# METAHEURISTICS

**INF273** 

#5: Local Search

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#### AGENDA

- Neighborhood
- Local Search Operators
  - ✓ Flip
  - ✓ Swap (2-exchange)
  - ✓ 2-Opt, 3-Opt, ... k-opt
  - **√** ...
- Local Search and Hill Climbing

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#### LOCAL SEARCH

- Given a Solution: How to Find a Better One?
- Modification of a given solution gives a "neighbor solution"
- A certain set of operations on a solution gives a set of neighbor solutions, a neighborhood
- Evaluations of neighbors
  - Objective function value
  - Feasibility?

- Also known as Neighborhood Operator
- Neighborhoods are most often defined by a given operation on a solution: N(s)
- Often simple operations
  - Remove an element
  - Add an element
  - Interchange two or more elements of a solution

- Example: Knapsack Problem (10 items, capacity: 100)
- Given solution 0010100000, is feasible. Value = 73

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0

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- Natural operator: "Flip" a bit, i.e.
  - If the item is in the knapsack, take it out
  - If the item is not in the knapsack, include it

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0

• Given solution: 0010100000, value = 73, size = 70

- Natural operator: "Flip" a bit, i.e.
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Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Neighbor Solution	0	1	1	0	1	0	0	0	0	0

- Given solution: 0010100000, value = 73, size = 70
- Neighbor solution: 0110100000, value = 105, size = 96

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Neighbor Solution	0	1	1	0	1	0	0	0	0	0

- Some Neighbors:
  - 0110100000, value = 105, size = 96

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Neighbor Solution	0	1	1	0	1	0	0	0	0	0
Another Neighbor	1	0	1	0	1	0	0	0	0	0

## • Some Neighbors:

- 0110100000, value = 105, size = 96

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Neighbor Solution	0	1	1	0	1	0	0	0	0	0
Another Neighbor	1	0	1	0	1	0	0	0	0	0

- 0110100000, value = 105, size = 96
- 1010100000 , value = 152, not feasible

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Neighbor Solution	0	1	1	0	1	0	0	0	0	0
Neighbor Solution	1	0	1	0	1	0	0	0	0	0

- 0110100000, value = 105, size = 96
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	1	2	3	4	5	6	7	8	9	10
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Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Neighbor Solution	0	1	1	0	1	0	0	0	0	0
Neighbor Solution	1	0	1	0	1	0	0	0	0	0
Neighbor Solution	0	0	1	0	0	0	0	0	0	0

- 0110100000, value = 105, size = 96
- 1010100000 , value = 152, not feasible
- 0010000000, value = 47, size = 48

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Neighbor Solution	0	1	1	0	1	0	0	0	0	0
Neighbor Solution	1	0	1	0	1	0	0	0	0	0
Neighbor Solution	0	0	1	0	0	0	0	0	0	0

- 0110100000, value = 105, size = 96
- 1010100000 , value = 152, not feasible
- 0010000000, value = 47, size = 48
- How many neighbors?

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Neighbor Solution	0	1	1	0	1	0	0	0	0	0
Neighbor Solution	1	0	1	0	1	0	0	0	0	0
Current Solution	0	0	1	0	0	0	0	0	0	0

- 0010100000, value = 73, size = 70
- 0110100000, value = 105, size = 96
- 10101000000, value = 152, not feasible
- 0010000000, value = 47, size = 48

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Current Solution	0	0	1	0	0	0	0	0	0	0

- 0010100000, value = 73, size = 70
- 0010000000, value = 47, size = 48

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Current Solution	0	0	1	0	0	0	0	0	0	0
Current Solution	0	0	1	0	0	0	0	0	0	1

- 0010100000, value = 73, size = 70
- 0010000000, value = 47, size = 48
- 0010000001, value = 106, size = 100

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Current Solution	0	0	1	0	0	0	0	0	0	0
Current Solution	0	0	1	0	0	0	0	0	0	1

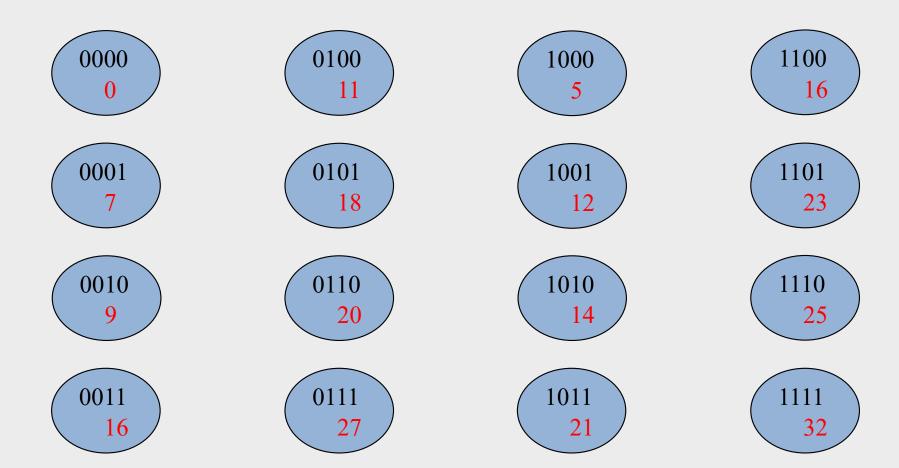
- 0010100000, value = 73, size = 70
- 0010000000, value = 47, size = 48
- 0010000001, value = 106, size = 100
- The global best?

