



CIS 2101

Priority Queue

Priority Queue

- A set ADT with operations
 1. Insert
 2. Deletemin
- Priority Queue includes operations
Initialize and Makenull
- Application: Process Scheduling in
a time-shared computing system

Priority Queue

* Operations

1. **Insert(x,A)** - adds element x in set A , if the element x is not yet a member of the set.
2. **Deletemin(A)** - removes and **returns** the smallest element in set A if the set is not empty; otherwise the operation fails

Priority Queue Implementations

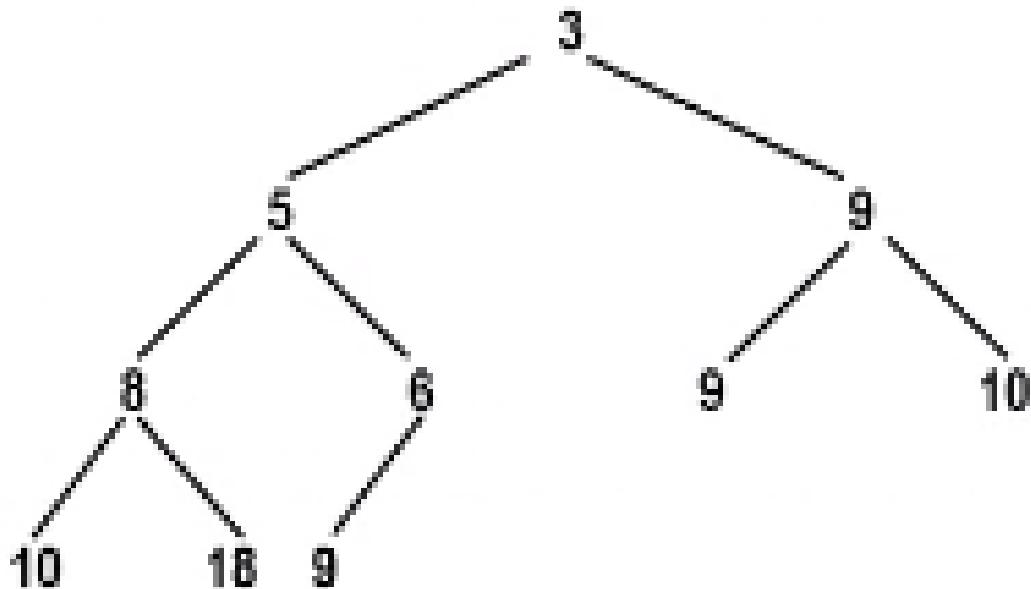
1. Bit-vector
2. Linked list
3. Array
4. Cursor based
5. Partially Ordered Tree (P.O.T.)



Prepared by chisp

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P.O.T. Illustration



Prepared by chesp

Characteristics of P.O.T.

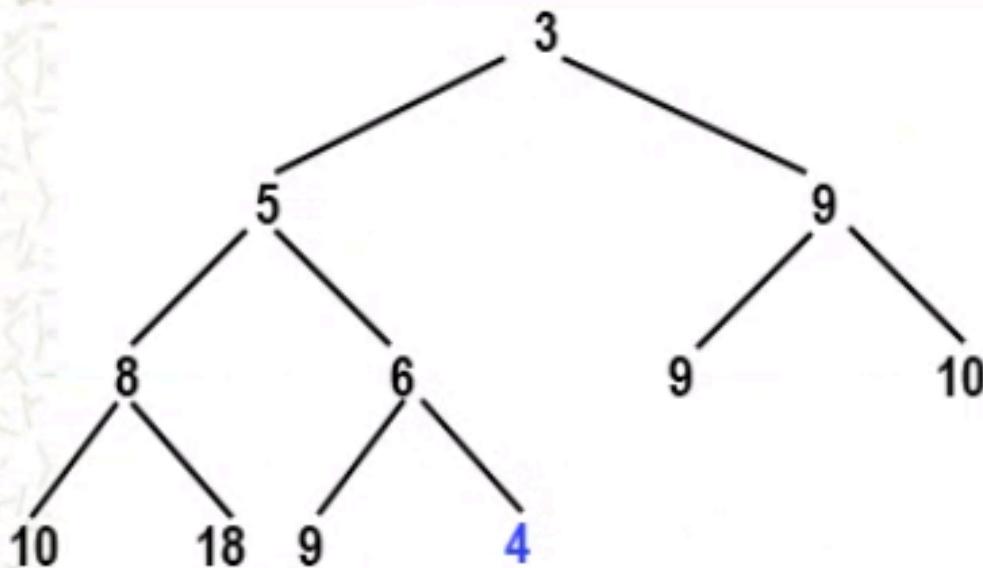
1. **Binary Tree** – a tree in which each node has either a left child, or a right child or both left and right children
2. **Balanced Tree** – a tree in which the height is a minimum for the # of nodes
3. At the lowest level, where the leaves may be missing, we require that all missing leaves are to the right of all the leaves present in the lowest level.
4. **POT Property** – the priority of the parent is less than or equal to the priority of the children



