

1. *Additional reading for this problem: J. Demmel, Applied Numerical Linear Algebra, available online via the UMCP library.*

The goal of this exercise is to understand how one can compute a QR decomposition using *Householder reflections*.

- (a) Let  $u$  be a unit vector in  $\mathbb{R}^n$ , i.e.,  $\|u\|_2 = 1$ . Let  $P = I - 2uu^\top$ . This matrix performs reflection with respect to the hyperplane orthogonal to the vector  $u$ . Show that  $P = P^\top$  and  $P^2 = I$ .

**Solution:** By the definition of  $P$  and the linearity of transpose,

$$P^\top = (I - 2uu^\top)^\top = I^\top - 2(uu^\top)^\top.$$

But then since  $I^\top = I$  and  $(uu^\top)^\top = (u^\top)^\top u^\top = uu^\top$ ,

$$P^\top = I^\top - 2(uu^\top)^\top = I - 2uu^\top = P$$

and hence  $P = P^\top$  as desired.

Now observe that since  $I^2 = I$  and  $I(uu^\top) = (uu^\top)I = uu^\top$ ,

$$P^2 = (I - 2uu^\top)(I - 2uu^\top) = I - 4uu^\top + 4(uu^\top)^2.$$

But then because  $\|u\|_2 = 1$ ,

$$(uu^\top)^2 = u(u^\top u)u^\top = u\|u\|_2^2 u^\top = uu^\top$$

and hence

$$P^2 = I - 4uu^\top + 4(uu^\top)^2 = I - 4uu^\top + 4uu^\top = I.$$

- (b) Let  $x \in \mathbb{R}^n$  be any vector,  $x = [x_1, \dots, x_n]^\top$ . Let  $u$  be defined as follows:

$$\tilde{u} := \begin{bmatrix} x_1 + \text{sign}(x_1)\|x\|_2 \\ x_2 \\ \vdots \\ x_n \end{bmatrix} \equiv x + \text{sign}(x_1)\|x\|_2 e_1, \quad u = \frac{\tilde{u}}{\|\tilde{u}\|_2}, \quad (1)$$

where  $e_1 = [1, 0, \dots, 0]^\top$ . The matrix with the vector  $u$  constructed according to (1) will be denoted  $\text{House}(x)$ :

$$P = I - 2uu^\top \equiv I - 2 \frac{\tilde{u}\tilde{u}^\top}{\tilde{u}^\top \tilde{u}} \equiv \text{House}(x).$$

Calculate  $Px$ .

**Solution:** We claim first that  $2\tilde{u}^\top x = \tilde{u}^\top \tilde{u}$ . To see this, note that

$$\begin{aligned} 2\tilde{u}^\top x &= 2(x + \text{sign}(x_1)\|x\|_2 e_1)^\top x = 2(x^\top + \text{sign}(x_1)\|x\|_2 e_1^\top)x \\ &= 2x^\top x + 2\text{sign}(x_1)\|x\|_2 e_1^\top x = 2x^\top x + 2\text{sign}(x_1)\|x\|_2 x_1 = 2x^\top x + 2|x_1|\|x\|_2 \end{aligned}$$

and also that since  $x^\top e_1 = e_1^\top x = x_1$  and  $e_1^\top e_1 = 1$  and  $\text{sign}(x_1)^2 = 1$ ,

$$\begin{aligned}\tilde{u}^\top \tilde{u} &= (x + \text{sign}(x_1)\|x\|_2 e_1)^\top (x + \text{sign}(x_1)\|x\|_2 e_1) \\ &= (x^\top + \text{sign}(x_1)\|x\|_2 e_1^\top)(x + \text{sign}(x_1)\|x\|_2 e_1) = x^\top x + 2\text{sign}(x_1)\|x\|_2 x_1 + \|x\|_2^2 \\ &= x^\top x + 2|x_1|\|x\|_2 + x^\top x = 2x^\top x + 2|x_1|\|x\|_2\end{aligned}$$