	1	2	3	4	5	6	7	8	9	10
Value	79	32	47	18	26	85	33	40	45	59
Size	85	26	48	21	22	95	43	45	55	52
Given Solution	0	0	1	0	1	0	0	0	0	0
Current Solution	0	0	1	0	0	0	0	0	0	0
Current Solution	0	0	1	0	0	0	0	0	0	1
Best Solution	0	1	0	0	1	0	0	0	0	1

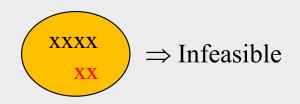
- 0010100000, value = 73, size = 70
- 0010000000, value = 47, size = 48
- 0010000001, value = 106, size = 100
- 0100100001, value = 117, size = 100

#### SEARCH SPACE

• The search space is the set of solutions  $(2^N)$   $xxxx \Rightarrow Solution$  Obj. Fun. Value



## FEASIBLE/INFEASIBLE SPACE

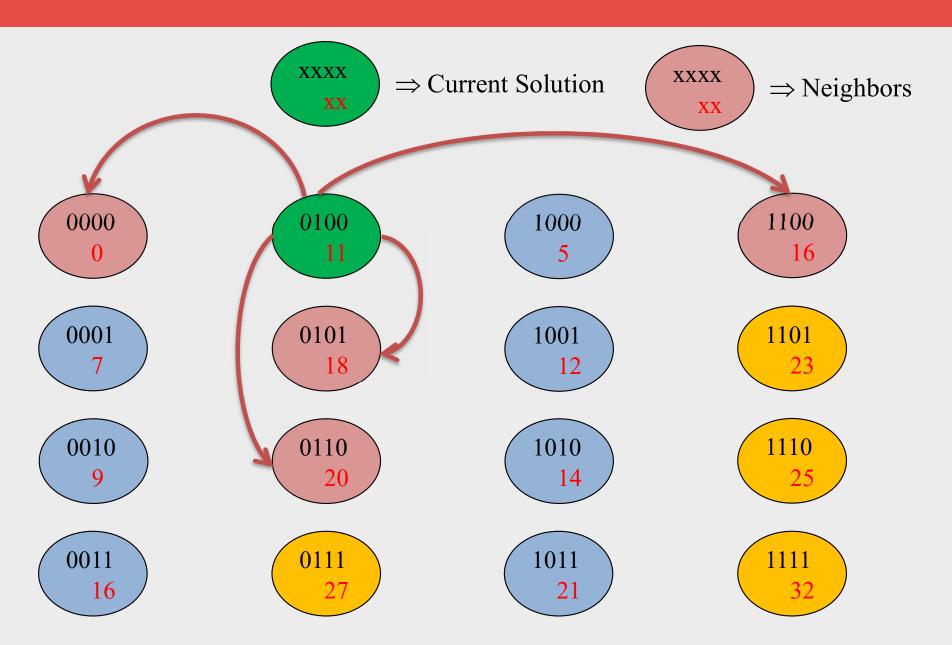


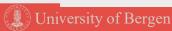
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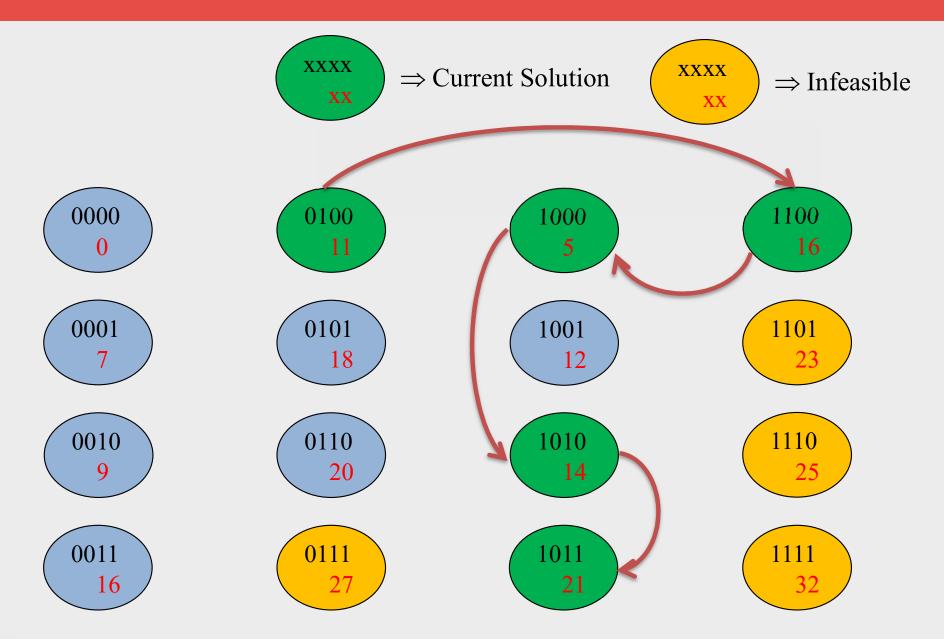
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## NEIGHBORS

