



UNIVERSITÀ
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Department of
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Automated Reasoning and **Formal Verification**

Laboratory 8

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Outline

1. Model Properties

Invariants

LTL

CTL

2. Fairness Constraints

3. Modelling a Program in nuXmv

4. Examples

5. Homework

Specifying Properties

- ▶ directly in the module

```
LTLSPEC G (req -> F sum = op1 + op2);
```

- ▶ or via nuXmv interactive shell

```
nuXmv > check_ltlspec -p "G (req -> F sum = op1 + op2)"
```

- ▶ `show_property` lists all properties specified (in the module or via `add_property`):
- ▶ properties can be verified one at a time using its **database index**

```
nuXmv > show_property
```

```
**** PROPERTY LIST [ Type, Status, Counter-example Number, Name ] ****
```

```
----- PROPERTY LIST -----  
000 :out < 2
```

[Invar	False	1	N/A]
--------	-------	---	------

```
001 : G ( F out = 1)
```

[LTL	Unchecked	N/A	N/A]
------	-----------	-----	------

```
nuXmv > check_ltlspec -n 1
```

Property verification:

- ▶ Each property is verified independently
- ▶ The result is either “TRUE” or “FALSE + counterexample”

Property types

Different kinds of properties are supported:

- Invariants** : properties on every reachable state;
- LTL** : properties on the computation paths;
- CTL** : properties on the computation tree.

- ▶ Invariant properties are specified via the keyword `INVARSPEC`:
`INVARSPEC <simple_expression>;`
- ▶ Invariants are checked via the `check_invar` command

Remark

When checking invariants, all the **fairness conditions** of the model are **ignored**.

Example: Modulo 4 Counter with Reset [1/2]

```
MODULE main
```

```
VAR  b0 : boolean; b1 : boolean;
     reset : boolean;
```

```
ASSIGN
```

```
  init(b0) := FALSE;
```

```
  next(b0) := case reset : FALSE;
                  !reset : !b0;
```

```
          esac;
```

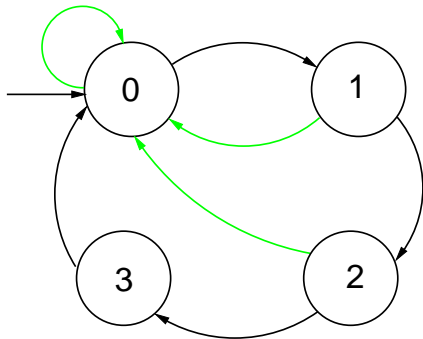
```
  init(b1) := FALSE;
```

```
  next(b1) := case reset : FALSE;
                  TRUE  : ((!b0 & b1) |
                           ( b0 & !b1));
```

```
          esac;
```

```
DEFINE out := toint(b0) + 2 * toint(b1);
```

```
INVARSPEC out < 2;
```



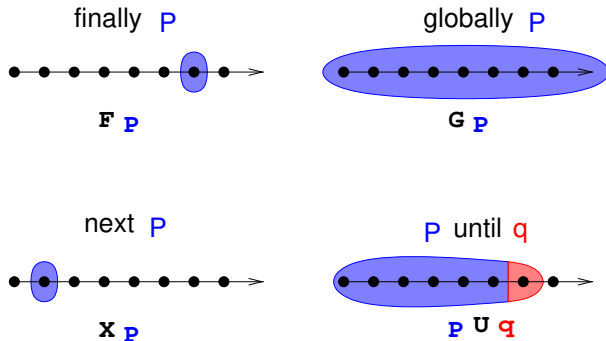
Example: Modulo 4 Counter with Reset [2/2]

► The invariant is **false**

```
nuXmv > read_model -i counter4reset.smv
nuXmv > go; check_invar
-- invariant out < 2 is false
...
-> State: 1.1 <-
  b0 = FALSE
  b1 = FALSE
  reset = FALSE
  out = 0
-> State: 1.2 <-
  b0 = TRUE
  out = 1
-> State: 1.3 <-
  b0 = FALSE
  b1 = TRUE
  out = 2
```

