

1)

$$a) \lim_{n \rightarrow \infty} \frac{(n^2 - 3n)^2}{5n^3 + n} = \frac{n^2(n-3)^2}{n(5n^2+1)} = \frac{n \cdot (n^2 - 6n + 9)}{5n^2+1} = \frac{n^3 - 6n^2 + 9n}{5n^2+1} \approx \frac{n^3}{n^2} \Rightarrow n \rightarrow \infty$$

$$\Rightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty \rightarrow f(n) = \Omega(g(n))$$

$$b) \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \frac{n^3}{\log_2 n^4} = \frac{n^3}{4 \cdot \log_2 n} \Rightarrow \text{L'Hospital} \Rightarrow \frac{3n^2}{4 \cdot \frac{1}{n \cdot \ln 2}} = \frac{3}{4} \cdot n^2 \cdot n \cdot \ln 2$$

$$\Rightarrow \frac{3}{4} \cdot \infty^2 \cdot \infty \cdot \ln 2 \Rightarrow \infty \Rightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty \rightarrow f(n) = \Omega(g(n))$$

$$c) \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \frac{5n \cdot \log_2^{4n}}{n \cdot \log_2^{5n}} = \frac{5n \cdot \log_2^{4n}}{n^2 \cdot \log_2^5} = \frac{5}{\log_2^5} \cdot \frac{\log_2^{4n}}{n} \Rightarrow \frac{\log_2^4 + \log_2^n}{n} = \frac{2 + \log_2^n}{n}$$

Constant

$$\Rightarrow \text{L'Hospital} \Rightarrow \frac{\frac{1}{n \cdot \ln 2}}{1} = \frac{1}{n \cdot \ln 2} \Rightarrow \frac{1}{\ln 2} \cdot \frac{1}{n} \Rightarrow \frac{1}{n} = \frac{1}{\infty} \Rightarrow 0$$

Constant

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \rightarrow f(n) = O(g(n))$$

$$d) \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \frac{n^n}{10^n} = \left(\frac{n}{10}\right)^n \Rightarrow n \text{ goes to infinity so it grows faster than } 10, \text{ so, } \frac{\infty}{\text{constant}} \rightarrow \infty$$

$$\Rightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty \rightarrow f(n) = \Omega(g(n))$$

$$e) \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \frac{8n \cdot \sqrt[5]{2n}}{n \cdot \sqrt[3]{n}} = 8 \cdot \frac{\sqrt[5]{2n}}{\sqrt[3]{n}} = 8 \cdot \frac{(2n)^{1/5}}{n^{1/3}} = 8 \cdot 2^{1/5} \cdot \frac{n^{1/5}}{n^{1/3}} = 8 \cdot 2^{1/5} \cdot n^{1/5 - 1/3} = 8 \cdot 2^{1/5} \cdot n^{-2/15} = \frac{1}{\sqrt[15]{n^2}}$$

Constant

$$\Rightarrow \frac{1}{\infty} \Rightarrow 0 \rightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \rightarrow f(n) = O(g(n))$$

2)  
a) Static void methodA (String str\_array[]) {  
for (int i=0; i<str\_array.length; i++)  
str\_array[i] = " ";

n times  $O(1) \Rightarrow \boxed{O(n)}$   
 $\left. \begin{array}{l} \text{str\_array}[0] = " " \rightarrow O(1) \\ \text{str\_array}[1] = " " \rightarrow O(1) \\ \vdots \\ \text{str\_array}[n] = " " \rightarrow O(1) \end{array} \right\} n \text{ times}$

b) Static void methodB (String str\_array[]) {  
for (int i=0; i<str\_array.length; i++)  
methodA (str\_array);  
for (int j=0; j<str\_array.length; j++)  
System.out.println (str\_array[j]);  
}

$$O(n^2) + O(n) = \boxed{O(n^2)}$$

$\left. \begin{array}{l} \text{assign} \\ \text{str\_array}[0] = " " \\ \vdots \\ \text{str\_array}[n] = " " \end{array} \right\} i=0 \rightarrow n-1 \left. \begin{array}{l} n \text{ times} \\ O(n) \\ \vdots \\ O(n^2) \end{array} \right\}$   
 $\left. \begin{array}{l} \text{print} \\ \text{str\_array}[0] \\ \vdots \\ \text{str\_array}[n] \end{array} \right\} n \text{ times } O(1) \Rightarrow O(n)$

c) Static void methodC (String str\_array[]) {  
for (int i=0; i<str\_array.length; i++)  
for (int j=0; j<str\_array.length; j++)  
methodB (str\_array);  
}

n times  $O(n^2) \Rightarrow O(n^3)$   
 $\left. \begin{array}{l} n \text{ times} \\ O(n^2) \end{array} \right\} \Rightarrow O(n^3) \Rightarrow \boxed{O(n^4)}$

d) Static void methodD (String str\_array[]) {  
for (int i=0; i<str\_array.length; i++)  
System.out.println (str\_array[i]);  
str\_array[i--] = " ";  
}

$\Rightarrow i \rightarrow 0$   
 $\text{str\_array}[0] \rightarrow \text{print}$   
 $\text{str\_array}[0] \rightarrow \text{assign}$   
 $i \rightarrow -1$   
 $\text{for } \Rightarrow i \rightarrow 0$   
 infinite loop

$\Rightarrow i$  cannot increase.  $[i--]$  decrease,  $\text{for}$  increase

$$\Rightarrow \boxed{O(\infty)}$$

e) Static void methodE (String str\_array[]) {  
for (int i=0; i<str\_array.length; i++)  
if (str\_array[i] = " ")  
break;  
}

$\Rightarrow$  If str\_array's any node assigned to " ", it goes until the end  
 so, it search n times to array

$$\rightarrow \boxed{O(n)}$$

3)

### a) Maximum Difference - Ascending Order

Algorithm max-difference (arr[])

min ← arr[0] -----  $O(1)$

max ← arr[arr.length - 1] ---  $O(1)$

return (max - min) - - - -  $O(1)$

// 1, 2, 3, 4, 5

↓  
arr[0]

↓  
arr[arr.length - 1]

=> Total Big-O is constant

=>  $O(1)$

### b) Maximum Difference - Not Sorted

Algorithm max-difference (arr[])

min ← ∞ -----  $O(1)$

max ← (-∞) -----  $O(1)$

// Assign max-value for initially  
// Assign min-value for initially

for (i = 0; i < arr.length; i++) -----  $O(n)$

if (arr[i] < min) -----  $O(1)$

min ← arr[i] -----  $O(1)$

if (arr[i] > max) -----  $O(1)$

max ← arr[i] -----  $O(1)$

return (max - min) - - - -  $O(1)$

// I write if-if instead of  
// if-else if. Because maybe list  
// have only one element, so  
// min and max will be the same.

// for area is n time  $O(1)$

-> So it will linear big-O.

->  $O(n)$

for A and B:

- Before calling functions, we need to check array length. List is empty or not.

Algorithm is-empty (arr[])

if (arr.length <= 0)

return true ----- // list is empty

else

return false ----- // list is not empty

// This method has  
 $O(1)$ . Big-O notation  
is constant.