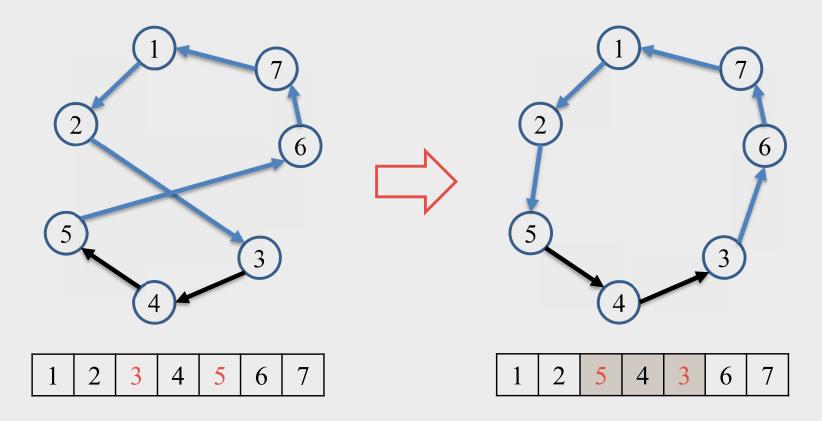




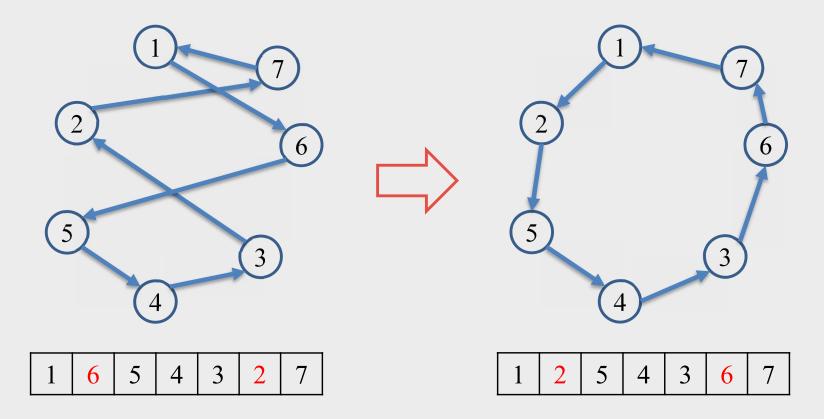
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• Swap (2-exchange)

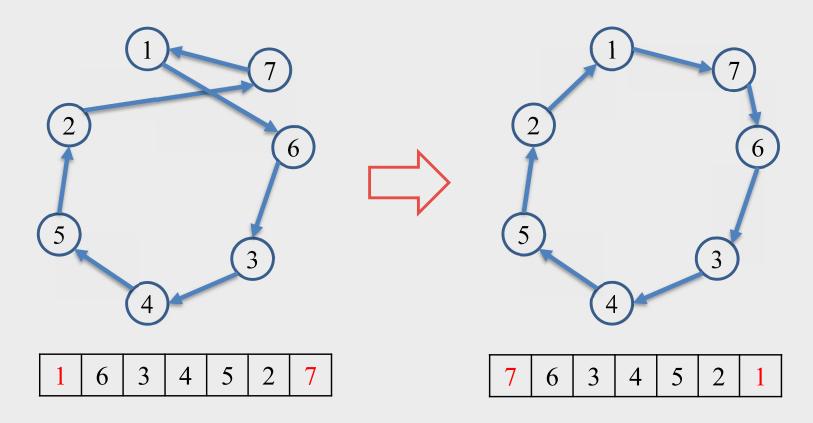


• Swap (2-exchange)



• Swap (2-exchange)

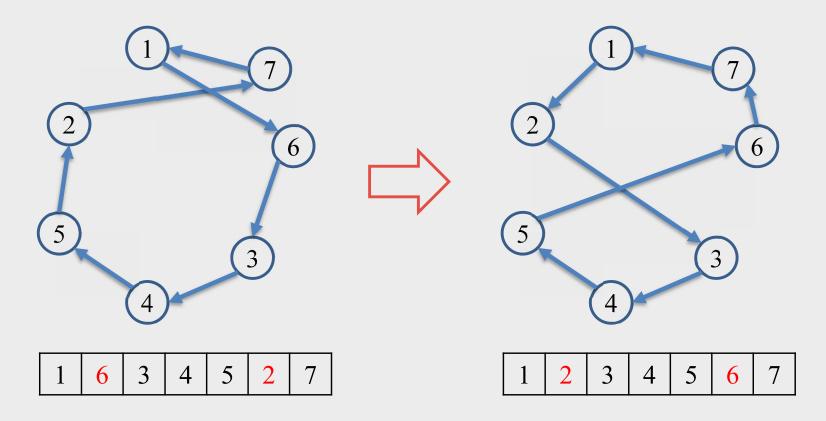
(Example: TSP)



