

# SQL DML, Physical Data Organization

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# Course Outcomes

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- ▶ CO5: Discuss and compare the aspects of Concurrency Control and Recovery in Database systems (Cognitive Knowledge Level: Apply)
- ▶ CO6: Explain various types of NoSQL databases (Cognitive Knowledge Level: Understand)



Company: Schema

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## EMPLOYEE

FNAME	MINIT	LNAME	<u>SSN</u>	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO
-------	-------	-------	------------	-------	---------	-----	--------	----------	-----

## DEPARTMENT

DNAME	<u>DNUMBER</u>	MGRSSN	MGRSTARTDATE
-------	----------------	--------	--------------

## DEPT\_LOCATIONS

<u>DNUMBER</u>	<u>DLOCATION</u>
----------------	------------------

## PROJECT

PNAME	<u>PNUMBER</u>	PLOCATION	DNUM
-------	----------------	-----------	------

## WORKS\_ON

<u>ESSN</u>	<u>PNO</u>	HOURS
-------------	------------	-------

## DEPENDENT

<u>ESSN</u>	<u>DEPENDENT_NAME</u>	SEX	BDATE	RELATIONSHIP
-------------	-----------------------	-----	-------	--------------

Figure: Company Schema

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EMPLOYEE	FNAME	MINIT	LNAME	<u>SSN</u>	BDATE	ADDRESS	SEX	SALARY	SUPERSSN	DNO
	John	B	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
	Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	M	40000	888665555	5
	Alicia	J	Zelaya	999887777	1968-07-19	3321 Castle, Spring, TX	F	25000	987654321	4
	Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
	Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	M	38000	333445555	5
	Joyce	A	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
	Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	M	25000	987654321	4
	James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	M	55000	null	1

					DEPT_LOCATIONS	DNUMBER	DLOCATION
						1	Houston
						4	Stafford
						5	Bellaire
						5	Sugarland
						5	Houston
DEPARTMENT	DNAME	DNUMBER	MGRSSN	MGRSTARTDATE			
	Research	5	333445555	1988-05-22			
	Administration	4	987654321	1995-01-01			
	Headquarters	1	888665555	1981-06-19			

WORKS_ON	<u>ESSN</u>	<u>PNO</u>	HOURS
	123456789	1	32.5
	123456789	2	7.5
	666884444	3	40.0
	453453453	1	20.0
	453453453	2	20.0
	333445555	2	10.0
	333445555	3	10.0
	333445555	10	10.0
	333445555	20	10.0
	999887777	30	30.0
	999887777	10	10.0
	987987987	10	35.0
	987987987	30	5.0
	987654321	30	20.0
	987654321	20	15.0
	888665555	30	10.0

PROJECT	PNAME	<u>PNUMBER</u>	PLOCATION	DNUM
	ProductX	1	Bellaire	5
	ProductY	2	Sugarland	5
	ProductZ	3	Houston	5
	Computerization	10	Stafford	4
	Reorganization	20	Houston	1
	Newbenefits	30	Stafford	4

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The general form of SQL

- ▶ SELECT column  
FROM table  
WHERE predicate on rows  
GROUP BY columns  
HAVING predicate on groups  
ORDER BY columns

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FROM EMPLOYEE  
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- ▶ Similar to a SELECT-PROJECT-JOIN of relational algebra operations:

- ▶ The SELECT clause specifies the projection attributes
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FROM EMPLOYEE, DEPARTMENT  
WHERE DNAME='Research' AND DNUMBER=DNO
```

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SELECT FNAME, LNAME, ADDRESS  
FROM EMPLOYEE, DEPARTMENT  
WHERE DNAME='Research' AND DNUMBER=DNO
```
- ▶ A query that refers to two or more attributes with the same name must qualify the attribute name with the relation name by prefixing the relation name to the attribute name
  - ▶ DNUMBER is both DEPARTMENT and DEPT\_LOCATIONS.
  - ▶ Qualify by relationname: DEPARTMENT.DNUMBER, DEPT\_LOCATIONS.DNUMBER

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```
SELECT PNUMBER, DNUM, LNAME, ADDRESS, BDATE  
FROM PROJECT,DEPARTMENT, EMPLOYEE,  
WHERE DNUM=DNUMBER AND MGSSSN=SSN AND  
PLOCATION ='STAFFORD'
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- ▶ The join condition `DNUM=DNUMBER` relates a project to its controlling department
- ▶ The join condition `MGRSSN=SSN` relates the controlling department to the employee who manages that department

