

Computer Architecture Final Project: Single Cycle CPU

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Announcement

- ◆ 1 ~ 3 people / group
- Please find a representative to fill out the google form before 8:00, 12/19 (Sun.)
 - https://forms.gle/NFBM5cAh66Jc3ckr6
 - TA will help you find group members if you can not find any partner
 - ◆ Select "徵隊友" in the form
- The final member list will be announced before 23:59, 12/21 (Tue.)
 - Those who do not response will be regarded as one people in one group

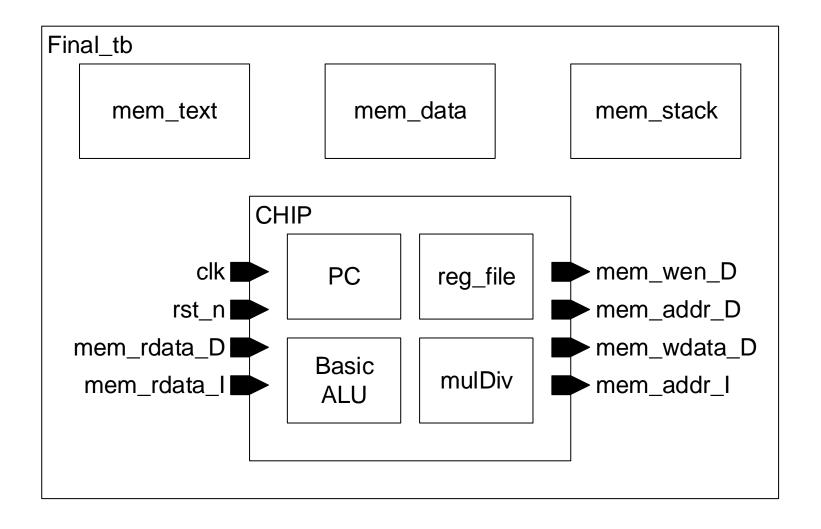


Goal

- Implement a single cycle CPU
- Add multiplication/division unit (mulDiv) to CPU (HW2)
- Handle multi-cycle operations
- Get more familiar with assembly and Verilog
- Run your own assembly in HW1-1 on your CPU (Bonus)



Specification





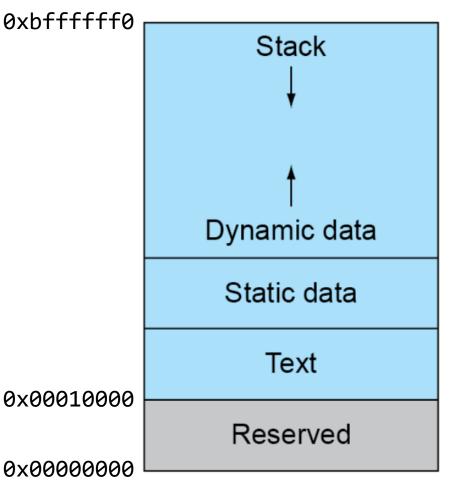
Port Definition

Name	I/O	Width	Description
clk	1	1	Positive edge-triggered clock
rst_n	I	1	Asynchronous negative edge reset
mem_wen_D	0	1	0: Read data from data/stack memory 1: Write data to data/stack memory
mem_addr_D	0	32	Address of data/stack memory
mem_wdata_D	0	32	Data written to data/stack memory
mem_rdata_D	I	32	Data read from data/stack memory
mem_addr_I	0	32	Address of instruction (text) memory
mem_rdata_I	I	32	Instruction read from instruction (text) memory



Memory Layout

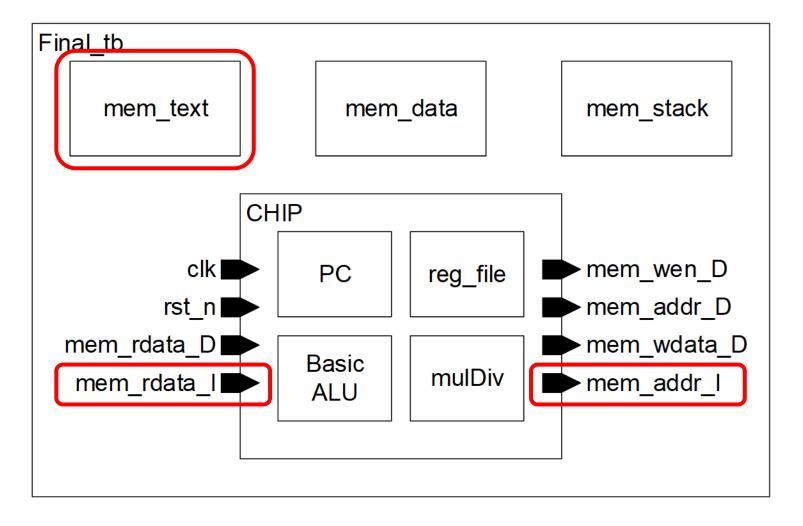
- In Jupiter simulator
- Text
 - Program code
- Data
 - Variables, arrays, etc.
- Stack
 - Automatic storage





