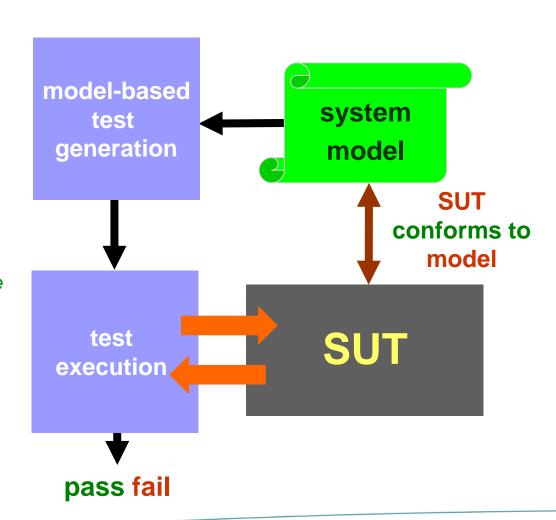
A Theory of Model-Based Testing with Labelled Transition Systems

The uioco Theory

MBT: Model-Based Testing

SUT conforms to model

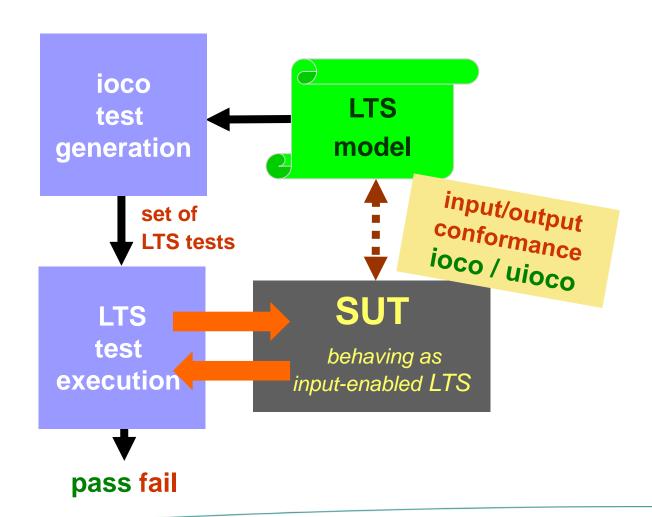
SUT passes tests



MBT: Labelled Transitions Systems

SUT uioco model

SUT passes tests

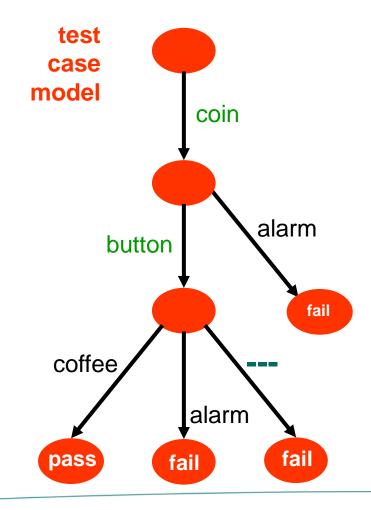


Models: Generation of Test Cases

specification model coin alarm button

button



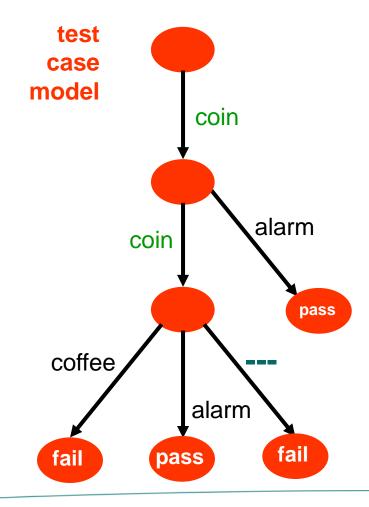


Models: Generation of Test Cases

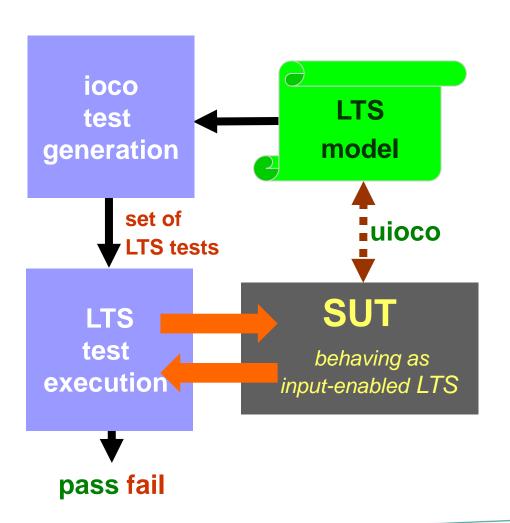
specification model coin alarm button

button





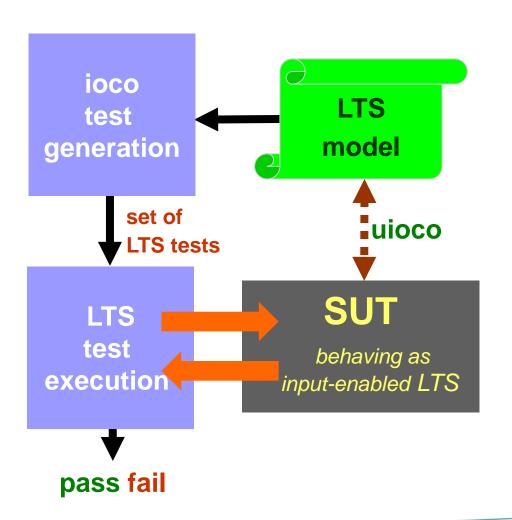
MBT: Labelled Transitions Systems



MBT with LTS topics:

- specification model
- implementation (SUT)
- implementation model
- conformance uioco
- test cases
- test generation
- test execution
- test result analysis
- sound & exhaustive

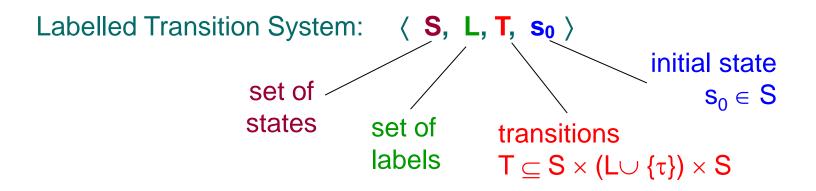
MBT: Labelled Transitions Systems



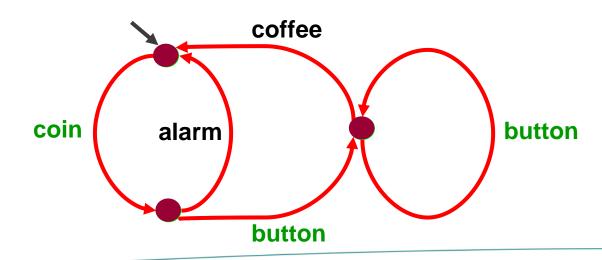
MBT with LTS topics:

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Models: Labelled Transition Systems







