

Question 2:

```
public class quest2 {  
  
    public static String question2(String word1, String word2) {  
  
        int w1L = word1.length();  
        int w2L = word2.length();  
        int[][] table = new int[w1L + 1][w2L + 1];  
        int maxLength = 0;  
        int endIndex = 0;  
  
        for (int i = 1; i <= w1L; i++) {  
            for (int j = 1; j <= w2L; j++) {  
                if (word1.charAt(i-1) == word2.charAt(j-1)) {  
                    table[i][j] = 1 + table[i-1][j-1];  
  
                    if (table[i][j] > maxLength) {  
                        maxLength = table[i][j];  
                        endIndex = i;  
                    }  
                }  
            }  
            else {  
                table[i][j] = 0;  
            }  
        }  
        if (maxLength == 0) {  
            return "No substring";  
        }  
    }  
}
```

```

        return word1.substring(endIndex - maxLength, endIndex);

    }

public static void main(String[] args) {
    System.out.println(question2("gears of war", "history of war"));
    System.out.println(question2("spy family", "pickle"));
    System.out.println(question2("hi", "talk"));

}

}

```

Question 6:

Algorithm 1:

Time Complexity: My first algorithm is Big-O($w1L \times w2L$) and Big- $\Omega(w1L \times w2L)$. Where $w1L$ is the length of word1 and $w2L$ is the length of word2. There is a nested for loop where the outer loop runs $w1L$ times and a inner loop that runs $w2L$ times. We go to each grid exactly once which is $w1L \times w2L$.

Space Complexity: We initialize a 2 dimensional array that is of size $w1L \times w2L$.

This means that our algorithm is Big-O($w1L \times w2L$) and Big- $\Omega(w1L \times w2L)$. Each cell gets written into so it is Big-O($w1L \times w2L$) and Big- $\Omega(w1L \times w2L)$.

Algorithm 2:

Time Complexity: This algorithm is also Big-O($w1L \times w2L$) and Big- $\Omega(w1L \times w2L)$. Where $w1L$ is the length of word1

and $w2L$ is the length of word2. There is a nested for loop where the outer loop runs $w1L$ times and a inner loop that runs $w2L$ times.

Space Complexity: This is once again the same and is Big-O($w1L \times w2L$) and Big- $\Omega(w1L \times w2L)$ since we are creating a

$w1L \times w2L$ array and filling each square.

Algorithm 3:

Time Complexity:

The time complexity of this will be Big-O(N) and Big- Ω (N). This is because there is one for loop that runs

numTerms where numTerms is the length of the sequence. It is followed by another loop that also runs numTerms where

numTerms is the length of the sequence. $N + N = 2N$ and we can just cut off the constants.

Space Complexity:

The space complexity is Big-O(N) and Big- Ω (N). This is because we initialize a long array that is the size of the sequence and fill it with constants. This is the biggest use of ram.

Algorithm 4:

Time Complexity:

The time complexity of this is Big-O(KL) and Big- Ω (K) where K is the index of the number that is equal or > than

the target because it will run until it gets to either one of those conditions.

Space Complexity:

The space complexity is Big-O(1) and Big- Ω (1) because we only create primitive data types and never

any array that would grow and not be constant. 1 is the constant space. we can factor out a 1 from how many times we intialize a primitive data type.

Algorithm 5:

Time Complexity:

The time complexity is Big-O(N) and Big- Ω (N) because inside of the logic we have a for loop that runs

nums.length where that is the length of the array (the input). It checks each number in the array which is why it happens

N times.

Space Complexity:

The space complexity is Big-O(1) and Big- Ω (1) because we are only ever writing into the array in the logic.

There is no point where we make an array in the logic it is only passed to us.