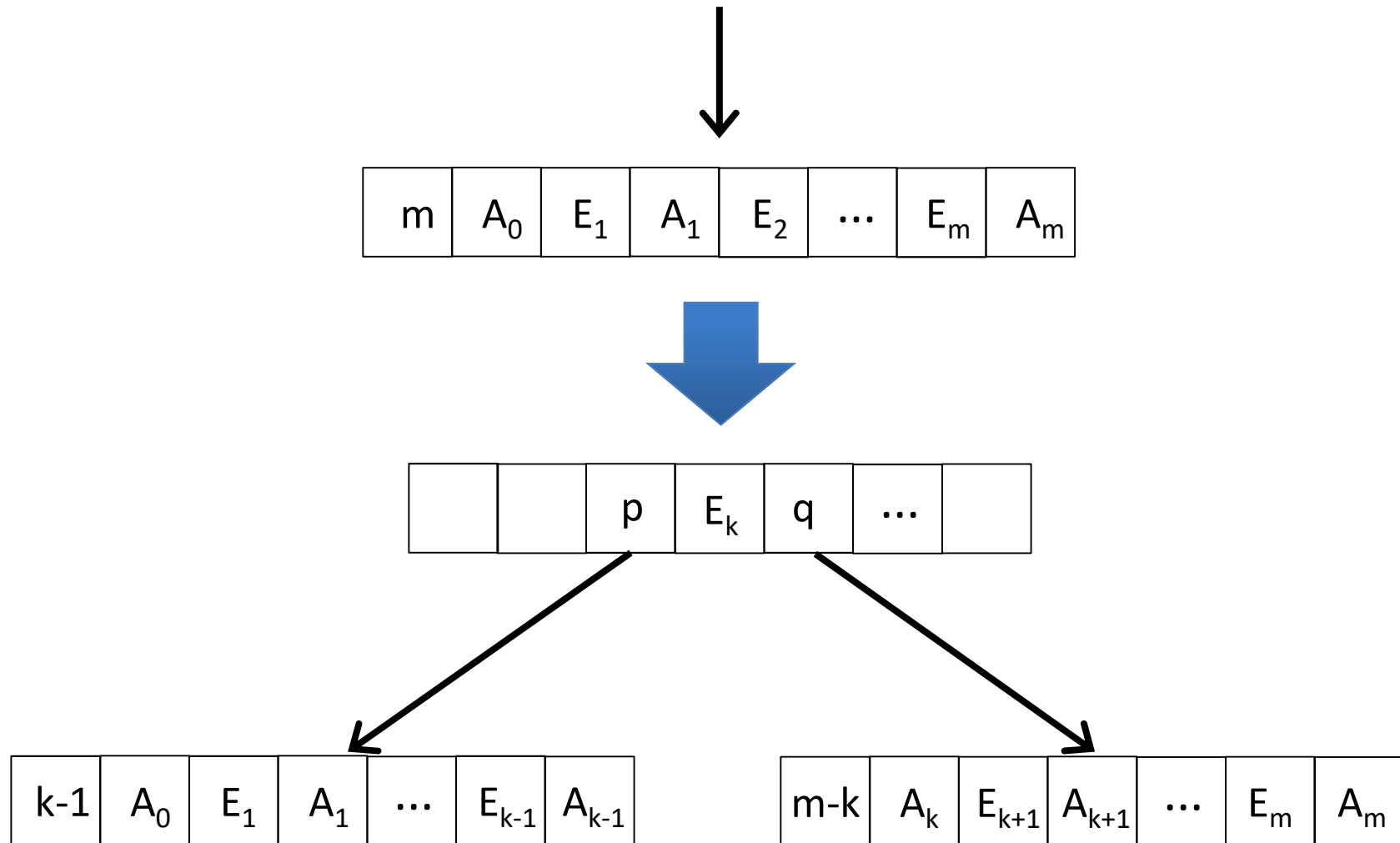


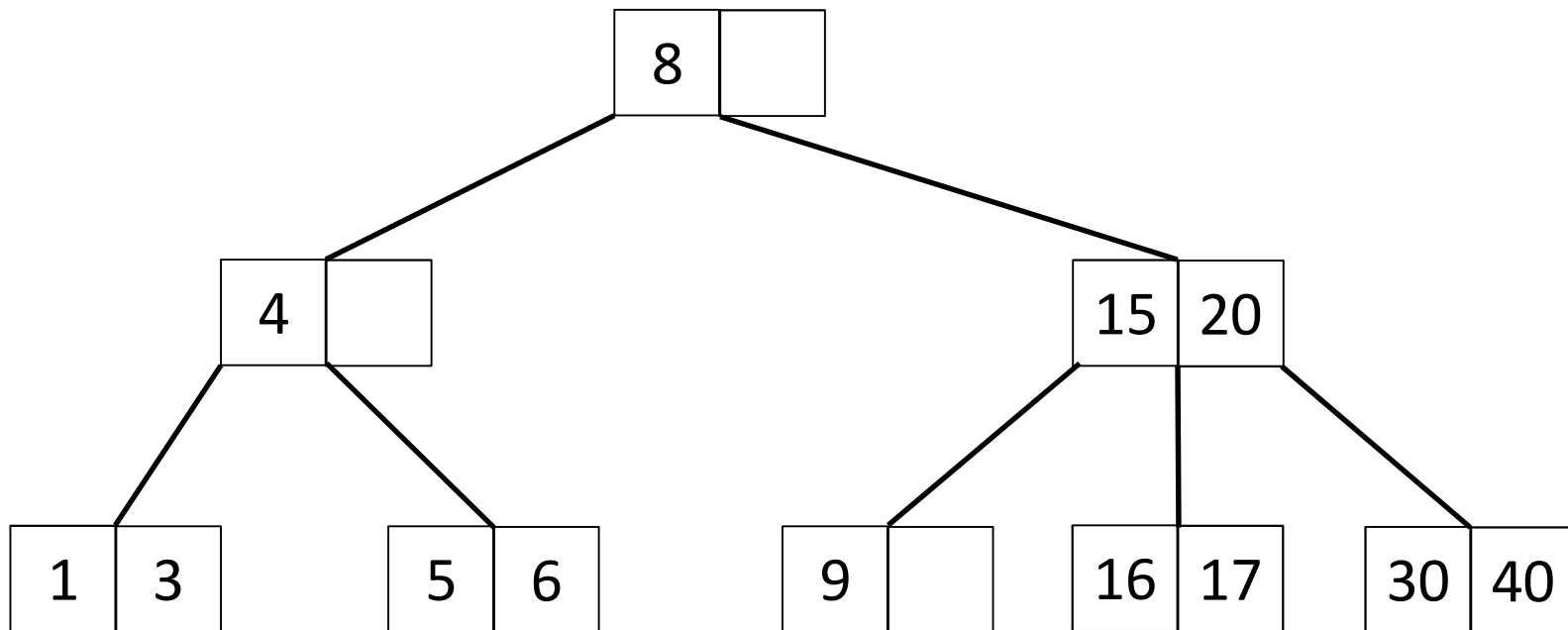
Insert

- If insertion results in overflow (already inserted with $m-1$ keys), split node
- Let node p have the format after insertion
 - $m, A_0, (E_1, A_1), \dots, (E_m, A_m)$
- p is split into two nodes p and q
 - Let $k = \text{ceil}(m/2)$
 - node p : $k-1, A_0, (E_1, A_1), \dots, (E_{k-1}, A_{k-1})$
 - node q : $m-k, A_k, (E_{k+1}, A_{k+1}), \dots, (E_m, A_m)$
 - (E_k, q) is inserted into the parent of p
- Splitting can propagate up to the root

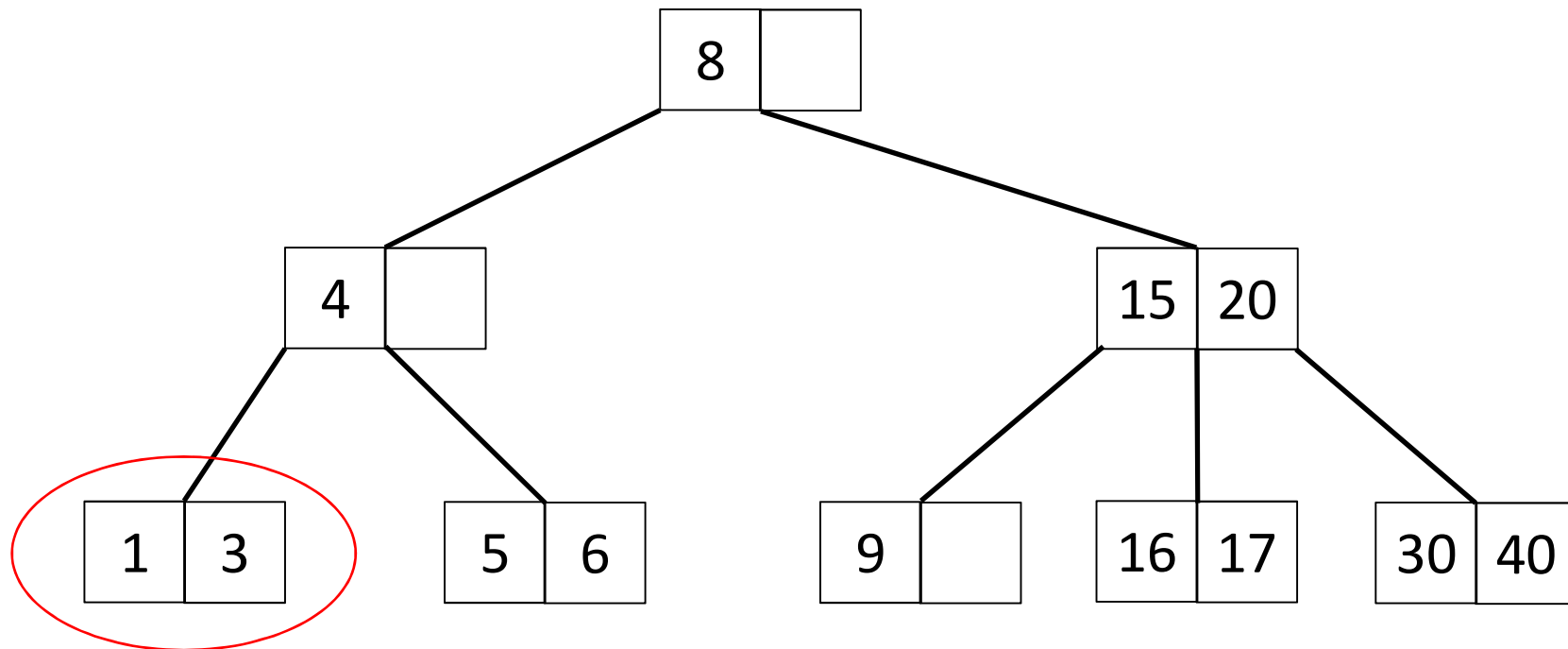
Split Node



Insert (3-way B-tree)