Hence, we have that  $2\tilde{u}^\top x = \tilde{u}^\top \tilde{u}$  (notably a scalar quantity). Then it follows that

$$2\frac{\tilde{u}\tilde{u}^\top}{\tilde{u}^\top \tilde{u}}x = \tilde{u} \cdot \frac{2\tilde{u}^\top x}{\tilde{u}^\top \tilde{u}} = \tilde{u}$$

and so

$$Px = \left(I - 2\frac{\tilde{u}\tilde{u}^\top}{\tilde{u}^\top \tilde{u}}\right)x = x - 2\frac{\tilde{u}\tilde{u}^\top}{\tilde{u}^\top \tilde{u}}x = x - \tilde{u} = -\text{sign}(x_1)\|x\|_2 e_1.$$

- (c) Let  $A$  be an  $m \times n$  matrix,  $m \geq n$ , with columns  $a_j$ ,  $j = 1, \dots, n$ . Let  $A_0 = A$ . Let  $P_1 = \text{House}(a_1)$ . Then  $A_1 := P_1 A_0$  has the first column with the first entry nonzero and the other entries being zero. Next, we define  $P_2$  as

$$P_2 = \begin{bmatrix} 1 & 0 \\ 0 & \tilde{P}_2 \end{bmatrix}$$

where the matrix  $\tilde{P}_2 = \text{House}(A_1(2:m, 2))$ . The notation  $A_1(2:m, 2)$  is Matlab's syntax indicating this is the vector formed by entries 2 through  $m$  of the 2nd column of  $A_1$ . Then we set  $A_2 = P_2 A_1$ . And so on.

This algorithm can be described as follows. Let  $A_0 = A$ . Then for  $j = 1, 2, \dots, n$  we set

$$P_j = \begin{bmatrix} I_{(j-1) \times (j-1)} & 0 \\ 0 & \tilde{P}_j \end{bmatrix}; \quad \tilde{P}_j = \text{House}(A_{j-1}(j:m, j)), \quad A_j = P_j A_{j-1}.$$

Check that the resulting matrix  $A_n$  is upper triangular, its entries  $(A_n)_{ij}$  are all zeros for  $i > j$ . Propose an **if**-statement in this algorithm that will guarantee that  $A_n$  has positive entries  $(A_n)_{jj}$ ,  $1 \leq j \leq n$ .

**Solution:** We proceed by induction on  $k$  to show that for  $0 \leq k \leq n$ ,  $A_k$  is upper triangular in columns  $1 \leq j \leq k$ . The base case of  $k = 0$  is trivially true, since the hypothesis asserts nothing about  $A_0$ ; there are no columns  $1 \leq j \leq 0$ .

Now for some  $0 \leq k \leq n - 1$ , suppose  $A_k$  is upper triangular in columns  $1 \leq j \leq k$ . Further, note that the entry

$$\begin{aligned}(A_{k+1})_{ij} &= (P_{k+1} A_k)_{ij} = \sum_{\ell=1}^m (P_{k+1})_{i\ell} (A_k)_{\ell j} \\ &= \sum_{\ell=1}^k (P_{k+1})_{i\ell} (A_k)_{\ell j} + \sum_{\ell=k+1}^m (P_{k+1})_{i\ell} (A_k)_{\ell j}.\end{aligned}$$

Further, recall that by definition,

$$P_{k+1} = \begin{bmatrix} I_{k \times k} & 0 \\ 0 & \tilde{P}_{k+1} \end{bmatrix}$$

and so it follows that

$$(P_{k+1})_{i\ell} = \begin{cases} (\tilde{P}_{k+1})_{i-k, \ell-k} & i, \ell > k \\ 1 & i = \ell \leq k \\ 0 & \text{otherwise} \end{cases} \quad (2)$$

Now focusing first on the left summation, note that because  $A_k$  is upper triangular in columns  $1 \leq j \leq k$ , and the first case of (2) is not possible since  $\ell \leq k$ , we have that

$$\sum_{\ell=1}^k (P_{k+1})_{i\ell} (A_k)_{\ell j} = \sum_{\ell=1}^j (P_{k+1})_{i\ell} (A_k)_{\ell j} = \sum_{\ell=1}^j (A_k)_{\ell j} \begin{cases} 1 & i = \ell \\ 0 & i \neq \ell \end{cases} = \begin{cases} (A_k)_{ij} & i \leq j \\ 0 & i > j \end{cases}.$$