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WHERE E.SUPERSSN=S.SSN



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- ▶ To retrieve all the attribute values of the selected tuples, a \* is used, which stands for all the attributes
  - ▶ SELECT \* FROM EMPLOYEE WHERE DNO=5
- ▶ To eliminate duplicate tuples in a query result, the keyword DISTINCT is used
  - ▶ SELECT DISTINCT SALARY  
FROM EMPLOYEE

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  - ▶ (SELECT PNUMBER, PNAME  
FROM PROJECT, DEPARTMENT EMPLOYEE  
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  - ▶ (SELECT PNUMBER, PNAME  
FROM PROJECT, DEPARTMENT EMPLOYEE  
WHERE DNUM=DNUMBER AND MGRSSN=SSN AND  
LNAME='Smith')  
UNION  
(SELECT PNUMBER, PNAME  
FROM PROJECT, WORKS\_ON, EMPLOYEE  
WHERE PNUMBER=PNO AND ESSN=SSN AND  
NAME='Smith')

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  - ▶ SELECT FNAME, LNAME, ADDRESS  
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WHERE DNO IN

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- ▶ The outer query select an EMPLOYEE tuple if its DNO value is in the result of the nested query
- ▶ The comparison operator **IN** compares a value **v** with a set (or multi-set) of values **V**, and evaluates to **TRUE** if **v** is one of the elements in **V**

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- ▶ **Evaluated once for each tuple in the outer query**

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```
SELECT FNAME, LNAME  
FROM EMPLOYEE  
WHERE SALARY ≥ ALL
```

## SQL: =ANY(or =SOME)

- ▶ In addition to **IN**, other operators like **=ANY(or =SOME)** can also be used for multiset comparison.
- ▶ The **=ANY(or =SOME)** operator returns TRUE if the value **v** is equal to some value in the set **V**, and hence equivalent to **IN**.
- ▶ **<, ≤, >, ≥, and <>** can be combined with **=ANY(or =SOME)**. The keyword **ALL** can also be combined with these operators.
- ▶ Retrieve names of all employees whose salary is greater than the salary of all employees in department 5
  - ▶ 

```
SELECT FNAME, LNAME
FROM EMPLOYEE
WHERE SALARY ≥ ALL
      (SELECT SALARY
       FROM EMPLOYEE
       WHERE DNO=5)
```

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  - ▶ SELECT E.FNAME, E.LNAME  
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WHERE EXISTS

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```
SELECT E.FNAME, E.LNAME
FROM EMPLOYEE AS E
WHERE EXISTS
    (SELECT *
     FROM DEPENDENT AS D
     WHERE E.SSN = D.ESSN AND E.SEX = D.SEX
     AND
     E.FNAME=D.DEPENDENT_NAME)
```



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WHERE EXISTS
      (SELECT *
       FROM DEPENDENT AS D
       WHERE E.SSN = D.ESSN AND E.SEX = D.SEX
       AND
       E.FNAME=D.DEPENDENT_NAME)
```
- ▶ For each employee tuple, evaluate the nested query, which retrieves **all DEPENDENT tuple with same SSN, sex and first name as the EMPLOYEE tuple**, if at least **one tuple EXISTS in the result** of the nested query, then select the EMPLOYEE tuple.

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FROM EMPLOYEE  
WHERE NOT EXISTS

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  - ▶ SELECT E.FNAME, E.LNAME  
FROM EMPLOYEE  
WHERE NOT EXISTS  
    (SELECT \*  
      FROM DEPENDENT  
      WHERE E.SSN = ESSN )

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  - ▶ SELECT E.FNAME, E.LNAME  
FROM EMPLOYEE  
WHERE NOT EXISTS  
(SELECT \*  
FROM DEPENDENT  
WHERE E.SSN = ESSN )
- ▶ For each EMPLOYEE tuple, the subquery **selects all DEPENDENT tuples whose ESSN value matches the EMPLOYEE SSN**. If subquery **results is empty**, **select the EMPLOYEE tuple and retrieve FNAME, LNAME**.



## SQL: NOT EXISTS

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AND

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```
SELECT E.FNAME, E.LNAME
FROM EMPLOYEE AS E
WHERE EXISTS
    (SELECT *
     FROM DEPENDENT
     WHERE E.SSN = ESSN )
AND
    EXISTS (SELECT *
           FROM DEPARTMENT
           WHERE SSN = MGRSSN);
```

