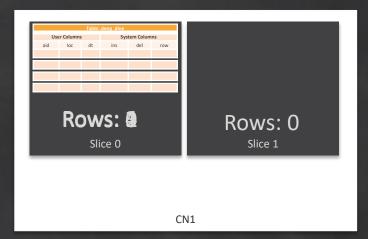


(3 User Columns + 3 System Columns) x (4 slices) = 24 Blocks (24 MB)

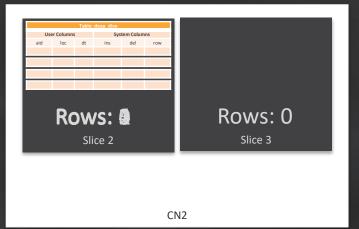


Data Distribution: ALL Example

```
CREATE TABLE deep_dive (
    aid INT --audience_id
    ,loc CHAR(3) --location
    ,dt DATE --date
) DISTSTYLE ALL;
```



```
INSERT INTO deep_dive VALUES
(1, 'SFO', '2016-09-01'),
(2, 'JFK', '2016-09-14'),
(3, 'SFO', '2017-04-01'),
(4, 'JFK', '2017-05-14');
```



(3 User Columns + 3 System Columns) x (2 slice) = 12 Blocks (12MB)



Terminology and Concepts: Data Distribution

- KEY
 - The key creates an even distribution of data
 - Joins are performed between large fact/dimension tables
 - Optimizing joins and group by
- ALL
 - Small and medium size dimension tables (< 2-3M)
- EVEN
 - When key cannot produce an even distribution

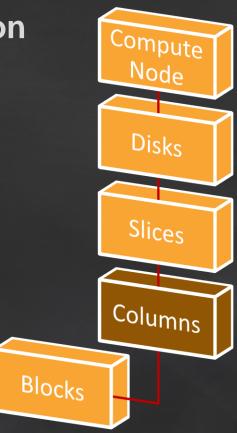




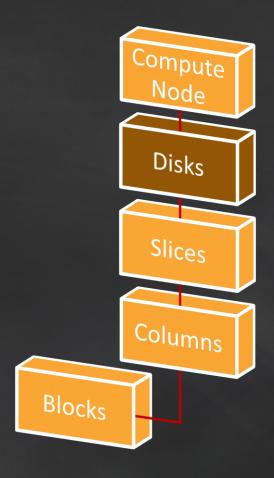
Table Design Considerations

- Denormalize/materialize heavily filtered columns into the fact table(s)
 - Timestamp/date should be in fact tables rather than using time dimension tables
- Add SORT KEYS to medium and large sized tables on the primary columns that are filtered on
 - Not necessary to add them to small tables
- Keep variable length data types as long as required
 - VARCHAR, CHAR and NUMERIC
- Add compression to tables
 - Optimal compression can be found using ANALYZE COMPRESSION



Storage Deep Dive: Disks

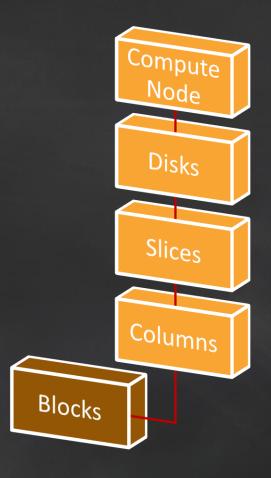
- Redshift utilizes locally attached storage devices
- Compute nodes have 2.5-3x the advertised storage capacity
- 1, 3, 8, or 24 disks depending on node type
- Each disk is split into two partitions
 - Local data storage, accessed by local CN
 - Mirrored data, accessed by remote CN
- Partitions are raw devices
 - Local storage devices are ephemeral in nature
 - Tolerant to multiple disk failures on a single node