We now focus on the right summation, in which the second case of (2) is not possible since  $\ell > k$ . Then

$$\begin{aligned} \sum_{\ell=k+1}^m (P_{k+1})_{i\ell} (A_k)_{\ell j} &= \sum_{\ell=k+1}^m (A_k)_{\ell j} \begin{cases} 0 & i \leq k \\ (\tilde{P}_{k+1})_{i-k, \ell-k} & i > k \end{cases} \\ &= \begin{cases} 0 & i \leq k \\ \sum_{\ell=k+1}^m (\tilde{P}_{k+1})_{i-k, \ell-k} (A_k)_{\ell j} & i > k \end{cases} \end{aligned}$$

and hence

$$(A_{k+1})_{ij} = \begin{cases} (A_k)_{ij} & i \leq j \\ 0 & i > j \end{cases} + \begin{cases} 0 & i \leq k \\ \sum_{\ell=k+1}^m (\tilde{P}_{k+1})_{i-k, \ell-k} (A_k)_{\ell j} & i > k \end{cases}.$$

From the above, note that for  $1 \leq j \leq k$ , if  $i > j$  then both summands will be 0 (the right will be 0 because  $\ell > k \geq j$  so  $(A_k)_{\ell j} = 0$  since  $j \leq k$ ). Hence, it suffices to show that  $A_{k+1}$  is upper triangular in column  $j = k + 1$ ;

The left summand cannot contribute any nonzero values if  $i > j$ , so we need only consider the right summand. Re-indexing to write it as a matrix multiplication using the Matlab notation and taking  $j = k + 1$ ,

$$\begin{aligned} \sum_{\ell=k+1}^m (\tilde{P}_{k+1})_{i-k, \ell-k} (A_k)_{\ell j} &= \sum_{\ell=1}^{m-k} (\tilde{P}_{k+1})_{i-k, \ell} (A_k)_{k+\ell, j} \\ &= (\tilde{P}_{k+1} A_k(k+1 : m, j))_{i-k} = (\tilde{P}_{k+1} A_k(k+1 : m, k+1))_{i-k}. \end{aligned}$$

But then by item (b),

$$\begin{aligned} (\tilde{P}_{k+1} A_k(k+1 : m, k+1))_{i-k} &= (-\text{sign}((A_k)_{k+1, k+1}) \|A_k(k+1 : m, k+1)\|_2 e_1)_{i-k} \\ &= -\text{sign}((A_k)_{k+1, k+1}) \|A_k(k+1 : m, k+1)\|_2 (e_1)_{i-k} \end{aligned}$$

which is 0 if  $i > j = k + 1$  since then  $i - k \geq 2$  so  $(e_1)_{i-k} = 0$ . Hence, we have that  $A_{k+1}$  is upper triangular in columns  $1 \leq j \leq k + 1$ , so by the principle of induction, for any  $0 \leq k \leq n$ ,  $A_k$  is upper triangular in columns  $1 \leq j \leq k$ . Letting  $k = n$ , we obtain the desired result:  $A_n$  is upper triangular in columns  $1 \leq j \leq n$ .

If we want to ensure that the diagonal entries  $(A_n)_{jj} \geq 0$  for  $1 \leq j \leq n$ , we could alter the definition of  $\tilde{u}$  such that

$$\tilde{u} \equiv x + \frac{|x_1|}{x_1} \|x\|_2 e_1,$$

that is, if  $A_{jj} < 0$  then use  $\tilde{u} = x - \text{sign}(x_1) \|x\|_2 e_1$ .

- (d) Extract the QR decomposition of  $A$  given the matrices  $P_j$ ,  $1 \leq j \leq n$ , and  $A_n$ .

**Solution:** From (c) we know that  $A_n$  is upper triangular. Hence, it suffices to construct an orthogonal matrix  $Q$  from the matrices  $P_j$ ,  $1 \leq j \leq n$  such that  $A = QA_n$ .

