## 1 C Memory Management

- 1. Match the items on the left with the memory segment in which they are stored. Answers may be used more than once, and more than one answer may be required.
  - 1. Static variables B
  - 2. Local variables D
  - 3. Global variables B
  - 4. Constants A, B, D
  - 5. Machine Instructions A
  - 6. Result of malloc() C
  - 7. String Literals B

- A. Code
- B. Static
- C. Heap
- D. Stack

4 Explained: With DEFINE, you can replace something in the code. With const declaration, it's stored in the static or stack depending on being declared in a function or not.

- 2. Write the code necessary to properly allocate memory (on the heap) in the following scenarios
  - 1. An array  $\operatorname{arr}$  of k integers

```
arr = (int *) malloc(sizeof(int) * k);
```

2. A string str containing p characters

```
str = (char *) malloc(sizeof(char) * (p + 1)); // Don't forget the null terminator!
```

3. An  $n \times m$  matrix mat of integers initialized to zeros

```
mat = (int *) calloc(n*m, sizeof(int));
```

Alternative solution. This might be needed if you wanted to efficiently permute the rows of the matrix.

```
mat = (int **) calloc(n, sizeof(int *));
for (int i = 0; i < n; i++) {
    mat[i] = (int *) calloc(m, sizeof(int));
}</pre>
```

3. What is wrong with the C code below?

```
int* pi = malloc(314 * sizeof(int));
if(!raspberry) {
   pi = malloc(1 * sizeof(int)); // Memory leak if not raspberry
}
return pi;
```

4. Write code to prepend (add to the start) to a linked list, and to free/empty the entire list. struct ll\_node { struct ll\_node\* next; int value; }

```
void prepend(struct ll_node** lst, int val) void free_ll(struct ll_node** lst)

struct ll_node* item = (struct ll_node*)
    malloc(sizeof(struct ll_node));
    item->value = val;
    item->next = *lst;
    *lst = item;
    *lst = NULL;
if(*lst) {
    free_ll(&((*lst)->next));
    free(*lst);
}

*lst = NULL;
```

Note: \*1st points to the first element of the list, or is NULL if the list is empty.

## 2 Data Structures in C

In this question, we will implement a array-based stack of integers in C. The stack will be represented by the struct below.

```
struct stack {
        int size;
                         // Number of element in stackArray
                        // Index of the array that is the top of the stack
        int topIndex;
        int *stackArray; // Array holding the elements of the stack
    }
Implement the functions below.
    // Create a new stack with the given array size
    struct stack *init_stack(int size) {
        struct stack *stk = (struct stack *) malloc(sizeof(struct stack));
        if (!stk) {
            return NULL;
        }
        stk->size = size;
        stk->topIndex = -1;
        stk->stackArray = (int *) malloc(sizeof(int) * size);
        if (!(stk->stackArray)) {
            free(stk);
            return NULL;
        return stk;
   }
    // Add the given element to the stack. Resize by doubling the array size if full.
    // Return 1 on success, 0 on failure. stk should be unchanged on failure.
    int push(int x, struct stack* stk) {
        if (stk->topIndex == stk->size - 1) {
            int *stackArray = (int *) realloc(stk->stackArray, stk->size * 2)
            if (!stackArray) {
                return 0;
            stk->stackArray = stackArray;
            stk->size = stk->size * 2;
        }
        stk->topIndex++;
        stk->stackArray[stk->topIndex] = x;
   }
    // Remove the top element from the stack. Return 0 if empty.
    int pop(struct stack *stk) {
        if (stk->topIndex < 0) {</pre>
            return 0;
        }
        int x = stk->stackArray[stk->topIndex];
        // Setting to 0 is not strictly needed, but can be useful for
        // debugging/assertions
        stk->stackArray[stk->topIndex] = 0;
        stk->topIndex--;
        return x;
```

## 3 RISC-V Intro

1. Assume we have an array in memory that contains int\* arr = {1,2,3,4,5,6,0}. Let the value of arr be a multiple of 4 and stored in register s0. What do the snippets of RISC-V code do? Note that these snippets all run immediately after each other on the same core (snippet a runs, then snippet b and so on).

```
a) lw t0, 12(s0)
Sets register t0 equal to arr[3]
b) slli t1, t0, 2
  add t2, s0, t1
  lw t3, 0(t2)
  addi t3, t3, 1
  sw t3, 0(t2)
Increments the array element specified by t0 (i.e. arr[t0]) by 1
c) lw t0, 0(s0)
  xori t0, t0, 0xFFF; Immediates are sign-extended (0xFFF -> all 1s);
  addi t0, t0, 1
Sets the register t0 to the two's complement negation of arr[0]
```

2. What are the instructions to branch to label on each of the following conditions? The only branch instructions you may use are beq and bne.

s0 < s1	s0 <= s1	s0 > 1
slt t0, s0, s1	slt t0, s1, s0	sltiu t0, s0, 2
bne t0, 0, label	beq t0, 0, label	beq t0, 0, label

## 4 Translating between C and RSIC-V

Translate between the C and RISC-V code. You may want to use the RISC-V Reference Card for more information on the instruction set and syntax. In all of the C examples, we show you how the different variables map to registers – you don't have to worry about the stack or any memory-related issues. You may assume all registers are initialized to zero.

```
\mathbf{C}
                                                 RISC-V
// s0 -> a, s1 -> b
                                                 addi s0, x0, 4
// s2 -> c, s3 -> z
                                                 addi s1, x0, 5
int a = 4, b = 5, c = 6, z;
                                                 addi s2, x0, 6
z = a + b + c + 10;
                                                 add s3, s0, s1
                                                 add s3, s3, s2
                                                 addi s3, s3, 10
// s0 -> int * p = intArr;
                                                      x0, 0(s0)
                                                 SW
// s1 -> a;
                                                 addi s1, x0, 2
*p = 0;
                                                     s1, 4(s0)
                                                 SW
                                                 slli t0, s1, 2
int a = 2;
p[1] = p[a] = a;
                                                 add t0, t0, s0
                                                      s1, 0(t0)
// s0 -> a, s1 -> b
                                                     addi s0, x0, 5
                                                     addi s1, x0, 10
int a = 5, b = 10;
                                                     add t0, s0, s0
if(a + a == b) {
                                                     bne t0, s1, else
    a = 0;
                                                     xor s0, x0, x0
} else {
                                                     jal x0, exit
    b = a - 1;
                                                     addi s1, s0, -1
                                                 exit:
// computes s1 = 2^30
                                                     addi s0, x0, 0
s1 = 1;
                                                     addi s1, x0, 1
for(s0=0;s0<30;s++) {
                                                     addi t0, x0, 30
                                                 loop:
   s1 *= 2;
                                                     beq s0, t0, exit
                                                     add s1, s1, s1
                                                     addi s0, s0, 1
                                                     jal x0, loop
                                                 exit:
// s0 -> n, s1 -> sum
                                                     addi s1, s1, 0
// assume n > 0 to start
                                                 loop:
for(int sum = 0; n > 0; n--) {
                                                     beq s0, x0, exit
  sum += n;
                                                     add s1, s1, s0
                                                     add s0, s0, -1
                                                     jal x0, loop
                                                 exit:
```