Models: LTS with Inputs and Outputs

coin, button

from user to machine initiative with user machine cannot refuse

? inputs L_I

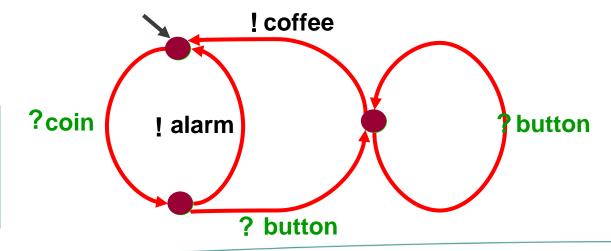
coffee, alarm

from machine to user initiative with machine user cannot refuse

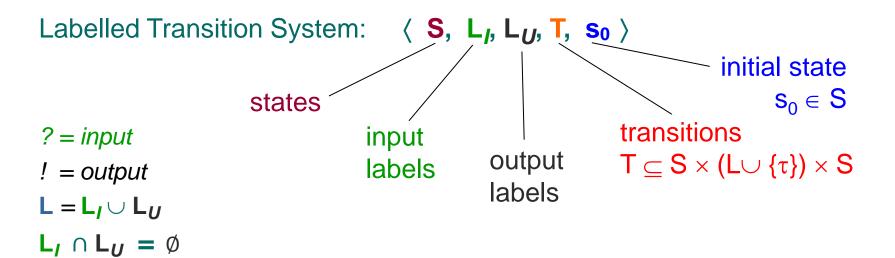
! outputs L_U

$$L_{I} \cap L_{U} = \emptyset$$

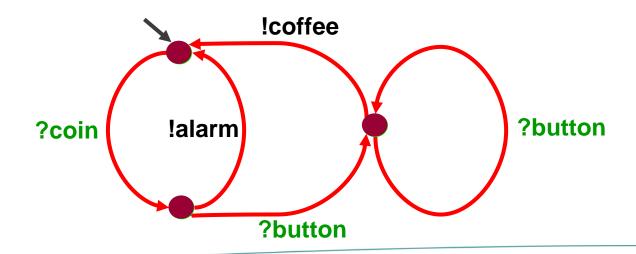
$$L_{I} \cup L_{U} = L$$



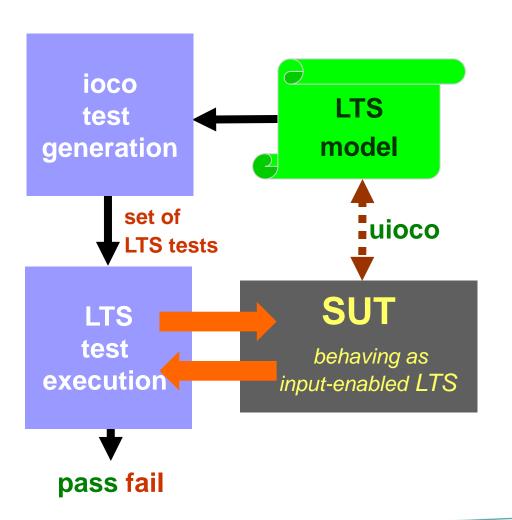
Models: LTS with Inputs and Outputs







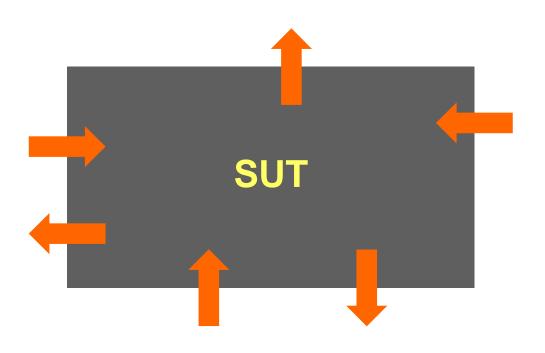
MBT: Labelled Transitions Systems



MBT with LTS topics:

- specification model
- implementation (SUT)
- implementation model
- conformance uioco
- * test cases
- test generation
- test execution
- test result analysis
- sound & exhaustive

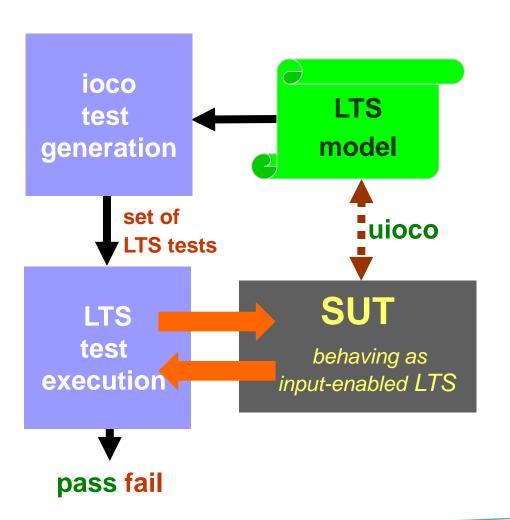
MBT: System Modelling