Towards this end, note that

$$A_n = P_n A_{n-1} = P_n P_{n-1} A_{n-2} = \cdots = P_n P_{n-1} \cdots P_1 A_0 = P_n P_{n-1} \cdots P_1 A$$

so  $Q$  should be defined such that

$$A = Q P_n P_{n-1} \cdots P_1 A$$

and so

$$Q = P_1^{-1} P_2^{-1} \cdots P_n^{-1}.$$

But finally, from (a) we know that each  $P_j$  is orthogonal ( $P_j P_j^\top = P_j^2 = I$ ) and hence  $P_j^{-1} = P_j$ . Hence, the QR decomposition of  $A$  is given by

$$Q = P_1 P_2 \cdots P_n, \quad R = A_n$$

where the orthogonality of  $Q$  is guaranteed as the product of orthogonal matrices.

2. Prove items (1)–(6) of Theorem 3 on page 14 of [LinearAlgebra.pdf](#).

**Solution:**

- (1) Let  $A$  be symmetric and  $A = U\Lambda U^\top$  is an eigendecomposition of  $A$ . Then since  $A$  is symmetric and  $U^\top U = I$ ,

$$U\Lambda^2 U^\top = (U\Lambda U^\top)(U\Lambda U^\top) = A^2 = AA^\top = (U\Sigma V^\top)(V\Sigma^\top U^\top) = U\Sigma^2 U^\top$$

and so  $\Sigma^2 = \Lambda^2$ ; in order to guarantee  $\sigma_i \geq 0$ , it must be that  $\sigma_i = |\lambda_i|$ . Now

$$U\Lambda U^\top = A = U\Sigma V^\top = U|\Lambda|V^\top$$

and so left multiplying by  $\text{sign}(\Lambda)\Lambda^{-1}U^\top$ , we obtain

$$\text{sign}(\Lambda)U^\top = \text{sign}(\Lambda)\Lambda^{-1}|\Lambda|V^\top = \text{sign}(\Lambda)\text{sign}(\Lambda)V^\top = V^\top.$$

Finally, taking the transpose,

$$V = (\text{sign}(\Lambda)U^\top)^\top = U\text{sign}(\Lambda)^\top = U\text{sign}(\Lambda),$$

so  $v_i = u_i\text{sign}(\lambda_i)$ .

- (2) Since  $U$  and  $V$  are orthogonal (so in particular  $U^\top U = I$  and  $V^\top = V^{-1}$ ) and  $\Sigma$  is diagonal (so in particular  $\Sigma^\top \Sigma = \Sigma^2$ ), we have that

$$A^\top A = (U\Sigma V^\top)^\top (U\Sigma V^\top) = V\Sigma^\top U^\top U\Sigma V^\top = V\Sigma^\top \Sigma V^\top = V\Sigma^2 V^{-1}.$$

Hence,  $V\Sigma^2 V^{-1}$  is the eigendecomposition of  $A^\top A$ , so the eigenvalues of  $A^\top A$  are the diagonal entries of  $\Sigma^2$ , those being  $\sigma_i^2$ . Further, the eigenvectors of  $A^\top A$  are the columns of  $V$ , the right singular vectors.

- (3) Since  $U$  and  $V$  are orthogonal (so in particular  $U^\top = U^{-1}$  and  $V^\top V = I$ ) and  $\Sigma$  is diagonal (so in particular  $\Sigma\Sigma^\top = \Sigma^2$ ), we have that

$$AA^\top = (U\Sigma V^\top)(U\Sigma V^\top)^\top = U\Sigma V^\top V\Sigma^\top U^\top = U\Sigma\Sigma^\top U^\top = U\Sigma^2 U^\top.$$

Note that  $U$  and  $\Sigma$  are  $m \times n$  and  $n \times n$  respectively; the eigendecomposition of  $AA^\top$  is obtained by padding  $U$  and  $\Sigma$  with zeros to make each  $m \times m$ . Hence, the eigenvalues of  $AA^\top$  are the  $\sigma_i^2$  from original  $\Sigma^2$  and the  $m - n$  padded zeros.

Further, the eigenvectors of  $AA^\top$  corresponding to the  $\sigma_i^2$  are the columns of  $U$ , the left singular vectors, and the eigenvectors corresponding to 0 can be any  $m - n$  vectors orthogonal to each and the columns of  $U$ .