Illustrations of P.O.T. Operations



Insert 4

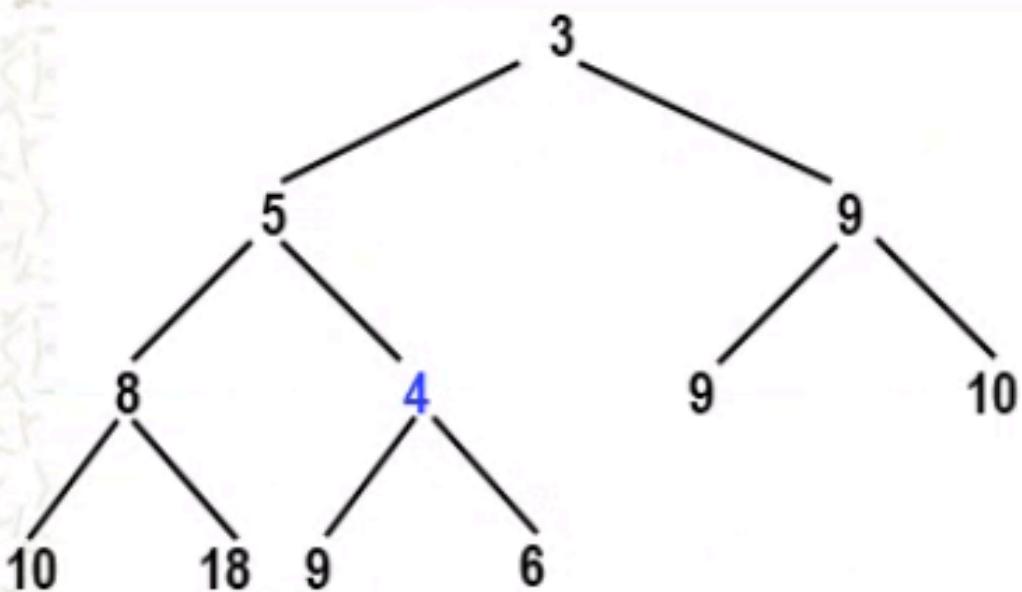


Is P.O.T. property satisfied? **NO!**

What is to be done? **SWAP 6 & 4 !**

Insert 4

(Swap 6 & 4)

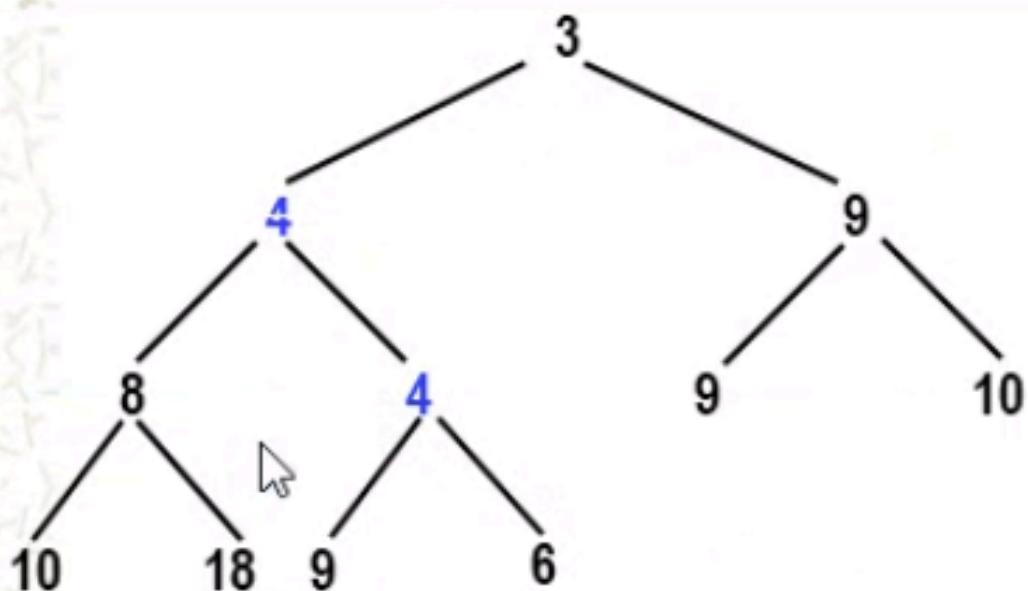


Is P.O.T. property satisfied? **NO!**

What is to be done? **SWAP 5 & 4 !**

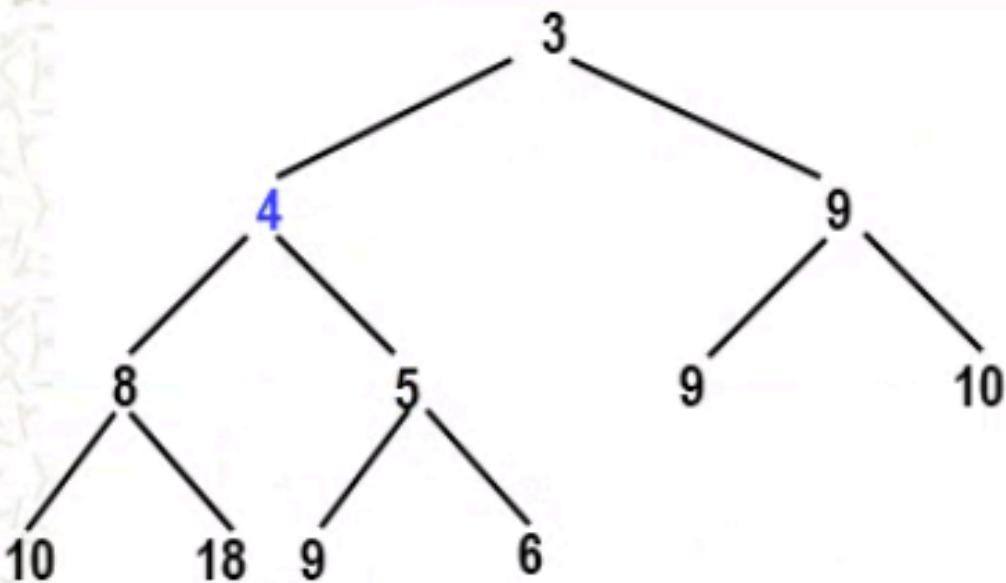
Insert 4

(Swap 5 & 4)



Insert 4

(Swap 5 & 4)