- ▶ LTL properties are specified via the keyword **LTLSPEC**:

LTLSPEC <ltl_expression>;



- ▶ LTL properties are checked via the **check_ltlspec** command



Specifications Examples:

- ▶ A state in which $\text{out}=3$ is eventually reached

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 $\text{LTLSPEC } F \text{ out} = 3;$
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LTLSPEC $F \text{ out} = 3;$

- ▶ Condition $\text{out}=0$ holds until reset becomes false

LTLSPEC $(\text{out} = 0) \ U \ (!\text{reset});$ -- False: reset can be true forever

LTLSPEC $(!\text{reset}) \ V \ (\text{out} = 0);$ -- True (V stands for "release")

- ▶ Every time a state with $\text{out}=2$ is reached, a state with $\text{out}=3$ is reached afterward

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LTLSPEC $G (\text{out} = 2 \rightarrow F \text{ out} = 3);$

All the previous specifications are false:

```
nuXmv > check_ltlspec
-- specification  F out = 3 is false ...
-- loop starts here --
-> State 1.1 <-
    b0 = FALSE
    b1 = FALSE
    reset = TRUE
    out = 0
-> State 1.2 <-
-- specification (out = 0 U (!reset)) is false ...
-- loop starts here --
-> State 2.1 <-
    b0 = FALSE
    b1 = FALSE
    reset = TRUE
    out = 0
-> State 2.2 <-
-- specification  G (out = 2 ->  F out = 3) is false ...
```

Q: Why?

CTL specifications

- CTL properties are specified via the keyword **CTLSPEC**:

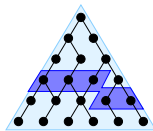
CTLSPEC <ctl_expression>;

finally **P**

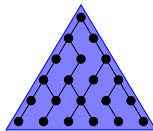
globally **P**

next **P**

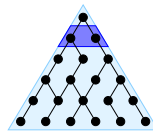
P until **q**



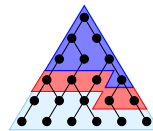
AF P



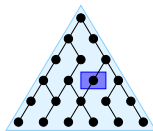
AG P



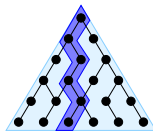
AX P



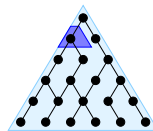
A[P U q]



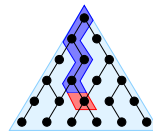
EF P



EG P



EX P



E[P U q]

- CTL properties are checked via the **check_ctlspec** command

1. Model Properties



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- ▶ It is always possible to reach a state in which $\text{out}=3$

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- ▶ It is inevitable that $\text{out}=3$ is eventually reached
`CTLSPEC AF out = 3;`
- ▶ It is always possible to reach a state in which $\text{out}=3$
`CTLSPEC AG EF out = 3;`
- ▶ Every time a state with $\text{out}=2$ is reached, a state with $\text{out}=3$ is reached afterward

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- ▶ The reset operation is correct

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- ▶ Every time a state with $\text{out}=2$ is reached, a state with $\text{out}=3$ is reached afterward
`CTLSPEC AG (out = 2 -> AF out = 3);`
- ▶ The reset operation is correct
`CTLSPEC AG (reset -> AX out = 0);`



Outline

1. Model Properties
2. Fairness Constraints
3. Modelling a Program in nuXmv
4. Examples
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The need for Fairness Constraints

The specification $F \text{ out} = 3$; is not verified

- ▶ On the path where **reset** is always **1**, the system loops on a state where **out** = **0**:
reset = **TRUE**, **TRUE**, **TRUE**, **TRUE**, **TRUE**, ...
out = 0, 0, 0, 0, 0, ...

Similar considerations for other properties:

- ▶ $F \text{ out} = 1$;
- ▶ $F \text{ out} = 2$;
- ▶ $G (\text{out} = 2 \rightarrow F \text{ out} = 3)$;
- ▶ ...



