

Question 2:

```
public class quest2 {

    public static String question2(String word1, String word2) {

        int w1L = word1.length();
        int w2L = word2.length();

        int[][] table = new int[w1L + 1][w2L + 1];

        int maxLength = 0;
        int endIndex = 0;

        for (int i = 1; i <= w1L; i++) {
            for (int j = 1; j <= w2L; j++) {
                if (word1.charAt(i-1) == word2.charAt(j-1)) {
                    table[i][j] = 1 + table[i-1][j-1];

                    if (table[i][j] > maxLength) {
                        maxLength = table[i][j];
                        endIndex = i;
                    }
                }
            }
        }

        if (maxLength == 0) {
            return "No substring";
        }
    }
}
```

```

        return word1.substring(endIndex - maxLength, endIndex);
    }

    public static void main(String[] args) {
        System.out.println(question2("gears of war", "history of war"));
        System.out.println(question2("spy family", "pickle"));
        System.out.println(question2("hi", "talk"));

    }
}

```

Question 6:

Algorithm 1:

Time Complexity: My first algorithm is $\text{Big-O}(w1L \times w2L)$ and $\text{Big-}\Omega(w1L \times w2L)$. Where $w1L$ is the length of word1 and $w2L$ is the length of word2. There is a nested for loop where the outer loop runs $w1L$ times and a inner loop that runs $w2L$ times. We go to each grid exactly once which is $w1L \times w2L$.

Space Complexity: We initialize a 2 dimensionial array that is of size $w1L \times w2L$.

This means that our algorithm is $\text{Big-O}(w1L \times w2L)$ and $\text{Big-}\Omega(w1L \times w2L)$. Each cell gets written into so itis $\text{Big-O}(w1L \times w2L)$ and $\text{Big-}\Omega(w1L \times w2L)$.

Algorithm 2:

Time Complexity: This algorithm is also $\text{Big-O}(w1L \times w2L)$ and $\text{Big-}\Omega(w1L \times w2L)$. Where $w1L$ is the length of word1

and $w2L$ is the length of word2. There is a nested for loop where the outer loop runs $w1L$ times and a inner loop that runs $w2L$ times.

Space Complexity: This is once again the same and is $\text{Big-O}(w1L \times w2L)$ and $\text{Big-}\Omega(w1L \times w2L)$ since we are creating a

$w1L \times w2L$ array and filling each square.

Algorithm 3:

Time Complexity:

The time complexity of this will be $\text{Big-O}(N)$ and $\text{Big-}\Omega(N)$. This is because there is one for loop that runs

`numTerms` where `numTerms` is the length of the sequence. It is followed by another loop that also runs `numTerms` where

`numTerms` is the length of the sequence. $N + N = 2N$ and we can just cut off the constants.

Space Complexity:

The space complexity is $\text{Big-O}(N)$ and $\text{Big-}\Omega(N)$. This is because we initialize a long array that is the size of the sequence and fill it with constants. This is the biggest use of ram.

Algorithm 4:

Time Complexity:

The time complexity of this is $\text{Big-O}(KL)$ and $\text{Big-}\Omega(K)$ where K is the index of the number that is equal or $>$ than

the target because it will run until it gets to either one of those conditions.

Space Complexity:

The space complexity is $\text{Big-O}(1)$ and $\text{Big-}\Omega(1)$ because we only create primitive data types and never

any array that would grow and not be constant. 1 is the constant space. we can factor out a 1 from how many times we initialize a primitive data type.

Algorithm 5:

Time Complexity:

The time complexity is $\text{Big-O}(N)$ and $\text{Big-}\Omega(N)$ because inside of the logic we have a for loop that runs

`nums.length` where that is the length of the array (the input). It checks each number in the array which is why it happens

N times.

Space Complexity:

The space complexity is $\text{Big-O}(1)$ and $\text{Big-}\Omega(1)$ because we are only ever writing into the array in the logic.

There is no point where we make an array in the logic it is only passed to us.