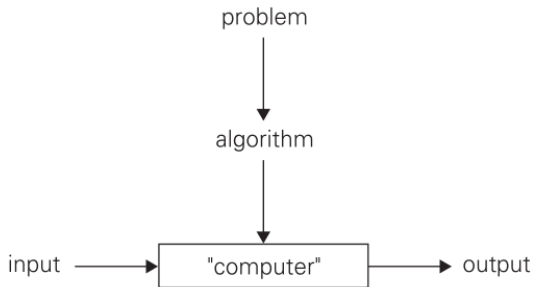

CS142 Wrapped

(Design and Analysis of Algorithms)

*Department of Mathematics and Computer Science
College of Science
University of the Philippines Baguio*

Algorithms

- ▶ The notion of an algorithm constitutes the **cornerstone of computer science**.
- ▶ Algorithms lie at the heart of **practical computing**.



Classifying Algorithms

Although algorithms can be classified in numerous ways, two of them are particularly important:

- **By Underlying Design Technique**
- **By Efficiency**

These two principal dimensions reflect the **needs of computing practice**.

General Design Techniques

1. Brute Force/ Exhaustive Search
2. Decrease-and-Conquer
3. Divide-and-Conquer
4. Transform-and-Conquer
5. Space-Time Trade-offs
6. Dynamic programming
7. Greedy technique
8. Backtracking
9. Branch-and-bound

There are other areas of algorithmics:

- ▶ **randomized algorithms** – makes random choices during its execution
- ▶ **parallel algorithms** – takes advantage of the capability of some newer computers to execute operations concurrently

Basic Efficiency Classes

The analysis framework classifies algorithms by the **order of growth** of their running time as a function of input size.

Class	Notation	Important examples
<i>constant time</i>	$\Theta(1)$	hashing (on average)
<i>logarithmic</i>	$\Theta(\log n)$	binary search (worst and average cases)
<i>linear</i>	$\Theta(n)$	sequential search (worst and average cases)
<i>linearithmic</i>	$\Theta(n \log n)$	advanced sorting algorithms
<i>quadratic</i>	$\Theta(n^2)$	elementary sorting algorithms
<i>cubic</i>	$\Theta(n^3)$	Gaussian elimination
<i>exponential</i>	$\Omega(a^n)$	combinatorial problems

Lower Bound Arguments and Decision Trees

Given a class of algorithms for solving a particular problem, a **lower bound** indicates the *best possible efficiency* any algorithm from this class can have.

- **Information—theoretic lower bound** - usually obtained through a mechanism of **decision trees**. This technique is particularly useful for comparison-based algorithms for sorting and searching. Specifically,
 - ▷ Any general comparison-based sorting algorithm must perform at least $\lceil \log_2 n! \rceil \approx n \log_2 n$ key comparisons in the worst case.
 - ▷ Any general comparison-based algorithm for searching a sorted array must perform at least $\lceil \log_2(n + 1) \rceil$ key comparisons in the worst case.

Lower Bound Arguments and Decision Trees

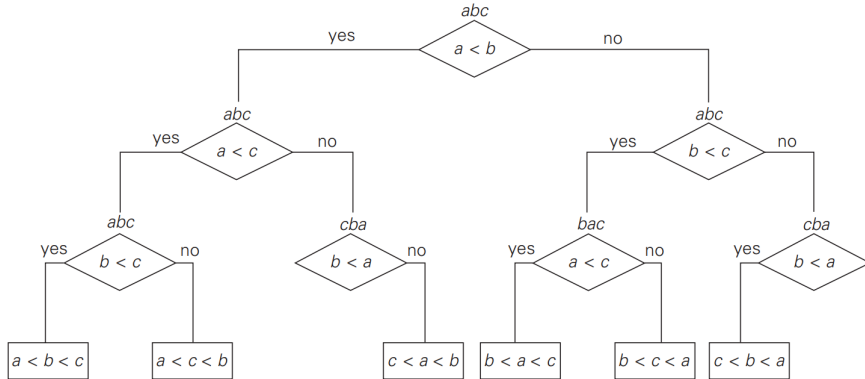


Figure: Decision tree for the three-element selection sort.