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    EXISTS (SELECT *
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           WHERE SSN = MGRSSN);
```

- ▶ Two nested correlated subqueries:
  - ▶ 1. Select all DEPENDENT tuples related to EMPLOYEE
  - ▶ 2. Select all DEPARTMENT tuples managed by the EMPLOYEE
  - ▶ If at least one tuple each first and second subquery exists, then that EMPLOYEE tuple is selected.

## SQL: NOT EXISTS

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- ▶ Retrieve names of each employee who works on all the projects controlled by department 5



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(

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```
SELECT E.FNAME, E.LNAME
FROM EMPLOYEE
WHERE NOT EXISTS
    (
      (SELECT PNUMBER
       FROM PROJECT
       WHERE DNUM = 5 )
    EXCEPT
```

# SQL: NOT EXISTS

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FROM EMPLOYEE
WHERE NOT EXISTS
    (
      (SELECT PNUMBER
       FROM PROJECT
       WHERE DNUM = 5 )
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      (SELECT PNO
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       WHERE SSN = ESSN);
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FROM PROJECT  
WHERE DNUM = 5 )  
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(SELECT PNO  
FROM WORKS\_ON  
WHERE SSN = ESSN);

- ▶ Two nested correlated subqueries:
  - ▶ 1. Select all projects controlled by department 5
  - ▶ 2. Select all projects that the particular employee being works on
  - ▶ If the set difference of first minus second subquery is empty, then that employee works on all projects

## Joined Tables in SQL and Outer Joins

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```
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- ▶ Here the FROM clause contains a single **joined** table.
- ▶ There are several types of join
  1. NATURAL JOIN: No join conditions are specified between the relations **R** and **S**; an *implicit EQUIJOIN condition* for each pair of attributes with the same name from **R** and **S** are created
  2. OUTER JOIN: When the user wants to keep all tuples of **R**, or all those tuples in **S**, or all those in both relations in the result of the JOIN, regardless of whether or not they have matching tuples in other relations

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```
SELECT FNAME, LNAME, ADDRESS  
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      (DEPARTMENT AS DEPT(DName, DNO, MSsn, MSdate)))  
WHERE DNAME ='Research'
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```
SELECT E.FNAME, E.LNAME, S.FNAME, S.LNAME  
FROM (EMPLOYEE AS E LEFT OUTER JOIN EMPLOYEE  
      AS S  
      ON E.SUPERSSN=S.SSN)
```

Joined Tables: LEFT OUTER, RIGHT OUTER, FULL  
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  - ▶ For every project located in 'Stafford', list the project number, the controlling department number, and the department manager's last name, address, and birthdate.



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- ▶ Multiway JOIN
  - ▶ For every project located in 'Stafford', list the project number, the controlling department number, and the department manager's last name, address, and birthdate.
  - ▶ 

```
SELECT PNUMBER, DNUM, LNAME, ADDRESS, BDATE  
FROM ((PROJECT JOIN DEPARTMENT ON DNUM  
=DNUMBER)  
JOIN EMPLOYEE ON MGRSSN = SSN)  
WHERE PLOCATION ='STAFFORD'
```

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```
SELECT SUM(SALARY), MAX(SALARY), MIN(SALARY),  
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```
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  - ▶ 

```
SELECT SUM(SALARY), MAX(SALARY), MIN(SALARY),  
      AVG(SALARY)  
FROM (EMPLOYEE JOIN DEPARTMENT ON DNO  
      =DNUMBER)  
WHERE DNAME ='RESEARCH'
```



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- ▶ Count the number of distinct salary values in the database

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  - ▶ `SELECT LNAME, FNAME`  
`FROM EMPLOYEE`  
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# Aggregate Functions in SQL

- ▶ Find the number of employees in Research department
  - ▶ `SELECT COUNT(*)`  
`FROM (EMPLOYEE, DEPARTMENT`  
`WHERE DNO=DNUMBER AND DNAME ='RESEARCH')`
- ▶ Count the number of distinct salary values in the database
  - ▶ `SELECT (DISTINCT SALARY)`  
`FROM EMPLOYEE`
- ▶ Retrieve names of employees who have two or more dependents
  - ▶ `SELECT LNAME, FNAME`  
`FROM EMPLOYEE`  
`WHERE ( SELECT COUNT(*)`  
`FROM DEPENDENT WHERE SSN=ESSN)  $\geq$  2`
- ▶ The above correlated query retrieves number of dependents of the employee, which used in the WHERE clause of outer query.



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```
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FROM EMPLOYEE
GROUP BY DNO
```