SUT

- black box
- inputs on interfaces
- outputs
 on interfaces

MBT: Labelled Transitions Systems



MBT with LTS topics:

- specification model
- implementation (SUT)
- implementation model
- conformance uioco
- test cases
- test generation
- test execution
- test result analysis
- sound & exhaustive

Models: Input-Output Transition Systems

In many systems, inputs are always enabled: input-enabled transition systems

= Input-Output Transition Systems

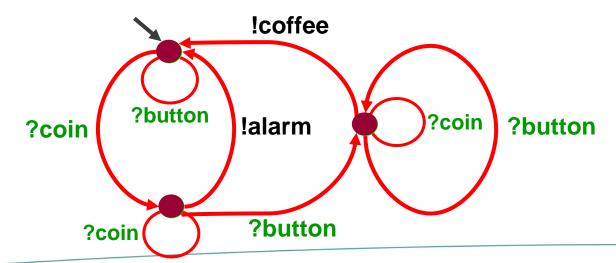
$$IOTS(L_I, L_U) \subseteq LTS(L_I \cup L_U)$$

input enabled:

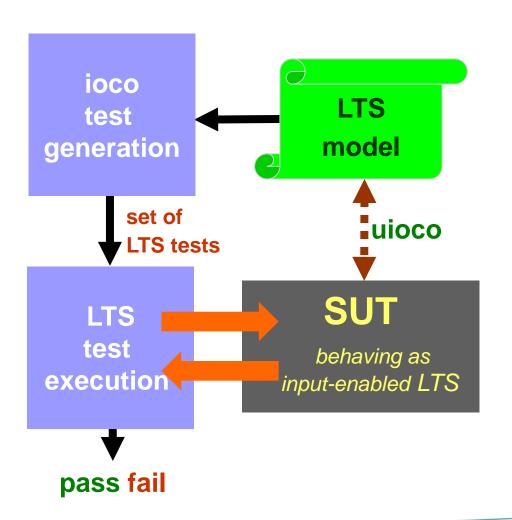
for all states **s**,

for all inputs $?a \in L_I$:

$$s \stackrel{?a}{\Longrightarrow}$$



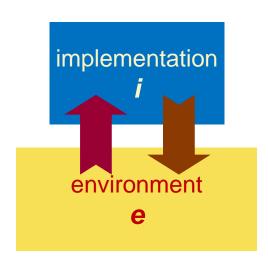
MBT: Labelled Transitions Systems



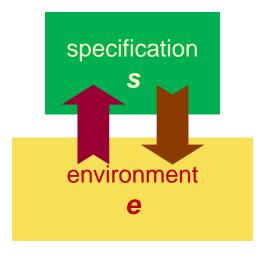
MBT with LTS topics:

- specification model
- implementation (SUT)
- implementation model
- conformance uioco
- * test cases
- test generation
- test execution
- test result analysis
- sound & exhaustive

Implementation Relations for Input-Output Transition Systems



imp



$$i \in IOTS(L_I, L_U)$$

$$s \in LTS(L_I, L_U)$$

imp
$$\subseteq IOTS(L_{I},L_{U}) \times LTS(L_{I},L_{U})$$

i imp s

i is a conforming implementation of s

```
i uioco s = def \forall \sigma \in Utraces(s): out(i after <math>\sigma) \subseteq out(s after \sigma)
```

s is a Labelled Transition System

i is (assumed to be) an input-enabled LTS

Alternative: ioco:

(see Lecture Notes)

```
i ioco s =<sub>def</sub> \forall \sigma \in Straces(s): out(i after \sigma) \subseteq out(s after \sigma)
```

```
i uioco s = def \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)
                                                             s is a Labelled Transition System
                                               i is (assumed to be) an input-enabled LTS
   p \xrightarrow{\delta} p \iff \forall !x \in L_U \cup \{\tau\}. p \xrightarrow{!x} \iff p \text{ refuses } L_U
    Straces (s) = { \sigma \in (L \cup \{\delta\})^* \mid s \stackrel{\sigma}{\Longrightarrow} \}
     Utraces(s) = \{ \sigma \in Straces(s) \mid
               \forall \sigma_1 ? b \sigma_2 = \sigma : not(s after \sigma_1 refuses {?b})
   out(P) = { |\mathbf{x} \in L_{II}| \exists \mathbf{p} \in \mathbf{P} : \mathbf{p} \xrightarrow{|\mathbf{x}|}  \cup \{\delta \mid \exists \mathbf{p} \in \mathbf{P} : \mathbf{p} \xrightarrow{\delta} \mathbf{p} \}
```

```
i uioco s = def \forall \sigma \in Utraces(s): out(i after <math>\sigma) \subseteq out(s after \sigma)
```

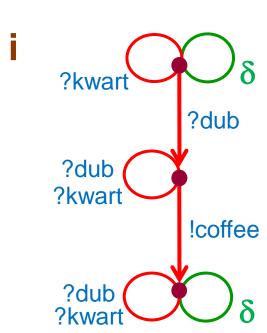
s is a Labelled Transition System

i is (assumed to be) an input-enabled LTS

Intuition:

- i uioco-conforms to s, iff
- if i produces output x after U-trace σ,
 then s can produce x after σ
- if i cannot produce any output after U-trace σ , then s cannot produce any output after σ (called *quiescence* δ)

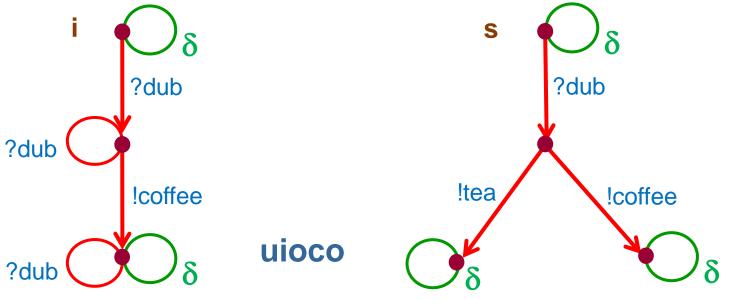
```
i uioco s =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)
out(i after \varepsilon) = \{\delta\}
```



```
{!coffee}
out(i after ?dub)
                                {!coffee}
out(i after ?dub ?dub)
                               \{\delta\}
out(i after ?dub!coffee) =
                                {δ}
out(i after ?kwart)
out(i after !coffee)
                                 Ø
out(i after ?dub !tea)
                                δ
out (i after \delta)
out ( i after \delta \delta?dub)
                                {!coffee}
```

```
i uioco s = def \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)
                                         uioco
                  = \{\delta\}
out (i after \varepsilon)
                                               out (s after \varepsilon)
                                                                           = \{\delta\}
out(i after ?dub) = {!coffee} out(s after ?dub)
                                                                           = {!coffee }
out (i after ?dub !coffee) = \{\delta\}
                                               out (s after ?dub !coffee) = \{\delta\}
```

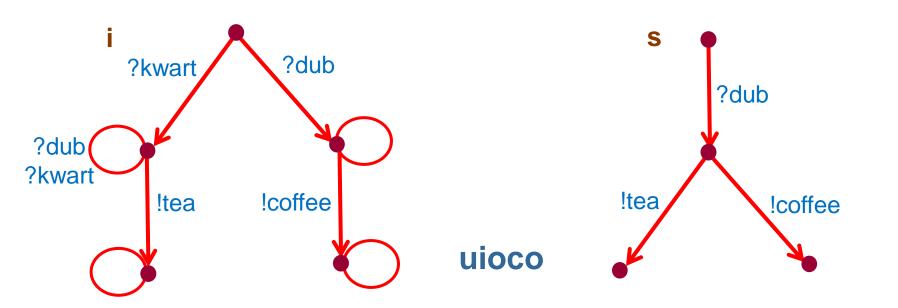
i uioco s $=_{def} \forall \sigma \in Utraces(s): out(i after <math>\sigma) \subseteq out(s after \sigma)$



i uioco s = def $\forall \sigma \in Utraces(s)$: out (i after σ) \subseteq out (s after σ)