Relate Memory to Testbench (1/4)

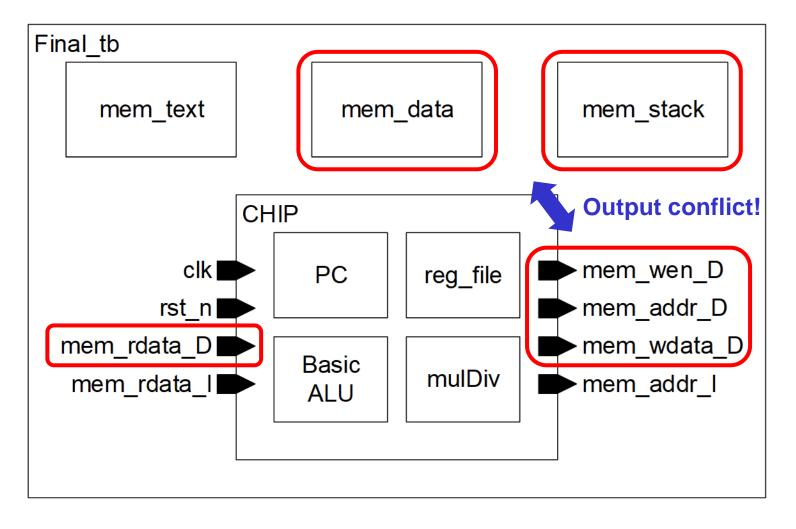
Instruction (text) memory





Relate Memory to Testbench (2/4)

Data/stack memory





define SIZE_TEXT 32

Relate Memory to Testbench (3/4)

 Reduce size of memory blocks to improve simulation speed

 Define offset address for each memory block

 Define high impedance to avoid output conflict

Not synthesizable coding style!

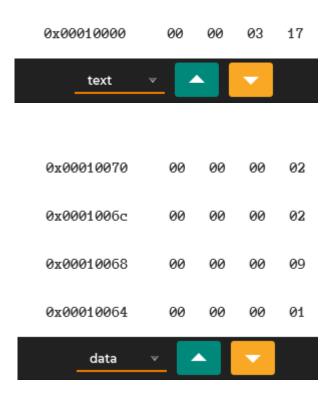
```
define SIZE_DATA 32
             define SIZE_STACK 32
module memory #(
        parameter BITS = 32,
        parameter word_depth = 32
        clk,
        rst n.
        wen,
        offset
```

```
always @(*) begin
    q = {(BITS-1){1'bz}};
    for (i=0; i<word_depth; i=i+1) begin
        if (mem_addr[i] == a)
          q = mem[i];
    end
end</pre>
```

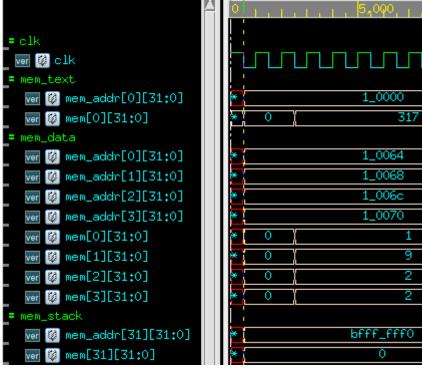


Relate Memory to Testbench (4/4)

In Jupiter



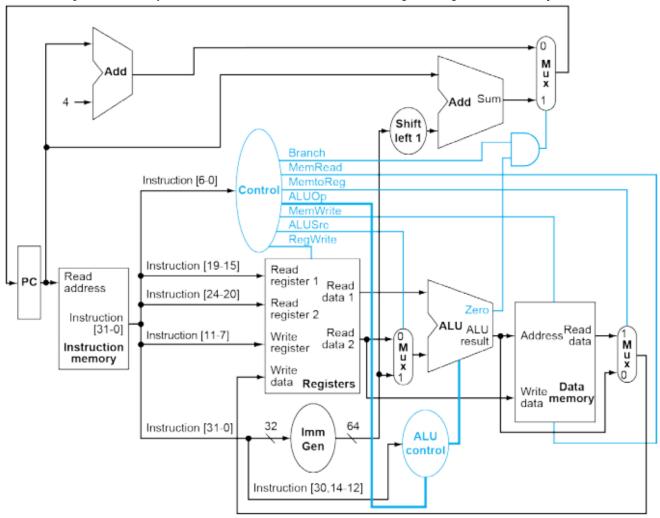
In Testbench





Architecture

Not complete (does not include jal, jalr, ...)