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out = 0, 0, 0, 0, 0, ...

Similar considerations for other properties:

- ▶ $F \text{ out} = 1$;
- ▶ $F \text{ out} = 2$;
- ▶ $G (\text{out} = 2 \rightarrow F \text{ out} = 3)$;
- ▶ ...

Fairness

It would be **fair** to consider only paths in which the **counter** is **not reset** with such a high frequency, so as to hinder its desired functionality.



Fairness Constraints

nuXmv supports both **justice** and **compassion** fairness constraints:

Fairness: **JUSTICE** $\langle p \rangle$ consider only the executions that satisfy p **infinitely often**

Strong Fairness: **COMPASSION** $(\langle p \rangle, \langle q \rangle)$ consider only those executions that either satisfy p **finitely often** or satisfy q **infinitely often**
(i.e., p is true infinitely often $\implies q$ is true infinitely often)

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Remarks

- ▶ For **verification**, properties must hold only on **fair paths**
- ▶ When checking invariants, all the fairness conditions are **ignored**.
- ▶ Currently, compassion constraints have some limitations, since they are supported only for BDD-based LTL model checking.

Example: Modulo 4 Counter with Reset

Add the following fairness constraint to the model:

JUSTICE out = 3;

We consider only paths in which the counter reaches value 3 infinitely often

All the properties are now verified:

```
nuXmv > reset
nuXmv > read_model -i counter4reset.smv
nuXmv > go
nuXmv > check_ltlspec
-- specification F out = 1 is true
-- specification G (out = 2 -> F out = 3) is true
-- specification G (reset -> F out = 0) is true
```



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Q: Given the following **code**, how can we **model** and **verify** it with nuXmv?

```
def gcd(a: int, b: int) -> int:
  while a != b:
    if a > b:
      a = a-b
    else:
      b = b-a
  return a; # at this point: GCD=a=b
```



Main idea

- ▶ We define a program counter `pc` that stores the current status of the execution (i.e., the line we reached).
- ▶ According to cycle and the conditional instructions, the program counter and the variables (when required) will change.

Step 1: label the **entry point** and the **exit point** of every block with the line number

```
def gcd(a: int, b: int) -> int:
  while a != b: # 11
    if a > b:    # 12
      a = a-b   # 13
    else:
      b = b-a   # 14
  return a;     # 15 -- at this point: GCD=a=b
```

Example: model programs in nuXmv [3/4]

Step 2: encode the transition system with the assign style

MODULE main

VAR

a: 0..100; b: 0..100;

pc: {11, 12, 13, 14, 15};

ASSIGN

init(pc) := 11;

next(pc) := case

pc = 11 & a != b : 12;

pc = 11 & a = b : 15;

pc = 12 & a > b : 13;

pc = 12 & a <= b : 14;

pc = 13 | pc = 14 : 11;

pc = 15 : 15;

esac;

next(a) := case

pc = 13 & a > b : a - b;

TRUE : a;

esac;

next(b) := case

pc = 14 & b >= a : b-a;

TRUE : b;

esac;

Step 2: (alternative): use the constraint style

```
MODULE main
```

```
VAR
```

```
  a : 0..100;  b : 0..100;  pc : {11, 12, 13, 14, 15};
```

```
INIT pc = 11
```

```
TRANS pc = 11 -> (((a != b & next(pc) = 12) | (a = b & next(pc) = 15)) &  
                  next(a) = a & next(b) = b);
```

```
TRANS pc = 12 -> (((a > b & next(pc) = 13) | (a < b & next(pc) = 14)) &  
                  next(a) = a & next(b) = b);
```

```
TRANS pc = 13 -> (next(pc) = 11 & next(a) = (a - b) & next(b) = b);
```

```
TRANS pc = 14 -> (next(pc) = 11 & next(b) = (b - a) & next(a) = a);
```

```
TRANS pc = 15 -> (next(pc) = 15 & next(a) = a & next(b) = b);
```


