

Spring 2021

Instructor: Dr. Stephen Krebsbach

Exam II

100 points Due: April 27th 6:00 PM (CST)

Record Your answers using the Quiz for EXAM II

Please do the following.

- 1) (6 pts) If we take a system and change its page/frame size to a smaller size then
 - a) The TLB reach will be smaller
 - b) The page table will be smaller
 - c) The Logical address register will need to get bigger
 - d) The Physical address register will need to get smaller

- 2) (6 pts) If the size of the “delta” (Δ) as part of the Working Set model is set smaller
 - a) Its working set will also get smaller
 - b) Its working set will get bigger
 - c) The locality of Reference will get smaller
 - d) It may lead to Thrashing

- 3) (6 pts) Assuming a 1-KB (1024 bytes) page size, what is the page number (p) and offset(d) for the following address reference (values are in decimal) - Address : **3265**
 - a) p = 3 d = 1
 - b) p = 1 d = 2433
 - c) p = 3 d = 193
 - d) p = 2 d = 24

- 4) (6 pts) Consider the logical address space of 512 pages of 1024 words each, mapped into a Physical memory of 64 frames. How many **bits** are there in the **Physical** address?
 - a) 18
 - b) 16
 - c) (64,1024)
 - d) 10

5) (6 pts) Consider the following segment table

Segment	Base	Length
0	200	350
1	2200	150
2	50	180

What is the physical address of the following logical address? Address : (2, 110)

- a) 160
- b) 290
- c) There is an error in the segment table!
- d) Segmentation Fault !

6) (6 pts) Consider the following page reference string: **1,2,1,2,3,1,3,2,3,1,2,3,1**

Assuming **pure** demand paging and **3 frames** allocation, how many times will the Page Replacement Algorithm need to be run?

Use the First-in-First-Out (**FIFO**) algorithm.

- a) 3
- b) 2
- c) 1
- d) 0

7) (6 pts) Consider the following page reference string: **1,2,1,3,4,3,2,3,4,3,3,1**

Assuming **pure** demand paging and **4 frames** allocation, how many page faults will the Least Recently Used (**LRU**) algorithm produce?

- a) 5
- b) 4
- c) 3
- d) 0

8) (6 pts) As the quantum (q) time for the Round Robin (RR) short term Schedule is set larger and larger

- a) It will become just like First-Come-First-Served Algorithm
- b) No process will ever run
- c) It would increase the number of Context switches
- d) The degree of Multi-programming will increase

9) (6 pts) The Scheduling Criteria “Waiting Time” for a process is

- a) the total time it waits in the “Wait” state
- b) the total time the user waits for it to get done
- c) the total time of the context switching it waits for
- d) None of the others

10) (6 pts) Using an External Priority Preemptive short-term scheduling algorithm

- a) can lead to deadlocks
- b) can lead to starvation
- c) will lead to the convoy effect
- d) is provably Optimal for large numbers of processes

11) (11 pts) Given the following information what will be the **AVERAGE WAIT TIME** for **ROUND ROBIN** (RR)? [If needed quantum = 5 and order in the queue is 1 -> 2 -> 3 -> 4]

Job #	Burst Rate	Arrival Time
1	24	0
2	17	0
3	33	0
4	100	0

- a) 63.00
- b) 43.50
- c) 61.75
- d) 210

12) (11 pts) Given the following information what will be the **AVERAGE WAIT TIME** for **SHORTEST REMINING TIME FIRST (SRJF)**? [If needed, quantum = 5 and order in the queue is 1 -> 2 -> 3 -> 4]

Job #	Burst Rate	Arrival Time
1	24	10
2	17	6
3	33	14
4	100	0

- a) 61.75
- b) 38.00
- c) 30.00
- d) 21.75

13) (6 pts) The Second-Chance page replacement algorithm when choosing a victim

- a) Gives a page only 1 second chance and will be the victim the next time it chooses one
- b) Gives a page a second chance when ever its reference bit is set to 1
- c) Only gives pages with their reference bit set to 0 a second chance
- d) Will only chance one page's reference bit

14) (6 pts) The COPY-ON_WRITE technique can be used when a parent forks a child process. In this technique

- a) new child pages are only created to copy the parent's code pages.
- b) the child shares the parent's pages using the same page table
- c) new child pages are only created when the shared page is going to be modified
- d) a copy of all pages is written out to the backing store and written back in when needed

15) (6 pts) The Page-Fault-Frequency (PFF) approach

- a) counts the sum of all page faults in the system
- b) allows for a bigger working set and less page faults
- c) attempts to determine the correct range of page faults to be allowed
- d) Never needs to swap out a process