Supporting Instructions

- Your design must <u>at least</u> support
 - ◆ auipc, jal, jalr
 - ♦ beq, lw, sw
 - addi, slti, add, sub
 - mul
- For bonus challengers
 - srai, slli, ... (observe which instructions do you use)
- See "Instruction_Set_Listings.pdf" for more information of machine code



Supplement: Instruction "auipc"

31	12 11	7 6	0
imm[31:12]	r	d	opcode
20		5	7
U-immediate[31:12]	$d\epsilon$	est	AUIPC

- Add upper immediate to PC, and store the result to rd
 - auipc rd, U-immediate
- ◆ Example: auipc x5, 1 (PC = 0x0001001c)
 - \bullet 0x0001001c + 0x00001000 = 0x0001101c
 - Store 0x0001101c in x5



Supplement: Instruction "mul"

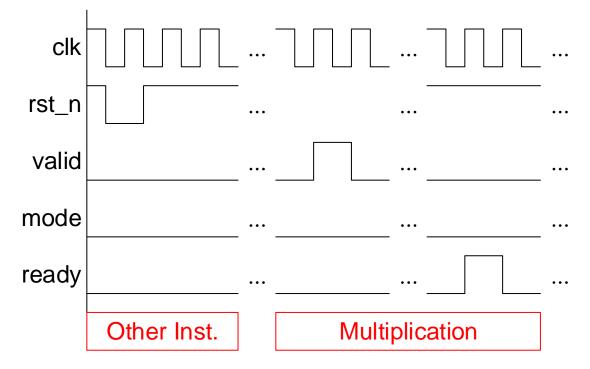
3	1 25	24 20) 19 15	14 12	11 7	6 0
	funct7	rs2	rs1	funct3	rd	opcode
	7	5	5	3	5	7
	MULDIV	multiplier	multiplicand	MUL/MULH[[S]	[U] dest	OP

- Not included in RV32I
- Store the lower 32-b result (rs1 x rs2) to rd
- ◆ Example: mul x10, x10, x6
 - \star x10 = 0x00000001, x6 = 0x00000002
 - \bullet 0x00000001 \times 0x00000002 = 0x00000002
 - Store 0x00000002 in x10
- Your mulDiv can support this instruction!



Multi-Cycle Operation

- Once CPU decodes mul operation, issue valid to your mulDiv
- Once CPU receives ready, store the lower 32-b result to rd
- You might have to design FSM in your CPU

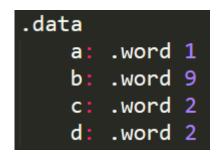


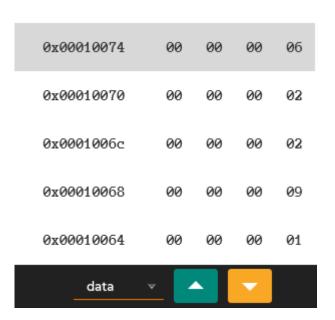


Test Pattern 1: Leaf Example

- Modified from lecture slides
- The procedure loads a,b,c,d from 0x00010064— 0x00010070, and stores the result to 0x00010074
- Run simulation:
 - \$ ncverilog Final_tb.v +define+leaf +access+r

```
def leaf(a,b,c,d):
    f = (a+b) - (c+d)
    return f
```



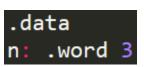


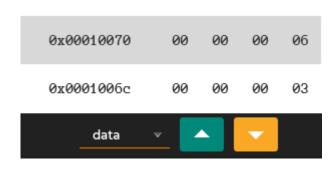


Test Pattern 2: Fact

- Modified from lecture slides
- ◆ The procedure loads n from 0x0001006c, and stores the result to 0x00010070
- Run simulation:
 - \$ ncverilog Final_tb.v +define+fact +access+r

```
def fact(n):
    if n < 1:
        return 1
    else:
        return n*fact(n-1)</pre>
```







(Bonus) Test Pattern 3: HW1 (1/3)

Design your assembly first (hw1.s)

•
$$T(n) = \begin{cases} 2T\left(\left\lfloor \frac{n}{2}\right\rfloor\right) + 8n + 5, & n \ge 2\\ 4, & n = 1 \end{cases}$$

- Example: T(10) = 291, T(30) = 1035
- Use recursive function(./Assembly/hw1.s)

```
# Todo: define your own function in HW1

# Do NOT modify this part!!!
__start:
    la t0, n
    lw x10, 0(t0)
    jal x1,FUNCTION
    la t0, n
    sw x10, 4(t0)
    addi a0,x0,10
    ecall
```



(Bonus) Test Pattern 3: HW1 (2/3)