- (4) By the non-negativity of norms,

$$\min_x \|Ax - b\| \geq 0.$$

Now take  $x = V\Sigma^{-1}U^\top b$ , in which case

$$\min_x \|Ax - b\| \leq \|(U\Sigma V^\top)(V\Sigma^{-1}U^\top b) - b\| = \|b - b\| = \|0\| = 0$$

and hence  $\min_x \|Ax - b\|$  attains its minimum value at  $x = V\Sigma^{-1}U^\top b$ . Further, since  $A$  is full rank, this solution is unique.

- (5) Recall we have seen a property relating the 2-norm of  $A$  to  $\rho(A^\top A)$ , the largest eigenvalue of  $A^\top A$ :

$$\|A\|_2 = \sqrt{\rho(A^\top A)}.$$

Further, by item (2), we know that the eigenvalues of  $A^\top A$  are  $\sigma_i^2$ , the largest of which is  $\sigma_1^2$ . Hence,

$$\|A\|_2 = \sqrt{\rho(A^\top A)} = \sqrt{\sigma_1^2} = \sigma_1.$$

Now suppose  $A$  is square and nonsingular; then  $A^{-1} = V\Sigma^{-1}U^\top$ , since

$$AA^{-1} = U\Sigma V^\top V\Sigma^{-1}U^\top = U\Sigma\Sigma^{-1}U^\top = UU^\top = I$$

and hence  $V\Sigma^{-1}U^\top$  is a singular value decomposition of  $A^{-1}$ . Note that the diagonal entries of  $\Sigma^{-1}$  are  $1/\sigma_i$ , and so the largest singular value of  $A^{-1}$  is the reciprocal of the smallest singular value of  $A$ ,  $\sigma_n$ . Hence, we have that

$$\|A^{-1}\|_2 = \frac{1}{\sigma_n}.$$

(6) Note that for  $1 \leq k \leq n$ ,

$$Av_k = U\Sigma V^\top v_k = U\Sigma e_k = U\sigma_k = \sigma_k u_k$$

and  $\sigma_1 \geq \dots \geq \sigma_r > \sigma_{r+1} = \dots = \sigma_n = 0$  by assumption. In particular, this means that  $Av_k \neq 0$  for  $1 \leq k \leq r$  and  $Av_k = 0$  for  $r+1 \leq k \leq n$ , and so we have both that

$$\text{span}\left(\frac{1}{\sigma_1}Av_1, \dots, \frac{1}{\sigma_r}Av_r\right) = \text{span}(u_1, \dots, u_r) \subseteq \text{range}(A)$$

and

$$\text{span}(v_{r+1}, \dots, v_n) \subseteq \text{null}(A)$$

so since  $v_1, \dots, v_n$  are orthogonal (and hence linearly independent) then  $\text{rank}(A) \geq r$  and  $\text{nullity}(A) \geq n - r$ . Now towards a contradiction suppose  $\text{rank}(A) > r$ . Then

$$\text{rank}(A) + \text{nullity}(A) \geq \text{rank}(A) + n - r > r + n - r = n$$

which contradicts the rank-nullity theorem. Hence, it must be that  $\text{rank}(A) = r$ , so

$$\text{null}(A) = \text{span}(v_{r+1}, \dots, v_n)$$

and

$$\text{range}(A) = \text{span}(u_1, \dots, u_r)$$

as desired.

3. Let  $A$  be an  $m \times n$  matrix where  $m < n$  and rows of  $A$  are linearly independent. Then the system of linear equations  $Ax = b$  is underdetermined, i.e., infinitely many solutions. Among them, we want to find the one that has the minimum 2-norm. Check that the minimum 2-norm solution is given by

$$x^* = A^\top(AA^\top)^{-1}b.$$

*Hint: One way to solve this problem is the following. Check that  $x^*$  is a solution to  $Ax = b$ . Show that if  $x^* + y$  is also a solution of  $Ax = b$  then  $Ay = 0$ . Then check that the 2-norm of  $x^* + y$  is minimal if  $y = 0$ .*

**Solution:** We observe first that  $x^*$  as defined is indeed a solution to  $Ax = b$  since

$$Ax^* = A(A^\top(AA^\top)^{-1}b) = (AA^\top)(AA^\top)^{-1}b = b.$$

Now suppose that  $x^* + y$  is also a solution to  $Ax = b$  (noting that every solution to  $Ax = b$

can be written in this manner). Then it follows that

$$b = A(x^* + y) = Ax^* + Ay = b + Ay$$

and hence  $Ay = 0$ .