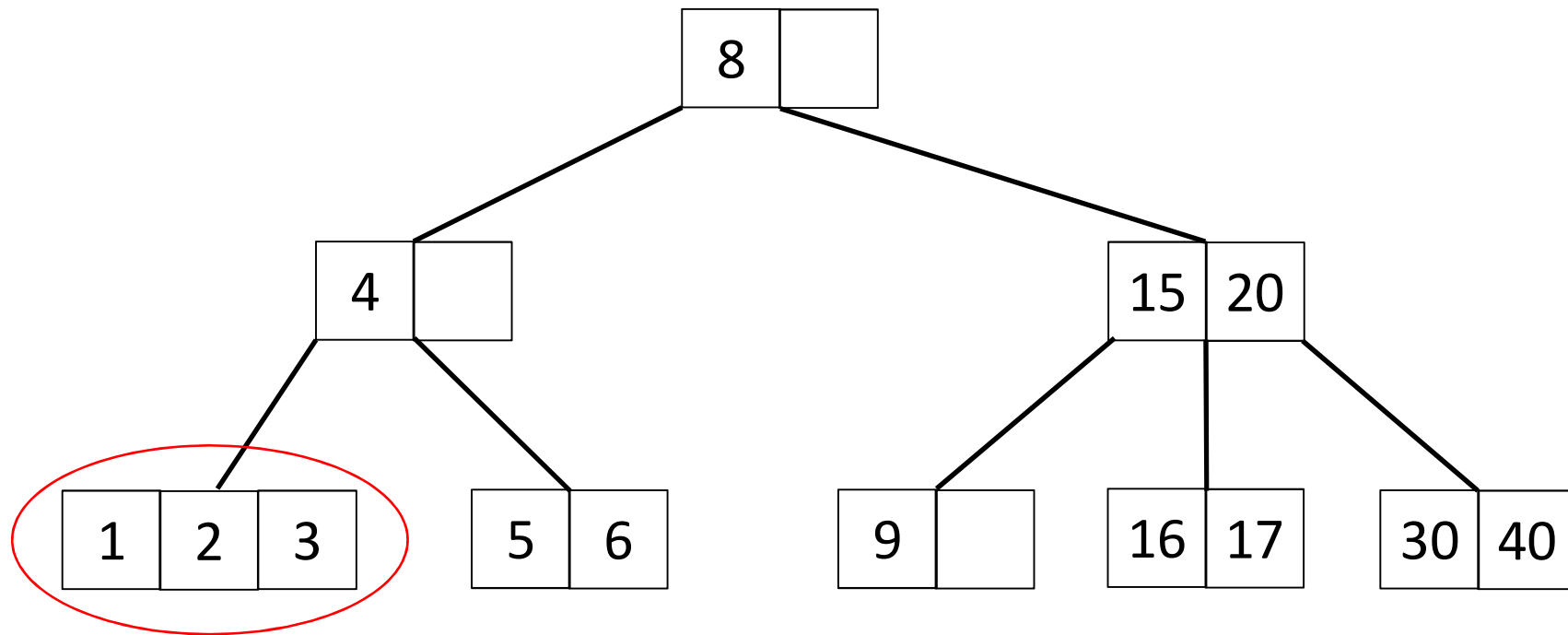


Insert (3-way B-tree)



Insert 2

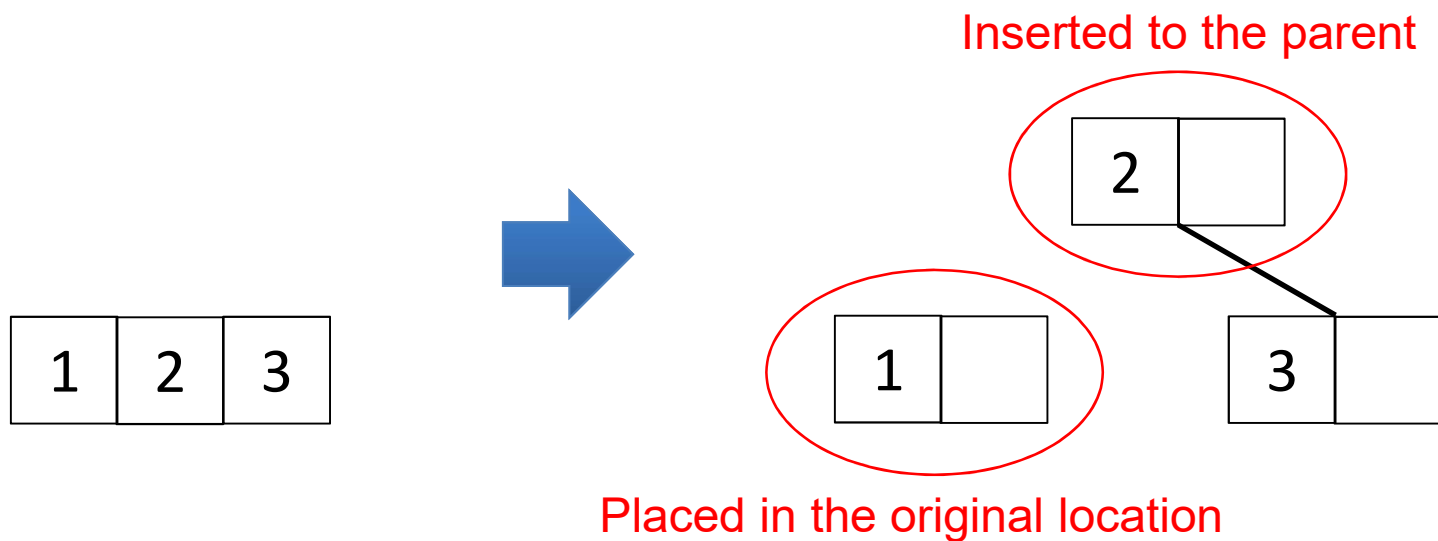
Insert (3-way B-tree)



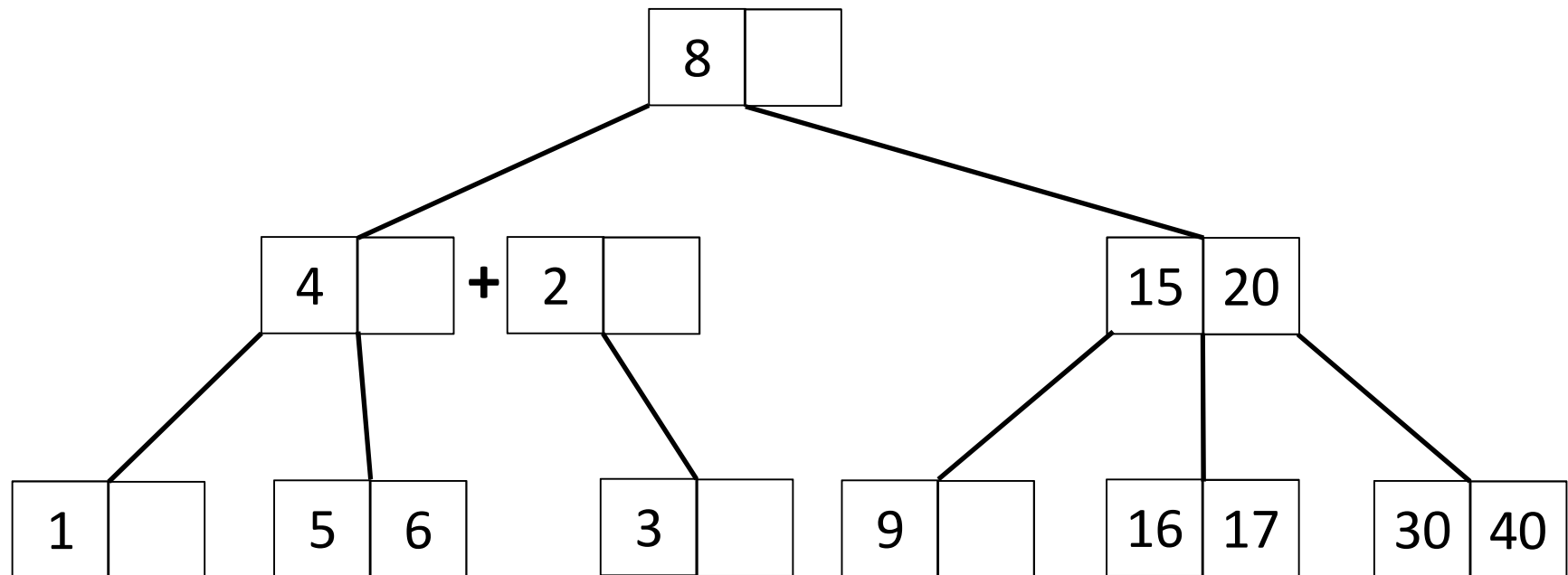
need split!