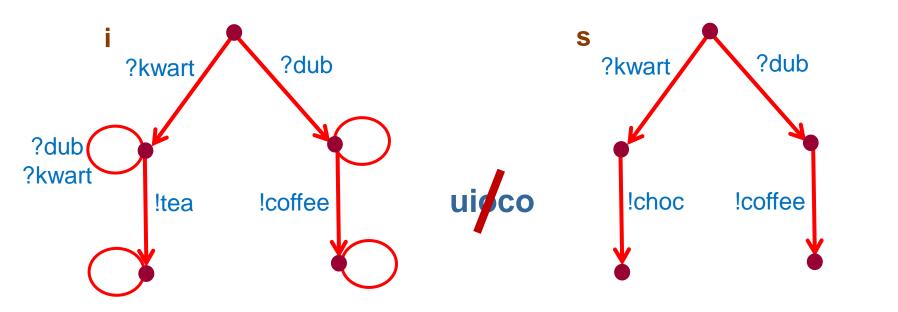
i uioco s = def $\forall \sigma \in Utraces(s)$: out(i after σ) \subseteq out(s after σ)



```
out(i 	ext{ after } ?dub) = \{ !coffee \} out(s 	ext{ after } ?dub) = \{ !coffee, !tea \} out(i 	ext{ after } ?kwart) = \{ !tea \} out(s 	ext{ after } ?kwart) = \emptyset
```

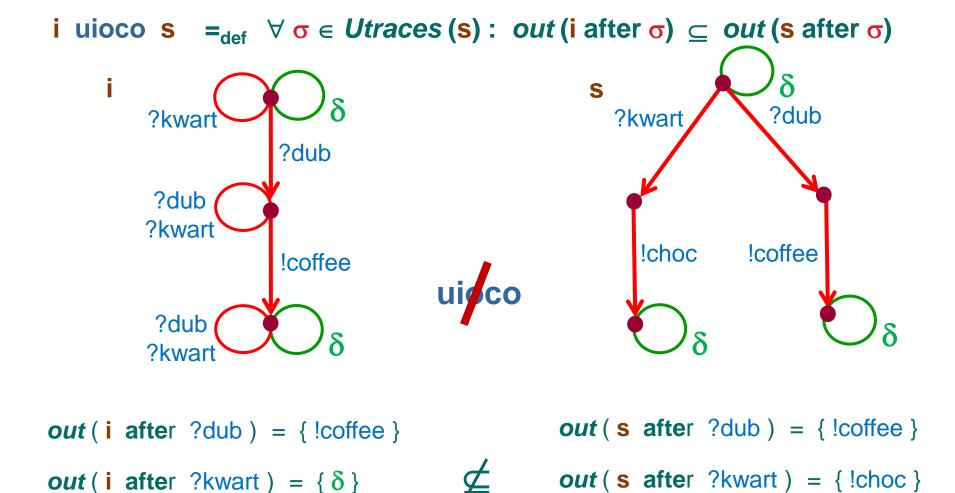
but ?kwart ∉ Utraces (s)

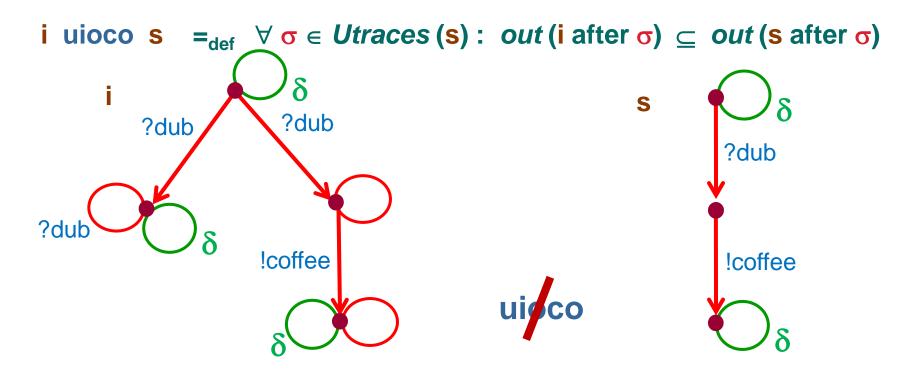
i uioco s $=_{def} \forall \sigma \in Utraces(s): out(i after <math>\sigma) \subseteq out(s after \sigma)$





