#### **HPCSE II**

**ScaLAPACK** and TBB

# Libraries for hybrid and multithreaded programming

### PLASMA

- Parallel Linear Algebra Software for Multicore Architectures
- Multi-threaded rewrite of LAPACK and BLAS functionality
- Optimizes for tiling/blocking and cache reuse on NUMA architectures
- Available from <a href="http://icl.cs.utk.edu/plasma">http://icl.cs.utk.edu/plasma</a>

### MAGMA

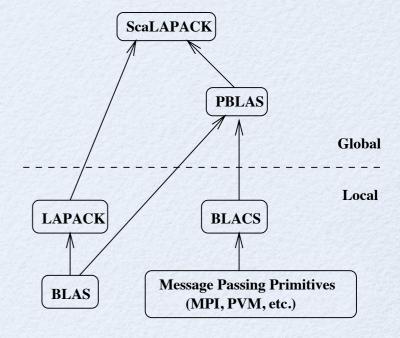
- Matrix Algebra on GPU and Multicore Architectures
- Hybrid CPU/GPU rewrite of LAPACK and BLAS functionality
- Available from <a href="http://icl.cs.utk.edu/magma">http://icl.cs.utk.edu/magma</a>

#### **ScaLAPACK**

- is a distributed memory extension of BLAS and LAPACK
- available from <a href="http://www.netlib.org/scalapack">http://www.netlib.org/scalapack</a>
- Builds upon
  - BLACS for communication abstraction over MPI, PVM or shared memory
  - BLAS and LAPACK for local computations
  - PBLAS

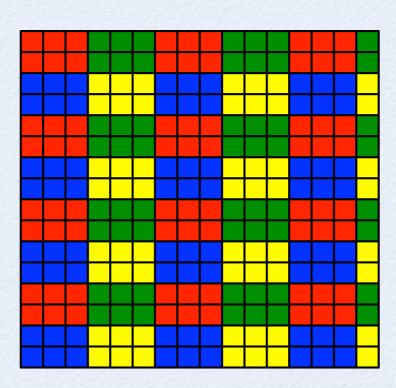
     a distributed BLAS using block-cyclic distributions

#### **ScaLAPACK Software Hierarchy**



#### Recall: block-cyclic distribution

- Block cyclic distribution
  - Example: 3x2 blocks, 2x2 process array
- Ideal block-sizes machine dependent
  - 32x32 or 64x64 are good starting guesses



#### Initializing the BLACS layer

 We do not need to initialize MPI ourselves but use the BLACS which typically builds on top of MPI

```
int main(int argc, char** argv)
{
  int rank;
  int nprocs;

  // initialize MPI using BLACS. Using MPI this is similar to the three statements below
  // MPI_Init(&argc, &argv);
  // MPI_Comm_size(MPI_COMM_WORLD, &nprocs);
  // MPI_Comm_rank(MPI_COMM_WORLD, &rank);
  Cblacs_pinfo(&rank,&nprocs);
  ...

  // we are done: call MPI_Finalize()
  Cblacs_exit(0);
  return 0;
}
```

#### Defining the process grid

Next initialize the process grid and find out who I am

```
// we want to use 2x3 processes
 int nprow=2;
 int npcol=3;
 // get the system context
 int ctxt;
 Cblacs_get(0,0,&ctxt);
// initialize a 2x3 process grid
 Cblacs gridinit(&ctxt, "Row-major", nprow, npcol);
 // get my coordinates in the process grid
 int myrow, mycol;
 Cblacs gridinfo(ctxt,&nprow,&npcol,&myrow,&mycol);
 // continue only if this rank is actually part of the grid
 if (myrow >= 0) {
   std::cout << "Rank " << rank << " has coordinates " << myrow << " " << mycol << "\n";</pre>
  . . .
else
   std::cout << "Rank " << rank << " is not used \n";</pre>
```

#### **Block-cyclic storage**

- We need to find the size of our local blocks
  - We will solve A\*X=B for a matrix A and a vector B, and

```
// we want a 100x 100 matrix
int n=100:
// and we want to solve with 1 right hand side
int nrhs = 1:
// we will use 32x32 size blocks in the block-cyclic layout
int nb=32;
// now intialize the matrix
// first calculate how many rows (np) and columns (ng) we store locally
int np = numroc_(n,nb,myrow,0,nprow);
int ng = numroc (n,nb,mycol,0,npcol);
int ngrhs = numroc_(nrhs, nb, mycol, 0, npcol);
// allocate local storage
hpc12::matrix<double> A(np.ng);
hpc12::matrix<double> B(np,ngrhs);
// create descriptors for the matrix and right hand side
int descA[9], descB[9], info;
descinit (descA, n, n , nb, nb, 0, 0, ctxt, A.leading dimension(), &info);
descinit_( descB, n, nrhs, nb, nb, 0, 0, ctxt, B.leading_dimension(), &info );
```

#### Finally call the solvers

- Solve the equation using pdgesv
- then use pdgemm to calculate the residual
- and finally use pdlange to get its 1-norm
  - pdlange needs work space. The size is taken from the manual.