Insert (3-way B-tree)

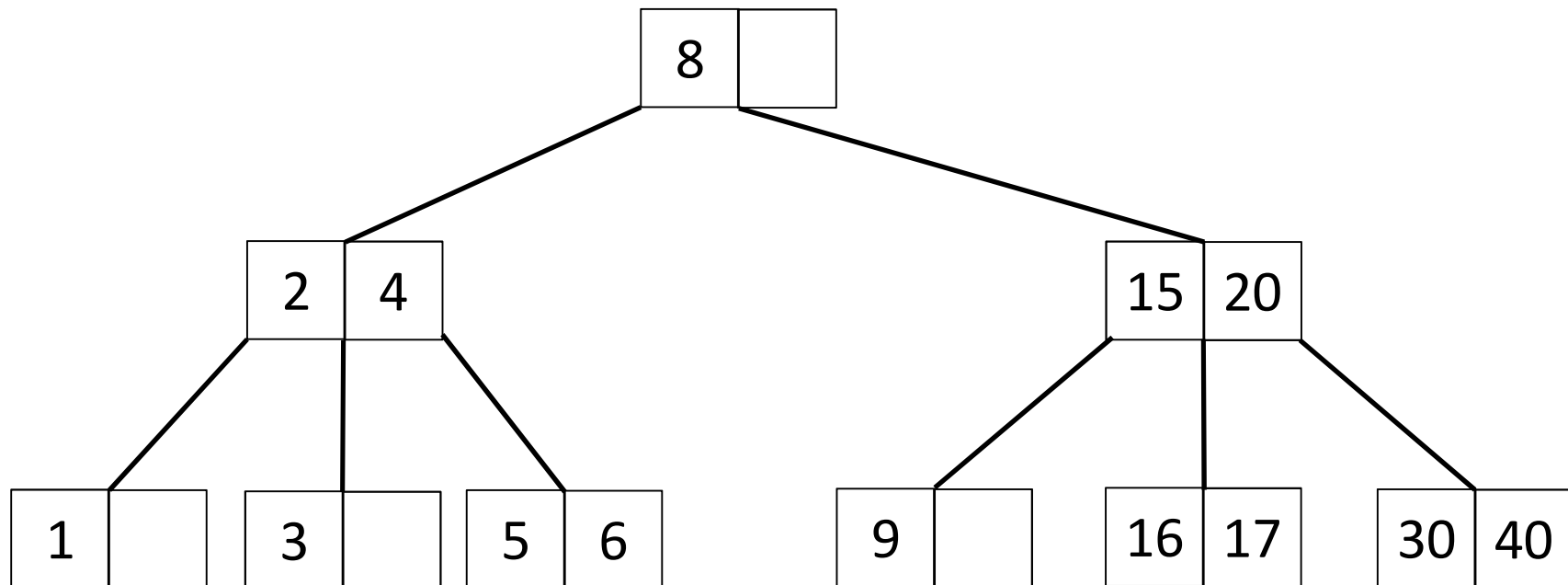
- Split overflowed node around middle key
- Insert middle key to its parent



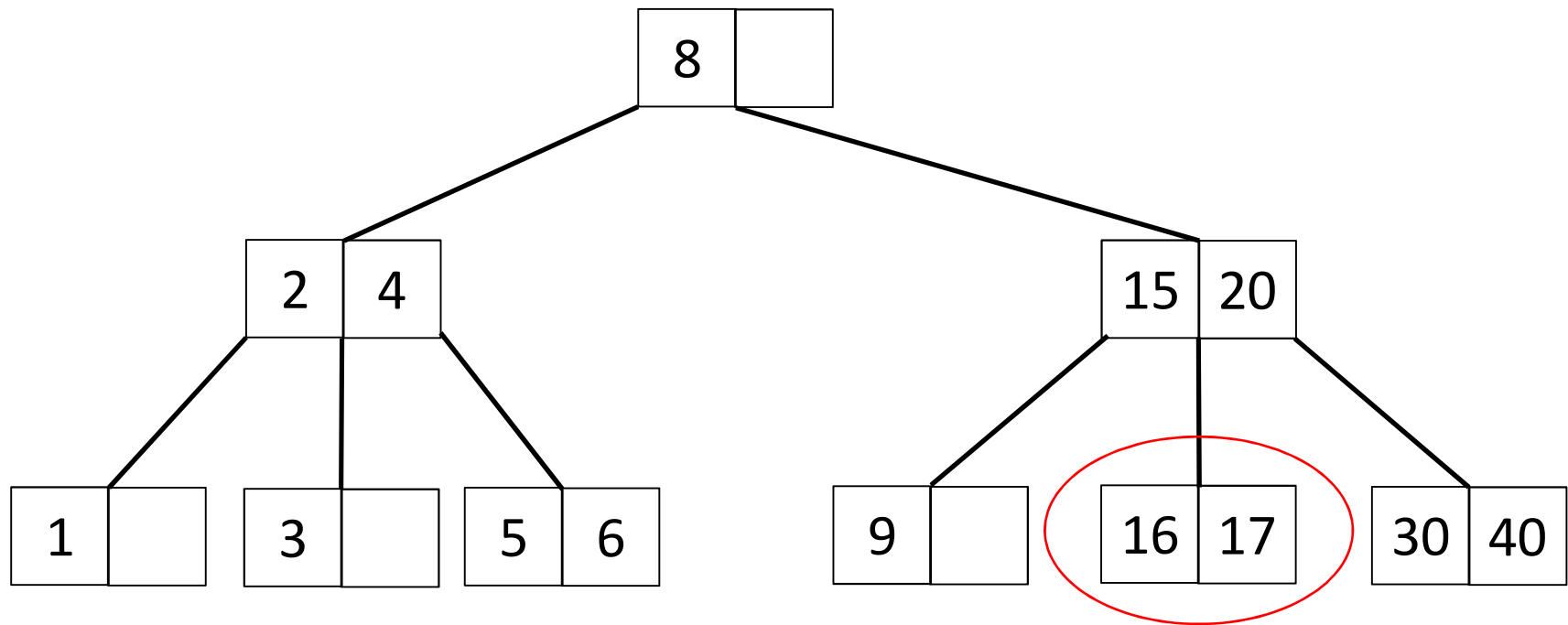
Insert (3-way B-tree)



Insert (3-way B-tree)

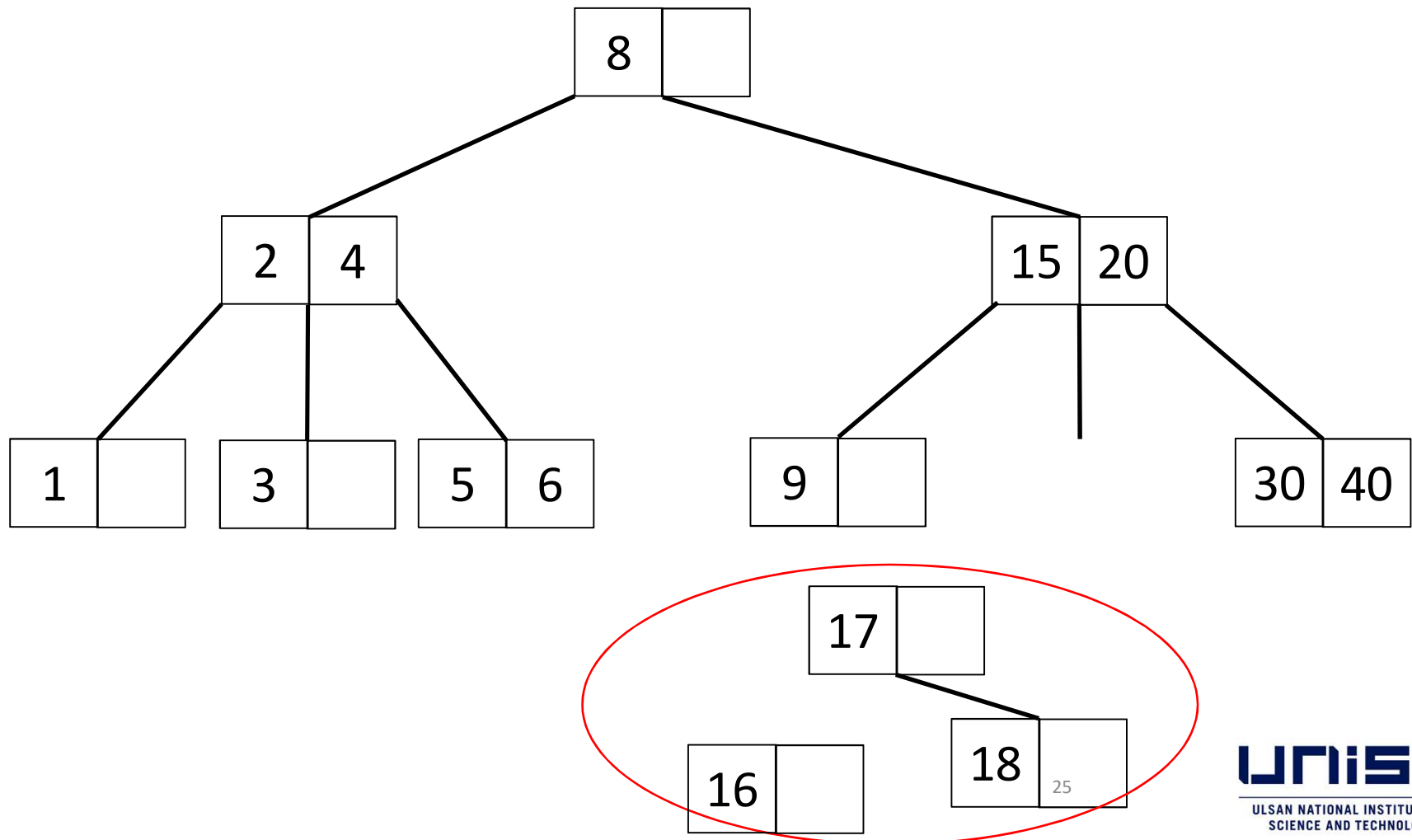


Insert (3-way B-tree)

