COMP90048: Workshop 7

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Logic Programming, but better

Intro

 In Prolog, it can be easy to accidentally write very inefficient code.

There are two tricks: tail recursion and asymptotic complexity.

What happens when you call a function in C?

```
int n = 2;
int x;
x = factorial(n);
// x = 2

int factorial(int n) {
    return n <= 1
        ? 1
        : (n * factorial(n - 1);
}</pre>
```

The stack

What happens when you call a function in C?

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int factorial(int n) {
    return n <= 1
        ? 1
        : (n * factorial(n - 1);
}</pre>
```

n (4 bytes)	

^{*} it grows downwards by convention

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n	(4	bytes)
Х	(4	bytes)

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    return n <= 1
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}</pre>
```

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n (4 bytes)
```

What happens when you call a function in C?

```
int n = 2;
int x;

x = factorial(n);
// x = 2

the calling function's
stack frame => extra
memory used
? 1
: (n * factorial(n - 1);
}
```

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n (4 bytes)
```

What happens when you call a function in C?

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int n = 2;
int x;
x = factorial(n);
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}</pre>
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```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n (4 bytes)
```

What happens when you call a function in C?

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n (4 bytes)

return address (8 bytes)

n - 1 (4 bytes)
```

What happens when you call a function in C?

```
int n = 2;
int x;
x = factorial(n);
// x = 2

int factorial(int n) {
    return n <= 1
        ? 1
        : (n * factorial(n - 1);
}</pre>
```

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n (4 bytes)

return address (8 bytes)

n - 1 (4 bytes)
```

What happens when you call a function in C?

```
int n = 2;
int x;
x = factorial(n);
// x = 2

int factorial(int n) {
    return n <= 1
    ? 1
    : (n * factorial(n - 1);
}</pre>
```

The stack

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n (4 bytes)

return address (8 bytes)

n 1 (4 bytes)
```

popped

What happens when you call a function in C?

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n * 1 (4 bytes)
```

What happens when you call a function in C?

```
int n = 2;
int x;
x = factorial(n);
// x = 2

int factorial(int n) {
    return n <= 1
    ? 1
    : (n * factorial(n - 1);
}</pre>
```

n (4 bytes)
x (4 bytes)
return address (8 bytes)
n * 1 (4 bytes)

What happens when you call a function in C?

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int n = 2;
int x;
x = factorial(n);

// x = 2

int factorial(int n) {
    return n <= 1
          ? 1
          : (n * factorial(n - 1);
}</pre>
```

n	(4	bytes)
Х	(4	bytes)

We've got two stack entries per call...

```
int n = 2;
int x;
x = factorial(n);
// x = 2

int factorial(int n) {
    return n <= 1
        ? 1
        : (n * factorial(n - 1);
}</pre>
```

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n (4 bytes)

return address (8 bytes)

n - 1 (4 bytes)
```

Now with tail recursion!

```
int n = 2;
int x;
x = factorial(n);
// x = 2

int factorial(int n) {
    return n <= 1
        ? 1
        : (n * factorial(n - 1);
}</pre>
```

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n (4 bytes)
```

Now with tail recursion!

```
int n = 2;
int x;
x = factorial(n);
// x = 2

int factorial(int n) {
    return n <= 1
    ? 1
    : (n * factorial(n - 1);
}</pre>
```

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n (4 bytes)

n - 1 (4 bytes)
```

Now with tail recursion!

```
n (4 bytes)
x (4 bytes)
return address (8 bytes)
n (4 bytes)
n - 1 (4 bytes)
```

Now with tail recursion!

```
int n = 2;
int x;
x = factorial(n);
// x = 2

int factorial(int n) {
    return n <= 1
          ? 1
          : (n * factorial(n - 1);
}</pre>
```

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n * 1 (4 bytes)

n - 1 (4 bytes)
```

Now with tail recursion!

```
int n = 2;
int x;
x = factorial(n);
// x = 2

int factorial(int n) {
    return n <= 1
        ? 1
        : (n * factorial(n - 1);
}</pre>
```

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n * 1 (4 bytes)

n 1 (4 bytes)
```

By reducing stack use, we save both memory and time.

```
int n = 2;
int x;
x = factorial(n);

// x = 2

int factorial(int n) {
    return n <= 1
        ? 1
        : (n * factorial(n - 1);
}</pre>
```

n	(4	bytes)
X	(4	bytes)

By reducing stack use, we save both memory and time.

The stack (no tail recursion)

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n (4 bytes)

return address (8 bytes)

n - 1 (4 bytes)
```

The stack (tail recursion)

```
n (4 bytes)

x (4 bytes)

return address (8 bytes)

n * 1 (4 bytes)

n - 1 (4 bytes)
```

We can do this because recursion is the last step.

```
int factorial(int n) {
   return n <= 1
   ? 1
   : (n * factorial(n - 1);
}</pre>
```

Tail recursion in Prolog

• We introduce an **accumulator** (and a helper predicate).

```
prodlist(List, Prod):-
    prodlist_acc(List, 1, Prod).

prodlist_acc([], Prod, Prod).
% to be continued
```

 The idea: prodlist_acc holds when its accumulator is the product.

