

**CSE221**

# Lecture 2: Arrays

2021 Fall

Young-ri Choi

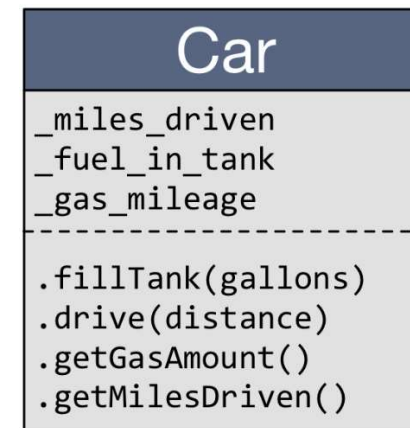
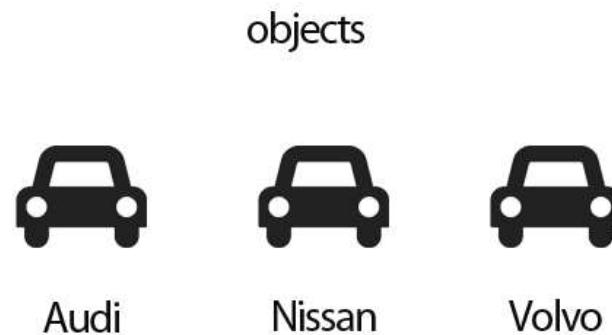
Acknowledgment: The content of this file is based on the slides of the textbook as well as the slides used in former lectures at UNIST.

# Outline

- Basic concepts
- Arrays
- Examples
  - Sparse matrices
  - Selection sort

# Object

- Entity that performs *computations* and has a *local state*
  - Procedural elements + variables
- An instance of a class



# Object-oriented Design

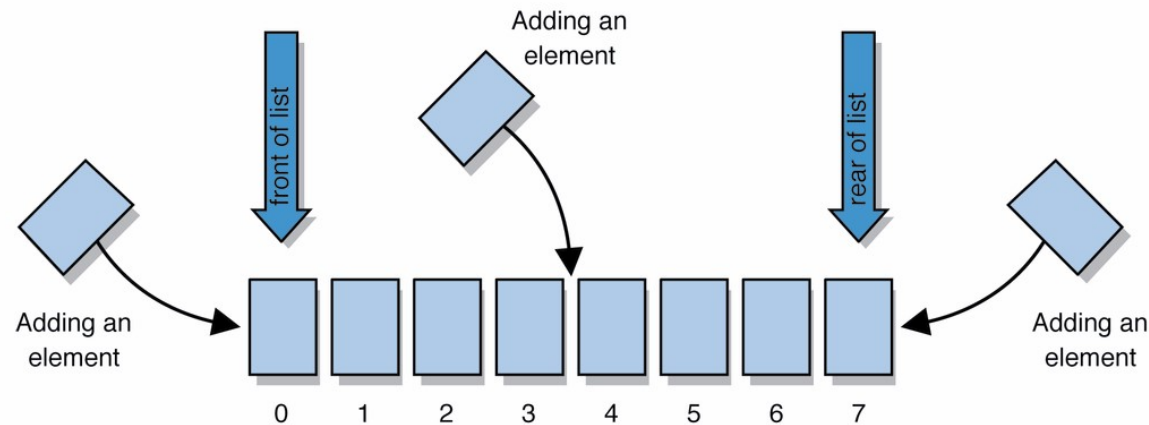
- Object as fundamental building block
  - View software as a set of interacting objects that model entities
  - Use *divide and conquer*
- Benefits
  - Encourage code reusability and flexibility
  - Localize changes: related objects *only*

# Data Encapsulation and Abstraction

- Data encapsulation
  - Concealing of the implementation details of a data object from the outside
- Data abstraction
  - Separation between the **specification (what)** of a data object and its **implementation (how)**

# Abstract Data Type (ADT)

- Definition for expected operations and behavior
- Example: list
  - A collection storing an ordered sequence of elements
  - Operations:



# Data Structure and ADT

- Data structure
  - How we organize/store data points
  - Determines how each operation would behave
- One ADT can be implemented in many ways
- One data structure can be used for multiple ADTs

# ADT First, Then Data Structure

- Wrong way:
  - I will represent my data A as an array
- Right way:
  - I will be frequently accessing a random data point of my data A, so I need operation []
  - I will therefore use an array



# Benefits of Abstraction and Encapsulation

- Simplification of software development
  - Enables parallel development of subtasks
- Testing and debugging
  - Easier to identify problems by testing subtasks
- Reusability
  - Easier to take out a part and use elsewhere
- Modification of data type
  - Implementation can be modified without changing interfaces

# Outline

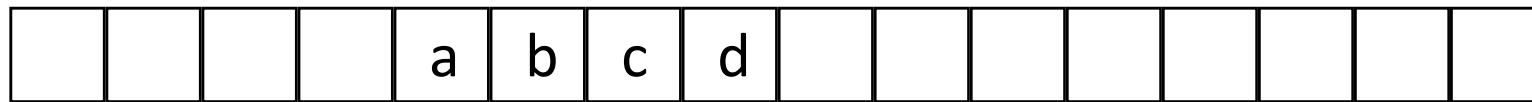
- Basic concepts
- Arrays
- Examples
  - Sparse matrices
  - Selection sort

# Arrays in C/C++

- **Continuous** memory space with the **same data type**
- `type name [elements];`
- **Example:**
  - `int a[100];`
  - `float b[10][20];`
  - `mytype c[5];` // user-defined data type
- Each element accessed by index