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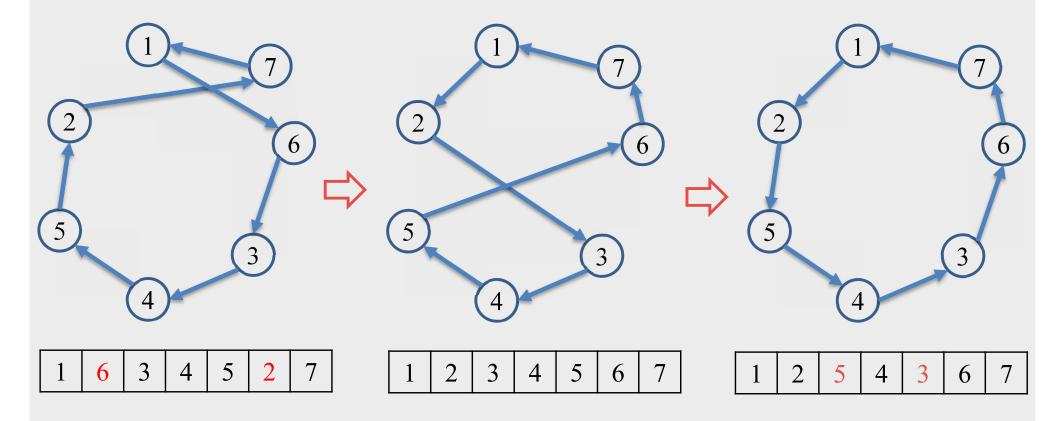
• Swap (2-exchange)

(Example: TSP)

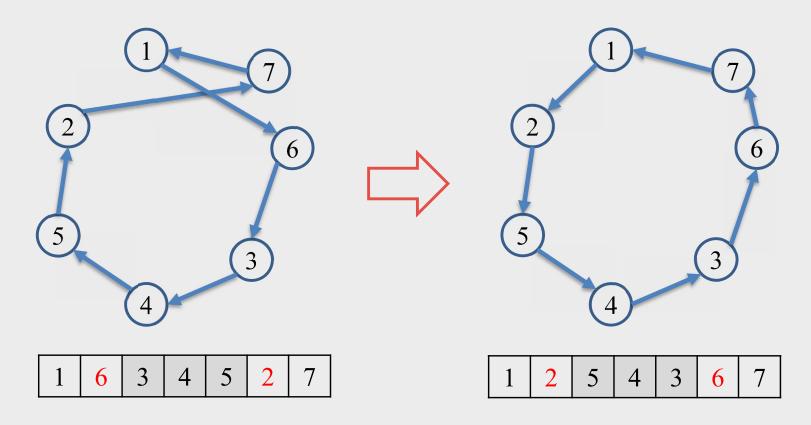


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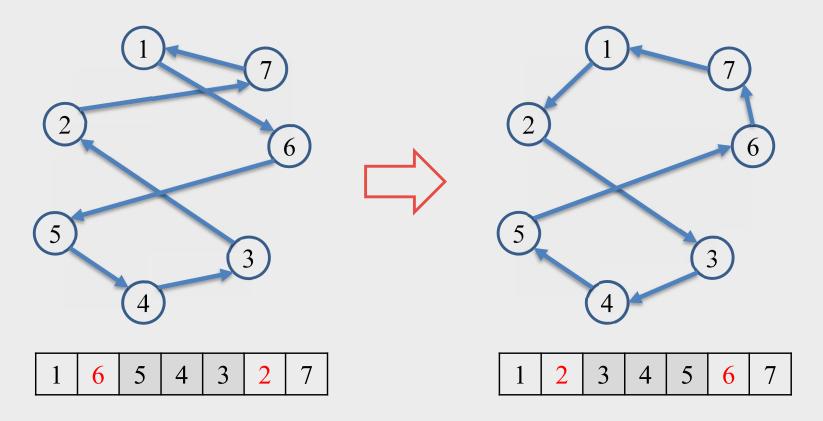
• Swap (2-exchange)