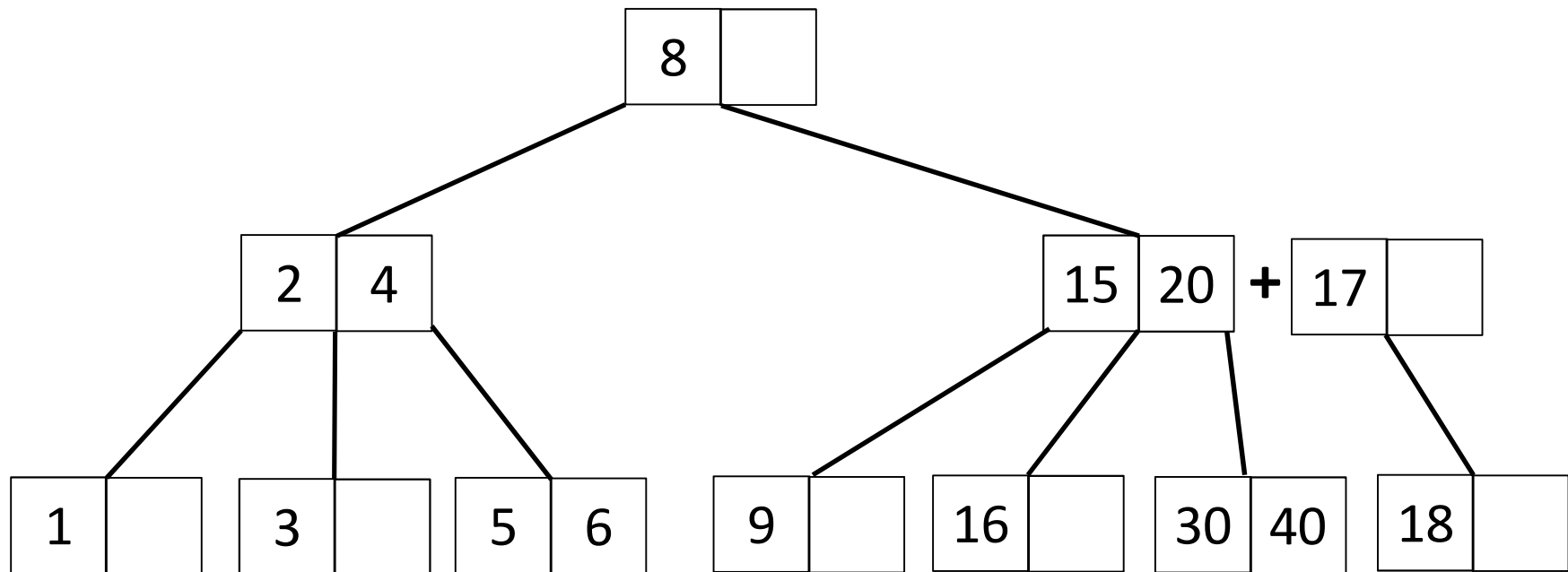


Insert 18

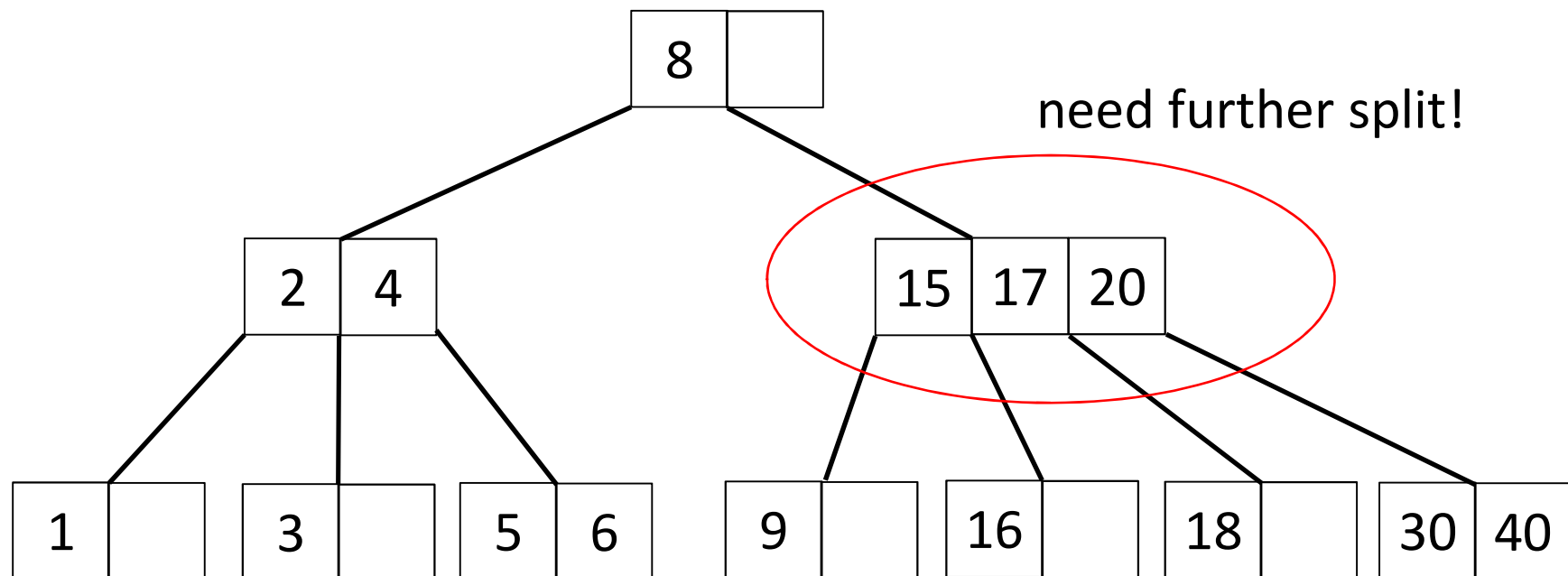
Insert (3-way B-tree)



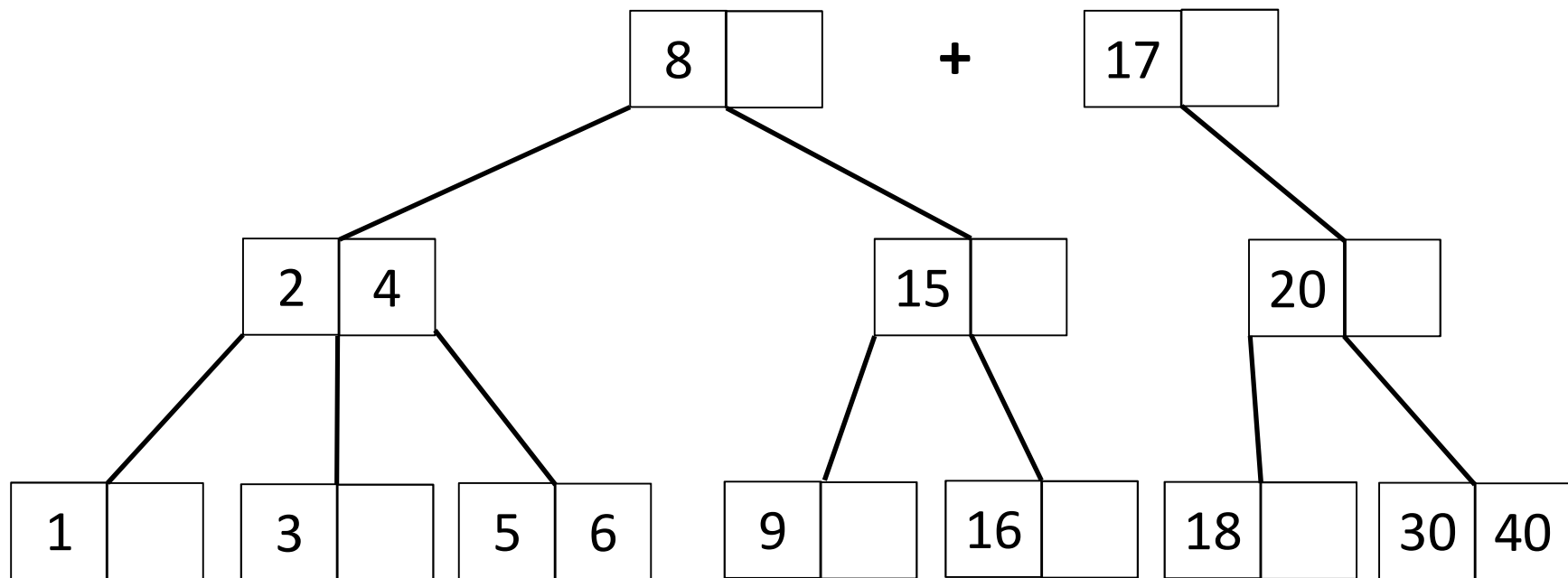
Insert (3-way B-tree)



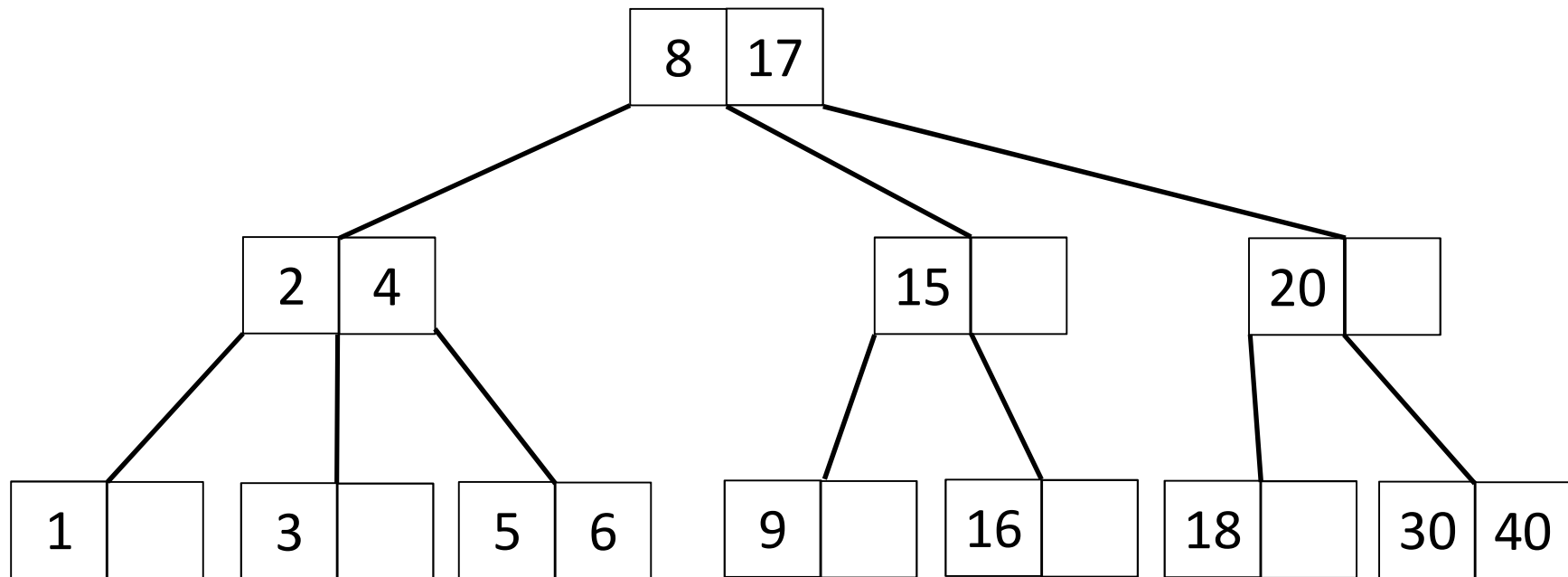
Insert (3-way B-tree)