Finally, we claim that if  $x^* + y$  is a solution to  $Ax = b$ , then  $\|x^* + y\|_2^2 = \|x^*\|_2^2 + \|y\|_2^2$ . To see this, by definition of the 2-norm and the fact that vector addition is element-wise,

$$\|x^* + y\|_2^2 = \sum_{i=1}^n ((x^* + y)_i)^2 = \sum_{i=1}^n (x_i^* + y_i)^2 = \sum_{i=1}^n (x_i^*)^2 + 2(x_i^*)(y_i) + (y_i)^2.$$

But then splitting the summation, we find

$$\|x^* + y\|_2^2 = \sum_{i=1}^n (x_i^*)^2 + \sum_{i=1}^n 2(x_i^*)(y_i) + \sum_{i=1}^n (y_i)^2 = \|x^*\|_2^2 + 2 \sum_{i=1}^n (x_i^*)(y_i) + \|y\|_2^2$$

and further, the remaining sum is simply  $(x^*)^\top y$ , meaning

$$\|x^* + y\|_2^2 = \|x^*\|_2^2 + 2(x^*)^\top y + \|y\|_2^2.$$

But now note that

$$(x^*)^\top y = (A^\top (AA^\top)^{-1} b)^\top y = b^\top ((AA^\top)^{-1})^\top Ay = b^\top ((AA^\top)^\top)^{-1} Ay = b^\top (AA^\top)^{-1} Ay$$

and recall that earlier we showed  $Ay = 0$ , hence  $(x^*)^\top y = 0$  and so

$$\|x^* + y\|_2^2 = \|x^*\|_2^2 + \|y\|_2^2 \tag{3}$$

when  $x^* + y$  is a solution to  $Ax = b$ . Therefore, by the non-negativity of norms,

$$\|x^* + y\|_2^2 = \|x^*\|_2^2 + \|y\|_2^2 \geq \|x^*\|_2^2$$

and so taking the square root of both sides yields  $\|x^* + y\|_2 \geq \|x^*\|_2$ . Further, note that equality is attained if and only if  $\|y\|_2 = 0$ , or equivalently,  $y = 0$ .

Hence, we have shown that  $\|x^* + y\|_2$  is minimal exactly when  $y = 0$ , and since every solution to  $Ax = b$  can be written in this form, this completes the proof.

4. Let  $A$  be a  $3 \times 3$  matrix, and let  $T$  be its Schur form, i.e., there is a unitary matrix  $Q$  (i.e.,  $Q^*Q = QQ^* = I$  where  $Q^*$  denotes the transpose and complex conjugate of  $Q$ ) such that

$$A = QTQ^*, \text{ where } T = \begin{bmatrix} \lambda_1 & t_{12} & t_{13} \\ 0 & \lambda_2 & t_{23} \\ 0 & 0 & \lambda_3 \end{bmatrix}.$$

Assume that  $\lambda_j, j = 1, 2, 3$  are all distinct.

- (a) Show that if  $v$  is an eigenvector of  $T$  then  $Qv$  is the eigenvector of  $A$  corresponding to the same eigenvalue.

**Solution:** Let  $v$  be an eigenvector of  $T$  with corresponding eigenvalue  $\lambda$ , such that  $Tv = \lambda v$ . Then since  $Q^*Q = I$ , we have that

$$A(Qv) = (QTQ^*)Qv = QT(Q^*Q)v = QTv = Q\lambda v = \lambda(Qv),$$

and so  $Qv$  is an eigenvector of  $A$  also corresponding to  $\lambda$ .