Step 3: check the properties of the program

- ▶ Let's check if, given $a = 16$ and $b = 12$, then we will eventually get as a result 4.
LTLSPEC $(a = 16 \ \& \ b = 12) \rightarrow \mathbf{F} \ (a = 4 \ \& \ b = 4);$
- ▶ Let's check if both numbers will never reach negative values:
INVARSPEC $a > 0 \ \& \ b > 0;$

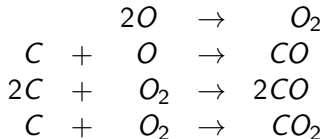


Outline

1. Model Properties
2. Fairness Constraints
3. Modelling a Program in nuXmv
4. Examples
 - Chemical reactions
 - The snail dungeon
5. Homework

Exercise 8.1

Assume the following chemical reactions hold:



Given 6 carbon atoms and 6 oxygen atoms, is there any way for the contents of this reaction vessel to progress to a state where it contains three molecules of CO₂?
Model the contents of the reaction vessel in nuXmv.



Science Modelling: variables [1/2]

- ▶ We can store the number of atoms and molecules in the current iteration with bounded integers.
- ▶ An enum variable can be used to define what reaction should be considered in the next step, ensuring non-determinism when necessary.

MODULE main

VAR

```
o          : 0..32;
o2         : 0..32;
c          : 0..32;
co         : 0..32;
co2        : 0..32;
r : {r1, r2, r3, r4, none};
```

ASSIGN

```
init(o)      := 6;
init(c)      := 6;
init(co)     := 0;
init(co2)    := 0;
init(o2)     := 0;
init(r) := none;
```

Transitions to define the next reaction that will take place on the next step.

TRANS

```
(next(o) < 2) -> (next(reaction) != r1);
```

TRANS

```
(next(o) < 1 | next(c) < 1) -> (next(reaction) != r2);
```

TRANS

```
(next(o2) < 1 | next(c) < 2) -> (next(reaction) != r3);
```

TRANS

```
(next(o2) < 1 | next(c) < 1) -> (next(reaction) != r4);
```

Science Modelling: transitions [2/2]

Transitions to define the new values for each molecule after a reaction took place.

TRANS

r=none \rightarrow (next(o) = o & next(o2) = o2 & next(c) = c &
next(co) = co & next(co2) = co2)

TRANS

r=r1 \rightarrow (next(o) = o - 2 & next(o2) = o2 + 1 & next(c) = c &
next(co) = co & next(co2) = co2);

TRANS

r=r2 \rightarrow (next(o) = o - 1 & next(o2) = o2 & next(c) = c - 1 &
next(co) = co + 1 & next(co2) = co2);

TRANS

r=r3 \rightarrow (next(o) = o & next(o2) = o2 - 1 & next(c) = c - 1 &
next(co) = co + 2 & next(co2) = co2);

TRANS

r=r4 \rightarrow (next(o) = o & next(o2) = o2 - 1 & next(c) = c - 1 &
next(co) = co & next(co2) = co2 + 1);

- ▶ If we are interested in knowing if there is a path that generates 3 CO_2 molecules, LTL apparently seems ineffective...
- ▶ ... but we can use it to search a valid **counter-example** corresponding to the desired execution.
- ▶ We try to verify that the number of CO_2 molecules **does not** reach 3 in any path:
LTLSPEC G (co2!=3)
- ▶ If the property is not satisfied, we get a sequence of steps reaching a state where $co2=3$.

Exercise 8.2

You want to simulate the gameplay of “**The Snail Dungeon**”:

- ▶ You have a path with 10 cells, with 2 good and 2 bad teleports. Use a variable **turn** whose value could be **{DICE, GOOD, BAD}**, and a variable **steps** counting how many times a dice has been thrown. Once set, the teleport positions remain fixed.
- ▶ Each turn you **throw a 3-valued dice** and **move forward** the designated number of cells. If the arrival cell is empty, you move on to the next turn. Notice that the dice **cannot get the same number from two consecutive throws**.
- ▶ If you get into a *good* teleport, the next value of turn will be **GOOD** and you will **move onward by 2 cells without increasing** steps.
- ▶ If you get into a *bad* teleport, the next value of turn will be **BAD** and you will **move back by 2 cells without increasing** steps.

