



Sorting Competition: Group13, 8th Place

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The Algorithm

- We focused on improving the speed at which we find prime factors.
 - At first we thought that identifying if N is prime would increase the speed of the algorithm so we used BigInteger's `isProbablePrime` to do this.
 - If N wasn't prime we just used a modified version of Group 0's old `productOfPrimeFactors()`!
 - Then we use `Arrays.sort()` to sort our program

Group 0's old Code

```
private static long productOfPrimeFactors(long n) {
    long prime1 = 1;
    long prime2 = 1;
    long bound = (long) Math.sqrt(n) + 1;

    for (int i = 2; i <= bound; ++i) {
        if ((n % i) == 0) { // the first found factor must be prime
            if (prime1 == 1) {
                prime1 = i;
            } else { // the second found factor is a prime or a power of the first one
                if (i % prime1 != 0) { // now we know it's a prime
                    prime2 = i;
                    break;
                }
            }
        }
    }

    // if we didn't find any prime factors, the number itself must be prime
    if (prime1 == 1 && prime2 == 1) {
        prime1 = n;
    } else if (prime2 == 1) { // if we have only one prime, the other one may be larger than the square
        // but only if it's not a power of the other prime
        long candidate = n / prime1;
        while (candidate % prime1 == 0) {
            candidate = candidate / prime1;
        }
        prime2 = candidate;
    }

    return prime1 * prime2;
}
```



Our Code

```
private static long productOfPrimeFactors(long n) {
    BigInteger primeChecker = new BigInteger(n + "");
    if (primeChecker.isProbablePrime(100)) {
        return n;
    }
    if (n == 1) {
        return 1;
    }

    long prime1 = 1;
    long prime2 = 1;

    if (n % 2 == 0) {
        prime1 = 2;
    }

    if (n % 3 == 0) {
        if (prime1 == 1) {
            prime1 = 3;
        } else {
            prime2 = 3;
            return prime1 * prime2;
        }
    }

    long limit = (long) Math.sqrt(n) + 1;
    long current = 5;
    long increment = 2;

    while (current <= limit) {
        if ((n % current) == 0) {
            if (prime1 == 1) {
                prime1 = current;
            } else if (current % prime1 != 0) {
                prime2 = current;
                break;
            }
        }

        current += increment;
        increment = 6 - increment;
    }

    if (prime2 == 1) {
        long candidate = n / prime1;
        while (candidate % prime1 == 0) {
            candidate = candidate / prime1;
        }
        prime2 = candidate;
    }

    return prime1 * prime2;
}
```



Data Representation

- We didn't change the original data representation
- We are still taking an array of longs
- Our functions still take longs



But..... is it correct?

- Our algorithm was evaluated by other classmates and was determined to sort correctly.
- Also, since our algorithm is just a slightly modified version of group0's old code, which we know sorted correctly, this algorithm should also sort correctly.



Theoretical stuff

- Since we are using TimSort to sort our data, our worst case for sorting is $\text{BigTheta}(n \log(n))$.
- We believe that the `isProbablePrime` method causes the biggest slow down in our program and its efficiency is unknown to us. Apparently it uses miller-rabin primality testing in some way but its efficiency is not mentioned anywhere. So our best estimation of the worst case is $\text{BigTheta}(n \log(n))$. It's possible it is larger if `isProbablePrime` is slower but we don't know that.



Things we tried

We tried to implement Pollard's Rho algorithm. We found that it was a lot less painful to just abandon this experiment and move on to something more manageable.



Our mistakes

Using BigInteger's `isProbablePrime()` method with a certainty of 100, was a mistake and probably caused the demise of our program.