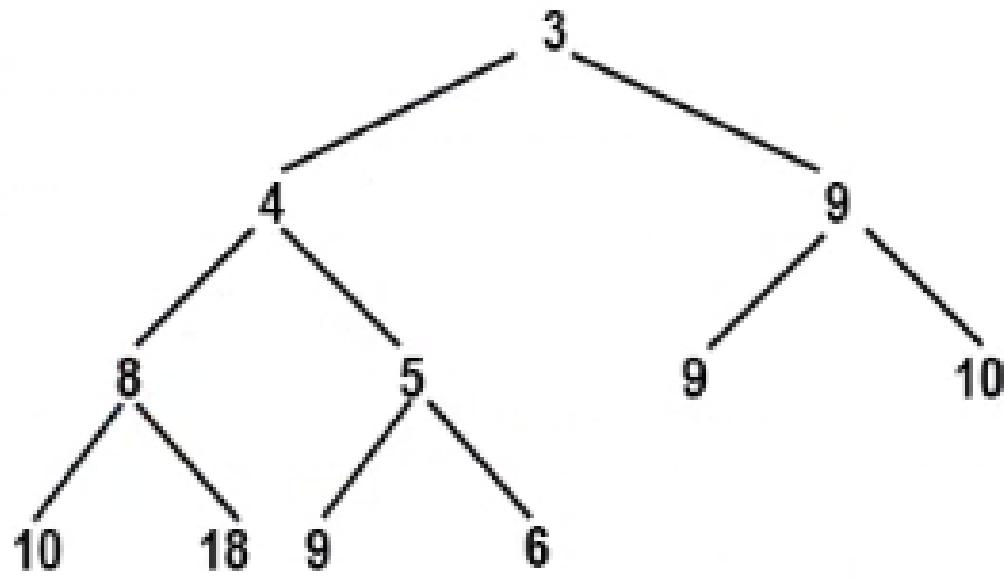
Is P.O.T. property satisfied? **YES!**

Steps in INSERT Operation

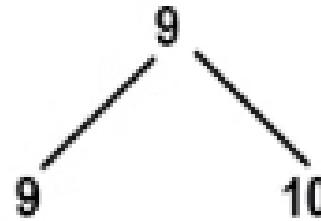
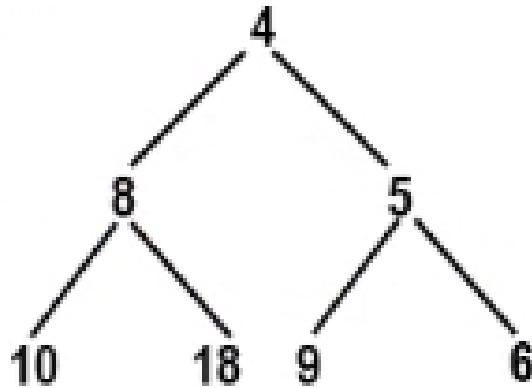
1. Place new element x at
the lowest level to the right of the leaves present
OR
next level if current level is full
2. while (x is not root and POT property is not satisfied)
SWAP(parent, x);



