

What is functions?

Set of instructions which we can use repeatedly to avoid Don not repeat yourself.

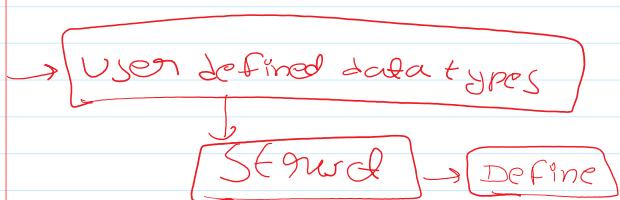
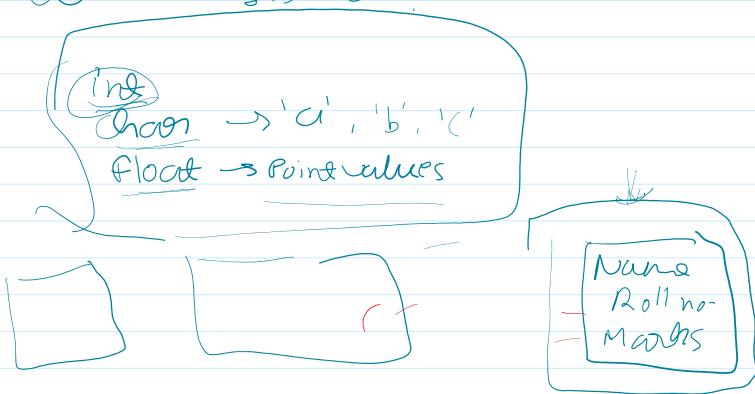
Programming Principle

Don't repeat code

Three parts:

- ① Function Declaration
- ② Function Definition
- ③ Function calling

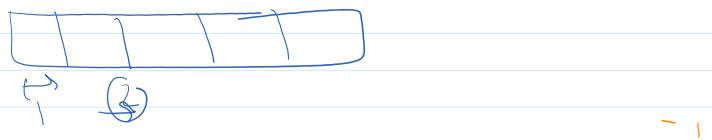
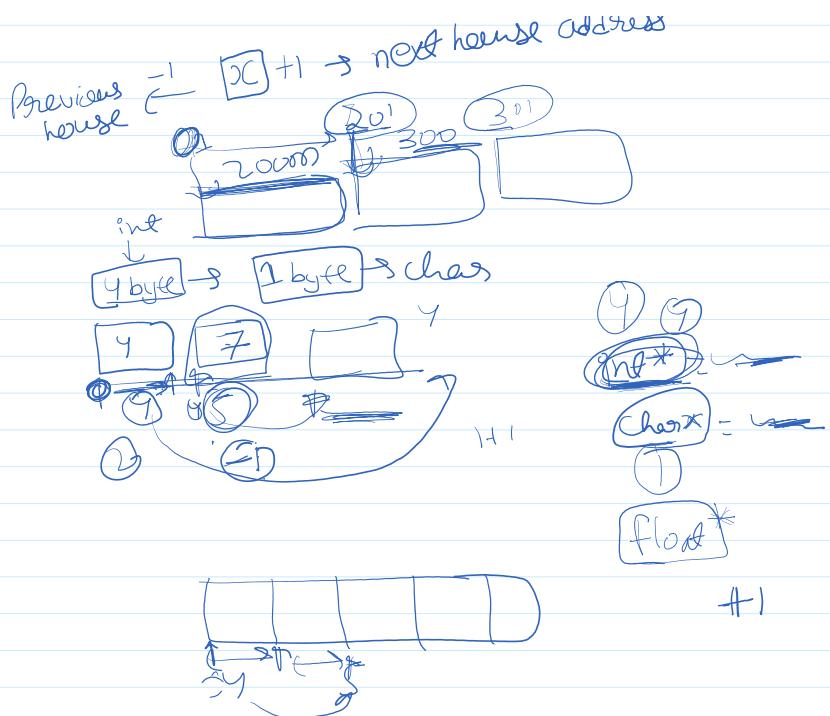
What is struct?



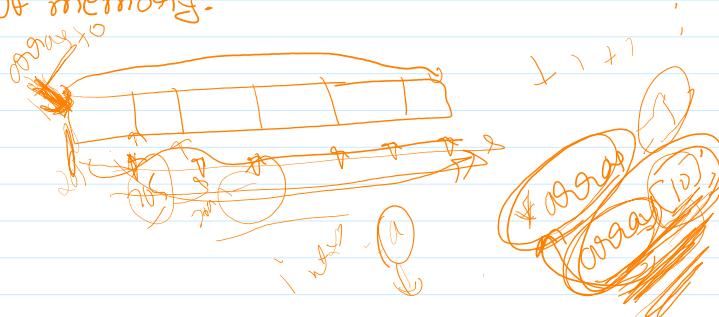
Pointers:

↳ Special type variable which is used to store address of another Data Obj.

Pointer Arithmetic:

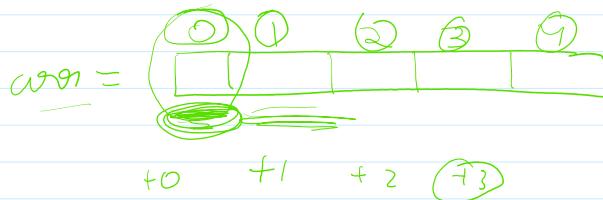


Array  $\rightarrow$  Array is a contiguous blocks  
of memory.



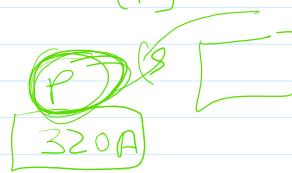
Index in arrays

↓



$$arr[0] = 10; \quad arr[1]$$

into



Int a = 5;

String in C :)

char  
Collection of characters.

Char c = 'a';  
Char b = 'b';

Rithik

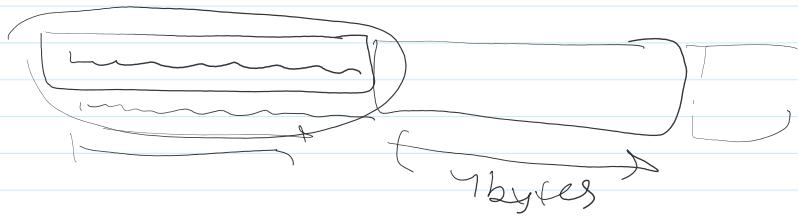
Binary → 01 → on/off

bits

Byte

1 Byte = 8 bits

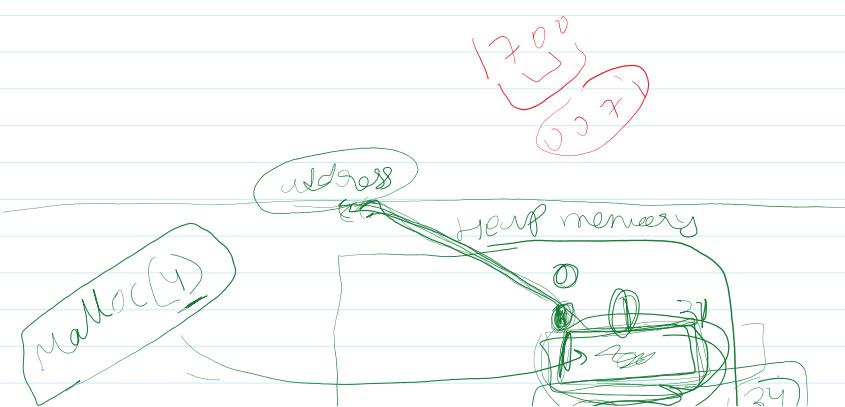
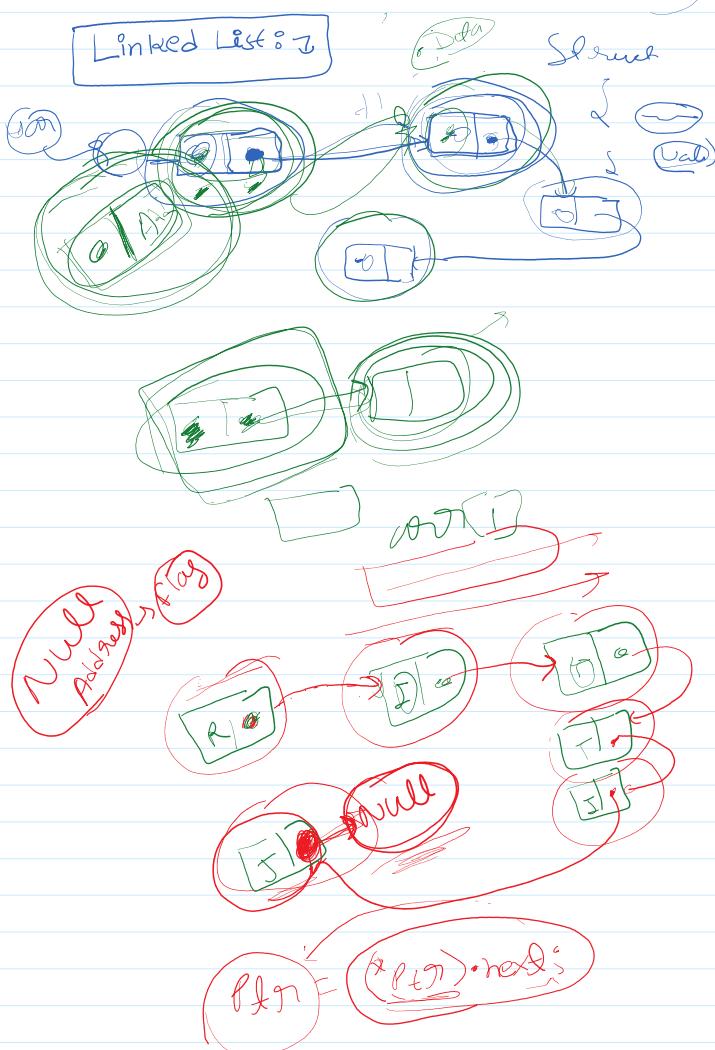
1 byte = 32 bits

