Assignment 2 proofs

<varlist>, <vardef>, <varname>, <method>, <ifstatemt>, <assignstatemt>, <factor>, <getvarref>

The two rules of Predictive Parsing are:

1. For every production A ::= α 1 | α 2 | α 3 | ... | α n, we must have

FIRST (
$$\alpha i$$
) \cap FIRST (αj) = \emptyset for each pair i, j, i $\neq j$

2. For every nonterminal A such that FIRST (A) contains λ , we must have

FIRST (A)
$$\cap$$
 FOLLOW (A) = \emptyset

<varname> ::= <letter> {<char>}

There are no methods (non-terminals) in the Java Class grammar where the FIRST(A) contains $\,\lambda$. The $\{\,\}$ metasymbol, which represents "0 or more" is never in the FIRST of any non-terminal. Thus, rule #2 is satisfied.

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<varlist>
<varlist> ::= <vardef> {, <vardef>}
<vardef> ::= <type> <varname> | <classname> <varref>
FIRST(\langle type \rangle) = \{I,S\} \cap FIRST(\langle classname \rangle) = \{C,D\} = \emptyset
Same can be proven with <vardef>.
<vardef> ::= <type> <varname> | <classname> <varref>
FIRST(\langle type \rangle = \{I,S\} \cap FIRST(\langle classname \rangle) = \{C,D\} = \emptyset
<varname>
<varname> ::= <letter> {<char>}
FIRST(<letter>) = { Y , Z } no intersection
<method>
<method> ::= <accessor> <type> <methodname> ([<varlist>]) B {<statemt>} <returnstatemt> E
FIRST(<accessor>) = { P , V } no intersection
<ifstatemt>
<ifstatemt> ::= F <cond> T B {<statemt>} E [L B {<statemt>} E]
FIRST<ifstatemt> is a terminal
<assignstatemt>
<assignstatemt> ::= <varname> = <mathexpr> | <varref> = <getvarref>
```