Tail recursion in Prolog

• We introduce an **accumulator** (and a helper predicate).

```
prodlist(List, Prod):-
        prodlist_acc(List, 1, Prod).

prodlist_acc([], Prod, Prod).
prodlist_acc([N|Ns], Acc, Prod):-
        Prod0 is Acc * N,
        prodlist_acc(Ns, Prod0, Prod).
```

• Similarly to Haskell, append/3 is O(n). We try to avoid it where we can. Consider:

```
reverse([], []).

reverse([A|BC], CBA):-
    reverse(BC, CB),
    append(CB, [A], CBA).
```

We can use the accumulator approach here too!

```
reverse(ABC, CBA):-
    reverse_acc(ABC, [], CBA).

reverse_acc([], A, A).

DCBAcc is the reversed list, with Acc on the end

reverse_acc([B|CD], Acc, DCBAcc):-
    reverse_acc(CD, [B|Acc], DCBAcc).
```

We can use the accumulator approach here too!

```
reverse(ABC, CBA):-
    reverse_acc(ABC, [], CBA).

reverse_acc([], A, A).

DCBAcc is the reversed list, with Acc on the end
reverse_acc([B|CD], Acc, DCBAcc):-
    reverse_acc(CD, [B|Acc], DCBAcc).

peel B from the front of the original list BCD and
```

stick it at the back of the new list Acc

We can use the accumulator approach here too!

```
reverse([], A, A).

reverse([B|CD], Acc, DCBAcc):-
    reverse(CD, [B|Acc], DCBAcc).
```

This is a general method for working with lists.

Using the Prolog debugger to understand reverse_acc/3

```
[trace] ?- reverse([a, b, c], List).
    Call: (8) reverse([a, b, c], _6824) ? creep
    Call: (9) reverse_acc([a, b, c], [], _6824) ? creep
    Call: (10) reverse_acc([b, c], [a], _6824) ? creep
    Call: (11) reverse_acc([c], [b, a], _6824) ? creep
    Call: (12) reverse_acc([], [c, b, a], _6824) ? creep
    Exit: (12) reverse_acc([], [c, b, a], [c, b, a]) ? creep
    Exit: (11) reverse_acc([c], [b, a], [c, b, a]) ? creep
    Exit: (10) reverse_acc([b, c], [a], [c, b, a]) ? creep
    Exit: (9) reverse([a, b, c], [], [c, b, a]) ? creep
    Exit: (8) reverse([a, b, c], [c, b, a]) ? creep
List = [c, b, a].
```

- peel a off the front

Using the Prolog debugger to understand reverse acc/3

```
[trace] ?- reverse([a, b, c], List).
    Call: (8) reverse([a, b, c], _6824) ? creep
    Call: (9) reverse_acc([a, b, c], [], _6824) ? creep
    Call: (10) reverse_acc([b, c], [a], _6824) ? creep
    Call: (11) reverse_acc([c], [b, a], _6824) ? creep
    Call: (12) reverse_acc([], [c, b, a], _6824) ? creep
    Exit: (12) reverse_acc([], [c, b, a], [c, b, a]) ? creep
    Exit: (11) reverse_acc([c], [b, a], [c, b, a]) ? creep
    Exit: (10) reverse_acc([b, c], [a], [c, b, a]) ? creep
    Exit: (9) reverse([a, b, c], [], [c, b, a]) ? creep
    Exit: (8) reverse([a, b, c], [c, b, a]) ? creep
List = [c, b, a].
```

← peel a off the front

accumulator done

Using the Prolog debugger to understand reverse acc/3

```
[trace] ?- reverse([a, b, c], List).
    Call: (8) reverse([a, b, c], _6824) ? creep
    Call: (9) reverse_acc([a, b, c], [], _6824) ? creep
    Call: (10) reverse_acc([b, c], [a], _6824) ? creep
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    Call: (12) reverse_acc([], [c, b, a], _6824) ? creep
    Exit: (12) reverse_acc([], [c, b, a], [c, b, a]) ? creep
    Exit: (11) reverse_acc([c], [b, a], [c, b, a]) ? creep
    Exit: (10) reverse_acc([b, c], [a], [c, b, a]) ? creep
    Exit: (9) reverse_acc([a, b, c], [], [c, b, a]) ? creep
    Exit: (8) reverse([a, b, c], [c, b, a]) ? creep
List = [c, b, a].
```

peel a off the front

accumulator done

unify accumulator with List

Continue with Grok Workshop 7 (Week 8).