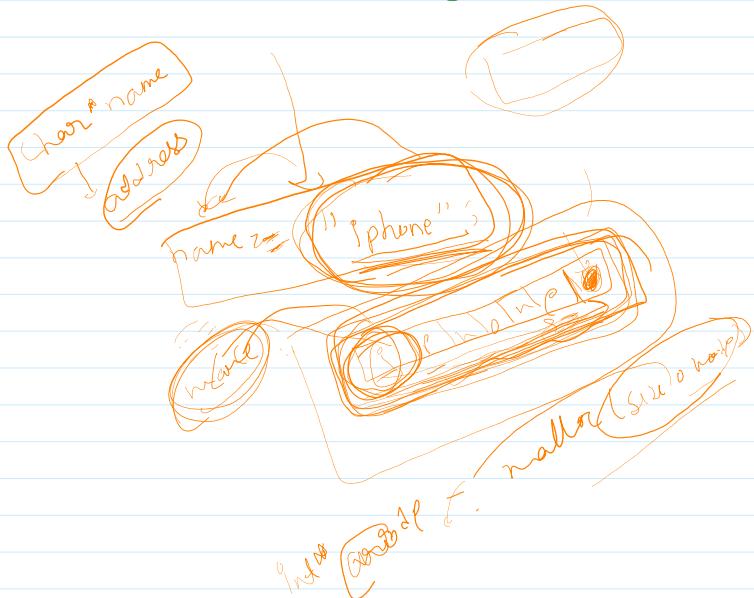
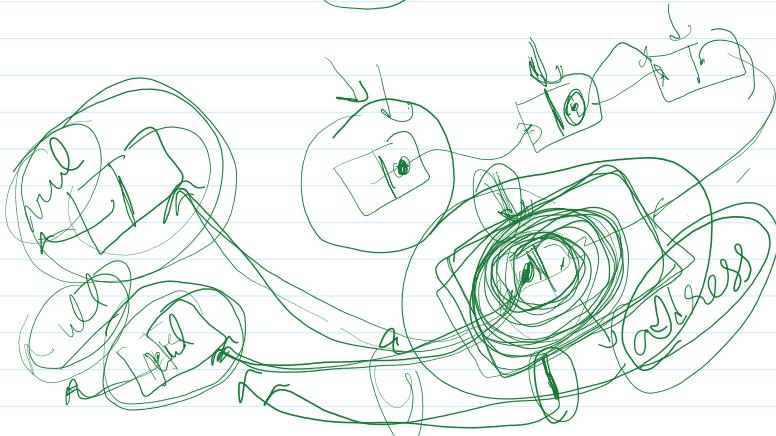
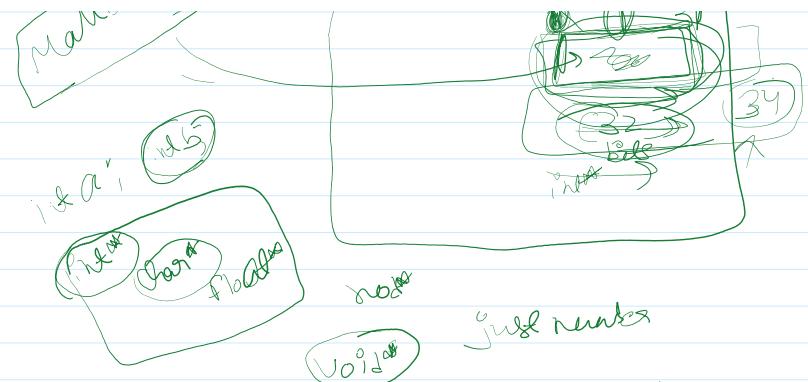


What is data structures?

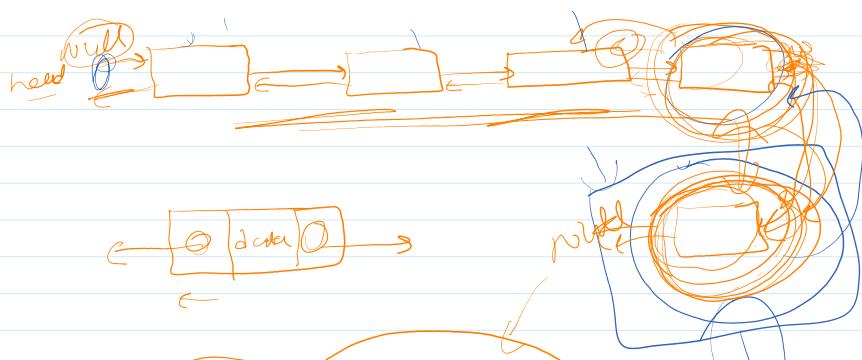
↪ The coarrangement of data in memory is called data structures.

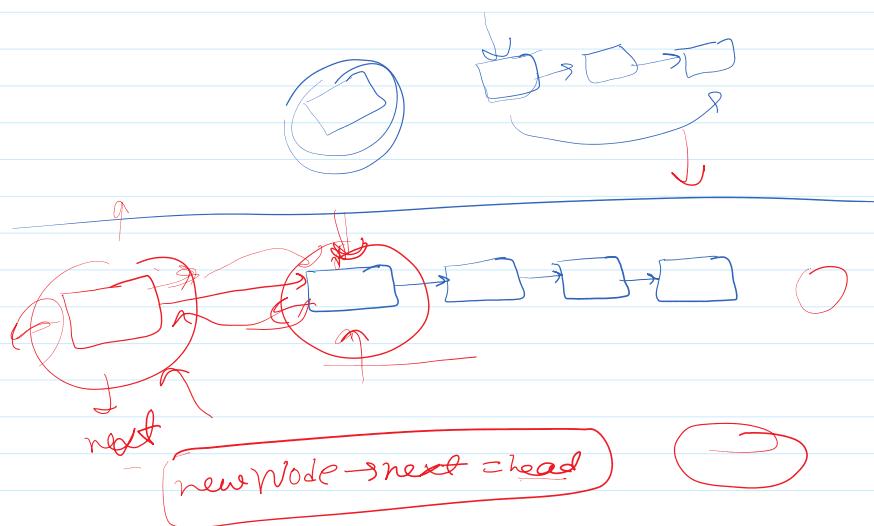
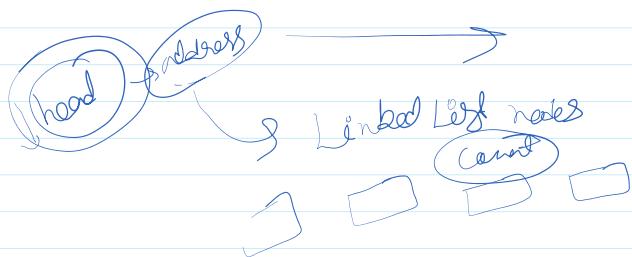
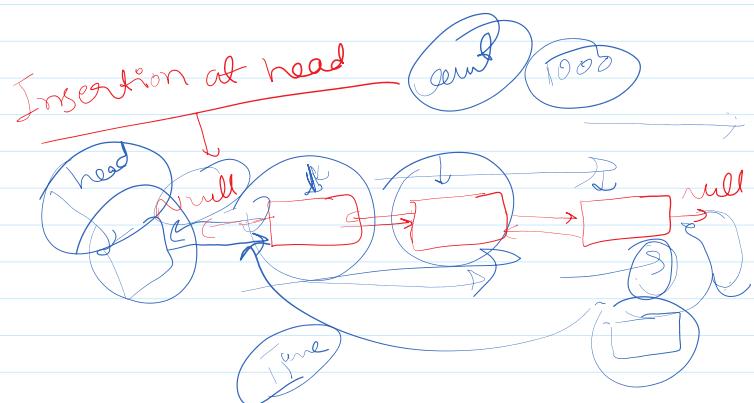
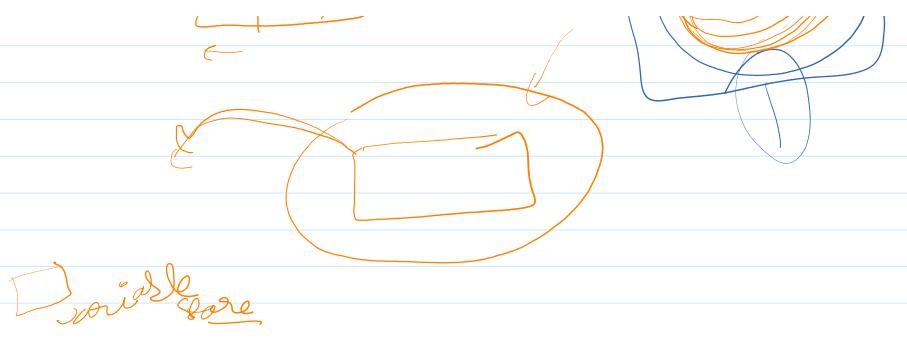
Linked List :-



