

# CIS 3223 TMQ 4

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Temple ID (last 4 digits:

1 (16 pts) Give a big- $\theta$  bound for the solutions of the following recurrence relations.

(a)  $T(n) = 2T(n/2) + n^{2/3}$

$a = 2, b = 2, d = \frac{2}{3}$

$\log_b a = \log_2 2 = 1$   
 $d = \frac{2}{3}$   
 $\left. \begin{array}{l} \log_b a = \log_2 2 = 1 \\ d = \frac{2}{3} \end{array} \right\} d < \log_b a, T(n) = O(n^{\log_2 a}) = \Theta(n)$

$\Theta(n)$

(b)  $T(n) = 9T(n/3) + 500n^2$

$a = 9, b = 3, d = 2$

$\log_b a = \log_3 9 = 2$   
 $d = 2$   
 $\left. \begin{array}{l} \log_b a = \log_3 9 = 2 \\ d = 2 \end{array} \right\} d = \log_b a, T(n) = \Theta(n^2 \log n)$

$\Theta(n^2 \log n)$

(c)  $T(n) = 7T(n/4) + 8n$

$a = 7, b = 4, d = 1$

$\log_b a = \log_4 7$   
 $d = 1$   
 $\left. \begin{array}{l} \log_b a = \log_4 7 \\ d = 1 \end{array} \right\} d < \log_b a, T(n) = \Theta(n^{\log_4 7})$

$\Theta(n^{\log_4 7})$

(d)  $T(n) = T(n-1) + n^2$

$$\Theta(n^3)$$

$$a = 1 \quad b = 1 \quad d = 2$$

$$T(n) = \Theta(n^{d+1}) = \Theta(n^3)$$

2 (4 pts) Algorithm Z solves problems by dividing them into four subproblems of half the size, recursively solving each subproblem, and combining the solution in linear time. What is the runtime of algorithm Z (in big- $O$  notation)?

$$a = 4 \quad b = 2 \quad d = 1$$

$$\left. \begin{array}{l} \log_b a = \log_2 4 = 2 \\ d = 1 \end{array} \right\} d < \log_b a \quad T(n) = \Theta(n^{\log_2 4}) = \Theta(n^2)$$

$$\Theta(n^2)$$

3 (extra credit, 2 pts) Let  $N = 2221$ . Find a power of 3 in the interval  $[N, 3N]$ .

8

Is there more than one power of 3 in the interval?

yes no

$$2221 = 1201000^1_3 \leq 10000000^8_3 = 3^8 < 3 \times 2221 = 6663$$

8 digits

Note :

$$\begin{array}{cccccccc} 1 & & & & & & & \\ 2221 & \rightarrow & -740 & \rightarrow & 246 & \rightarrow & 82 & \rightarrow & 27 & \rightarrow & 9 & - & 3 & - & 1 \\ & & 1 & & 2 & & 0 & & 1 & & 0 & & 0 & & 1 \end{array}$$