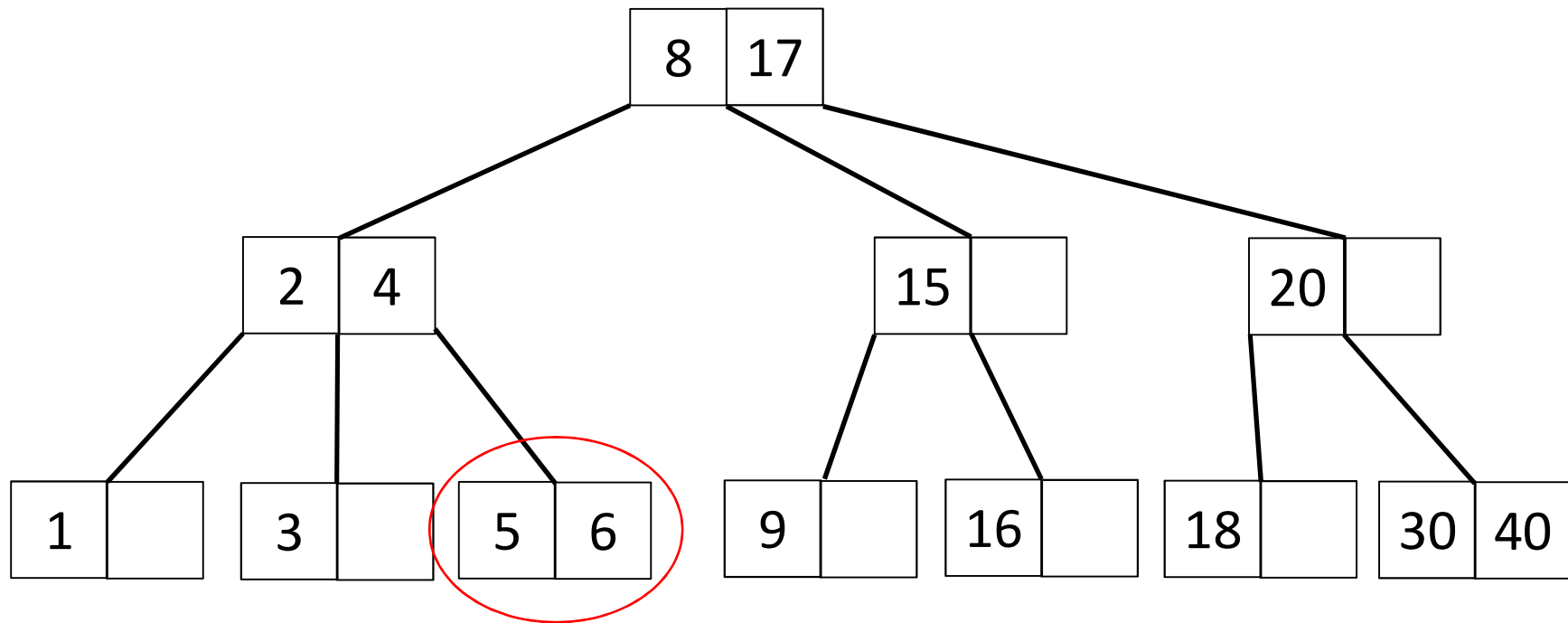
Insert (3-way B-tree)



Insert (3-way B-tree)

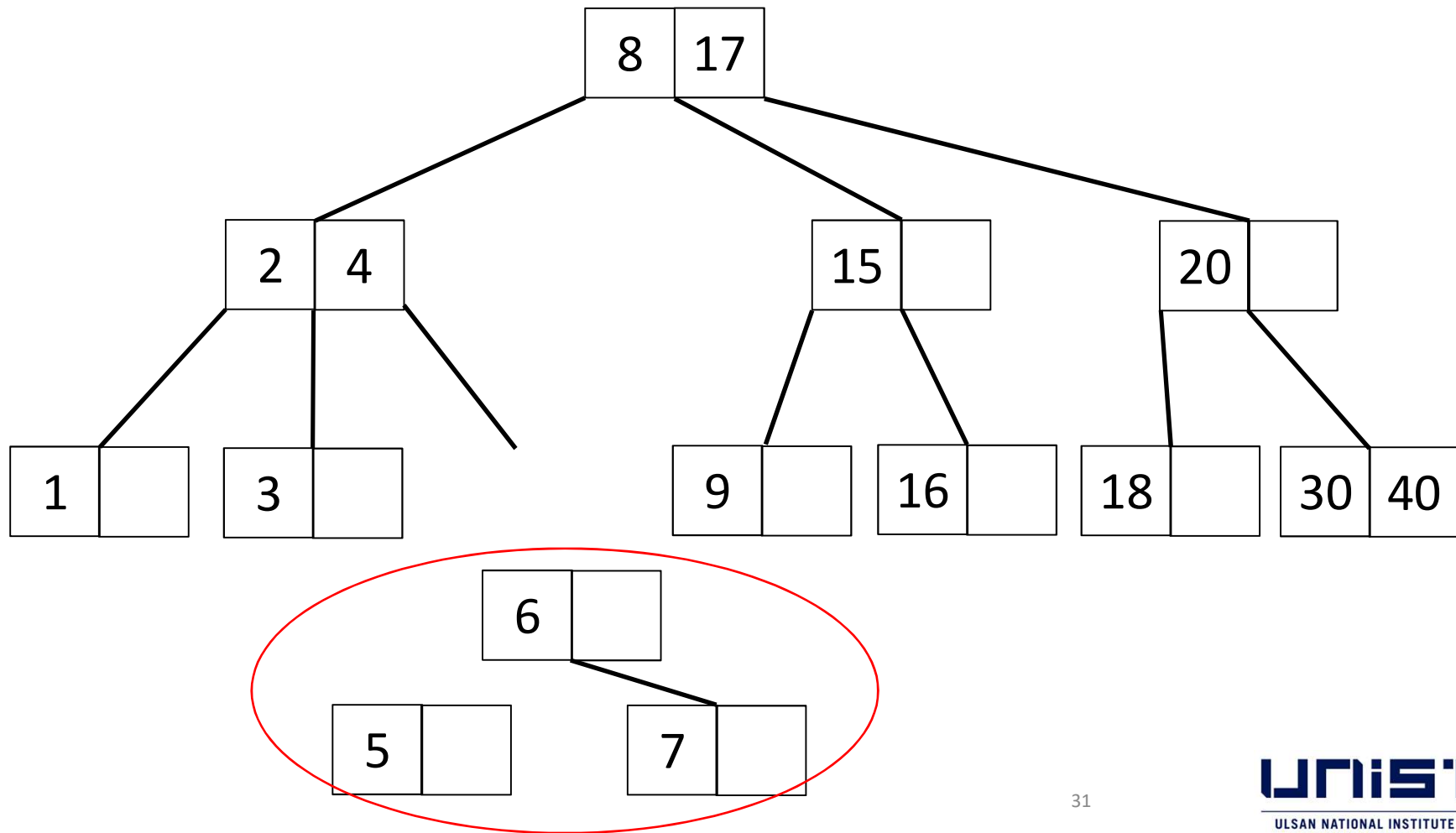


Insert (3-way B-tree)