Lower Bound Arguments and Decision Trees

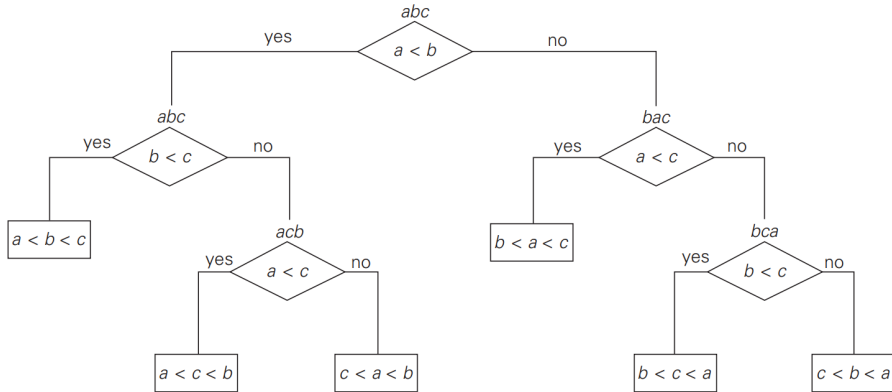


Figure: Decision tree for the three-element insertion sort.

Complexity Theory

Complexity theory seeks to classify problems according to their computational complexity: **tractable** and **intractable problems** – problems that can and cannot be solved in polynomial time, respectively.

Class P - class of all decision problems that can be **solved** in polynomial time

Class NP - class of all decision problems whose randomly guessed solutions can be **verified** in polynomial time.

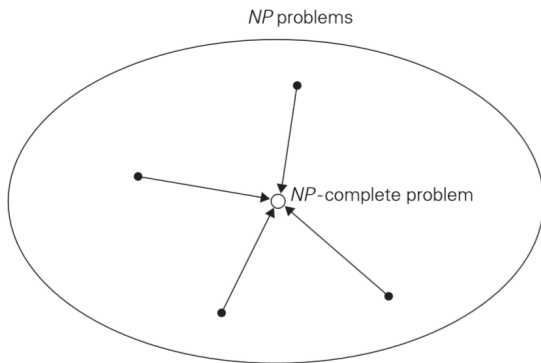
Is $P = NP$?

Most important unresolved issue in theoretical computer science

NP – Complete Problems

A decision problem D is said to be NP –**complete** if:

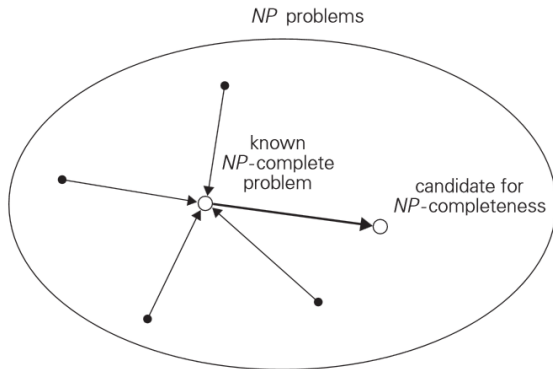
- (i) it belongs to class NP
- (ii) every problem in NP is **polynomially reducible** to D



NP – Complete Problems

A decision problem D is said to be NP –**complete** if:

- (i) it belongs to class NP
- (ii) every problem in NP is **polynomially reducible** to D



NP –**completeness** implies that if there exists a deterministic polynomial-time algorithm for just one NP –complete problem, then $P = NP$.

“...whichever direction you take in your future journey through the land of algorithms in your studies and your career, the road ahead is as exciting as it has ever been.”
-Anany Levitin

The **field of algorithms** is **constantly evolving**.
Stay curious and keep building your understanding.

**Thank you for stopping by
at the CS142 Place!**