Storage Deep Dive: Blocks

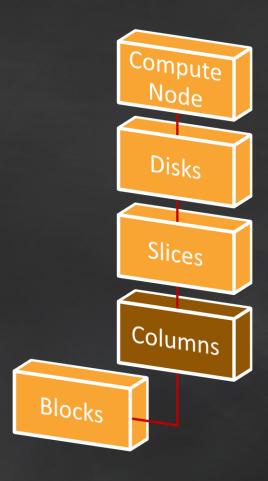
- Column data is persisted to 1MB immutable blocks
- Each block contains in-memory metadata:
 - Zone Maps (MIN/MAX value)
 - Location of previous/next block
- Blocks are individually compressed with 1 of 10 encodings
- A full block contains between 16 and 8.4 million values





Storage Deep Dive: Columns

- Column: Logical structure accessible via SQL
- Physical structure is a doubly linked list of blocks
- These blockchains exist on each slice for each column
- All sorted & unsorted blockchains compose a column
- Column properties include:
 - Distribution Key
 - Sort Key
 - Compression Encoding
- Columns shrink and grow independently, 1 block at a time
- Three system columns per table-per slice for MVCC





Block Properties: Design Considerations

- Small writes:
 - Batch processing system, optimized for processing massive amounts of data
 - 1MB size + immutable blocks means that we clone blocks on write so as not to introduce fragmentation
 - Small write (~1-10 rows) has similar cost to a larger write (~100 K rows)
- UPDATE and DELETE:
 - Immutable blocks means that we only logically delete rows on UPDATE or DELETE
 - Must VACUUM or DEEP COPY to remove ghost rows from table



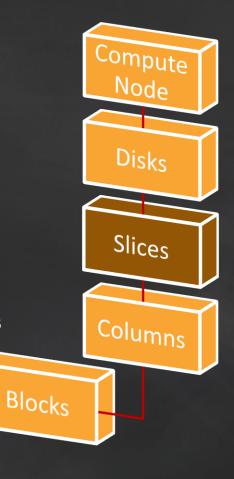
Column Properties: Design Considerations

- Compression:
 - COPY automatically analyzes and compresses data when loading into empty tables
 - ANALYZE COMPRESSION checks existing tables and proposes optimal compression algorithms for each column
 - Changing column encoding requires a table rebuild
- DISTKEY and SORTKEY significantly influence performance (orders of magnitude)
- Distribution Keys:
 - A poor DISTKEY can introduce data skew and an unbalanced workload
 - A query completes only as fast as the slowest slice completes
- Sort Keys:
 - A sortkey is only effective as the data profile allows it to be
 - Selectivity needs to be considered



Storage Deep Dive: Slices

- Each compute node has either 2, 16, or 32 slices
- A slice can be thought of like a "virtual compute node"
 - Unit of data partitioning
 - Parallel query processing
- Facts about slices:
 - Table rows are distributed to slices.
 - A slice processes only its own data
 - Within a compute node all slices read from and write to all disks





- Parser & Rewriter
- Planner & Optimizer
- Code Generator
 - <u>Input</u>: Optimized plan
 - Output: >=1 C++ functions
- Compiler

- Task Scheduler
- WLM
 - Admission
 - Scheduling
- PostgreSQL Catalog Tables
- Redshift System Tables (STV)



Leader Node



- Parser & Rewriter
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Leader Node



Query Execution Terminology

- Step: An individual operation needed during query execution. Steps need to be combined to allow compute nodes to perform a join. Examples: scan, sort, hash, aggr
- Segment: A combination of several steps that can be done by a single process.
 The smallest compilation unit executable by a slice. Segments within a stream run in parallel.
- Stream: A collection of combined segments which output to the next stream or SQL client.