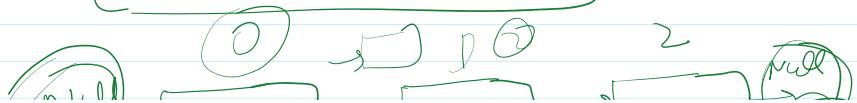


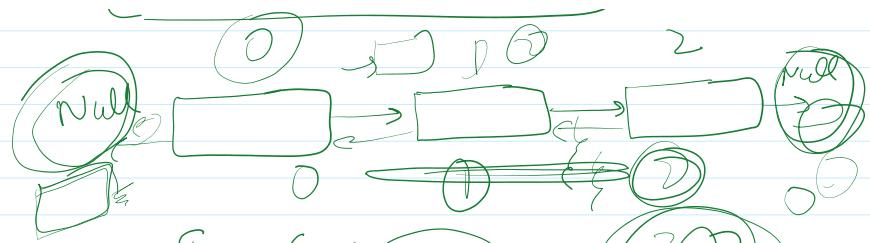
Doubly linked list :-





Insert at a given index

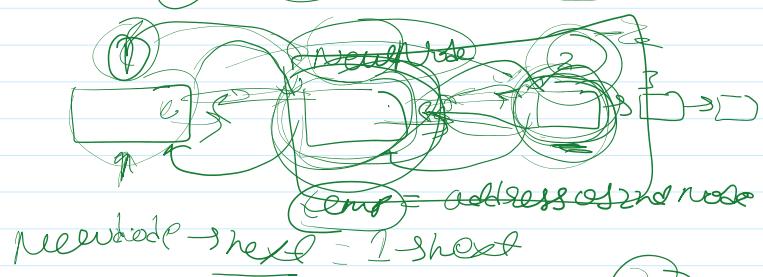




`Func(int index)`



①  
②

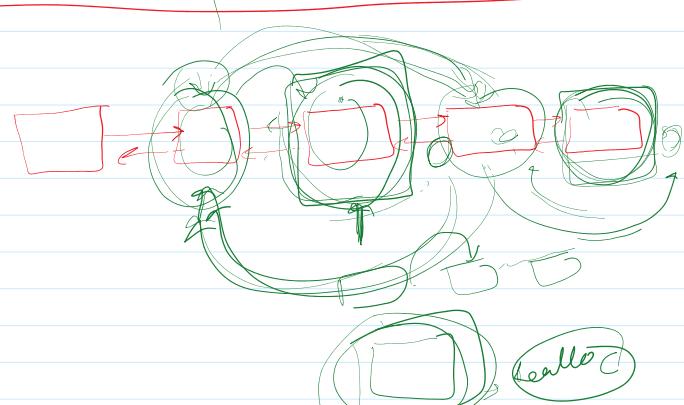
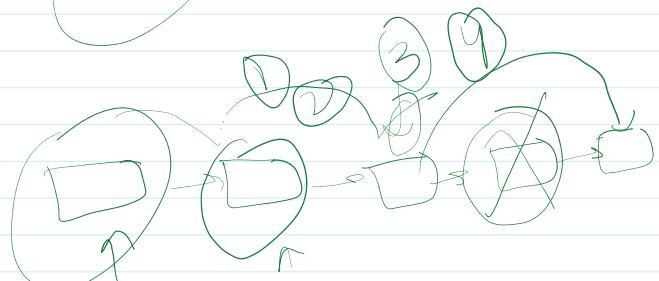
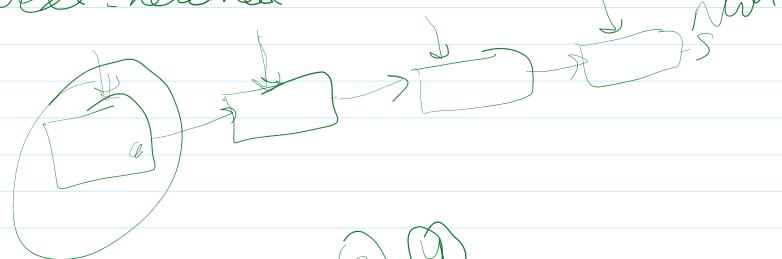


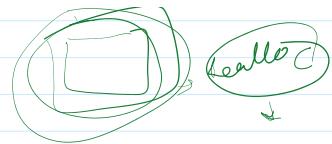
newNode → next = 1 - sheet

1 - sheet → prev = newNode

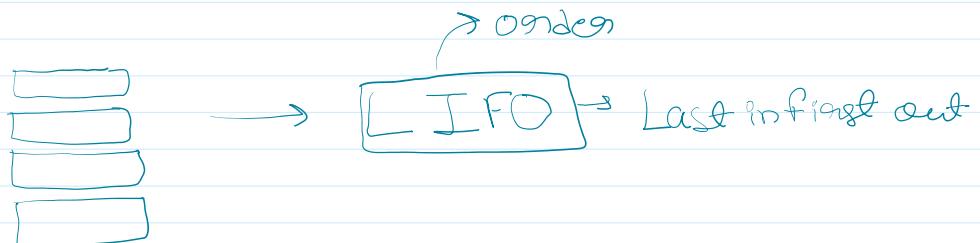
`newPrev = 1`

`1 - next = newNext`

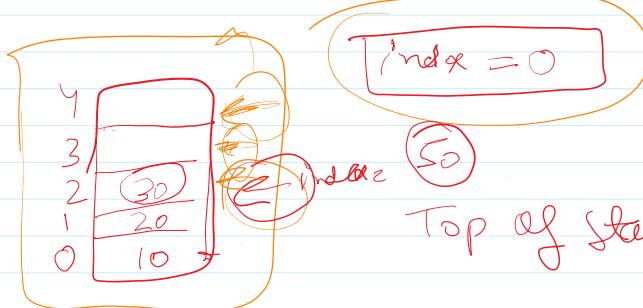
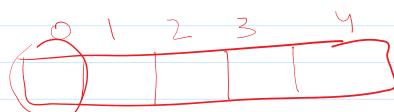
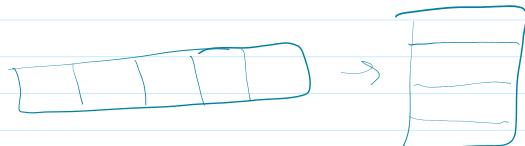