# 1D Array Representation In C

1D Memory Space

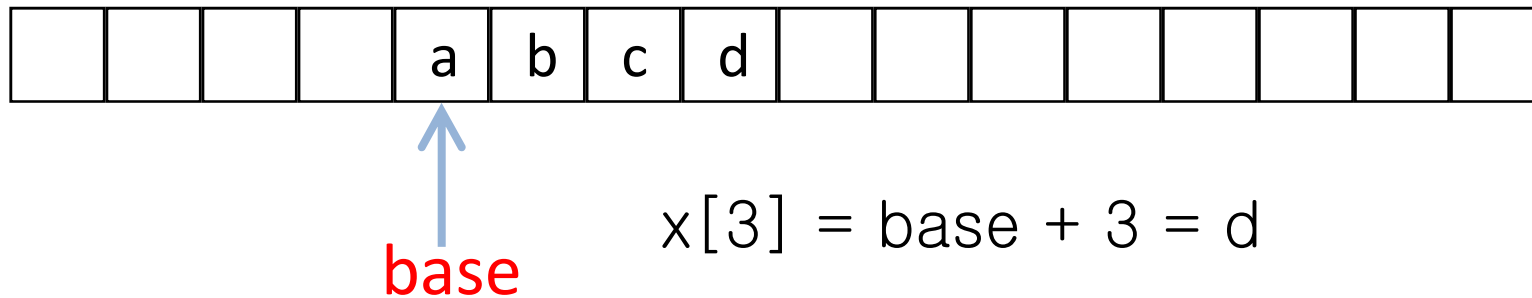


$$\text{base} \quad x[0] = \text{base} + 0 = \text{a}$$

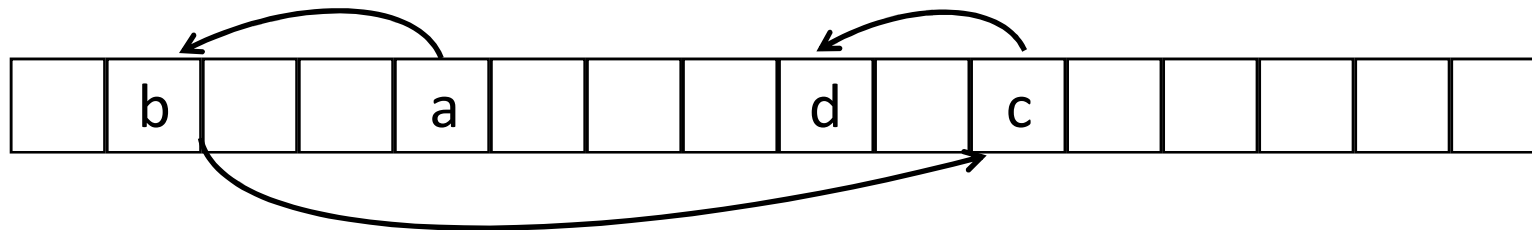
- 1-dimensional array  $x = [a, b, c, d]$
- map into contiguous memory locations
- $\text{location}(x[i]) = \text{base} + i$

# Array v.s. Pointer-based Linked List

Array



Linked list



# Array v.s. Pointer-based Linked List

- Arrays
  - Good for random and sequential access
    - Indexing
    - Modern architecture reads data at a chunk
  - Good when the size is quite fixed
- Linked list
  - Good for frequent inserting and deleting
  - Resizing is easier

# 2D Arrays

- The elements of a 2-dimensional array `x` declared as:
- `int x[3][4];`
- may be shown as a table

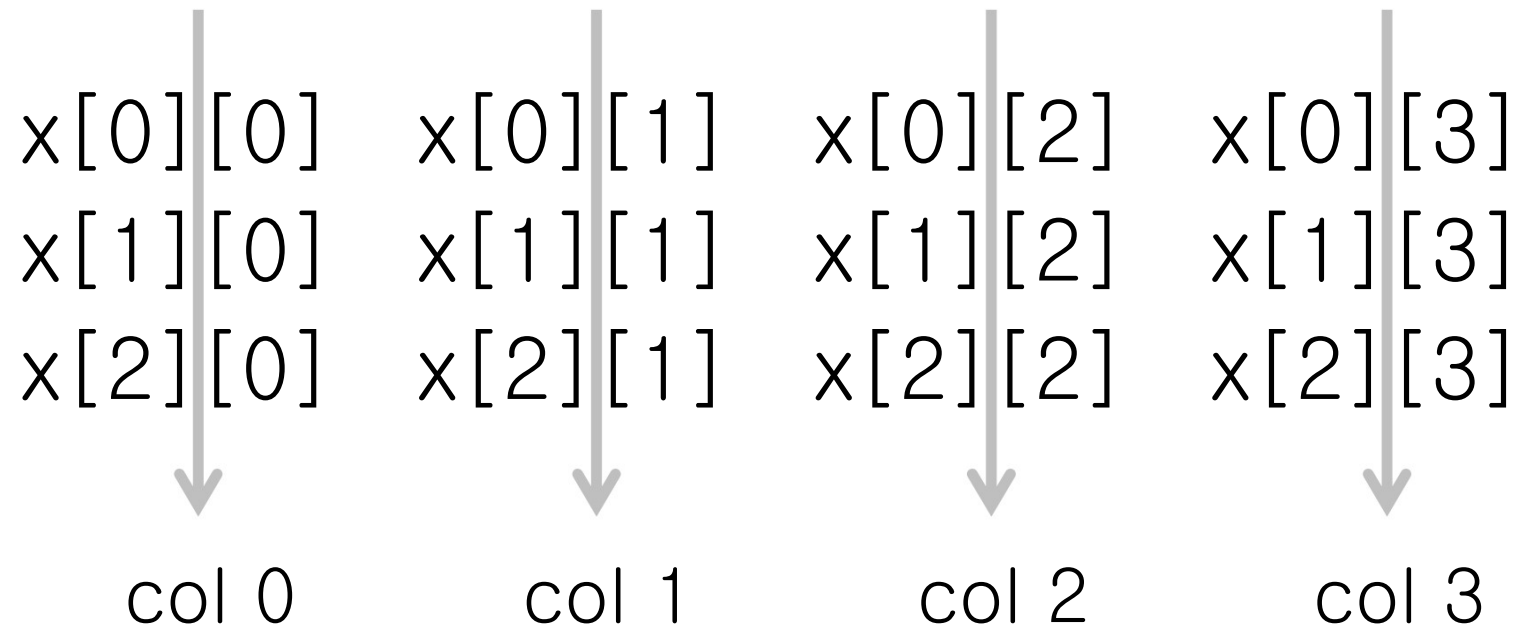
<code>x[0][0]</code>	<code>x[0][1]</code>	<code>x[0][2]</code>	<code>x[0][3]</code>
<code>x[1][0]</code>	<code>x[1][1]</code>	<code>x[1][2]</code>	<code>x[1][3]</code>
<code>x[2][0]</code>	<code>x[2][1]</code>	<code>x[2][2]</code>	<code>x[2][3]</code>