- (b) Find the eigenvectors of  $T$ . *Hint: Check that  $v_1 = [1, 0, 0]^\top$ . Look for  $v_2$  of the form  $v_2 = [a, 1, 0]^\top$ , and then for  $v_3$  of the form  $v_3 = [b, c, 1]^\top$ , where  $a, b, c$  are to be expressed via the entries of the matrix  $T$ .*

**Solution:** Since  $T$  is (upper) triangular, its eigenvalues are its diagonal entries  $\lambda_1$ ,  $\lambda_2$ , and  $\lambda_3$ . Hence, we want to find vectors  $v_1$ ,  $v_2$ , and  $v_3$  such that

$$Tv_1 = \lambda_1 v_1, \quad Tv_2 = \lambda_2 v_2, \quad Tv_3 = \lambda_3 v_3.$$

That is, for each  $1 \leq i \leq 3$ , we want to find  $a_i$ ,  $b_i$ , and  $c_i$  such that

$$\begin{aligned} \begin{bmatrix} \lambda_1 & t_{12} & t_{13} \\ 0 & \lambda_2 & t_{23} \\ 0 & 0 & \lambda_3 \end{bmatrix} \begin{bmatrix} a_i \\ b_i \\ c_i \end{bmatrix} - \lambda_i \begin{bmatrix} a_i \\ b_i \\ c_i \end{bmatrix} &= \begin{bmatrix} \lambda_1 a_i + t_{12} b_i + t_{13} c_i \\ \lambda_2 b_i + t_{23} c_i \\ \lambda_3 c_i \end{bmatrix} - \begin{bmatrix} \lambda_i a_i \\ \lambda_i b_i \\ \lambda_i c_i \end{bmatrix} \\ &= \begin{bmatrix} (\lambda_1 - \lambda_i) a_i + t_{12} b_i + t_{13} c_i \\ (\lambda_2 - \lambda_i) b_i + t_{23} c_i \\ (\lambda_3 - \lambda_i) c_i \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}. \end{aligned}$$

Then for  $i = 1$ ,  $c_1 = 0$  since  $(\lambda_3 - \lambda_1)c_1 = 0$ . Further, then  $(\lambda_2 - \lambda_1)b_1 + t_{23}c_1 = 0$ , it must be that  $b_1 = 0$ . Finally, we see that  $a_1$  is free, and hence an eigenvector corresponding to  $\lambda_1$  is

$$v_1 = [1, 0, 0]^\top.$$

Now repeating the same procedure for  $i = 2$ , again we find that  $c_2 = 0$ , but now  $b_2$  is free since any choice will satisfy  $(\lambda_2 - \lambda_2)b_2 + t_{23}c_2 = 0$ . Hence, from the first equation it follows that

$$a_2 = \frac{t_{12}b_2}{\lambda_2 - \lambda_1}$$

and so taking  $b_2 = 1$  yields that an eigenvector corresponding to  $\lambda_2$  is

$$v_2 = \begin{bmatrix} \frac{t_{12}}{\lambda_2 - \lambda_1} & 1 & 0 \end{bmatrix}^\top.$$

Finally, for  $i = 3$ , we see that  $c_3$  is free since any choice will satisfy the third equation  $(\lambda_3 - \lambda_3)c_3 = 0$ . Taking  $c_3 = 1$ , from the second equation we have that

$$b_3 = \frac{t_{23}}{\lambda_3 - \lambda_2}$$

and hence from the first equation, we have that

$$a_3 = \frac{t_{12}b_3 + t_{13}}{\lambda_3 - \lambda_1} = \frac{t_{12}t_{23}}{(\lambda_3 - \lambda_1)(\lambda_3 - \lambda_2)} + \frac{t_{13}}{\lambda_3 - \lambda_1}$$



and hence an eigenvector corresponding to  $\lambda_3$  is

$$v_3 = \left[ \frac{t_{12}t_{23}}{(\lambda_3 - \lambda_1)(\lambda_3 - \lambda_2)} + \frac{t_{13}}{\lambda_3 - \lambda_1} \quad \frac{t_{23}}{\lambda_3 - \lambda_2} \quad 1 \right]^\top.$$

- (c) Write out the eigenvectors of  $A$  in terms of the found eigenvectors of  $T$  and the columns of  $Q$ :  $Q = [q_1, q_2, q_3]$ .