out(s after ?dub) = {!coffee}
out(s after ?kwart) = {!choc}



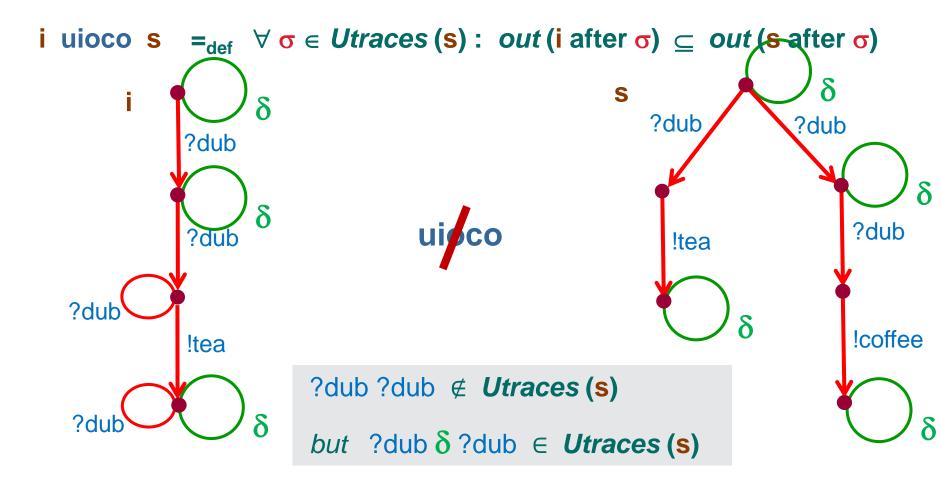


```
out(i after ?dub) = \{!coffee, \delta\}   \not\subseteq out(s after ?dub) = \{!coffee\}
```

i uioco s $=_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$? dub
? dub $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$ $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$ $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$ $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$ $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$ $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$ $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$ $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$ $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$ $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$ $\delta =_{def} \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)$

 $out(i after ?dub) = \{!coffee, \delta\}$ $out(s after ?dub) = \{!coffee, \delta\}$

```
i uioco s = def \forall \sigma \in Utraces(s): out (i after \sigma) \subseteq out (s after \sigma)
                                                                               ?dub
                                       uioco
               ?dub
                                                                            !coffee
               !tea
                      ?dub ?dub \notin Utraces (s) = { \epsilon, ?dub, \delta, ... }
                      because s after ?dub refuses { ?dub }
```



 $out(i after ?dub \delta ?dub) = \{ !tea \} \not = out(s after ?dub \delta ?dub) = \{ !coffee \}$

i uioco s = def $\forall \sigma \in Utraces(s)$: out (i after σ) \subseteq out (s after σ) ?dub ?dub uioco ?dub !tea ?dub !coffee !tea ?dub ?dub ∉ Utraces (s) but $?dub \delta ?dub \in Utraces (s)$

 $out(i after ?dub \delta ?dub) = \emptyset \subseteq out(s after ?dub \delta ?dub) = \{!coffee\}$

```
i uioco s = def \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)
                                                             s is a Labelled Transition System
                                               i is (assumed to be) an input-enabled LTS
   p \xrightarrow{\delta} p \iff \forall !x \in L_U \cup \{\tau\}. p \xrightarrow{!x} \iff p \text{ refuses } L_U
    Straces (s) = { \sigma \in (L \cup \{\delta\})^* \mid s \stackrel{\sigma}{\Longrightarrow} \}
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               \forall \sigma_1 ? b \sigma_2 = \sigma : not(s after \sigma_1 refuses {?b})
   out(P) = { |\mathbf{x} \in L_{II}| \exists \mathbf{p} \in \mathbf{P} : \mathbf{p} \xrightarrow{|\mathbf{x}|}  \cup \{\delta \mid \exists \mathbf{p} \in \mathbf{P} : \mathbf{p} \xrightarrow{\delta} \mathbf{p} \}
```

```
i uioco s = def \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)
                                                                              ?dub
                                       uioco
               ?dub
                                                                           !coffee
               !tea
                      ?dub ?dub \notin Utraces (s) = { \epsilon, ?dub, \delta, ... }
                      because s after ?dub refuses { ?dub }
```

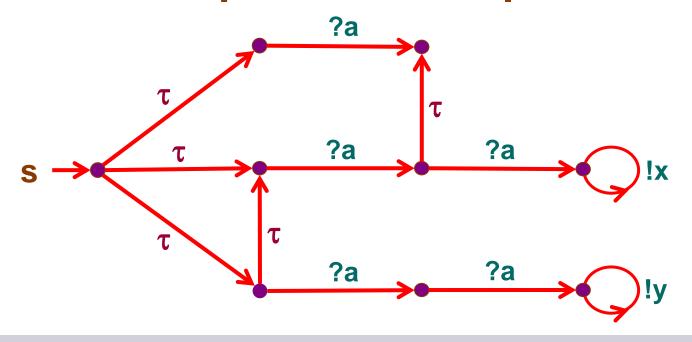
Alternative Implementation Relation ioco