# Rows of a 2D Array

$x[0][0]$   $x[0][1]$   $x[0][2]$   $x[0][3]$  → row 0  
 $x[1][0]$   $x[1][1]$   $x[1][2]$   $x[1][3]$  → row 1  
 $x[2][0]$   $x[2][1]$   $x[2][2]$   $x[2][3]$  → row 2



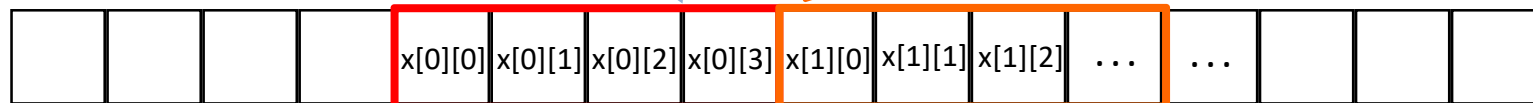
# Columns of a 2D Array



# 1D Indexing of 2D Array

$x[0][0]$	$x[0][1]$	$x[0][2]$	$x[0][3]$
$x[1][0]$	$x[1][1]$	$x[1][2]$	$x[1][3]$
$x[2][0]$	$x[2][1]$	$x[2][2]$	$x[2][3]$

1D Memory Space

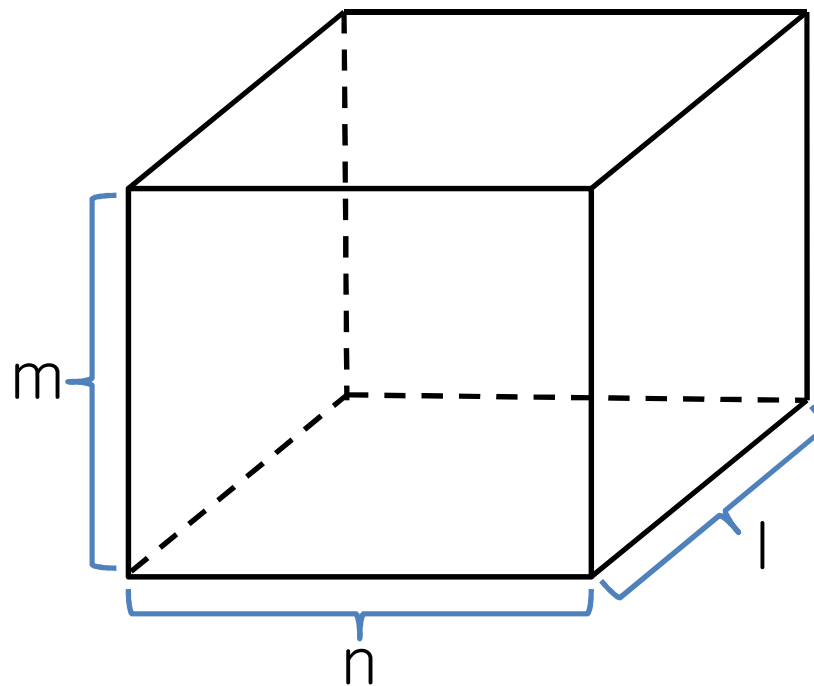


base

$(i,j)$  in  $a[m][n]$  :  $\text{base} + i*n + j$   
 $m$  : # rows,  $n$  : # columns (size of a row)  
 $i$  : row index,  $j$  : column index