Deletemin



Deletemin



To be returned

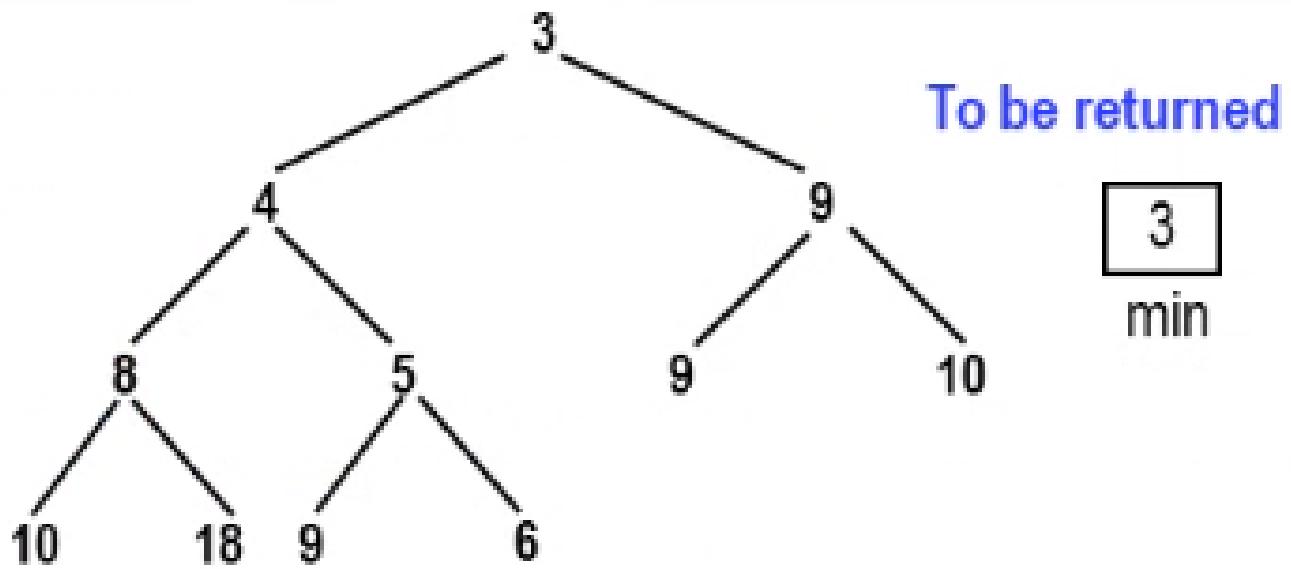
3

min

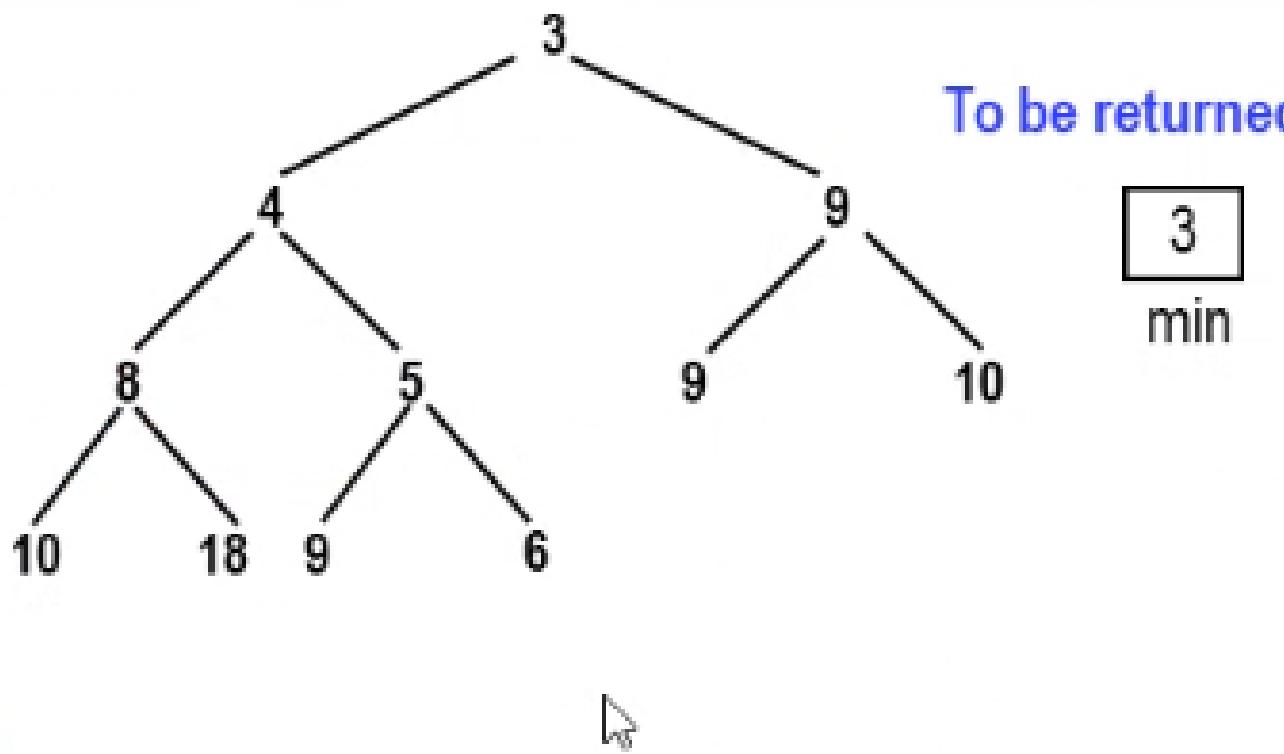
What is the result? A Forest!