i ioco s = def $\forall \sigma \in Straces(s)$: out (i after σ) \subseteq out (s after σ) ?dub ?dub !coffee !tea ?dub ?dub ∈ Straces (s) Straces (s) = { $\sigma \in (L \cup \{\delta\})^* \mid s$

```
i uioco s = def \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)
      i ioco s = def \forall \sigma \in Straces(s) : out(i after \sigma) \subseteq out(s after \sigma)
                        Straces (s) = { \sigma \in (L \cup \{\delta\})^* \mid s \xrightarrow{\sigma} \}
?dime
             ?dime
                         Utraces(s) = \{ \sigma \in Straces(s) \mid
                              \forall \sigma_1 ? b \sigma_2 = \sigma : not(s after \sigma_1 refuses {?b})
   ?dime
                       ?dime ?dime \in Straces (s)
                                                                          ioco ⊂ uioco
                       ?dime ?dime ∉ Utraces (s)
```

```
i uioco s = def \forall \sigma \in Utraces(s): out(i after \sigma) \subseteq out(s after \sigma)
i ioco s = def \forall \sigma \in Straces(s) : out(i after \sigma) \subseteq out(s after \sigma)
                                   i ioco s
                                   i uioco s
                                                                            ?b
```

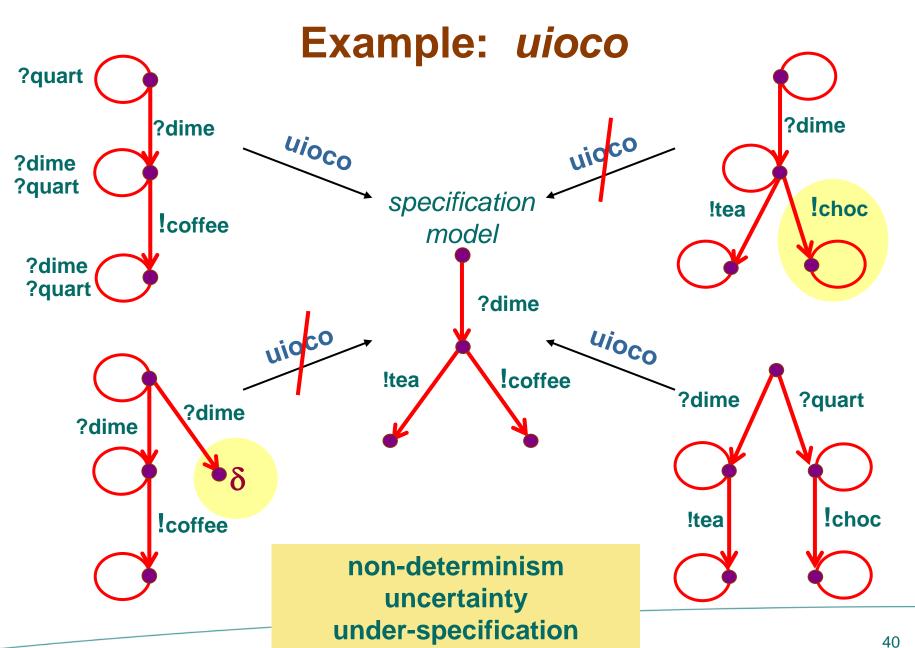
A non-ioco-implementable specification



There is **no** implementation i; what would be **out** (i after δ ?a δ ?a) ?

let $out(i \text{ after } \delta?a\delta?a) = X \text{ then}$

- **X** ≠ Ø
- $X \subseteq out(i \text{ after } \delta?a?a) \subseteq out(s \text{ after } \delta?a?a) = \{ !x \}$
- $X \subseteq out(i \text{ after } ?a\delta?a) \subseteq out(s \text{ after } ?a\delta?a) = \{ !y \}$