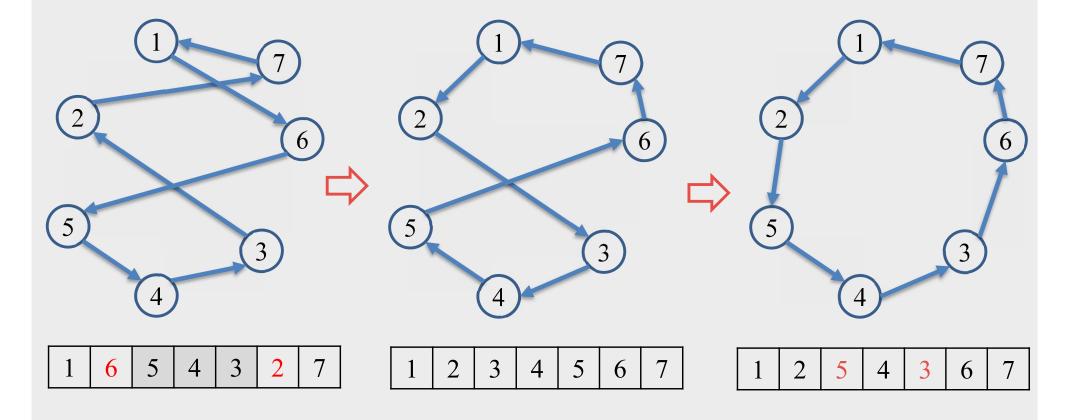
• Swap (2-exchange) with reversing



• Swap (2-exchange) with reversing



• Swap (2-exchange) with reversing (Example: TSP)

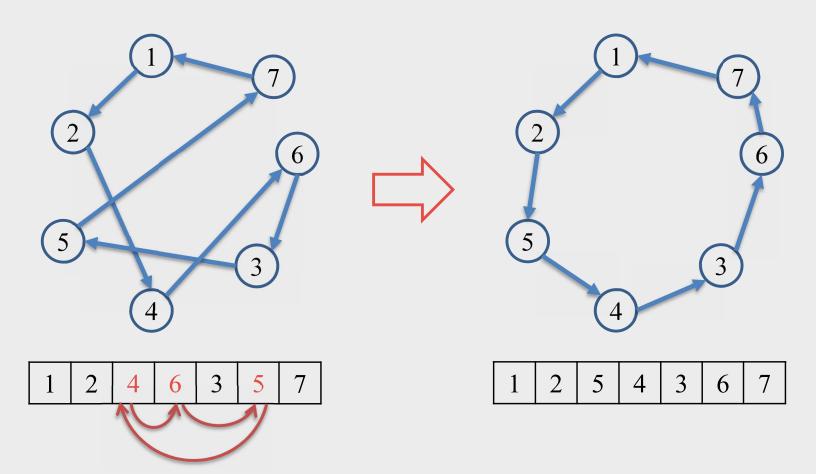


• 2-opt

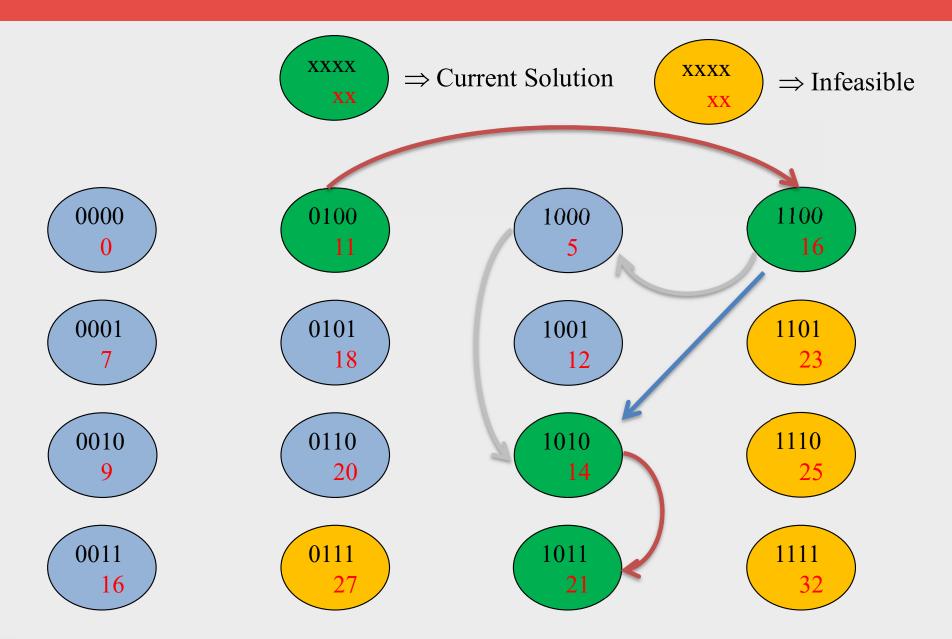
- Note: a finite sequence of "swap (2-exchange) with or without reversing" can generate any tour (for TSP), including the optimum tour
- <u>Strategy</u>:
  - Select the best swap among N\*(N-1)/2 neighbors (2-move neighborhood)
  - Repeat this process until no improvement can be made

• 3-exchange  $\rightarrow$  3-opt

(Example: TSP)



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#### LOCAL SEARCH-ADVANTAGES

- For many problems, it is quite easy to design a local search (LS can be applied to almost any problem)
- The idea of improving a solution by making small changes is easy to understand
- The use of neigborhoods sometimes makes the optimal solution seem "close", e.g.:
  - A knapsack has n items
  - The search space has  $2^n$  members
  - From any solution, no more than n flips are required to reach an optimal solution!