# 1D Indexing of 3D Array

- $(i,j,k)$  in  $x[l][m][n]$  :  $\text{base} + i*m*n + j*n + k$



# Array Abstract Data Type

- More general than array in C/C++
  - Contiguous memory location is more of a implementation issue
- Properties:
  - Mapping between index and value
    - Do not need to be implemented with pure array!
  - Operators to retrieve and store value

```

class GeneralArray {
// A set of pairs <index, value> where for each
// value of index in index set, there is a value of
// type float
public:
    // Constructor.
    // j : dimension, list : range of each dimension.
    // initValue : initial value
    GeneralArray(int j,
                  RangeList list,
                  float initValue);

    // Return the float associate with i if i is in
    // the index set. Otherwise throw an exception
    float Retrieve(index i);

    // Replace the old value associate with i if i is
    // in the index set. Otherwise throw an exception
    void Store(index i, float x)
}

```

# Outline

- Recap OOP: programming concepts
- Arrays
- Examples
  - Sparse matrices
  - Selection sort

# Matrices

- $A[m][n]$ 
  - $m \times n$  matrix  $A$
  - $m$  : # of rows
  - $n$  : # of column
  - $m*n$  : # of elements
- Accessing an element at  $i$ -th row and  $j$ -th col
  - $A[i][j]$
- 2D array is often used to represent matrices → Dense matrix

# Sparse Matrices

- Mostly zero

$$\begin{array}{c} \begin{array}{ccc} & 0 & 1 & 2 \\ \begin{array}{c} 0 \\ 1 \\ 2 \\ 3 \\ 4 \end{array} & \left[ \begin{array}{ccc} -27 & 3 & 4 \\ 6 & 82 & -2 \\ 109 & -64 & 11 \\ 12 & 8 & 9 \\ 48 & 27 & 47 \end{array} \right] \end{array}\end{array}$$

(a) Dense Matrix

$$\begin{array}{c} \begin{array}{cccccc} & 0 & 1 & 2 & 3 & 4 & 5 \\ \begin{array}{c} 0 \\ 1 \\ 2 \\ 3 \\ 4 \\ 5 \end{array} & \left[ \begin{array}{cccccc} 15 & 0 & 0 & 22 & 0 & -15 \\ 0 & 11 & 3 & 0 & 0 & 0 \\ 0 & 0 & 0 & -6 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 91 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 28 & 0 & 0 & 0 \end{array} \right] \end{array}\end{array}$$

(b) Sparse Matrix



# Sparse Matrix Representation

- Array of triples <row, col, value> for nonzero elements
  - # elements to store (**space** requirement ) for nxn matrix
    - Dense :  $n^2$
    - Sparse diagonal matrix :  $3*n$
- } Meaningful when n is very large!

Tridiagonal  
Matrix

$$f_n = \begin{vmatrix} a_1 & b_1 & & & \\ c_1 & a_2 & b_2 & & \\ & c_2 & \ddots & \ddots & \\ & & \ddots & \ddots & b_{n-1} \\ & & & c_{n-1} & a_n \end{vmatrix}.$$

# Row Major vs. Column Major

- 2D->1D layout
- Row major
  - Store each row in contiguous memory
- Column major
  - Store each column in contiguous memory

# Sparse Matrix


- Array of triples <row, col, value> for nonzero elements

```
class SparseMatrix;  
    class MatrixTerm {  
        friend class SparseMatrix;  
        private:  
            int row, col, value;  
    };  
  
    private:  
        int rows, cols, terms, capacity;  
        MatrixTerm * smArray;
```