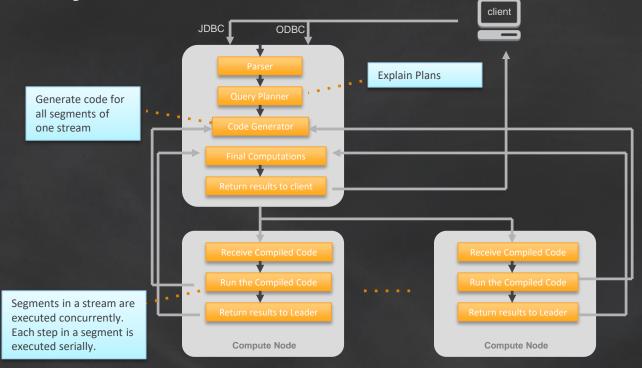


Visualizing Streams, Segments, and Steps





Query Lifecycle



Query Execution Deep Dive: Leader Node

- 1. The leader node receives the query and parses the SQL.
- 2. The parser produces a logical representation of the original query.
- 3. This query tree is input into the query optimizer (volt).
- 4. Volt rewrites the query to maximize its efficiency. Sometimes a single query will be rewritten as several dependent statements in the background.
- 5. The rewritten query is sent to the planner which generates >= 1 query plans for the execution with the best estimated performance.
- 6. The query plan is sent to the execution engine, where it's translated into **steps**, **segments**, and **streams**.
- This translated plan is sent to the code generator, which generates a C++ function for each segment.
- 8. This generated C++ is compiled with gcc to a .o file and distributed to the compute nodes.



Query Execution Deep Dive: Compute Nodes

- Slices execute the query segments in parallel.
- Executable segments are created for one stream at a time. When the segments
 of that stream are complete, the engine generates the segments for the next
 stream.
- When the compute nodes are done, they return the query results to the leader node for final processing.
- The leader node merges the data into a single result set and addresses any needed sorting or aggregation.
- The leader node then returns the results to the client.

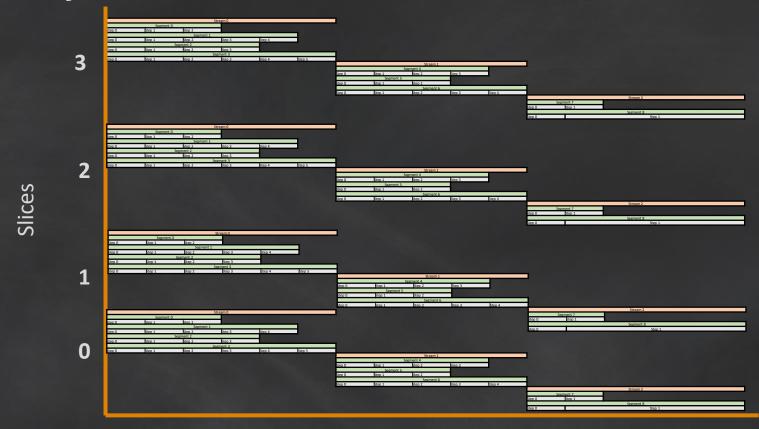


Visualizing Streams, Segments, and Steps





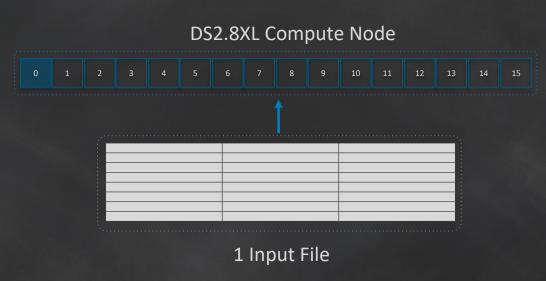
Query Execution





Data Ingestion: COPY Statement

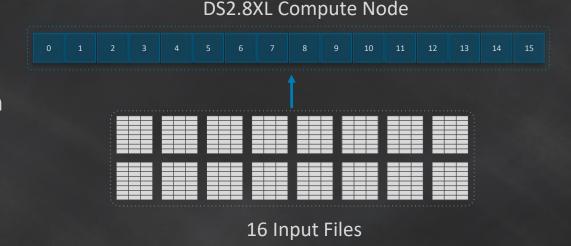
- Ingestion Throughput:
 - Each slice's query processors can load one file at a time:
 - Streaming decompression
 - Parse
 - Distribute
 - Write
- Realizing only partial node usage as 6.25% of slices are active





Data Ingestion: COPY Statement

- Number of input files should be a multiple of the number of slices
- Splitting the single file into 16 input files, all slices are working to maximize ingestion performance
- COPY continues to scale linearly as you add nodes