## LOCAL SEARCH

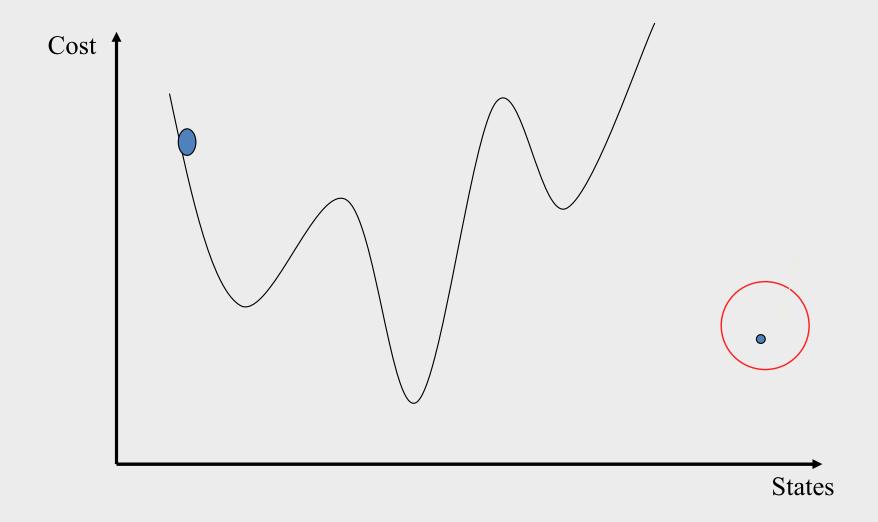
#### Local search 1: Best Accept

```
1: Input: initial solution (s_0)
 2: Input: neighborhood operator N, N(s) \rightarrow all neighbors for s
 3: Input: evaluation function f, f(s) \rightarrow the cost of s
 4: Current \leftarrow s_0
     done \Leftarrow false
     while done = false do
              Best\_neighbor \Leftarrow Current
 7:
 8:
              for each s \in N(Current) do
 9:
                      if f(s) < f(Best\_neighbor) then
10:
                               Best\_neighbor \Leftarrow s
11:
                       end if
              end for
12:
13:
              if Current = Best_neighbor then
14:
                       done \Leftarrow true
15:
              else
16:
                       Current \leftarrow Best\_neighbor
17:
              end if
18: end while
```

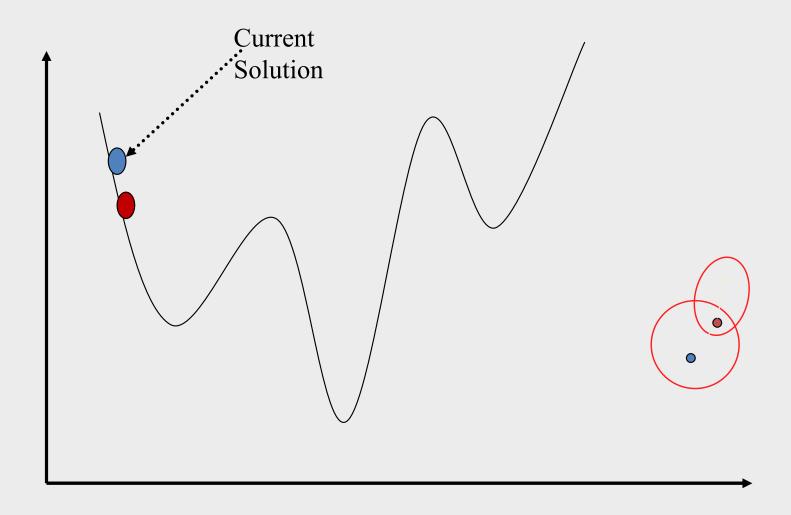
### Local search 2: First Accept

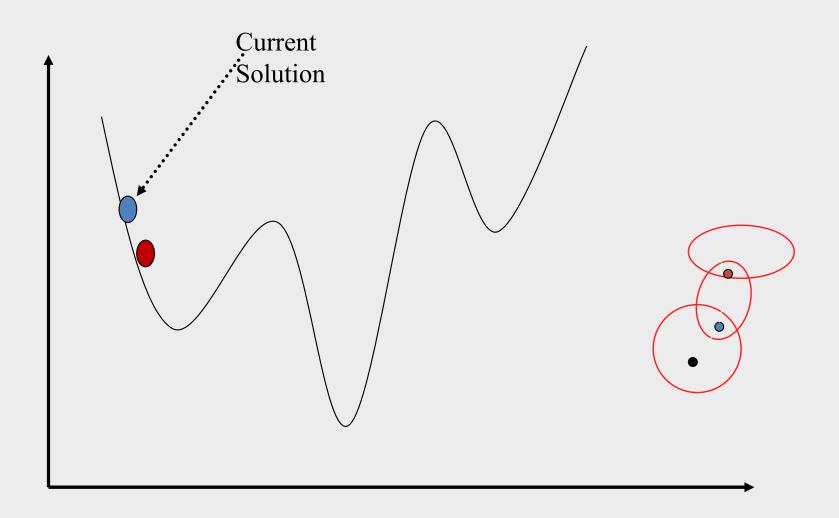
```
1: Input: initial solution (s_0)
     Input: neighborhood operator N, N(s) \rightarrow all neighbors for s
 3: Input: evaluation function f, f(s) \rightarrow the cost of s
 4: Current \leftarrow s_0
     done \Leftarrow false
     while done = false do
              Best\_neighbor \Leftarrow Current
 7:
 8:
              for each s \in N(Current) do
 9:
                       if f(s) < f(Best\_neighbor) then
                                Best\_neighbor \Leftarrow s
10:
11:
                                exit the for-loop
12:
                       end if
13:
              end for
14:
              if Current = Best_neighbor then
15:
                       done \Leftarrow true
16:
              else
17:
                       Current \leftarrow Best\_neighbor
18:
              end if
19
     end while
```