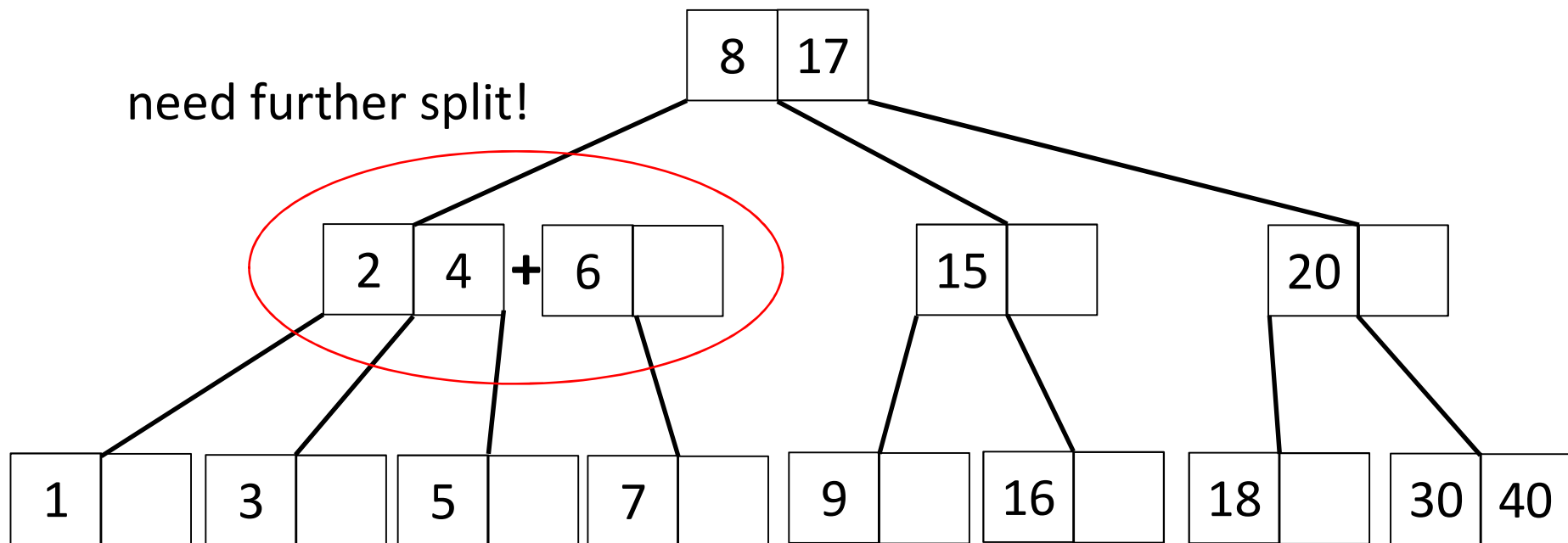


Insert 7

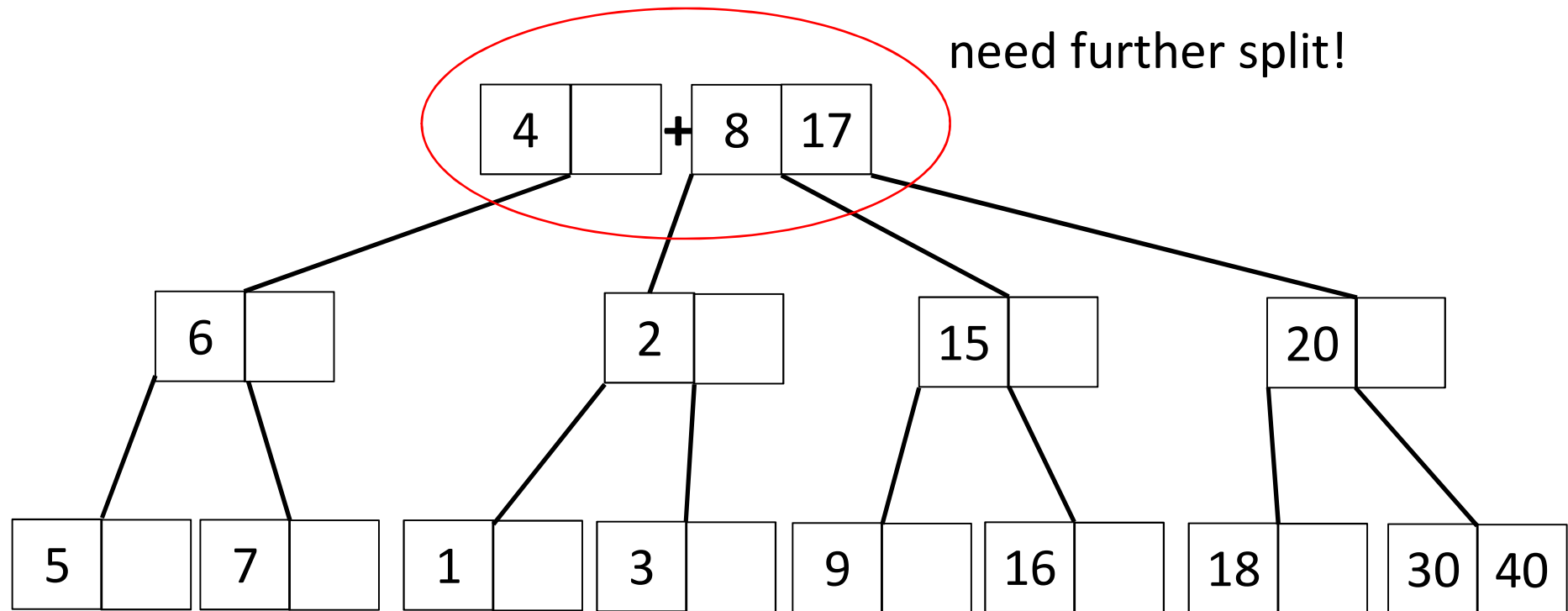
Insert (3-way B-tree)



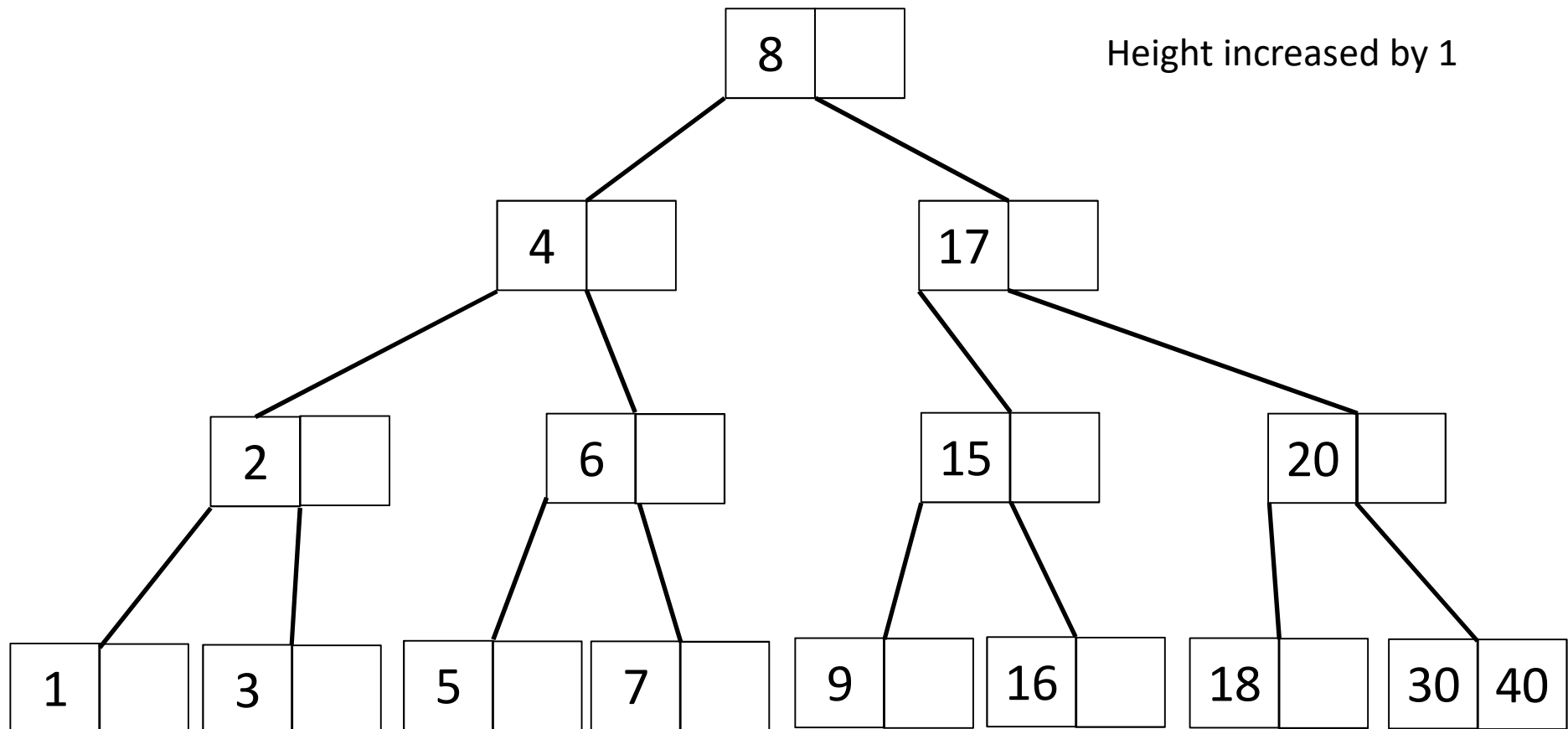
Insert (3-way B-tree)



Insert (3-way B-tree)



Insert (3-way B-tree)

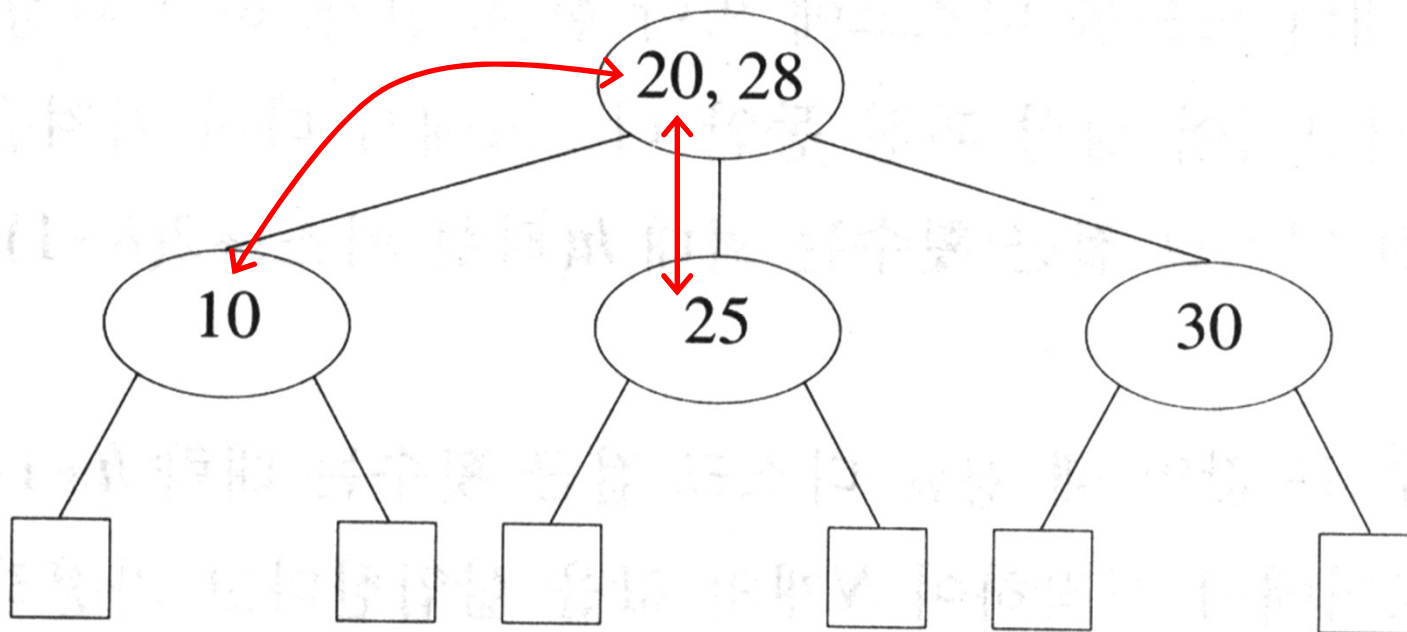


Deletion

- Delete from interior node
 - Replace with the largest in left subtree or the smallest in right subtree
 - The smallest/largest is in the leaf node
 - Deletion from an interior node is transformed into a **deletion from a leaf node**
 - If deletion results in less than $\text{ceil}(m/2)$ children, rotation or combine must be done