Solution? Replace the root node

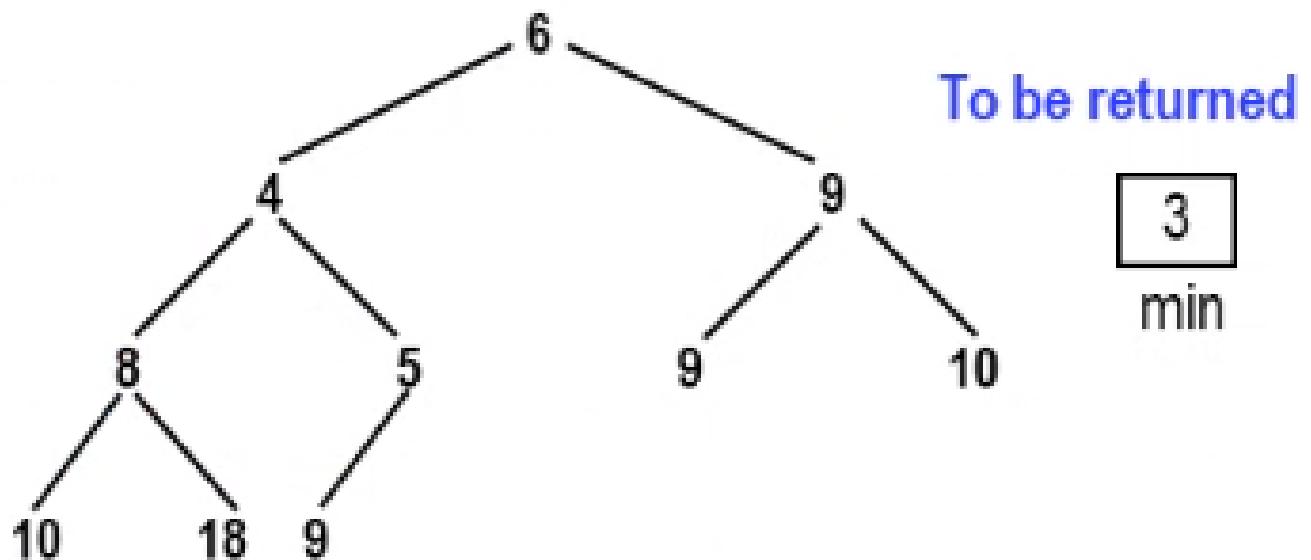
Deletemin



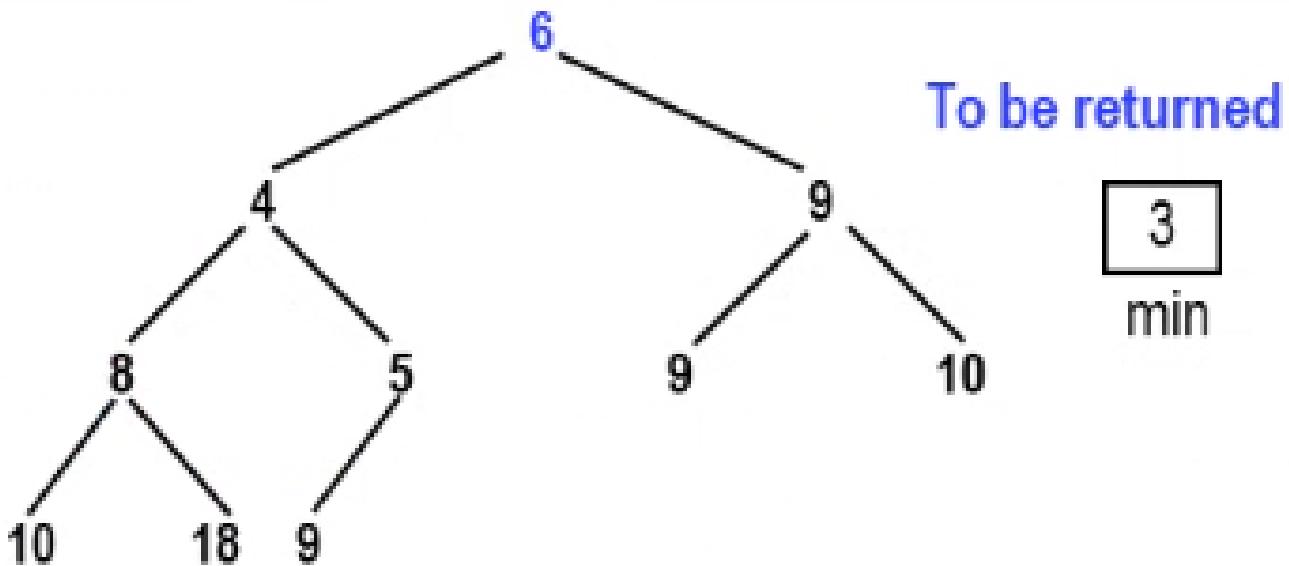
Deletemin



Deletemin



Deletemin



To be returned

3

min

Is P.O.T. property satisfied?

NO!

What is to be done?

Swap 6 with its smallest child!

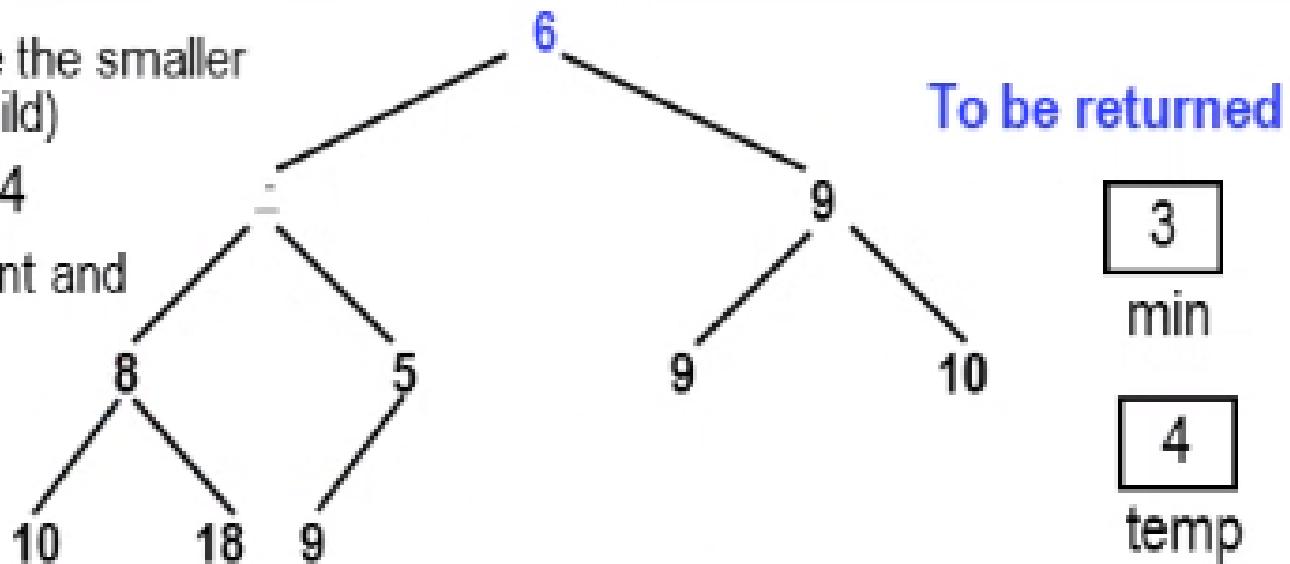
Deletemin

(Swap 4 & 6)