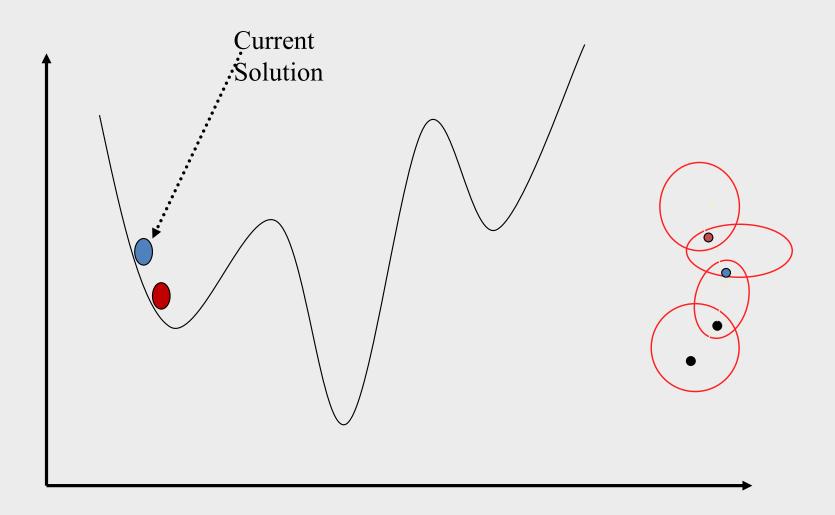
- Start with an *initial solution*
- Iteratively search in the neighborhood for better solutions
- Sequence of solutions  $s_{k+1} \in N(s_k)$
- Strategy for which solution in the neighborhood that will be accepted as the next solution
- Stopping Criteria
- What happens when the neighborhood does not contain a better solution?