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```
SELECT DNO, COUNT(*), AVG(SALARY)
FROM EMPLOYEE
GROUP BY DNO
```
- ▶ The SELECT clause include only the grouping attribute and aggregate functions to be applied on each group

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SELECT DNO, COUNT(*), AVG(SALARY)
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```
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- ▶ For each project, retrieve the project number, the project name, and the number of employees who work on this project

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```
- ▶ The SELECT clause include only the grouping attribute and aggregate functions to be applied on each group
- ▶ For each project, retrieve the project number, the project name, and the number of employees who work on this project
  - ▶ 

```
SELECT PNUMBER, PNAME, COUNT(*)
FROM PROJECT, WORK_ON
WHERE PNUMBER =PNO
GROUP BY PNUMBER, PNAME
```

## GROUP BY and HAVING



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- ▶ For each project on which more than two employees work, retrieve the project number, project name, and the number of employees works on the project.

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```
SELECT PNUMBER, PNAME, COUNT(*)  
FROM PROJECT, WORK_ON  
WHERE PNUMBER =PNO  
GROUP BY PNUMBER, PNAME  
HAVING COUNT(*) > 2
```

## GROUP BY and HAVING

- ▶ For each project on which more than two employees work, retrieve the project number, project name, and the number of employees works on the project.
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```
SELECT PNUMBER, PNAME, COUNT(*)  
FROM PROJECT, WORK_ON  
WHERE PNUMBER =PNO  
GROUP BY PNUMBER, PNAME  
HAVING COUNT(*) > 2
```
- ▶ For each project, retrieve the project number, the project name, and the number of employees from department 5 who work on the project

## GROUP BY and HAVING

- ▶ For each project on which more than two employees work, retrieve the project number, project name, and the number of employees works on the project.
  - ▶ 

```
SELECT PNUMBER, PNAME, COUNT(*)  
FROM PROJECT, WORK_ON  
WHERE PNUMBER =PNO  
GROUP BY PNUMBER, PNAME  
HAVING COUNT(*) > 2
```
- ▶ For each project, retrieve the project number, the project name, and the number of employees from department 5 who work on the project
  - ▶ 

```
SELECT PNUMBER, PNAME, COUNT(*)  
FROM PROJECT, WORK_ON, EMPLOYEE  
WHERE PNUMBER =PNO AND SSN=ESSN AND DNO=5  
GROUP BY PNUMBER, PNAME
```

## GROUP BY and HAVING

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- ▶ Count the total number of employees whose salary exceeds rs.40000 in each department, but only for departments where more than five employees work

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```
SELECT DNAME, COUNT(*)  
FROM DEPARTMENT, EMPLOYEE  
WHERE DNUMBER =DNO AND SALARY > 40000  
GROUP BY DNAME HAVING COUNT(*) > 5
```

## GROUP BY and HAVING

- ▶ Count the total number of employees whose salary exceeds rs.40000 in each department, but only for departments where more than five employees work
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```
SELECT DNAME, COUNT(*)  
FROM DEPARTMENT, EMPLOYEE  
WHERE DNUMBER =DNO AND SALARY > 40000  
GROUP BY DNAME HAVING COUNT(*) > 5
```
- ▶ This is incorrect, it will select only departments that have more than employees who can earn more than rs. 40000



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```
- ▶ This is incorrect, it will select only departments that have more than employees who can earn more than rs. 40000
- ▶ Here WHERE is executed first, but HAVING is applied later.

## GROUP BY and HAVING

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GROUP BY DNAME HAVING COUNT(*) > 5
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- ▶ This is incorrect, it will select only departments that have more than employees who can earn more than rs. 40000
- ▶ Here WHERE is executed first, but HAVING is applied later.
- ▶ The correct query is

## GROUP BY and HAVING

- ▶ Count the total number of employees whose salary exceeds rs.40000 in each department, but only for departments where more than five employees work

- ▶ 