1. Determine the smaller child (SChild)

SChild = 4

2. Swap parent and SChild



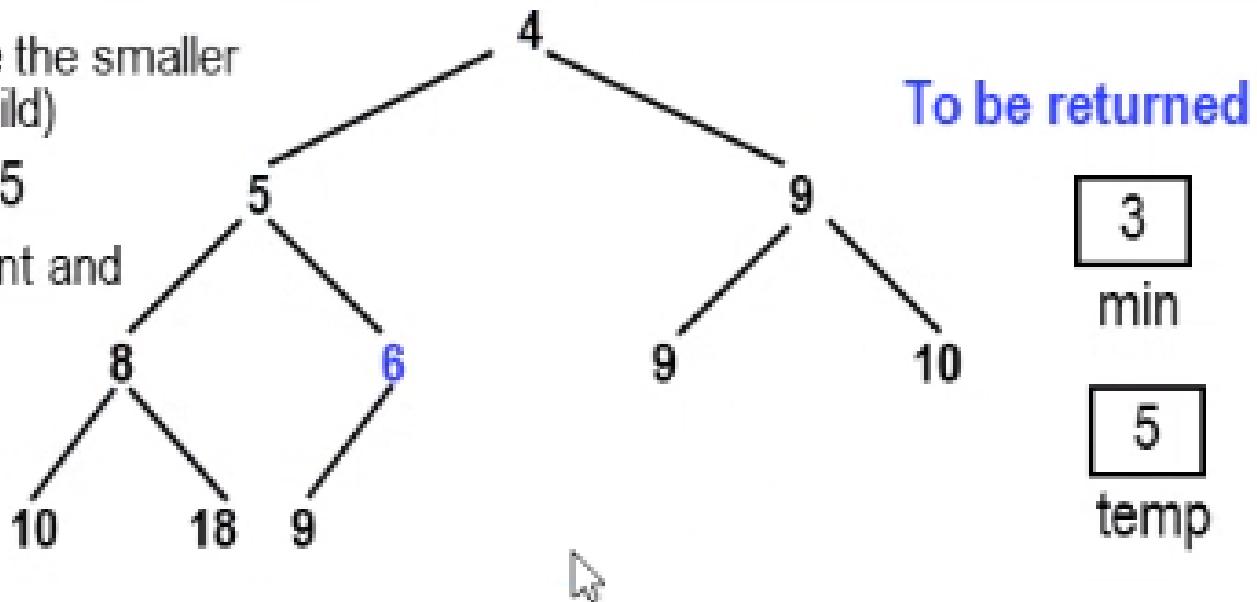
Deletemin

(Swap 5 & 6)

1. Determine the smaller child (SChild)

SChild = 5

2. Swap parent and SChild



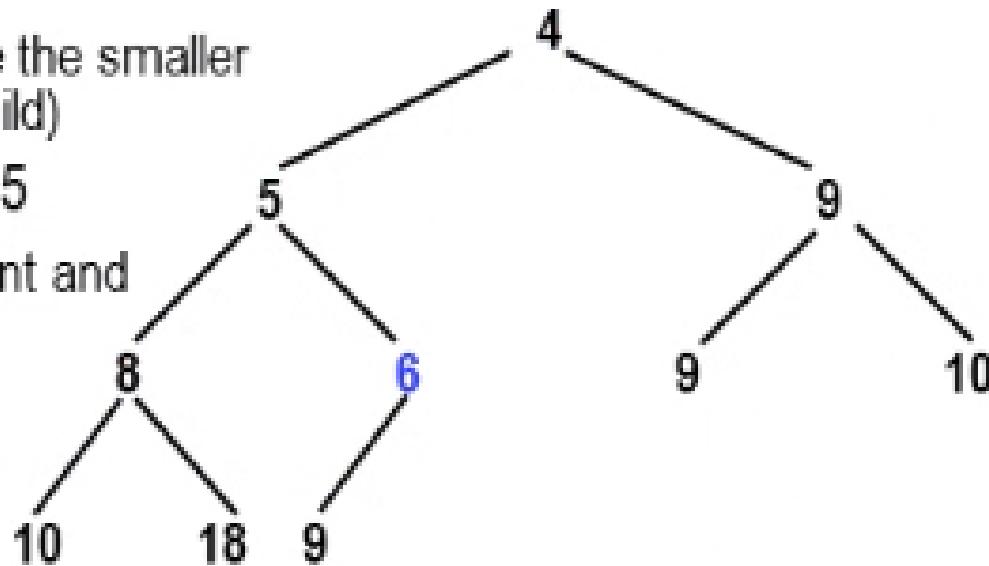
Deletemin

(Swap 5 & 6)

1. Determine the smaller child (SChild)

SChild = 5

2. Swap parent and SChild



Is P.O.T. property satisfied? **YES!**

Return Minimum Priority : **3**

Deletemin operation is **DONE!**

Prepared by chnsp

Steps in Deletemin Operation

1. min = root node
2. Replace **root** with the element **X** found at the lowest level far right
3. While ((**X** is not leaf) and NOT POT) {
 SChild = smaller child of parent **X**;
 SWAP(**X**, SChild);
}
4. return min;

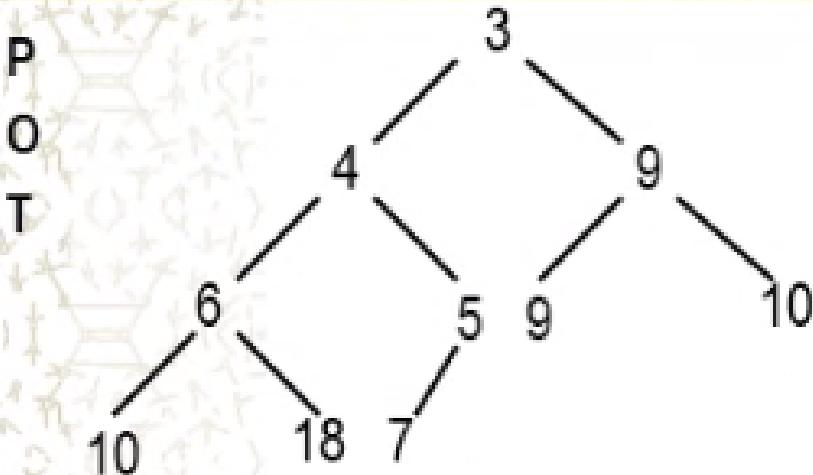


Implementation of P.O.T

Heap

- Array Implementation of P.O.T.