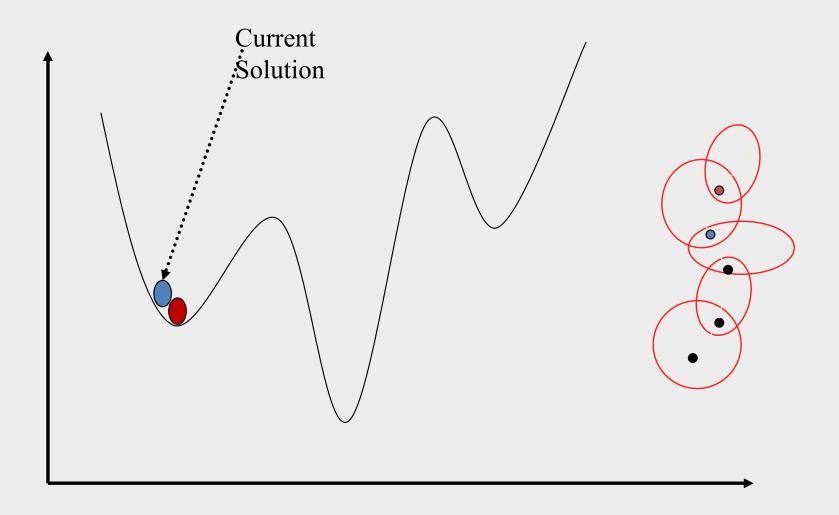




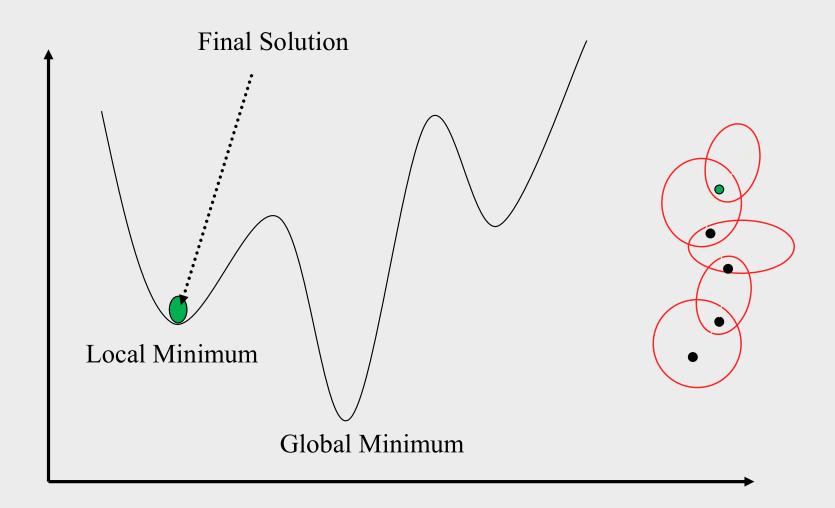






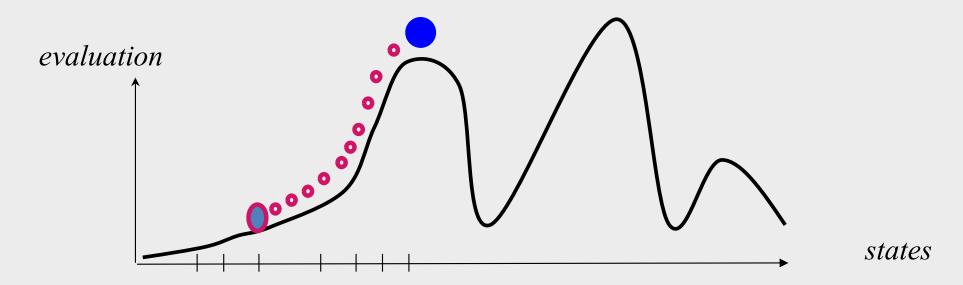






#### Local Search - Hill Climbing

- Hill climbing search algorithm (also known as greedy local search) uses a loop that continually moves in the direction of increasing values (that is uphill).
- It terminates when it reaches a peak where no neighbor has a higher value.



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- Local optimum:
  - If a solution x is "better" than all the solutions in its neighborhood, N(x), we say that x is a local optimum
  - We note that local optimality is defined relative to a particular neighborhood

In a local search we need the following:

- ✓ a Combinatorial Optimization Problem (COP)
- ✓ a starting solution (e.g. random)
- ✓ a defined search neighborhood (neighboring solutions)
- ✓ a move (e.g. changing a variable from  $0 \rightarrow 1$  or  $1 \rightarrow 0$ ), going from one solution to a neighboring solution
- ✓ a move evaluation function (a neighborhood evaluation)
- ✓ a move selection strategy
- $\checkmark$  a stopping criterion e.g. a local optimum

#### LOCAL SEARCH-DISADVANTAGES

- The search stops when no improvement can be found
- Restarting the search might help, but is often not very effective in itself
- Some neighborhoods can become very large (time consuming to examine all the neighbors)

#### Main Challenge in Local Search

How can we avoid the search stopping in a local optimum?

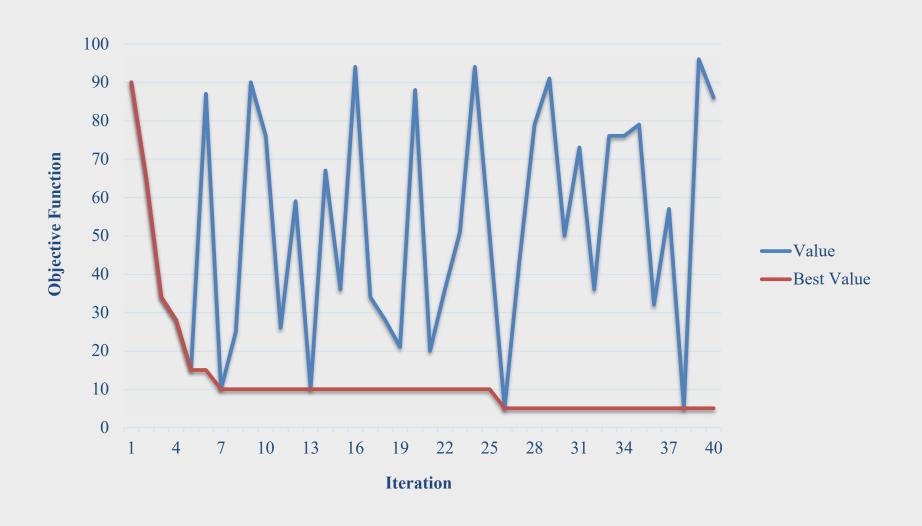
Diversification

#### HOW TO IMPROVE THE LOCAL SEARCH

Some well-known Metaheuristics:

- Simulated Annealing (SA)
- Tabu Search (TS)
- Genetic Algorithms (GA)

### Typical Search Trajectory





#### METAHEURISTICS AND LOCAL SEARCH

- In Local Search, we iteratively improve a solution by making small changes until we cannot make further improvements
- Metaheuristics can be used to guide a Local Search, and to help it to escape a local optimum
- Several metaheuristics are based on Local Search, but the mechanisms to escape local optima vary widely
  - We will look at Simulated Annealing and Tabu Search, as well as some others

## NEXT LECTURE

## LECTURE #6:

### SIMULATED ANNEALING