**Solution:** By item (a), we know that if  $v$  is an eigenvector of  $T$  then  $Qv$  is an eigenvector of  $A$  corresponding to the same eigenvalue. Hence, using the eigenvectors of  $T$  found in item (b),

$$v'_1 = Qv_1 = q_1$$

$$v'_2 = Qv_2 = \left( \frac{t_{12}}{\lambda_2 - \lambda_1} \right) q_1 + q_2$$

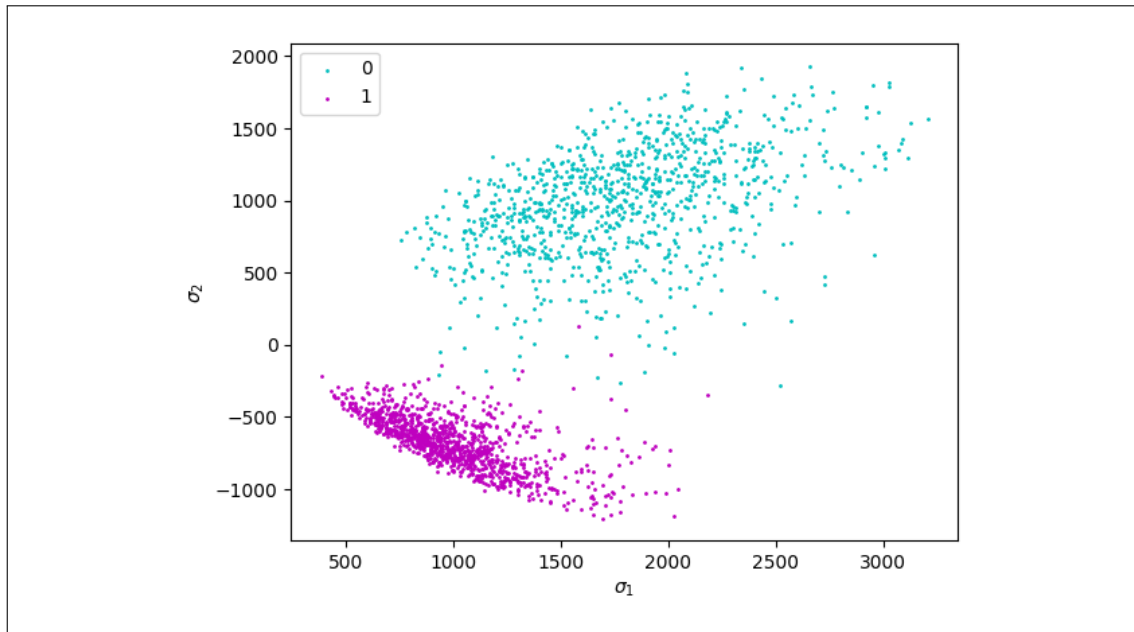
$$v'_3 = Qv_3 = \left( \frac{t_{12}t_{23}}{(\lambda_3 - \lambda_1)(\lambda_3 - \lambda_2)} + \frac{t_{13}}{\lambda_3 - \lambda_1} \right) q_1 + \left( \frac{t_{23}}{\lambda_3 - \lambda_2} \right) q_2 + q_3$$

are eigenvectors of  $A$  corresponding to eigenvalues  $\lambda_1$ ,  $\lambda_2$ , and  $\lambda_3$  respectively.

5. Download the **MNIST** dataset. For your convenience, I prepared it as the `mnist.mat` file. This file contains 60,000 training images and 10,000 test images of handwritten digits from 0 to 9, and labels for the training and test images. Each image is 28-by-28 pixels because I stripped off paddings with zeros. You can use Matlab or Python.

- (a) Convert the set of test images into a matrix  $A$  of size  $10^4 \times 400$ . Compute an SVD of this matrix. Project the data onto the space spanned by the first two right singular vectors  $v_1$  and  $v_2$ . Display only points corresponding to digits 0 and 1 in this 2D space and color the points corresponding to 1 and 0 in different colors. Check if 0s and 1s cluster in this 2D space.

**Solution:** The points corresponding to 0s and 1s indeed appear to cluster in this 2D space; there is a clear decision boundary between them.



- (b) Do this task for  $k = 10, 20$ , and  $50$ . Compute  $A_k = U_k \Sigma_k V_k^\top$ , where  $U_k$  ( $V_k$ ) is comprised of the first  $k$  left (right) singular vectors, and  $\Sigma_k = \text{diag}\{\sigma_1, \dots, \sigma_k\}$ . Then take the first four rows of  $A$  and  $A_k$ , reshape each of these rows back to  $20 \times 20$  images, and display them.

**Solution:** This diagram and the one above were produced by [this code](#).