Example

- Delete 20
 - Replace with 10 or 25



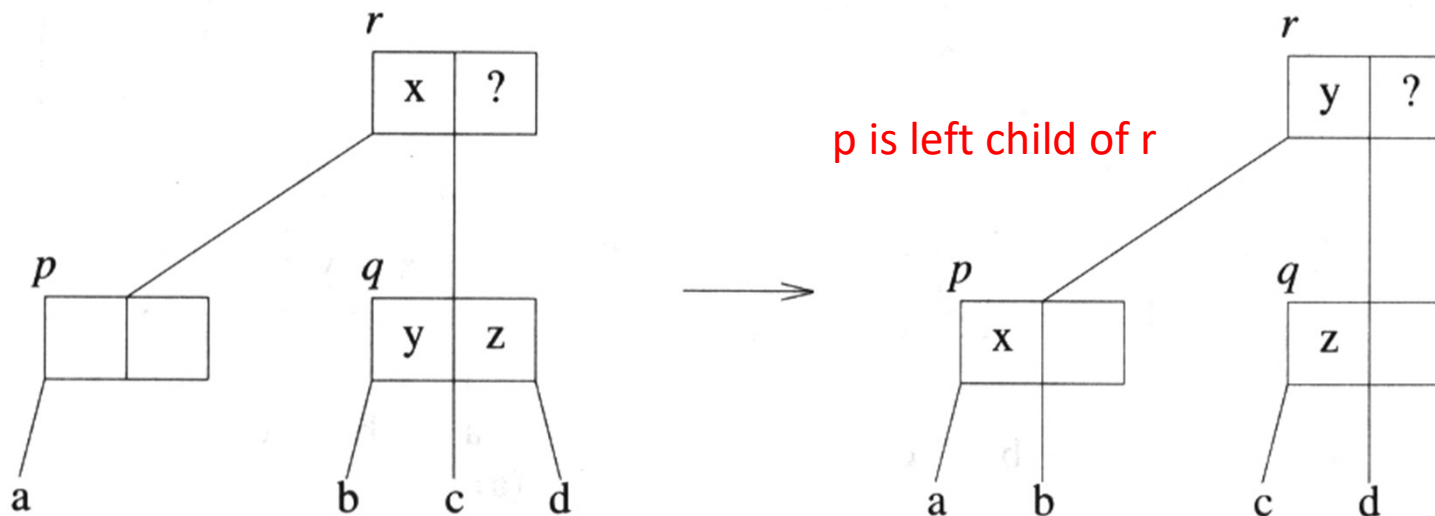
Deletion

Four cases when deleting an element from a leaf node p

1. p is root and left with at least one element after delete
 - OK: root is not empty = **at least two** children
2. p is internal and left with at least $\text{ceil}(m/2)-1$ elements after delete
 - OK: **$\text{ceil}(m/2)-1$ elements = $\text{ceil}(m/2)$ children**

Deletion

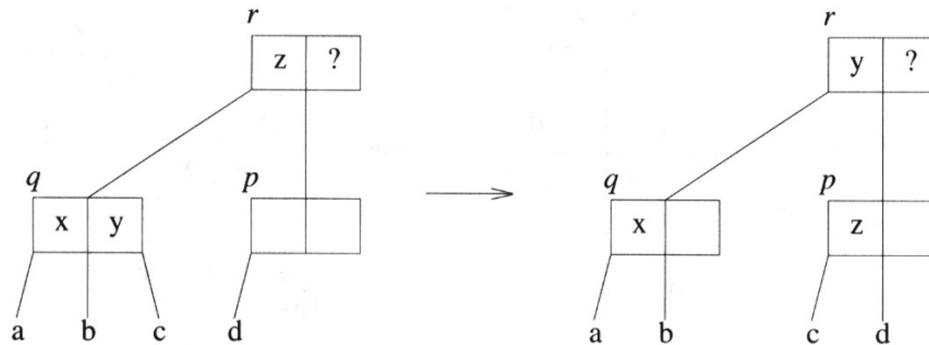
3. p has $\text{ceil}(m/2)-2$ elements and its sibling q has **at least $\text{ceil}(m/2)$** elements
–Rotation: **$\text{ceil}(m/2)-1$** elements in p



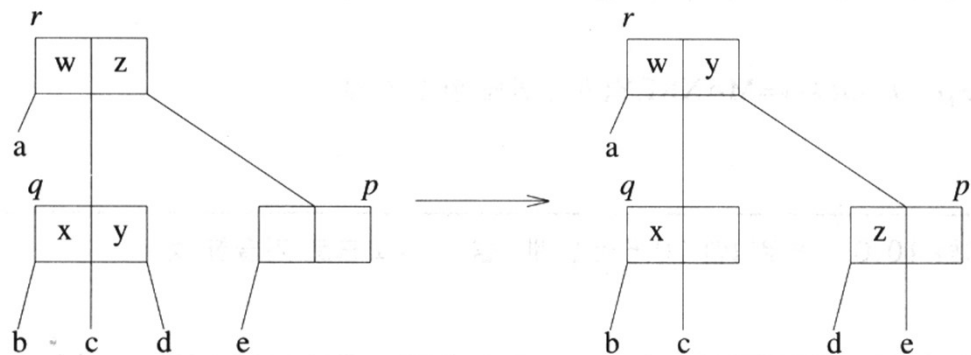
Order kept: $a < x < b < y < c < z < d$

Deletion

3. More rotation examples



p is middle child of r



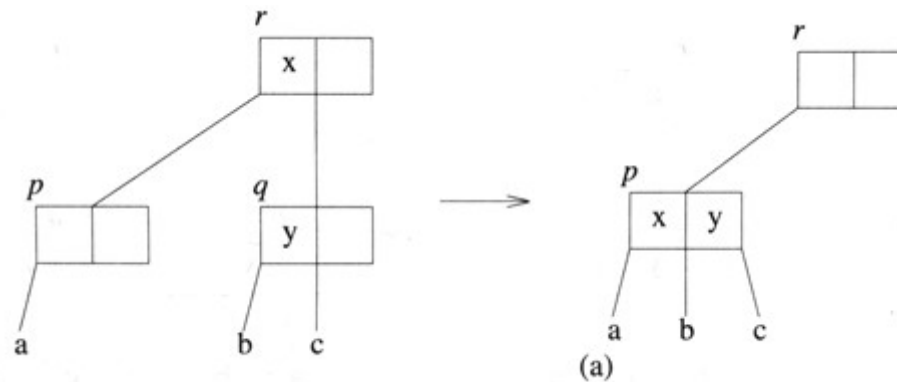
p is right child of r

Deletion

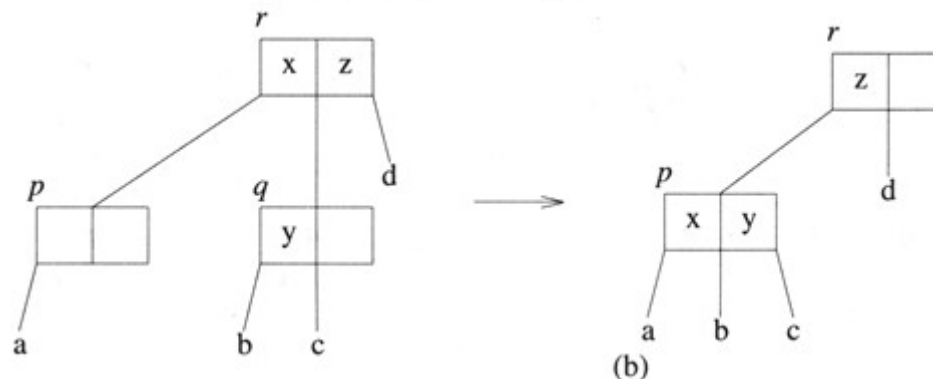
4. p has **$\text{ceil}(m/2)-2$** elements and its sibling q has **$\text{ceil}(m/2)-1$** elements
 - q has the minimum number of elements
 - Cannot rotate as we cannot reduce q's element
 - p and q are combined while borrowing a key (in-between element) from parent r
 - If r has $\text{ceil}(m/2)-2$ elements, rotation or combine is applied upward to the root

Deletion

4. p has $\text{ceil}(m/2)-2$ elements and its sibling q has $\text{ceil}(m/2)-1$ elements