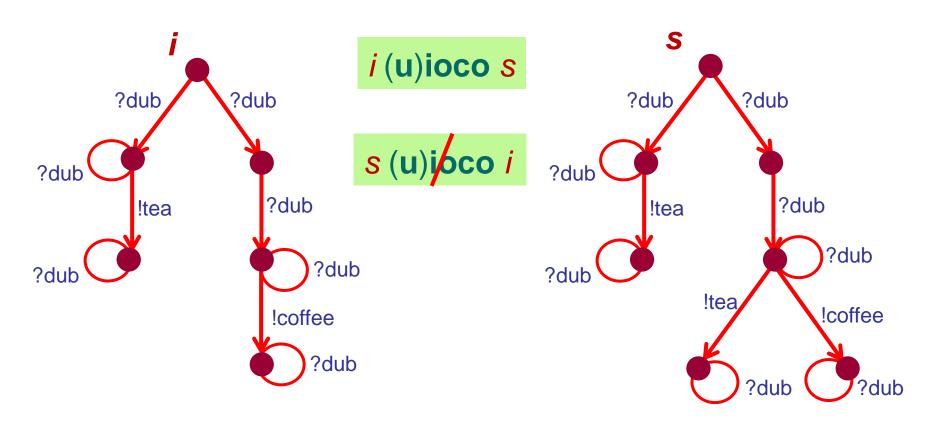


Example: ioco

```
i u/ioco s = _{def}

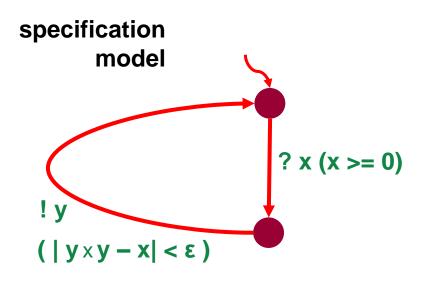
\forall \sigma \in U/Straces (s) :

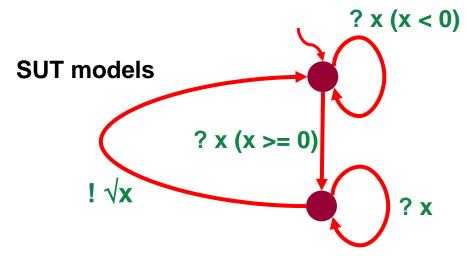
out (i after \sigma) \subseteq out (s after \sigma)
```



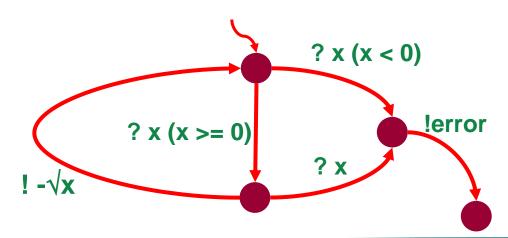
```
out(i 	ext{ after ?dub.?dub}) = out(s 	ext{ after ?dub.?dub}) = \{ !tea, !coffee \}  out(i 	ext{ after ?dub.} \delta.?dub) = \{ !coffee \} \neq out(s 	ext{ after ?dub.} \delta.?dub) = \{ !tea, !coffee \}
```

MBT: Nondeterminism, Underspecification

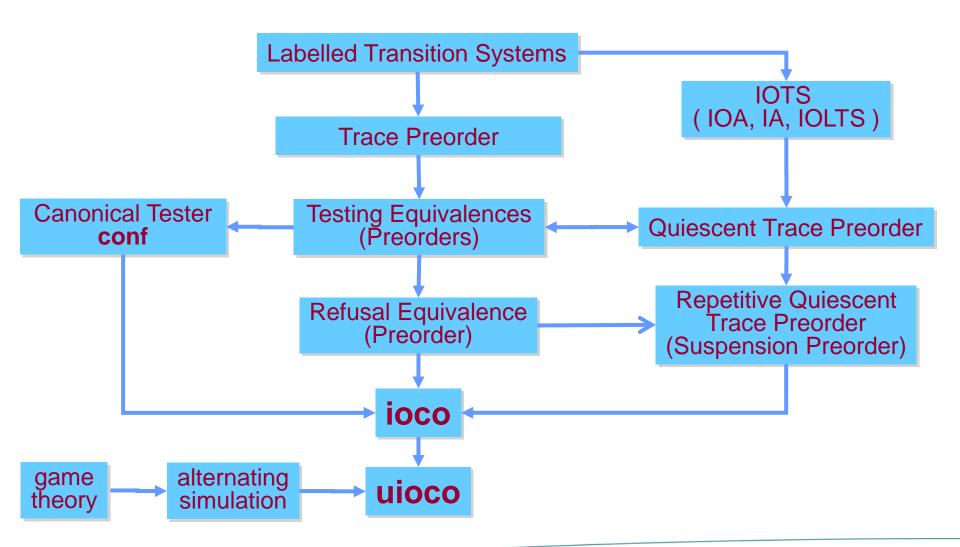




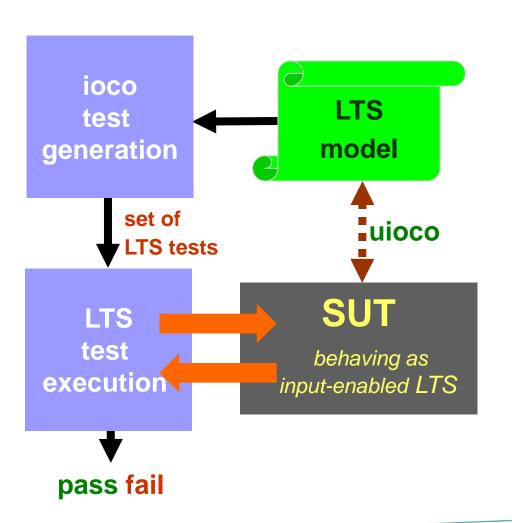
- non-determinism
- under-specification
- specification of properties
 rather than construction



Genealogy of u/ioco



MBT: Labelled Transitions Systems



MBT with LTS topics:

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- implementation model
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- sound & exhaustive