## **Intel TBB**

#### **Intel Thread Building Blocks**

- A threading library by Intel, based on a draft version of C++11
  - implements C++11 threads plus a few extra locks and mutexes
  - higher level algorithmic abstractions
    - parallel\_for
    - parallel\_do
    - parallel while
    - tasks
  - thread-local data
  - thread-safe data structures

#### **Parallel for**

Consider a simple serial loop:

```
int main() {
  const int n = 10000;
  std::vector<double> x(n);

for (int i=0; i<n; ++i)
   x[i]=std::sin(i*0.01);
}</pre>
```

We want to run it in parallel using tbb::parallel\_for

```
#include <tbb/parallel_for.h>
#include <tbb/blocked_range.h>

int main() {

  const int n = 10000;
  std::vector<double> x(n);

  tbb::parallel_for( tbb::blocked_range<int>(0, n),
    [&] (tbb::blocked_range<int> const& r) {
    for( int i=r.begin(); i!=r.end(); ++i )
        x[i]=std::sin(i*0.01);
  });
}
```

#### Parallel for (continued)

The parallel\_for

```
tbb::parallel_for( tbb::blocked_range<int>(0, n), ...);
```

 needs a parallel function object that takes a blocked\_range. We used a C++11 lambda function

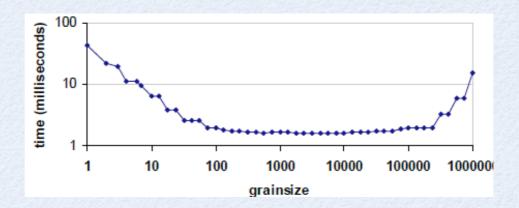
```
[&] (tbb::blocked_range<int> const& r) {
    for( int i=r.begin(); i!=r.end(); ++i )
        x[i]=std::sin(i*0.01);
}
```

TBB automatically launches threads and splits the range

#### **Controlling chunking**

- tbb::parallel\_for will be profitable when the loop runs for about 500µs
- Further optimization can be obtained by controlling the chunking:

```
tbb::parallel_for(tbb::blocked_range<int>(0,n,G), f, tbb::simple_partitioner());
```



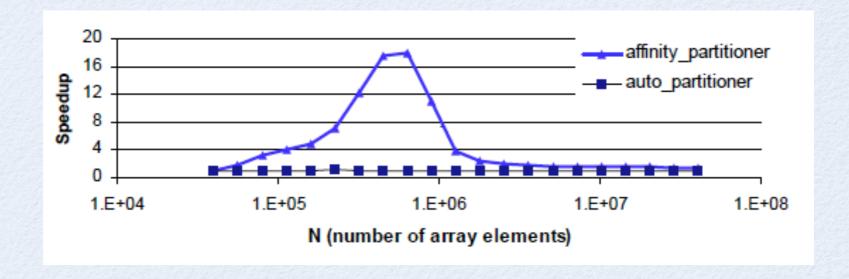
- The simple\_partitioner splits the loop into chunk sizes between G/2 and G
- The auto\_partitioner is the default
- The affinity\_partitioner can improve cache locality

#### The affinity partitioner

 When running a loop many times the affinity partitioner can improve cache affinity. The partitioner remembers which thread did which chunk

```
tbb::affinity_partitioner ap; // create a partitioner that is reused
for (int i=0; i< iterations; ++i)
  tbb::parallel_for(tbb::blocked_range<int>(0,n), f, ap);
```

• Intel shows substantial speedup is possible:



#### **Parallel reductions**

A reduction

```
double sum=0.;
for (int i=0; i<n; ++i)
  sum += x[i];</pre>
```

can be done by parallel\_reduce

```
sumit s(x);
tbb::parallel_reduce( tbb::blocked_range<int>(0, n), s);
```

 but now it needs a more complex function object and not just a lambda function

#### Parallel reduce (continued)

```
class sumit {
public:
  sumit(std::vector<double> const& x)
  : data (x)
  , sum_(0.)
  sumit(sumit const& x, tbb::split )
  : data (x.data )
  , sum_(0.)
  void join(sumit const& y ) {sum_ += y.sum_;}
  void operator()(tbb::blocked_range<int> const& r )
    double s = sum_; // remember to add to partial sum done locally
   for( int i=r.begin(); i!=r.end(); ++i )
      s += data [i];
   sum = s;
  double sum() const { return sum_;}
private:
  double sum ;
  std::vector<double> const& data_;
```

 The splitting constructor is used to make copies for other threads that should copy the data but not the partial sums

#### Thread-local data in TBB: combinable

```
template <typename T>
class combinable {
 public:
  combinable();
  template <typename Flnit>
  combinable(FInit finit);
                                    // constructed for each thread from return value
  void clear();
  T& local();
  T& local(bool & exists);
  template<typename FCombine>
  T combine(FCombine fcombine); // a reduction over all threads
  template<typename Func>
  void combine_each(Func f);
                                  // a function applied to each thread's value
};
```