# Space Cost Analysis


1 integer



	0	1	2	3	4	5
0	15	0	0	22	0	-15
1	0	11	3	0	0	0
2	0	0	0	-6	0	0
3	0	0	0	0	0	0
4	91	0	0	0	0	0
5	0	0	28	0	0	0

→

3 integers  
(3x more space)



	row	col	val
[0]	0	0	15
[1]	0	3	22
[2]	0	5	-15
[3]	1	1	11
[4]	1	2	3
[5]	2	3	-6
[6]	4	0	91
[7]	5	2	28

space saving = 12  
(36/24 integers in dense/sparse matrix)

- Sparse representation is more effective when there are many empty elements

# Sparse Matrix Transpose

- Exchange  $(i,j)$  to  $(j,i)$
  - Pseudo code:
    - for each row  $i$ , store  $(i,j)$  to  $(j,i)$
    - Problem
      - $(0, 0, 15) \rightarrow (0, 0, 15)$
      - $(0, 3, 22) \rightarrow (3, 0, 22)$
      - $(0, 5, -15) \rightarrow (5, 0, -15)$
      - $(1, 1, 11) \rightarrow (1, 1, 11)$  **x**
- Violate row-major order

# Sparse Matrix Transpose

- New approach
  - Use column major order traversal
    - After transpose, matrix is row-major order

for (all element in **column** j)  
store (i,j,value) of the original matrix as (j,i,value) of transposed matrix

# Example


		row	col	val			row	col	val
smArray	[0]	0	0	15	→ smArray →	[0]	0	0	15
	[1]	0	3	22		[1]	0	4	91
	[2]	0	5	-15		[2]			
	[3]	1	1	11		[3]			
	[4]	1	2	3		[4]			
	[5]	2	3	-6		[5]			
	[6]	4	0	91		[6]			
	[7]	5	2	28		[7]			

for Column 0



# Example

		row	col	val			row	col	val
smArray	[0]	0	0	15	smArray	[0]	0	0	15
	[1]	0	3	22		[1]	0	4	91
	[2]	0	5	-15		[2]	1	1	11
	[3]	1	1	11		[3]			
	[4]	1	2	3		[4]			
	[5]	2	3	-6		[5]			
	[6]	4	0	91		[6]			
	[7]	5	2	28		[7]			



for Column 1

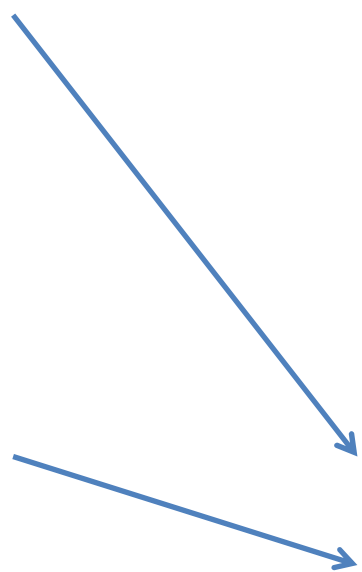
# Example

		row	col	val			row	col	val
smArray	[0]	0	0	15	smArray	[0]	0	0	15
	[1]	0	3	22		[1]	0	4	91
	[2]	0	5	-15		[2]	1	1	11
	[3]	1	1	11		[3]	2	1	3
	[4]	1	2	3		[4]	2	5	28
	[5]	2	3	-6		[5]			
	[6]	4	0	91		[6]			
	[7]	5	2	28		[7]			

for Column 2

# Example

		row	col	val			row	col	val
smArray	[0]	0	0	15	smArray	[0]	0	0	15
	[1]	0	3	22		[1]	0	4	91
	[2]	0	5	-15		[2]	1	1	11
	[3]	1	1	11		[3]	2	1	3
	[4]	1	2	3		[4]	2	5	28
	[5]	2	3	-6		[5]	3	0	22
	[6]	4	0	91		[6]	3	2	-6
	[7]	5	2	28		[7]			