Heap (Data Structure)

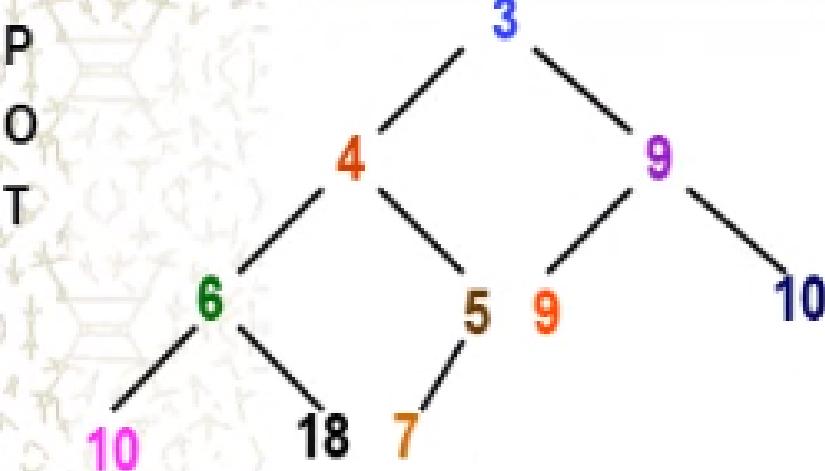


H
E
A
P

Elem	[]	[]	[]	[]	[]	[]	[]	[]	[]	[]	[]	[]	[]		
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
lastNdx	[]														

Heap

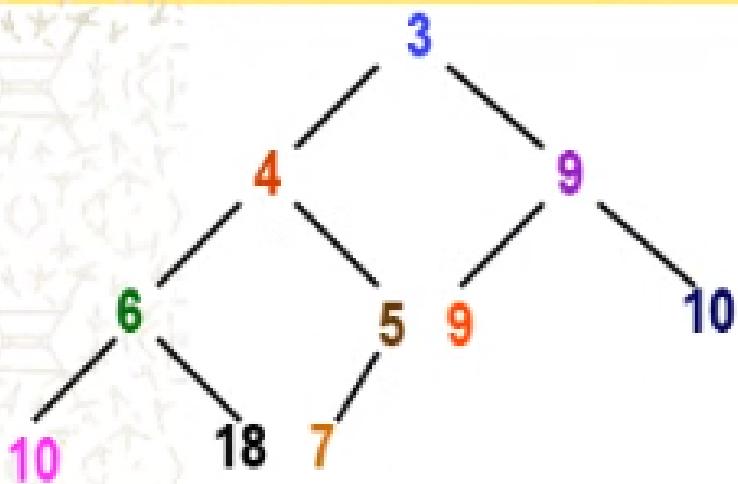
(How elements are stored)



H	E	A	P	Elem	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
				lastNdx															

Heap

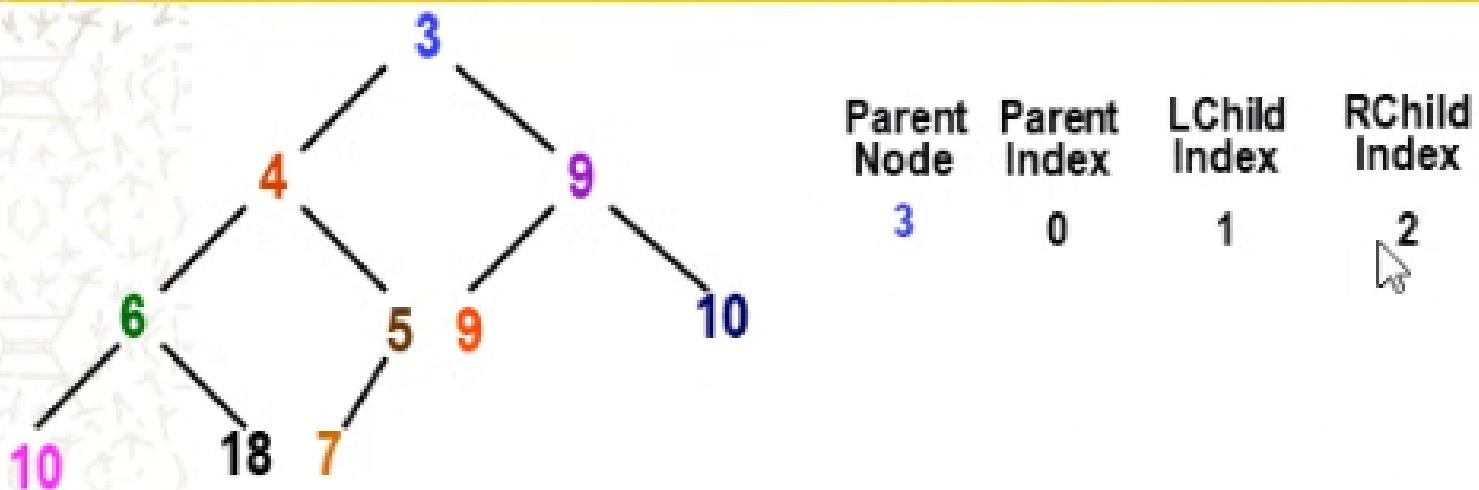
(How elements are stored)



H E A P	Elem	3	4	9	6	5	9	10	10	18	7					
		0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
lastNdx		9														

Heap

(How elements are stored)

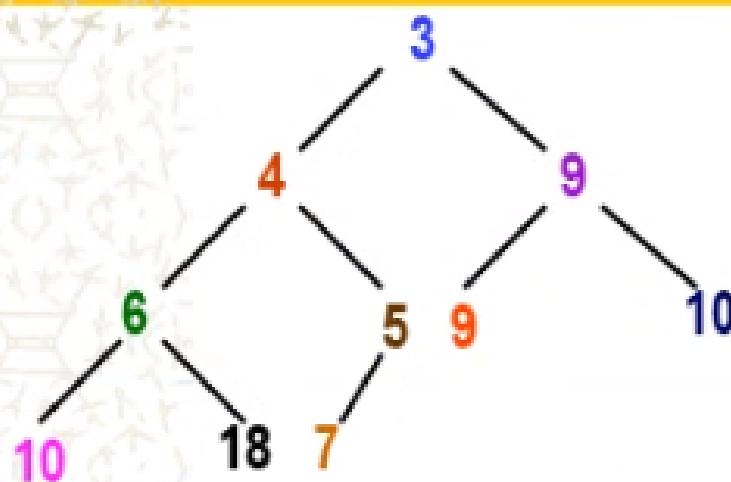


H E A P	Elem	3	4	9	6	5	9	10	18	7						
	lastNdx	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

Heap

(How elements are stored)

P
O
T



Parent Node	Parent Index	LChild Index	RChild Index
3	0	1	2
4	1	3	4
9	2	5	6
6	3	7	8

Can you see a pattern ?