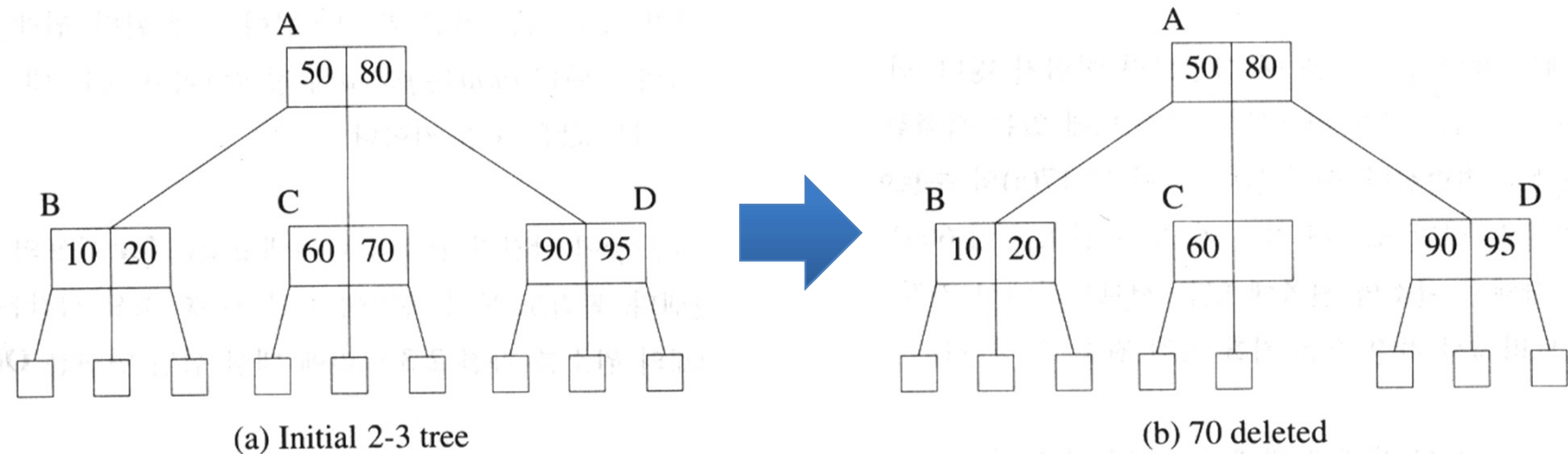


r has insufficient element
 \rightarrow rotation/combine is applied upward



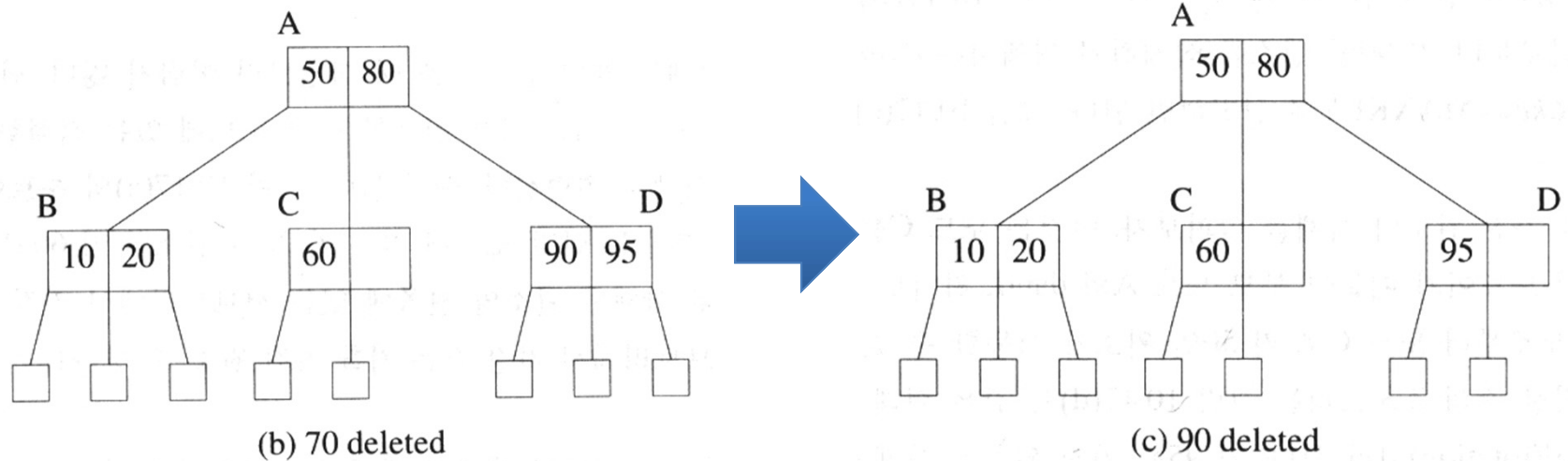
push x down
 \rightarrow p is left child of r

Example



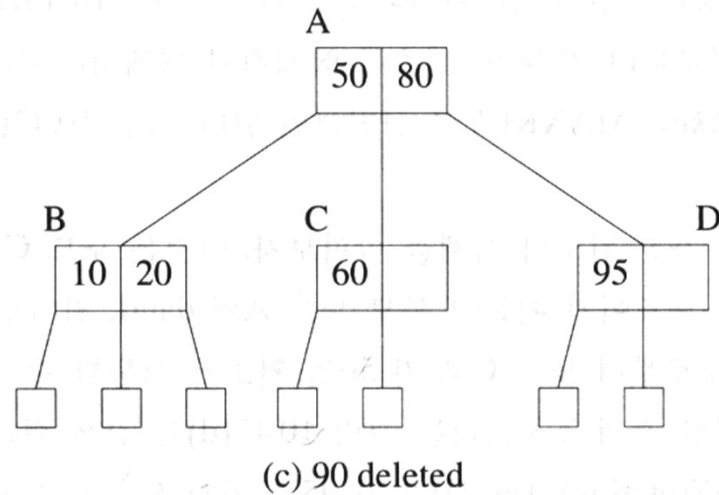
delete 70: node C has 1 element left, ok

Example

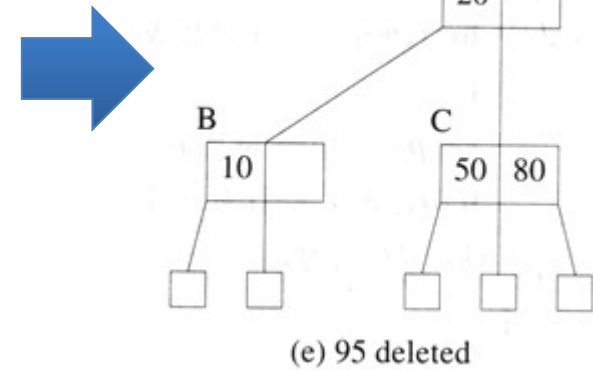
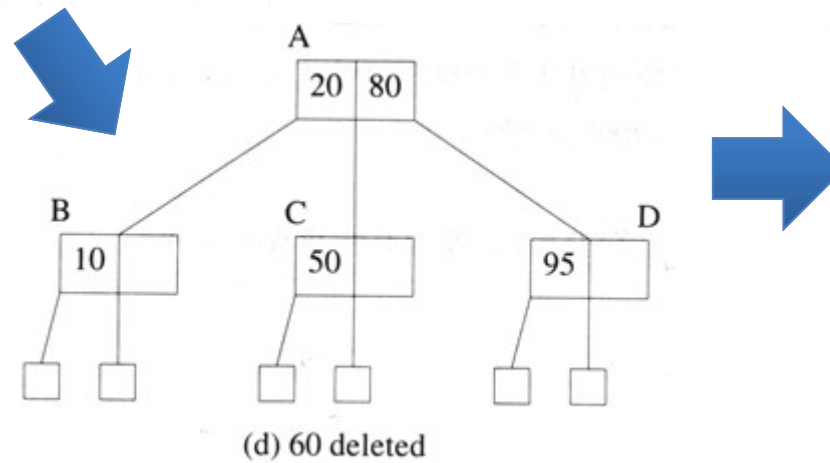


delete 90: node D has 1 element left, ok

Example

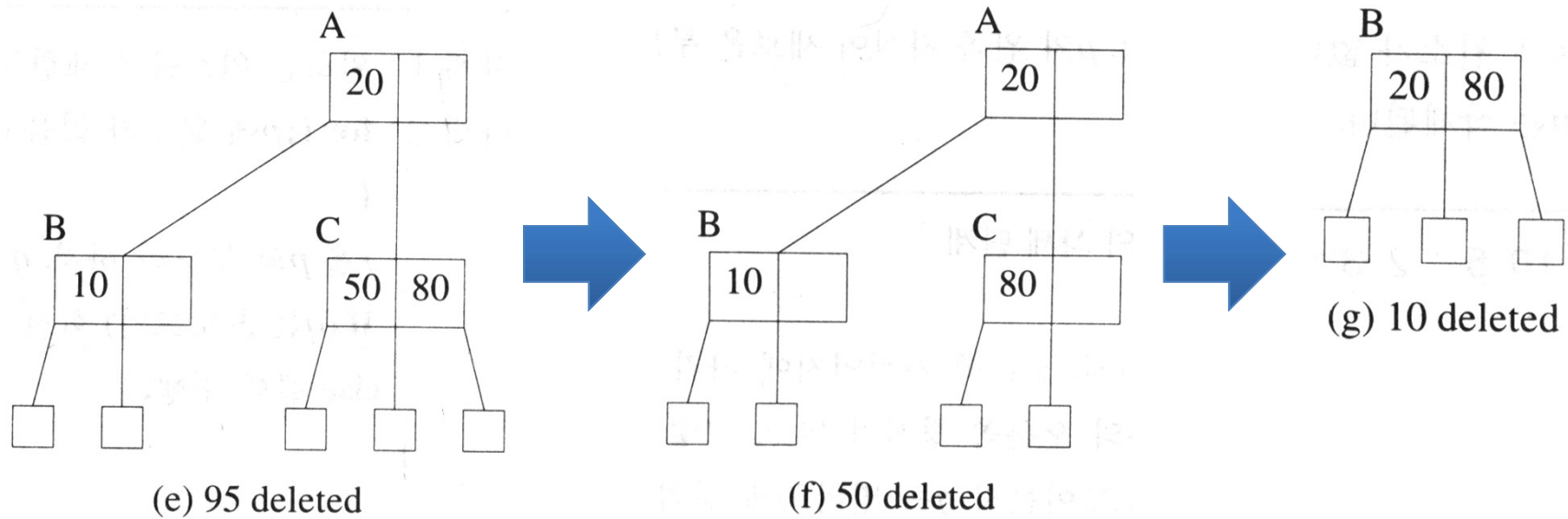


delete 95: combine
→ by pushing 80 down, we keep all
external nodes in the same level



delete 60: rotation right

Example



delete 50: node C has 1 element left, ok

delete 10: combine

Questions?