Recommendation is to use delimited files – 1MB to 1GB after GZIP compression



Best Practices: Data Ingestion

- Export Data from Source System
 - Delimited Files are recommend
 - Pick a simple delimiter '|' or ',' or tabs
 - Pick a simple NULL character (\N)
 - Use double quotes and an escape character ('\') for varchars
 - UTF-8 varchar columns take 4 bytes per char
 - Split Files so there is a multiple of the number of slices
 - Files sizes should be 1MB 1GB after gzip compression
- Useful COPY Options
 - MAXERRORS
 - ACCEPTINVCHARS
 - NULLAS



Data Ingestion: Deduplication / UPSERT

Table: deep_dive		
aid	loc	dt
1	SFO	2017-10-20
2	JFK	2017-10-20
3	SFO	2017-04-01
4	JFK	2017-05-14
5	SJC	2017-10-10
6	SEA	2017-11-29

s3://bucket/dd.csv		
aid	loc	dt
1	SFO	2017-10-20
2	JFK	2017-10-20
5	SJC	2017-10-10
6	SEA	2017-11-29



Data Ingestion: Deduplication / UPSERT

Steps:

1. Load CSV data into a staging table

2. Delete duplicate data from the production table

3. Insert (or append) data from the staging into the production table



Data Ingestion: Deduplication / UPSERT

```
BEGIN;
CREATE TEMP TABLE staging (LIKE deep dive);
COPY INTO staging : 'creds' COMPUPDATE OFF;
DELETE deep dive d
USING staging s WHERE d.aid = s.aid;
INSERT INTO deep dive SELECT * FROM staging;
DROP TABLE staging;
COMMIT;
```



Best Practices: Deduplication / UPSERT

- Wrap workflow/statements in an explicit transaction
- With Staging Tables:
 - Use temporary tables or "NO BACKUP" option
 - If possible use DISTSTYLE KEY on both the staging table and production table to speed up the INSERT/SELECT statement
 - Turn off automatic compression COMPUPDATE OFF
 - Copy compression settings from production table or use ANALYZE
 COMPRESSION statement
 - Use CREATE TABLE LIKE or write encodings into the DDL



Query Monitoring Rules (QMR)

- Goal Automatic handling of run-away (poorly written) queries
- Extension to WLM
- Rules applied onto a WLM Queue
 - Queries can be LOGGED, CANCELED or HOPPED



Query Monitoring Rules (QMR)

- Metrics with operators and values (e.g. query_cpu_time > 1000) create a predicate
- Multiple predicates can be AND-ed together to create a rule
- Multiple rules can be defined for a queue in WLM. These rules are OR-ed together

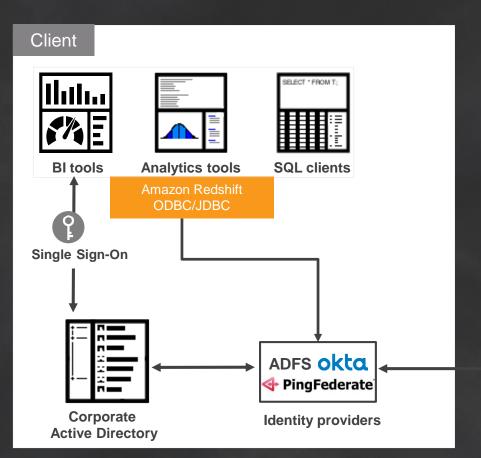
```
If { rule } then [action]
      { rule : metric operator value } e.g.: rows_scanned >
100000
```

- Metric cpu_time, query_blocks_read, rows scanned, query execution time, cpu & io skew per slice, join_row_count, etc.
- Operator: <, >, ==
- Value integer

[action] hop, log, abort



IAM Authentication





New Redshift ODBC/JDBC drivers. Grab the ticket (userid) and get a SAML assertion.