Stack :-



array



size

Index, arr →

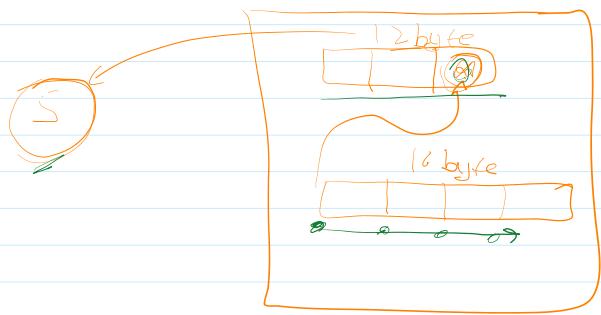
↓ ↓

Scansit one

↑

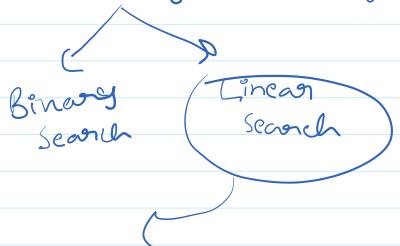
array, index

Presentation of stack in loop

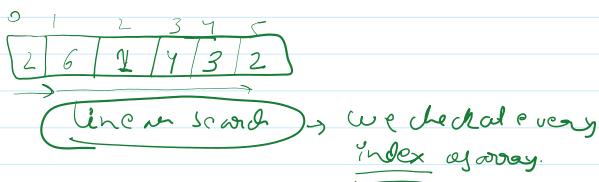


Int arr = (int\*) malloc(sizeof(int)\*4);

Searching :- (algorithm)

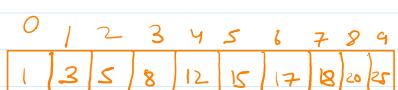
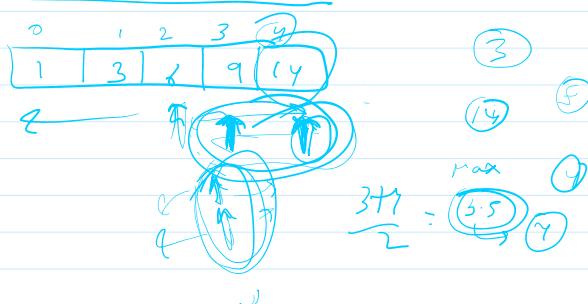


Linear search :-



Binary search

Numbers should be sorted



target

18

L M R

0 4 9

5 7 9

①

$\frac{5+9}{2} = 7$

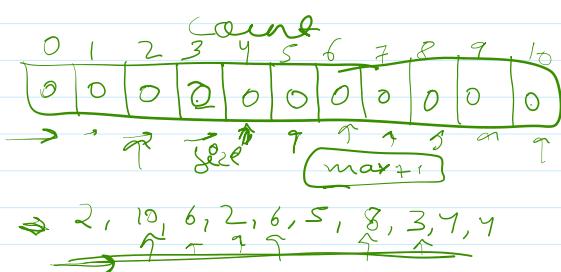
$\left(\frac{1+9}{2}\right) \text{, } n = 6 \text{, } \text{array} := \text{true}$

②

Sorting :-

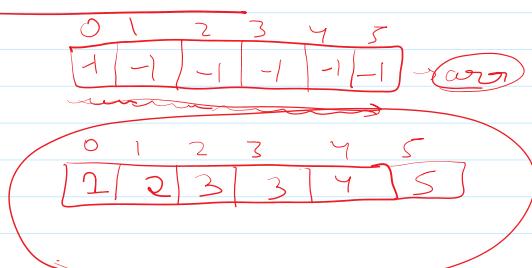
2, 4, 5, 7, 11 → ↘

Count sort :-



2 2 3 4 5 6 6 8 10

Selection sort :-



mini = ∞

10, 11, 8, 12

mini = 18

Bubble sort :-

array :- 5 13 6 7 6

5 6 7 7 13

(1) → no swap

if i > then i+1  
we swap

5 13 8 6 1

↑

5 8 13 6 1

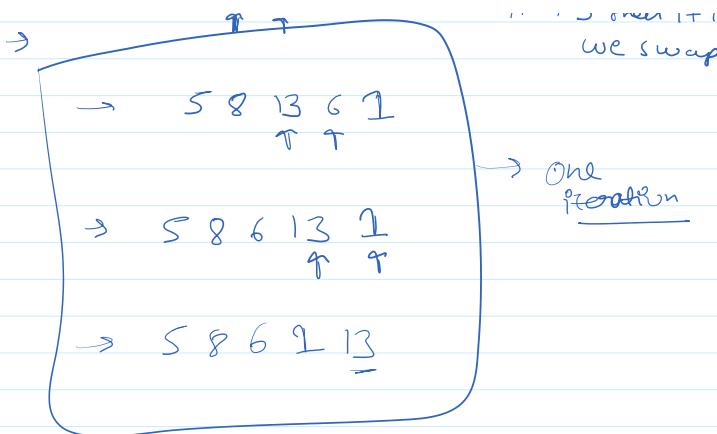
↑ ↑

5 8 6 13 1

↑ ↑

5 8 6 1 13

→ One iteration



→ one iteration  
 we swap

2nd iteration

$$\rightarrow \begin{matrix} 5 & 8 & 6 & 1 & 13 \\ & \downarrow & & & \\ & 1 & & & \end{matrix}$$

$$\begin{matrix} 5 & 8 & 6 & 1 & 13 \\ & \downarrow & \downarrow & & \end{matrix}$$

$$\begin{matrix} 5 & 6 & 8 & 1 & 13 \\ & \downarrow & \downarrow & & \end{matrix}$$

$$\rightarrow 5 6 1 \textcircled{8} 13$$

3rd iteration

$$\begin{matrix} 5 & 6 & 1 & 8 & 13 \\ & \downarrow & \downarrow & & \end{matrix}$$

$$5 1 \textcircled{6} 8 13$$

4th iteration

$$\begin{matrix} 5 & 1 & 6 & 8 & 13 \\ & \uparrow & & & \end{matrix}$$

1 5 6 8 13 sorted

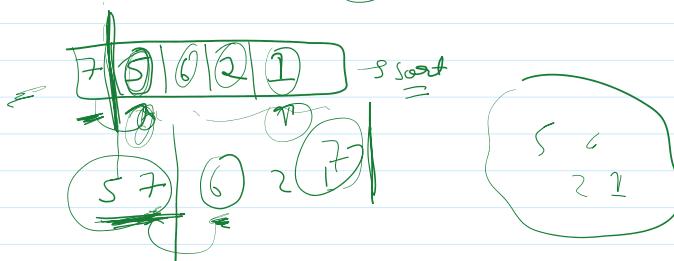
Insertion Sort :-

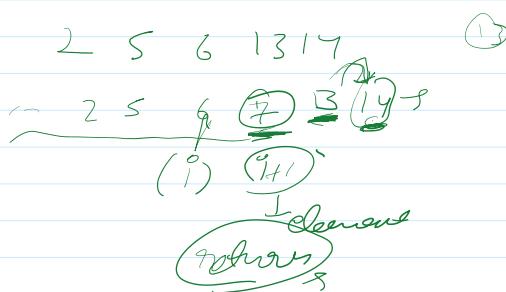
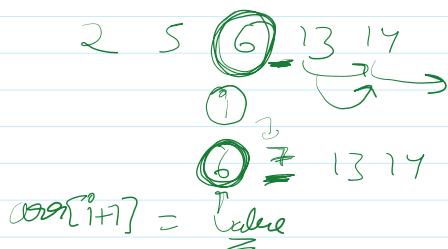
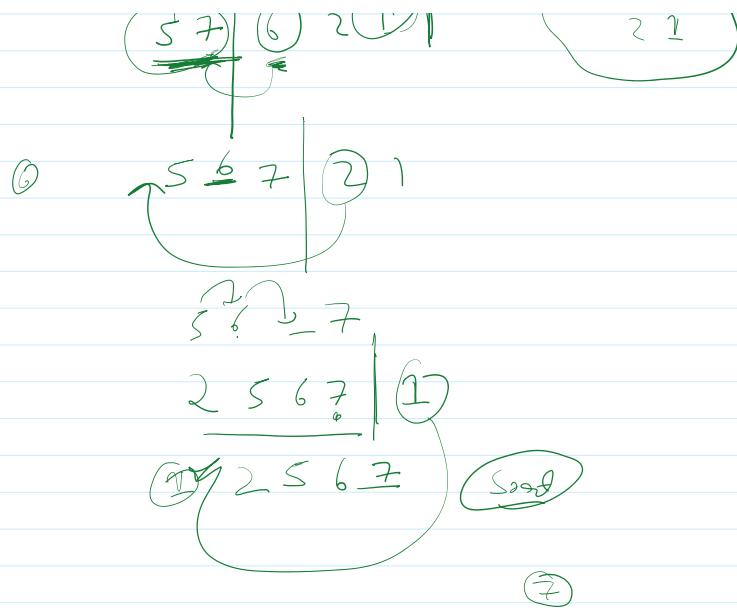