#### The reduction using combinable

Use thread-local storage and combine it at the end

```
tbb::combinable<double> sum(0.);
tbb::parallel_for( tbb::blocked_range<int>(0, n),
  [&] (tbb::blocked_range<int> const& r) {
    double s=0.;
    for( int i=r.begin(); i!=r.end(); ++i )
        s +=x[i];
    sum.local() += s;
});
std::cout << "The sum is " << sum.combine(std::plus<double>()) << "\n";</pre>
```

# Thread-local data in TBB: enumerable\_thread\_specific

 is a container with one element per thread. Easiest to see in use, again for the reduction

```
tbb::enumerable_thread_specific<double> sum(0.);
tbb::parallel_for( tbb::blocked_range<int>(0, n),
  [&] (tbb::blocked_range<int> const& r) {
    double s=0.;
    for( int i=r.begin(); i!=r.end(); ++i )
        s +=x[i];
    sum.local() += s;
});
std::cout << "The sum is " << std::accumulate(sum.begin(),sum.end(),0.) << "\n";</pre>
```

#### More parallel operations

- TBB contains further and more flexible parallelization constructs:
  - Arbitrary iteration spaces
  - parallel\_do
    - Iteration through lists
    - Iteration through trees, offloading children to other threads
  - parallel\_pipeline
    - for pipelining various steps done on a continuous stream of input data
- When should one use TBB?
  - If you need finer control than OpenMP but don't want to manage the threads yourself manually
- Is it worth it?
  - That depends on whether you prefer to read manuals or code yourself.

#### Thread safe containers

- Thread safe containers circumvent the need to always lock a data structure manually for parallel access. TBB contains
  - concurrent\_hash\_map
  - concurrent\_vector
  - concurrent\_queue and concurrent\_bounded\_queue

#### **Example: concurrent\_vector**

- is similar to std::vector, but
  - might not be contiguous in memory
  - allows concurrent insertion by
    - v.push\_back(x);
    - v.grow\_by(n);
    - v.grow\_to\_at\_least(n);
  - allows concurrent iteration and check for size()

#### However, watch out:

- only iteration and insertion is thread-safe
- access to an element still needs to be synchronized manually
- calling clear() is not thread-safe

#### concurrent\_vector example

All threads can push safely into the same vector

```
#include <tbb/concurrent vector.h>
#include <tbb/blocked_range.h>
#include <tbb/parallel for.h>
#include <random>
#include <iostream>
int main() {
  tbb::concurrent vector<double> x;
  const int n = 100000;
  tbb::parallel_for( tbb::blocked_range<int>(0, n),
    [&] (tbb::blocked_range<int> const& r) {
      std::mt19937 mt:
      std::uniform_real_distribution<double> ureal_d(0.,10.);
      for( int j = r.begin(); j !=r.end(); ++j )
        x.push back(ureal d(mt)); // concurrent push back is safe!
      });
  std::cout << " The size is " << x.size() << "\n";</pre>
  return 0:
```

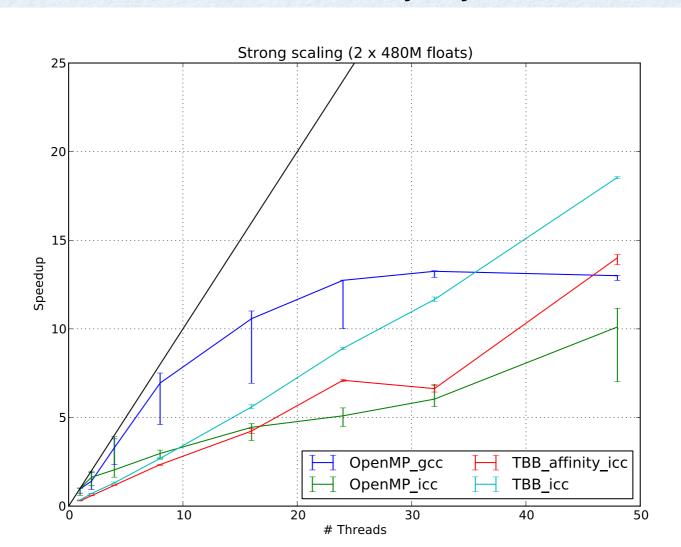
#### Task scheduler

- TBB provides tasks as an abstraction over thread
  - split a job into many tasks
  - TBB will schedule the tasks over the available threads
  - creating a task is 10-100 times faster than creating a thread

- How does it compare to OpenMP tasks?
  - Much harder to create and manage
  - But finer control over dependencies
- Is it worth it?

#### Some benchmarks by Andreas Hehn

- Multi-threaded vector multiplications, taking care of NUMA effects
- Lesson learned: benchmarks are the only way to tell what is fastest



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