| Quantum Factoring | Reduce Factoring to Order Finding | +(x) -> | Quantum | > period of f(x), P. | Computer | P: smallest, s.t. f(x+p) = f(x)

With this, how to factor N?

- @ generate random num a < N
- 2) If grd(a, N) = 1, done. (we've factored N)
- 3 Define:  $f(x) = a^x \mod N$ then, find period, P, of f(x):  $a^{x+p} \pmod N = a^x \pmod N$  $\Rightarrow a^p = 1 \pmod N$ 
  - G If P is odd, go back to Step  $\mathbb{O}$  (meaning re-pick a random "a"). If P is even, but  $\mathbb{A}^{\frac{p}{2}} = -1$  (mod N), goto  $\mathbb{O}$  (have to repick rand  $\mathbb{A}$ )
  - (5) grd  $(a^{\frac{p}{2}}+1, N)$  and grd  $(a^{\frac{p}{2}}-1, N)$  are factors of N

note: The random number "a" has a high (n \frac{1}{2})

probability of success.

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Proof:
    aP = 1 (mod N)
    Γ = a= (mod N) is a sqrt of a (mod N)
   Ir because P is even
  r + 1 because P is the period. (It I was I, then P is
                                     would be the period)
  r = -1 by construction
- At the end of the algorithm, we will find an r, s.t.
   \Gamma \neq 1 and \Gamma \neq -1 and \Gamma = A^{\frac{1}{2}} = \sqrt{1} (mod N)
- There exists such an - by CRT
 Claim: f = grd (r-1, N) is a factor of N.
          That is, f \ I and f \ N
| F + N = If f = N, then r-1 is a multiple of N
           \Rightarrow \Gamma-1 \equiv 0 \pmod{N} \Rightarrow \Gamma=1 \pmod{N}, (ontradiction
f + 1 If f=1 = g(d(r-1, N)
         then I = (\Gamma - I)u + NV,
 mul both sides by
     r+1 , r+1=(r2-1)u + N(r+1) V
    r2=1 (mod N), so r2-1=0 (mod N)
          => r+1=0 (mod N) => r=-1 (mod N)
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Contradiction

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Factor N ( Finding non-trivial sgrt of 1 (mod N) ( find u s.t. u^2 = 1 (mod N) and u \neq \pm 1 (mod N)
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## Background:

How many 39rts of 1 are there mod N?

$$\chi^2 \equiv 1 \pmod{p}$$

$$\Rightarrow \chi^2 - 1 \equiv 0 \pmod{p} \Rightarrow (x+1)(x-1) \equiv 0 \pmod{p}$$

$$\Rightarrow \chi = \pm 1 \pmod{p} \Rightarrow 2 \text{ Syrts}$$
degree n, at most n sulutions

How about :