$$\rightarrow [2, 4, \underline{5}, 6, 7]$$

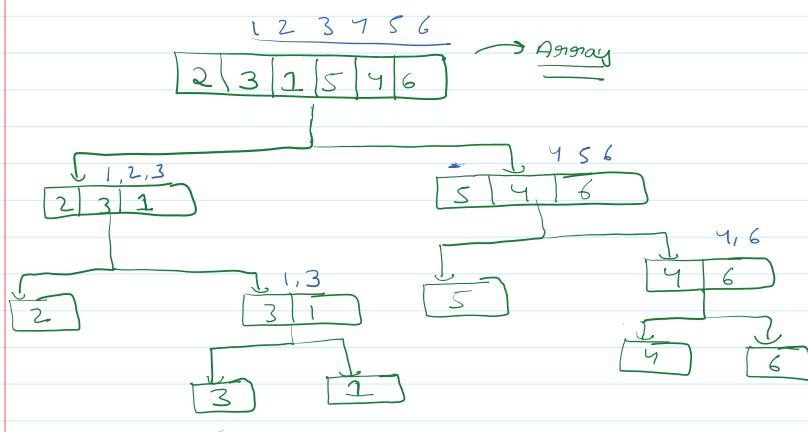
$$\begin{matrix} 4 & 2 & 7 & 6 \\ & \swarrow & & \end{matrix}$$

$$4 2 \textcircled{5} 7 6$$



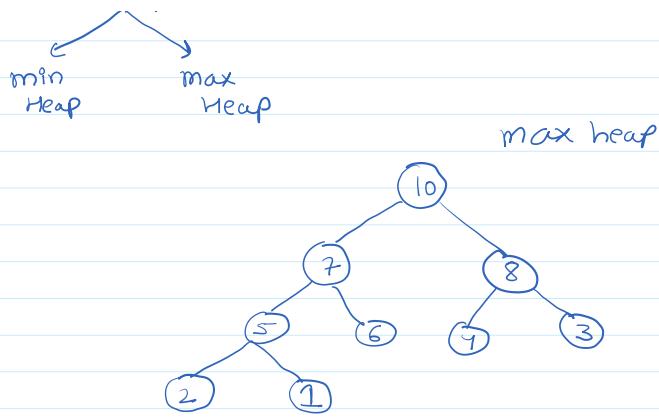


Merge Sort → Divide and Conquer



Heap Sort :-

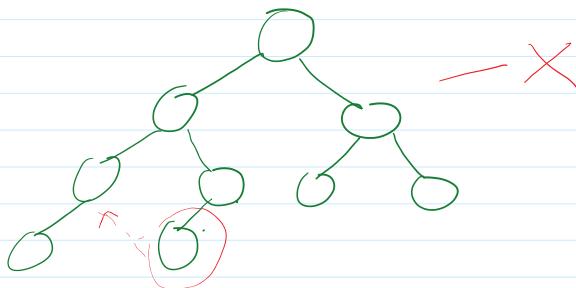
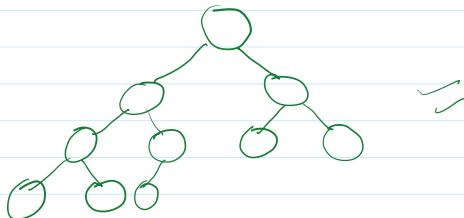




Each parent in heap should be greater than its children.

- Heap is Binary Tree
- Heap is special Binary tree which is Complete Binary tree.

e.g. ↴

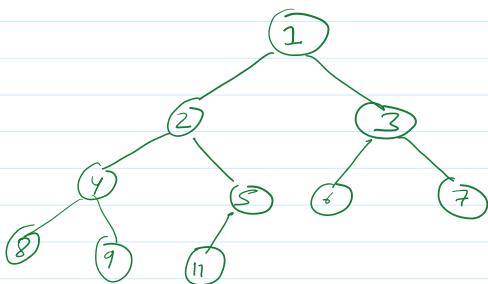


Complete Binary tree + Parent - Child Relationship

↓  
Heap

min-heap :-

It is also a complete binary tree.

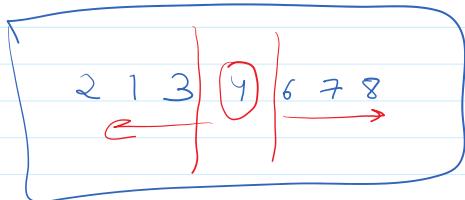


Root node is always greater than everyone in  
Max-heap.

Selection Sort + heap = heap Sort

Quick Sort :-

Divide and conquer



Dry Run

7 2 5 1 6  $\Rightarrow$  curr Pivot

temp[5];  
for (int i = 0; i < size; i++)

if (arr[i] < pivot)

{ temp[index] = arr[i];  
index++;

{

temp = 2 5 1

temp[index] = pivot;  
↓

temp = 2 5 1 6

for (int i = 0; i < size; i++)

{ if (arr[i] > pivot)  
{

temp[index] = arr[i];  
index++;

{

temp = 2 5 1 6 7  
← → i

→ new arr

2 5 ① → Pivot

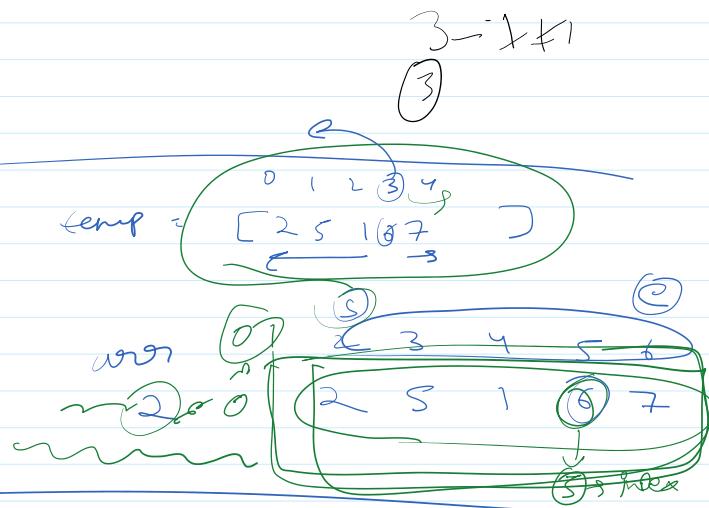
temp[0] = 2; Pivot

temp → 1

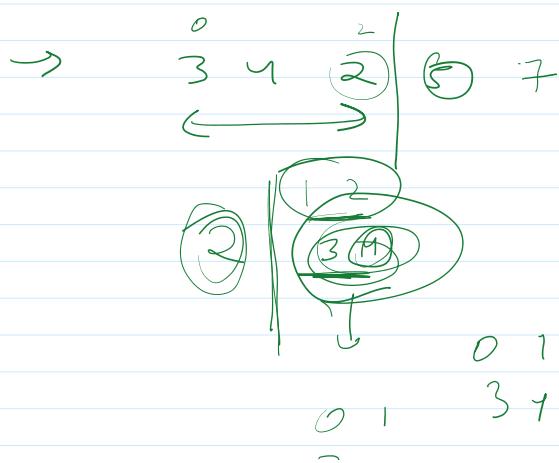
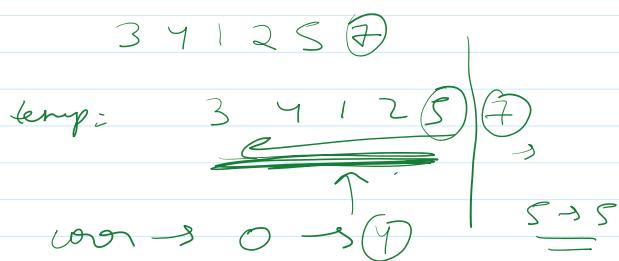
large numbers than pivot

temp → [1, 2, 5]