```
SELECT DNAME, COUNT(*)  
FROM DEPARTMENT, EMPLOYEE  
WHERE DNUMBER =DNO AND SALARY > 40000  
GROUP BY DNAME HAVING COUNT(*) > 5
```

- ▶ This is incorrect, it will select only departments that have more than employees who can earn more than rs. 40000
- ▶ Here WHERE is executed first, but HAVING is applied later.
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```
SELECT DNAME, COUNT(*)  
FROM DEPARTMENT, EMPLOYEE  
WHERE DNUMBER =DNO AND SALARY > 40000 AND  
      (SELECT DNO  
       FROM EMPLOYEE  
       GROUP BY DNAME  
       HAVING COUNT(*) > 5)
```

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  - ▶ The UNIQUE clause specifies alternate keys

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  - ▶ CREATE TABLE EMPLOYEE (  
.  
    ESSN CHAR(9),  
    DNO INTEGER DEFAULT 1,  
    SUPERSSN CHAR(9),  
    PRIMARY KEY (ESSN),  
    FOREIGN KEY (DNO) REFERENCES DEPARTMENT  
    ON DELETE SET DEFAULT  
    ON UPDATE CASCADE,  
    FOREIGN KEY (SUPERSSN) REFERENCES  
EMPLOYEE  
    ON DELETE SET NULL  
    ON UPDATE CASCADE);

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- ▶ Specify a constraint that **the salary of an employee must not be greater than the salary of the manager of the department that the employee works for**
- ▶ CREATE ASSERTION SALARY\_CONSTRAINT  
CHECK (NOT EXISTS(SELECT \*  
                    FROM EMPLOYEE E, EMPLOYEE M,  
                                DEPARTMENT D  
                    WHERE E.Salary > M.Salary  
                    AND E.Dno =D.DNumber  
                    AND D.Mgr\_ssn =M.Ssn))



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- ▶ For example, it is useful to specify a condition that, if violated, causes some users to be informed about the violation.
- ▶ We want to check **whenever employees salary is greater than the salary of his or her direct supervisor.** Several events can trigger this rule :
  - ▶ Insertion of new employee tuple
  - ▶ Changing an employee's salary
  - ▶ Changing an employee's supervisor

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  - ▶ Changing an employee's supervisor
- ▶ Let the action required is **call an external stored procedure**.

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  - ▶ CREATE TRIGGER SALARY\_VIOLATION  
BEFORE INSERT OR UPDATE OF  
SALARY, SUPERVISOR\_SSN ON EMPLOYEE  
FOR EACH ROW  
WHEN (NEW.SALARY > (SELECT SALARY  
FROM EMPLOYEE  
WHERE SSN=NEW.SUPERVISOR\_SSN))  
INFORM\_SUPERVISOR  
(NEW.SUPERVISOR\_SSN, NEW.SSN);



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- ▶ **Action** : Usually sequence of SQL statements; or it can be a stored procedure. Here `INFORM_SUPERVISOR`

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AFTER INSERT ON EMPLOYEE  
FOR EACH ROW  
    WHEN (NEW.Dno IS NOT NULL)  
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- ▶ The **statement-level trigger** is triggered once for triggering statement.

## Row-level trigger vs Statement-level trigger

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  - ▶ UPDATE EMPLOYEE  
SET Salary=1.1\*Salary  
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- ▶ This operation would be an event for that trigger the **Total\_Sal** trigger
- ▶ The above statement could update multiple records, hence a rule using **row-level semantics** is required to trigger for each row. A **statement-level semantics** is triggered only once, which is not suitable for the above update.

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```
CREATE VIEW DEPT_INFO(Dept_name, No_of_emps,  
Total_sal)  
AS SELECT Dname, COUNT(*), SUM(Salary)  
FROM DEPARTMENT, EMPLOYEE  
WHERE Dnumber =Dno  
GROUP BY Dname
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- ▶ To dispose a view: `DROP VIEW WORKS_ON`

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  - ▶ Views defined using grouping and aggregate functions are not updatable.