Amazon Redshift Spectrum

Run SQL queries directly against data in S3 using thousands of nodes





High concurrency:
Multiple clusters access
same data

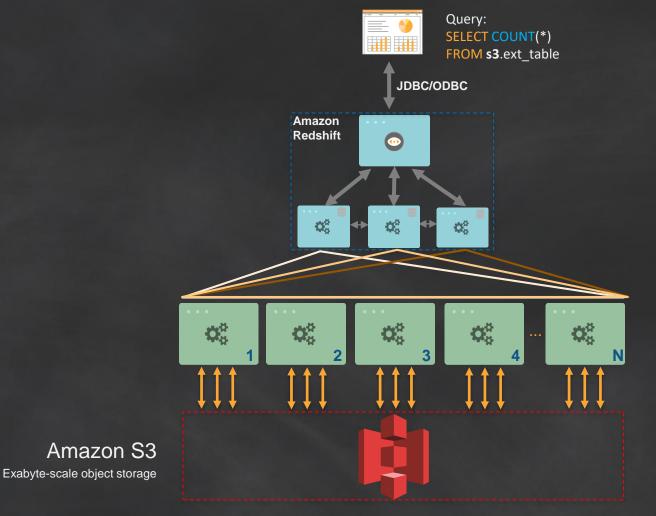




No ETL: Query data inplace using open file formats









Data Catalog
Apache Hive Metastore

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An Exabyte Query: Harry Potter

```
demo=# SELECT
       P.ASIN.
        P.TITLE,
demo-#
demo-#
        R.POSTAL_CODE.
        P.RELEASE_DATE,
demo-#
         SUM(D.QUANTITY * D.OUR_PRICE) AS SALES_sum
demo-#
demo-# FROM s3.d_customer_order_item_details D, asin_attributes A, products P, regions r
demo-# WHERE D.ASIN = P.ASIN AND
demo-#
            P.ASTN = A.ASTN AND
demo-#
            D.REGION_ID = R.REGION_ID AND
demo-#
            A.EDITION LIKE '%FIRST%' AND
            P.TITLE LIKE '%Potter%' AND
demo-#
            P.AUTHOR = 'J. K. Rowling' AND
demo-#
            R.COUNTRY\_CODE = 'US' AND
demo-#
demo-#
            R.CITY = 'Seattle' AND
demo-#
            R.STATE = 'WA' AND
            D.ORDER_DAY :: DATE >= P.RELEASE_DATE AND
demo-#
            D.ORDER_DAY :: DATE < dateadd(day, 3, P.RELEASE_DATE)</pre>
demo-#
demo-# GROUP BY P.ASIN, P.TITLE, R.POSTAL_CODE, P.RELEASE_DATE
demo-# ORDER BY sales_sum DESC
demo-# LIMIT 20;
```

Roughly 140 TB of customer item order detail records for each day over past 20 years.

190 million files across 15,000 partitions in S3. One partition per day for USA and rest of world.

Need a **billion-fold** reduction in data processed.

Running this query using a 1000 node Hive cluster would take over 5 years.*

- * Estimated using 20 node Hive cluster & 1.4TB, assume linear
- * Query used a 20 node DC1.8XLarge Amazon Redshift cluster
- * Not actual sales data generated for this demo based on data format used by Amazon Retail.

Additional Resources

AWS Labs on Github - Amazon Redshift

- https://github.com/awslabs/amazon-redshift-utils
- https://github.com/awslabs/amazon-redshift-monitoring
- https://github.com/awslabs/amazon-redshift-udfs
- Admin Scripts

Collection of utilities for running diagnostics on your cluster

- Admin Views
 - Collection of utilities for managing your cluster, generating schema DDL, etc.
- Analyze Vacuum Utility

Utility that can be scheduled to vacuum and analyze the tables within your Amazon Redshift cluster

- Column Encoding Utility
 - Utility that will apply optimal column encoding to an established schema with data already loaded



AWS Big Data Blog - Amazon Redshift

Amazon Redshift Engineering's Advanced Table Design Playbook

https://aws.amazon.com/blogs/big-data/amazon-redshift-engineerings-advanced-table-design-playbook-preamble-prerequisites-and-prioritization/

- Zach Christopherson

Top 10 Performance Tuning Techniques for Amazon Redshift

https://aws.amazon.com/blogs/big-data/top-10-performance-tuning-techniques-for-amazon-redshift/

- Ian Meyers and Zach Christopherson

10 Best Practices for Amazon Redshift Spectrum

https://aws.amazon.com/blogs/big-data/10-best-practices-for-amazon-redshift-spectrum/

- Po Hong and Peter Dalton





Thank you!

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