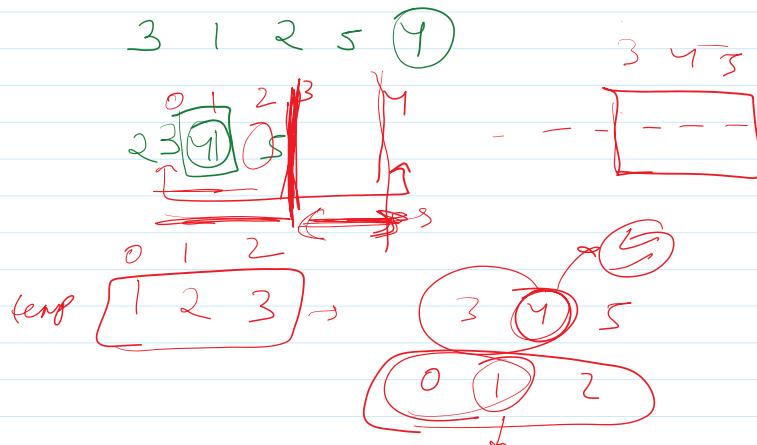
Results: → [1 2 5 6 7] Sorted



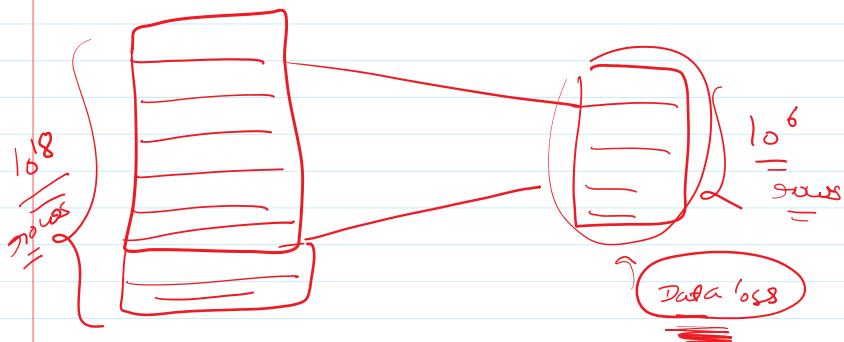
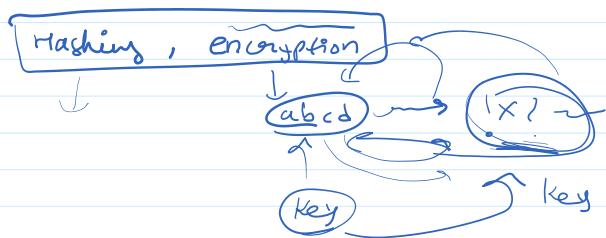
dry run : 1



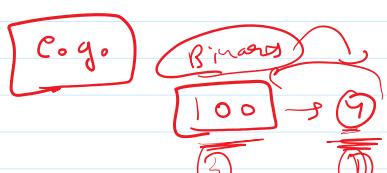
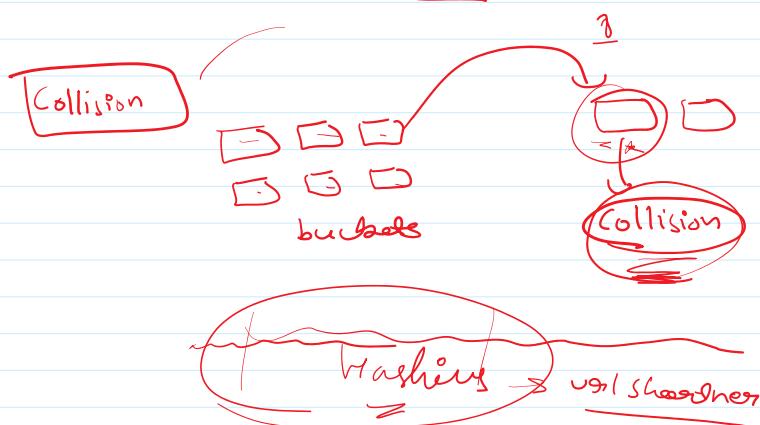
3 4

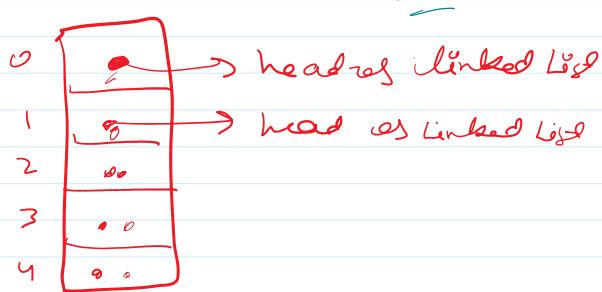
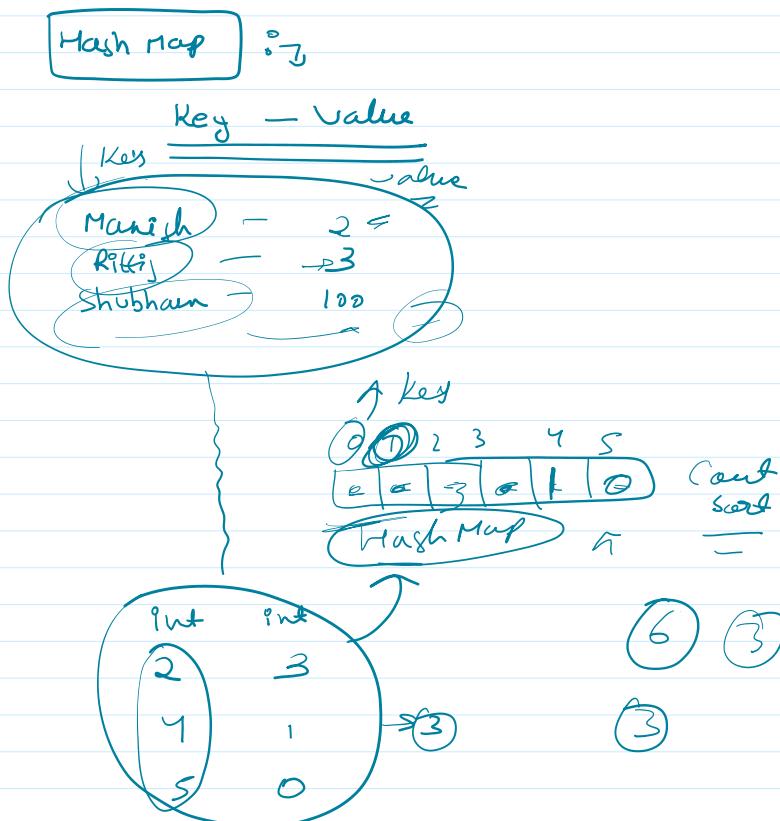
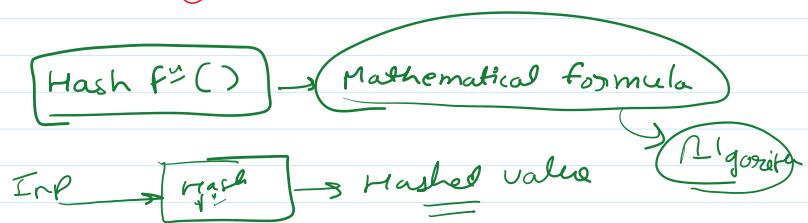