H	E	A	P	Elem	3	4	9	6	5	9	10	10	18	7							
				0	1	2	3	4	5	6	7	8	9	10	11	12	13	14			
				lastNdx	9																

Inserting in a Heap

3

Insert 2 in the Heap

H
E
A
P

Elem	3	4	9	6	5	9	10	10	18	7				
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
last	9													

Insert 2 in the Heap

H E A P	Elem	3	4	9	6	5	9	10	10	18	7	2				
		0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
	last	10														

Insert 2 in the Heap

HEAP	P	C													
Elem	3	4	9	6	5	9	10	10	18	7	2				
last	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

Insert 2 in the Heap

H	E	A	P	Elem												C		
			last	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
				3	4	9	6	5	9	10	10	18	7	2				

Is P.O.T. property satisfied ? NO!

What is to be done ? SWAP 2 & 5 !

Insert 2 in the Heap

(1st Iteration)

H E A P	P	C													
Elem	3	4	9	6	2	9	10	10	18	7	5				
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
last	10														

Is P.O.T. property satisfied ? NO!

What is to be done ? SWAP 2 & 4 !

Insert 2 in the Heap

(2nd Iteration)

H E A P	P	C													
Elem	3	2	9	6	4	9	10	18	7	5					
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
last	10														

Is P.O.T. property satisfied ? **NO!**

What is to be done ? **SWAP 2 & 3 !**

Insert 2 in the Heap

(3rd Iteration)

C	2	3	9	6	4	9	10	18	7	5			
H E A P	0	1	2	3	4	5	6	7	8	9	10	11	12
last	10											13	14



Insert 2 in the Heap

(3rd Iteration)

H	E	A	P	C	Elem	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
				last	10															

Is P.O.T. property satisfied ? YES !

Insert Operation is **DONE !**

Insert 2 in the Heap

(3rd Iteration)

Heap H

ELEM	2	3	9	6	4	9	10	10	18	7	5				
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
last	10														

- Given the illustration of the heap above, write an appropriate definition of datatype Heap and declaration of Heap H.
- Based on the definition and declaration, write the codes of the following operations:
 - Insert 
 - DeleteMin

Heapsort

- ★ A sorting technique which uses the concept of a P.O.T. implemented using an array

Heapsort

- ★ A sorting technique which uses the concept of a P.O.T. implemented using an array
- ★ Steps in HEAPSORT (**Descending order**)
 1. INSERT all elements to be sorted in an initially empty Partially Ordered Tree
 2. While (POT is not empty)
 - DELETEMIN and store the deleted minimum element in the position of the element which has replaced the root

Heapsort

Step 1: Insert all elements in an initially empty POT (heap)
OR
Heapify the list (i.e. make list into a heap)

Resulting Heap H

ELEM	2	3	9	6	4	9	10	18	7	5					
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
last	10														

Heapsort: Step 2a

DeleteMin: 2

2

10

temp

oldLast

Heap H

Elem	2	3	9	6	4	9	10	10	18	7	5				
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
last	10														

Heapsort: Step 2a

Delete min: 2

2

10

temp

oldLast

Push the new root down the Tree
until P.O.T. Property is Satisfied!!

Heap H

P	LC	RC														
Elem	5	3	9	6	4	9	10	10	18	7	2					
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	
last	9															



Heapsort: Step 2a

Deletemin: 2

5

10

temp

oldLast

Push the new root down the Tree
until P.O.T. Property is Satisfied!!

Heap H

Elem	3	4	9	6	5	9	10	10	18	7	2				
last	9														
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14

Heapsort: Step 2a

Deletemin: 2

5

10

temp

oldLast

Push the new root down the Tree
until P.O.T. Property Satisfied!!

Heap H

Elem	3	4	9	6	5	9	10	10	18	7	2				
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
last	9														

P

SC

Heapsort Result

last	Elem														
10	2	3	9	6	4	9	10	10	18	7	5				
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
9	3	4	9	6	5	9	10	10	18	7	2				
8	4	5	9	6	7	9	10	10	18	3	2				
7	5	6	9	10	7	9	10	18	4	3	2				
6	6	7	9	10	18	9	10	5	4	3	2				
5	7	10	9	10	18	9	6	5	4	3	2				
4	9	10	9	10	18	7	6	5	4	3	2				

Prepared by chrisp

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Heapsort Result

last	Elem	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
10	2	3	9	6	4	9	10	10	18	7	5					
4	9	10	9	10	18	7	6	5	4	3	2					
3	9	10	18	10	9	7	6	5	4	3	2					
2	10	10	18	9	9	7	6	5	4	3	2					
1	10	18	10	9	9	7	6	5	4	3	2					
0	18	10	10	9	9	7	6	5	4	3	2					
-1	18	10	10	9	9	7	6	5	4	3	2					

From Heap to Sorted List

10

oldLast

Elem	18	10	10	9	9	7	6	5	4	3	2				
	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
last	-1														