- Go to simulator
- Dump code → binary file

1	0x00000317
2	0x00830067
3	0x00000297
4	0x02428293
5	0x0002a503
6	0xff5ff0ef
7	0x00000297
8	0x01428293
9	0x00a2a223
10	0x00a00513
11	0x00000073

	D	н		dump code
Bkpt	ext address	Machine Code	used inst.	Source Code
	0x00010000	0x00000317	auipc x6, 0	auipc жб, 0
	0x00010004	0x00830067	jalr x0, x6, 8	jalr x0, x6, 8
	0x00010008	0x00000297	auipc x5, 0	la t0, n
	0x0001000c	0x02428293	addi x5, x5, 36	la t0, n
	0x00010010	0x0002a503	lw x10, x5, 0	lw x10, 0(t0)
	0x00010014	0xff5ff0ef	jal x1, −12	jal x1, FUNCTION
	0x00010018	0x00000297	auipc x5, 0	la t0, n
	0x0001001c	0x01428293	addi x5, x5, 2 0	la t0, n
	0x00010020	0x00a2a223	sw x5, x10, 4	sw x10, 4(t0)
	0x00010024	0x00a00513	addi x10, x0, 10	addi a0, x0, 10
	0x00010028	0x00000073	ecall	ecall



(Bonus) Test Pattern 3: HW1 (3/3)

- Modify the code and save as: ./Verilog/hw1/hw1_text.txt
- Test pattern generation: ./Verilog/hw1/hw1_gen.py
- Run simulation:
 - \$ ncverilog Final_tb.v +define+hw1 +access+r

Delete

1	0x <mark>00000317</mark>
2	0x <mark>00830067</mark>
3	0x <mark>0</mark> 0000297
4	0x <mark>02428293</mark>
5	0x <mark>0002a503</mark>
6	0xff5ff0ef
7	0x <mark>0</mark> 0000297
8	0x <mark>01428293</mark>
9	0x <mark>00a2a223</mark>
10	0x00a00513
11	0x00000073



1	00000317
2	00830067
3	00000297
4	02428293
5	0002a503
6	ff5ff0ef
7	00000297
8	01428293
9	00a2a223



Pattern Generation

- Three python codes provided:
 - leaf_gen.py
 - fact_gen.py
 - hw1_gen.py
- ◆ TA will change the variables in *_gen.py to generate new test patterns when testing your CPU design



Coding Style Check

Register Name	 Type	W	/idth	I	Bus	I	МВ	I	AR	I	AS	I	SR	I	ss	I	ST
alu_in_reg counter_reg shreg_reg state_reg	Flip-flop Flip-flop Flip-flop Flip-flop		32 5 64 2		Y Y Y		N N N		Y Y Y Y		N N N		N N N		N N N		N N N

- All sequential elements must be flip-flops
- Check by Design Compiler
- Command:
 - \$ dv -no_gui
 - design_vision> read_verilog CHIP.v
- Exit:
 - design_vision> exit



Report

- Briefly describe your CPU architecture
- Describe how you design the data path of instructions not referred in the lecture slides (jal, jalr, auipc, ...)
- Describe how you handle multi-cycle instructions (mul)
- Record total simulation time (CYCLE = 10 ns)
 - ◆ Leaf: a = 5, b = 6, c = 8, d = 0
 - ◆ Fact: n = 8
 - (Bonus) HW1: n = 10

Simulation complete via \$finish(1) at time 4795 NS + 0

- Describe your observation
- Snapshot the "Register table" in Design Compiler (p. 22)
- List a work distribution table



Submission

- ◆ Deadline: 01/17 (Mon.) 8:00 am
 - Late submission: 20 % reduction per day
- Upload Final_group_<group_id>.zip to ceiba
 - Final_group_<group_id>.zip
 - Final _group_<group_id>/
 - Final _group_<group_id>/CHIP.v
 - Final _group_<group_id>/hw1.s (bonus)
 - Final _group_<group_id>/hw1_text.txt (bonus)
 - Final _group_<group_id>/report.pdf

Example

```
[r08943003@eda1 ~]$ unzip Final_group_0.zip
Archive: Final_group_0.zip
    creating: Final_group_0/
    inflating: Final_group_0/CHIP.v
    inflating: Final_group_0/hw1.s
    inflating: Final_group_0/hw1_text.txt
    extracting: Final_group_0/report.pdf
```



Score

- ◆ Simulation: 70 % (+ bonus 20 %)
 - Leaf
 - Default: 15 %
 - Change test pattern: 15 %
 - Fact
 - Default: 20 %
 - > Change test pattern: 20 %
 - HW1 (bonus)
 - > Default: 10 %
 - Change test pattern: 10 %
- ◆ Report: 30 %
 - Content: 20 %
 - Snapshots: 5 %
 - Work distribution: 5 %