### Hashing:-



Mapping a large set of data into a small set  $\rightarrow$  Hashing

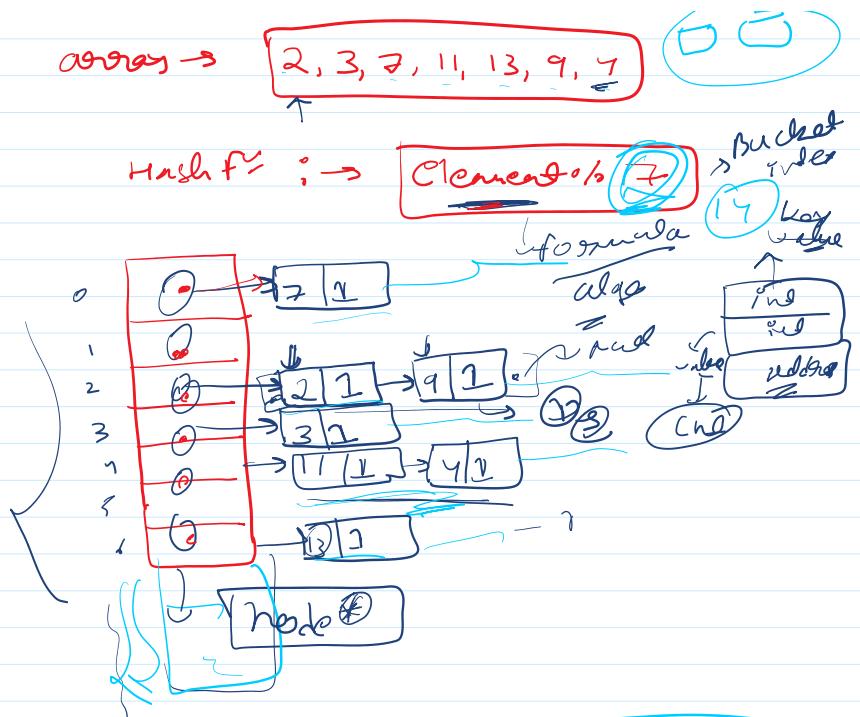




and **Separate chaining**



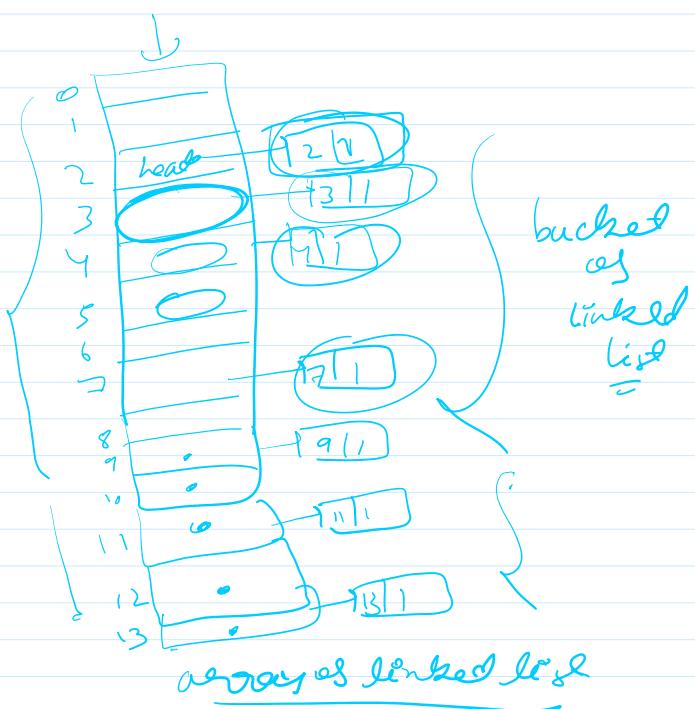
array → 2, 3, 7, 11, 13, 9, 7



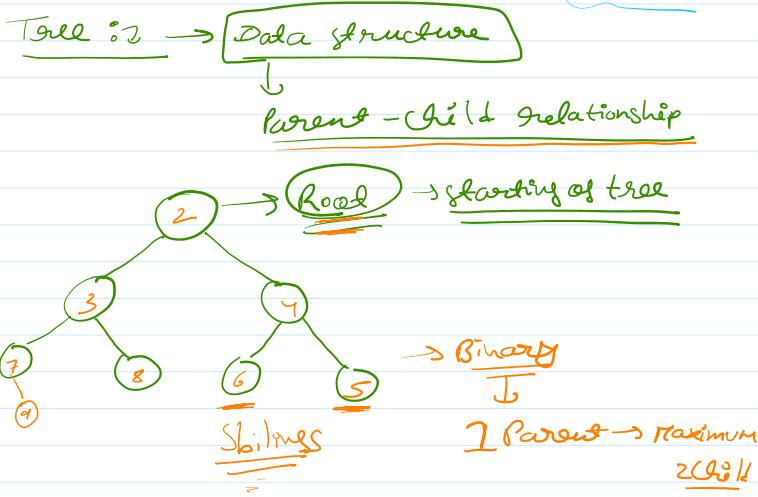
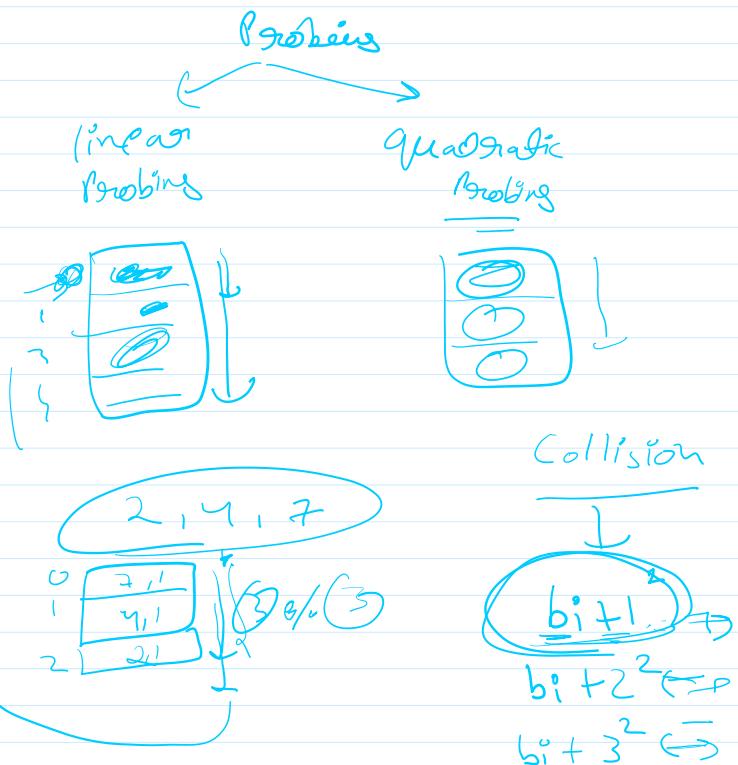
Load factor : → amount data in hash tab  
Size of array (7) =  
0.5 → Rehashing

Rehashing : → array size × 2 → copy

Hash fn = Element % ~~Size of array~~

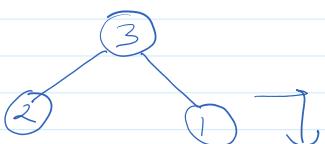


↗ array of linked list  
separate chaining → //  
Probing ↗ ~~3x~~



Traversal Techniques:

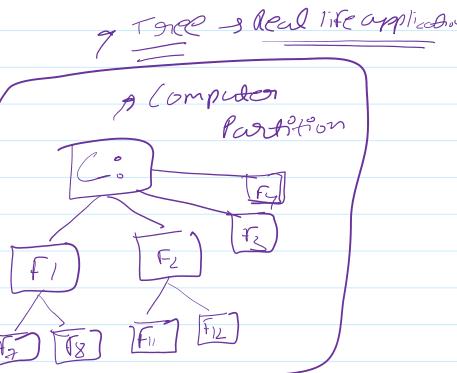
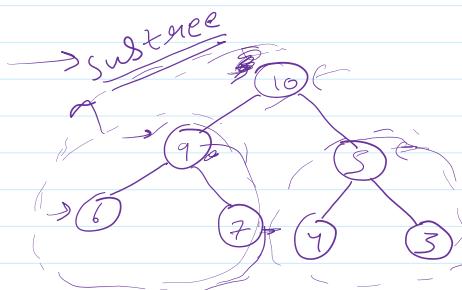
- (i) Preorder → Root → Left → Right
- (ii) Inorder → Left → Root → Right
- (iii) Postorder → Left → Right → Root



(i) Preorder:  $\underline{\underline{3, 2, 1}}$

(ii) Postorder:  $\underline{\underline{2, 1, 3}}$

(iii) Inorder:  $\underline{\underline{2, 3, 1}}$



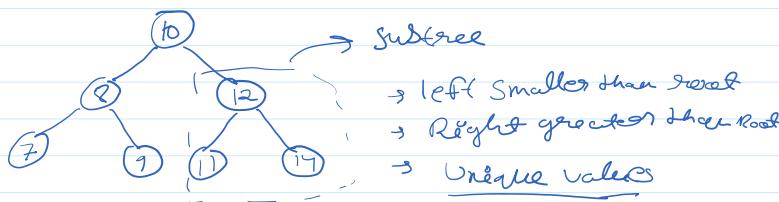
Postorder:  $\underline{\underline{6, 7, 9, 4, 3, 5, 10}}$

Preorder:  $\underline{\underline{10, 9, 6, 7, 5, 4, 3}}$

Inorder:  $\underline{\underline{6, 9, 7, 10, 4, 5, 3}}$

Binary Search Tree:  $\exists$

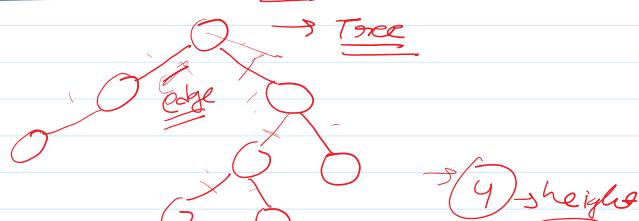
Binary Tree



Inorder:  $\underline{\underline{7, 8, 9, 10, 11, 12, 14}} \rightarrow$  sorted

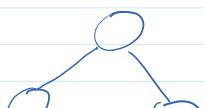
If we traverse in inorder fashion then we get sorted elements.