Encode the game in nuXmv and find a way to reach position 10 with less than 3 steps.

The Snail Dungeon: variables

- ▶ The text already indicates some variables. Notice that with the default **BDD-based engine** we can only use **bounded integers**.

VAR

```
pos    : 0..10;  
turn   : {DICE, GOOD, BAD};  
steps  : 0..100;  
dice   : 1..3;
```

- ▶ We can use two arrays of two integers in 1..10 to store the positions of teleports.

```
good : array 0..1 of 1..10; bad  : array 0..1 of 1..10;
```

- ▶ These arrays are **fixed**. **Instead of adding**

```
ASSIGN next(good[0]) := good[0]; next(good[1]) := good[1]; ...
```

we can **declare them as FROZENVAR**:

```
FROZENVAR good : array 0..1 of 1..10;  
          bad  : array 0..1 of 1..10;
```

The Snail Dungeon: initial values

- Initialize the game:

ASSIGN

```
init(pos)    := 1;  
init(steps)  := 0;  
init(turn)   := DICE;
```

- And constrain the teleport positions:

INIT

```
good[0] != good[1] & bad[0] != bad[1] &  
good[0] != bad[0]  & good[0] != bad[1] &  
good[1] != bad[0]  & good[1] != bad[1];
```



The Snail Dungeon: transitions [1/2]

We encode the moves' logic:

ASSIGN

```
next(turn) := case
  next(pos) = good[0] | next(pos) = good[1] : GOOD;
  next(pos) = bad[0]  | next(pos) = bad[1]  : BAD;
  TRUE: DICE;
esac;
next(steps) := case
  turn = DICE : min(steps + 1, 100);
  TRUE      : steps;
esac;
next(pos) := case
  turn = DICE : min(pos + dice, 10);
  turn = GOOD : min(pos + 2, 10);
  turn = BAD  : max(pos - 2, 0);
esac;
```

Notice: we must ensure that bounded integers do not exceed their limits!

4. Examples



The Snail Dungeon: transitions [2/2]

Finally, we encode the dice logic: it cannot give the same value twice in a row.

TRANS

```
(turn = DICE -> next(dice) != dice) &  
(turn != DICE -> next(dice) = dice);
```

The Snail Dungeon: properties

To get a trace where we reach position 10 with less than 3 steps, can negate the property and get a counter-example:

INVARSPEC pos=10 -> steps >= 3;

Trace Type: Counterexample

-> State: 1.1 <-

good[0] = 8

good[1] = 4

bad[0] = 10

bad[1] = 2

pos = 1

turn = DICE

steps = 0

dice = 3

-> State: 1.2 <-

pos = 4

turn = GOOD

steps = 1

dice = 2

-> State: 1.3 <-

pos = 6

turn = DICE

-> State: 1.4 <-

pos = 8

turn = GOOD

steps = 2

dice = 1

-> State: 1.5 <-

pos = 10

turn = BAD



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Homework 8.1: Bubblesort

Implement a transition system which sorts the following input array $\{4, 1, 3, 2, 5\}$ with increasing order. Verify the following properties:

- ▶ there exists no path in which the algorithm ends
- ▶ there exists no path in which the algorithm ends with a sorted array

Bubblesort

You might use the following *bubblesort* code as reference:

```
def bubbleSort(A : list[int]) -> None:
    n = len(A)
    swapped = True
    while swapped:
        swapped = False
        for i in range(1, n):
            if A[i-1] > A[i]: # if this pair is out of order
                # swap and remember something changed
                A[i-1], A[i] = A[i], A[i-1]
                swapped = true
```