for Column 3

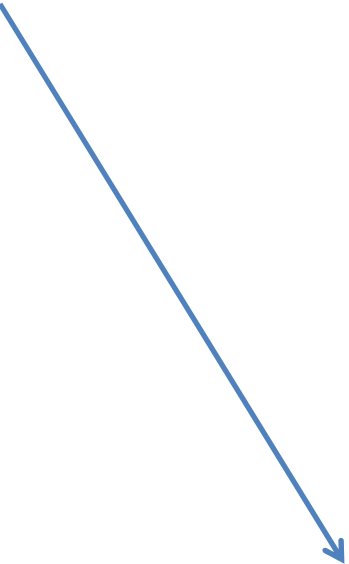
# Example

		row	col	val			row	col	val
smArray	[0]	0	0	15	smArray	[0]	0	0	15
	[1]	0	3	22		[1]	0	4	91
	[2]	0	5	-15		[2]	1	1	11
	[3]	1	1	11		[3]	2	1	3
	[4]	1	2	3		[4]	2	5	28
	[5]	2	3	-6		[5]	3	0	22
	[6]	4	0	91		[6]	3	2	-6
	[7]	5	2	28		[7]			

for Column 4 : No Match 36

# Example

		row	col	val			row	col	val
smArray	[0]	0	0	15	smArray	[0]	0	0	15
	[1]	0	3	22		[1]	0	4	91
	[2]	0	5	-15		[2]	1	1	11
	[3]	1	1	11		[3]	2	1	3
	[4]	1	2	3		[4]	2	5	28
	[5]	2	3	-6		[5]	3	0	22
	[6]	4	0	91		[6]	3	2	-6
	[7]	5	2	28		[7]	5	0	-15



for Column 5

```

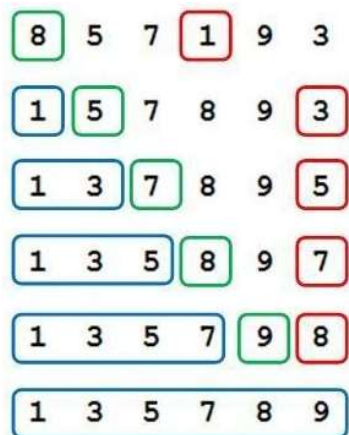
SparseMatrix SparseMatrix::Transpose(){
    SparseMatrix b;           // b : output (transposed matrix)
    b.Rows = Cols;            // b rows = a cols
    b.Cols = Rows;            // b cols = a rows
    b.Terms = Terms;          // # of elements is same

    if (Terms > 0){           // non empty matrix
        int CurrentB = 0;
        for (int c = 0; c < Cols; c++)    // per each column
            for (int i = 0; i < Terms; i++) // for all nonzero elements in sparse
                if (smArray[i].col == c) { // matrix
                    b.smArray[CurrentB].row = c;
                    b.smArray[CurrentB].col = smArray[i].row;
                    b.smArray[CurrentB].value = smArray[i].value;
                    CurrentB++;
                }
    }
    return b;
}

```

# Selection Sort

- *From those integers that are currently unsorted, find the smallest and place it next in the sorted list*



[i] [j]

```
for (int i = 0; i < n; i++) [  
    examine a[i] to a[n-1] and suppose the smallest is at a[j];  
    interchange a[i] and a[j];  
]
```

*This is algorithm!*

# Selection Sort

```
1    void sort(int *a, const int n)
2    // sort n integers a[0] to a[n-1] into nondecreasing order.
3    {
4        for (int i = 0; i < n; i++)
5        {
6            int j = i;
7            // find smallest integer in a[i] and a[n-1]
8            for (int k = i + 1; k < n; k++)
9                if (a[k] < a[j]) j = k;
10           // swap
11           int temp = a[i]; a[i] = a[j]; a[j] = temp;
12       }
13   }
```

*This is C++ program!*



# Questions?