Height of tree:  $\exists$   $\rightarrow$  Code

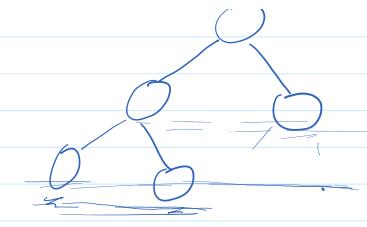


Connection b/w nodes is called edge.

Complete Binary Tree:  $\exists$



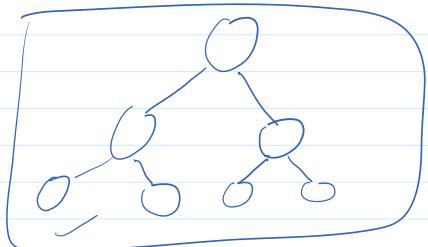
$\rightarrow$  Parent should have 0 or 2 children



↓ Complete Binary

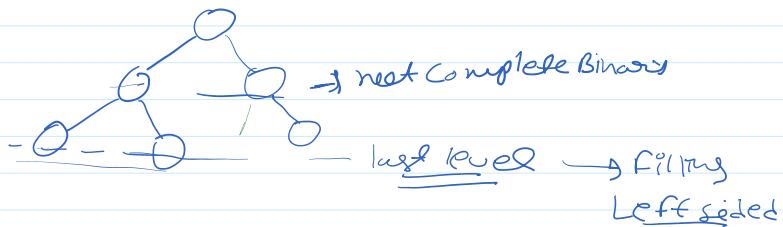
→ Parent should  
have 0 child  
or 2 children

→ Left shifted  
child as last level

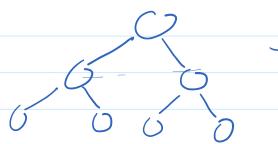


→ all the above  
level (except  
last)  
should be filled completely

→ Complete Binary Tree  
→ Perfect Binary Tree  
→

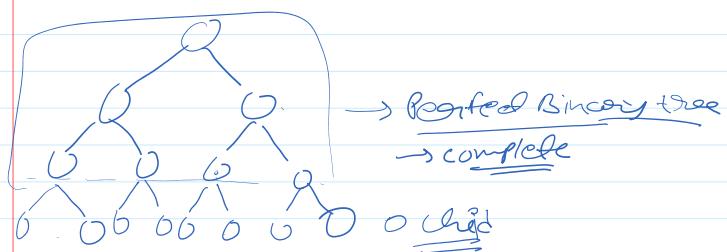


Perfect Binary Tree:



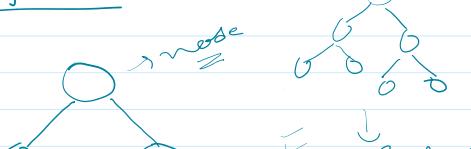
→ Last Level should have  
zero child.

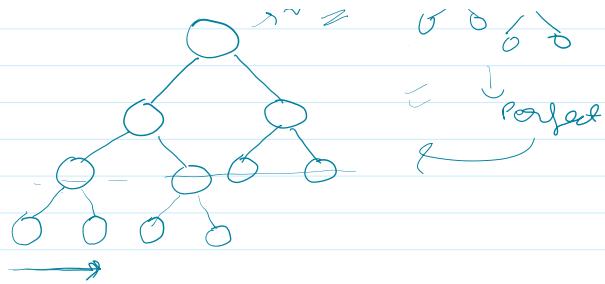
→ above level must  
have 2 children.



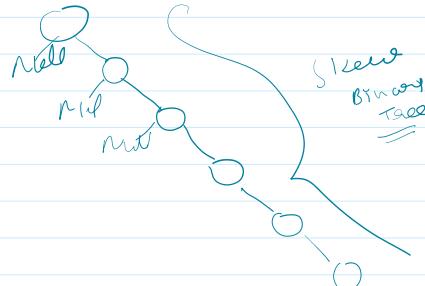
→ Perfect Binary Tree is complete Binary Tree  
but every complete tree may not be Perfect.

Complete Binary tree is





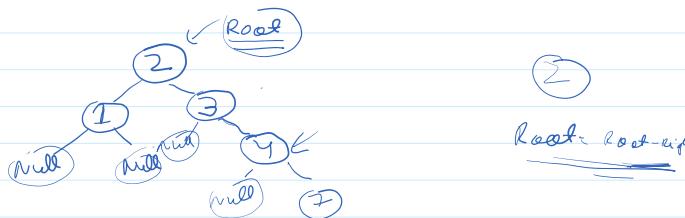
$\approx 5 \ 0 \ 6 \ 7$   
perfect



1. Structure of Binary Tree
2. Insertion in BST
3. Traversals in BST → Recursion
4. Recursion

Dry Run of insertion in BST (loop) :-

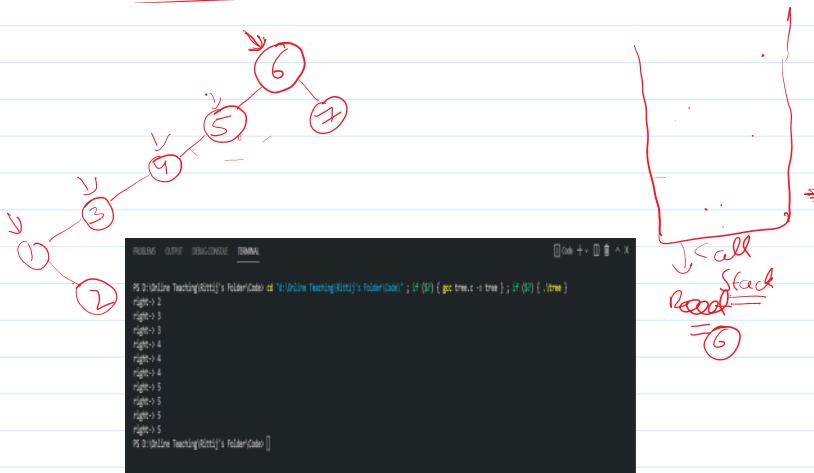
2, 3, 4, 1, 5, 7



2  
Root: Root-left

Recursive :-  
↓  
6, 5, 4, 7, 3, 1, 2

1



```

public class Solution {
    public void solve(TreeNode root) {
        if (root == null) return;
        solve(root.left);
        System.out.print(root.data + " ");
        solve(root.right);
    }
}

class TreeNode {
    int data;
    TreeNode left;
    TreeNode right;

    public TreeNode(int data) {
        this.data = data;
        left = null;
        right = null;
    }
}

```

6, 5, 4, 7, 3, 1, 2

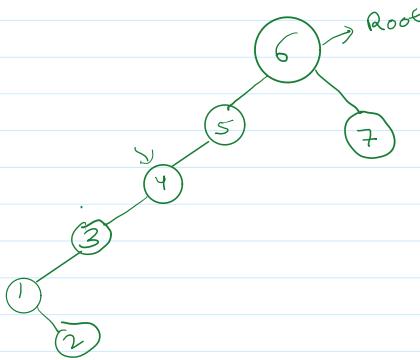
$\approx 6 \ 5 \ 4 \ 7 \ 3 \ 1 \ 2$

```

16 int g = 0;
17
18 struct node* insert(struct node* root, int value) // recursive
19 {
20     if(g == 0)
21     {
22         if (root == NULL)
23         {
24             return createnode(value);
25         }
26     }
27
28     if (value > root->value && root->right == NULL)
29     {
30         struct node* newnode = createnode(value);
31         root->right = newnode;
32         return root;
33     }
34
35     else if (value < root->value && root->left == NULL)
36     {
37         struct node* newnode = createnode(value);
38         root->left = newnode;
39         return root;
40     }
41
42     else if (value < root->value)
43     {
44         insert(root->left, value);
45         printf("right-> %d\n", g);
46     }
47     else if (value > root->value)
48     {
49         insert(root->right, value);
50         printf("right-> %d\n", g);
51     }
52 }
53
54 return root;
55

```

$$g = y$$



$R=1$	$63$	$g=5$
$R=3$	$77$	$g=4$
$R=4$	$77$	$g=3$
$R=5$	$27$	$g=2$
$R=6$	$77$	$g=1$
<code>main 18</code>		

↑  
2  
 $g = 5$

This is for  
insertion of 2