



The Network Layer

Lecture given by Emmanuel Lochin

ISAE-SUPAERO

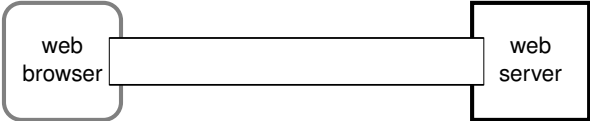
Original slides from A. Carzaniga (Univ. Lugano)
Extended/modified by E. Lochin (ISAE-SUPAERO) with author permission Textbook Chap. #4 Sections 4.1, 4.2.2, 4.3, 4.4.1

Outline

- Basic network-layer architecture of a datagram network
- Introduction to forwarding
- Introduction to routing
- General architecture of a router
- Switching fabric and queuing
- Internet network-layer protocol
- The Internet protocol (IP)
- Fragmentation

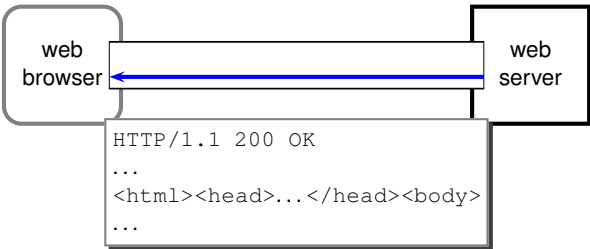
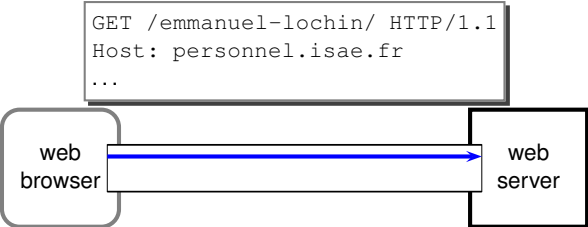
Application Level

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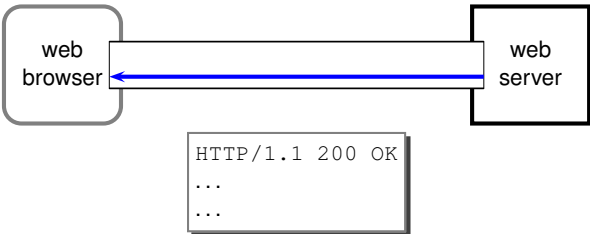
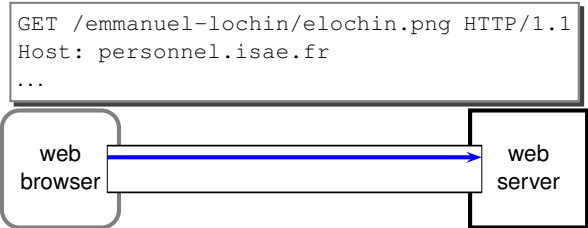
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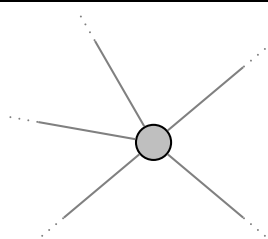
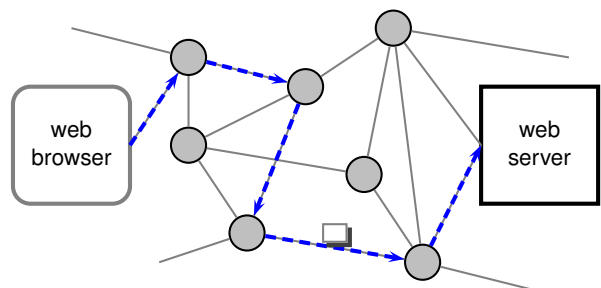
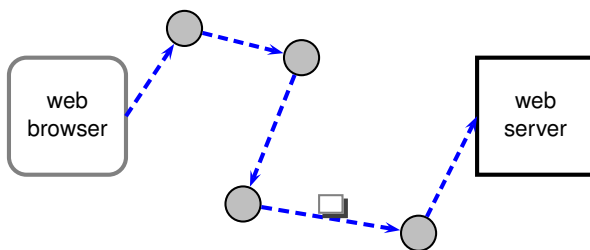
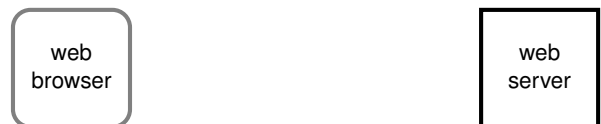
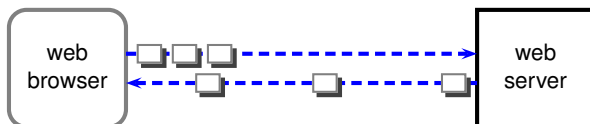
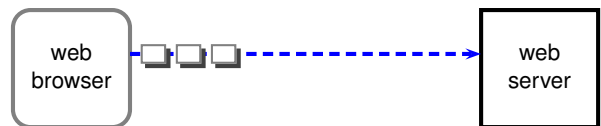
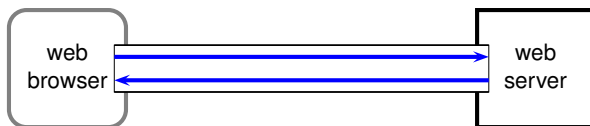
Application Level



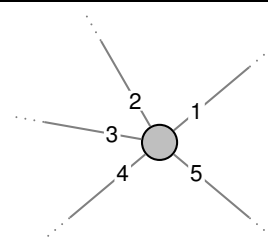
Application Level

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- Fundamental component of the network layer
- A node in a graph
- A finite set of input/output (physical) connections
 - a.k.a., **interfaces** or **ports**



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- **Packet-switched network**

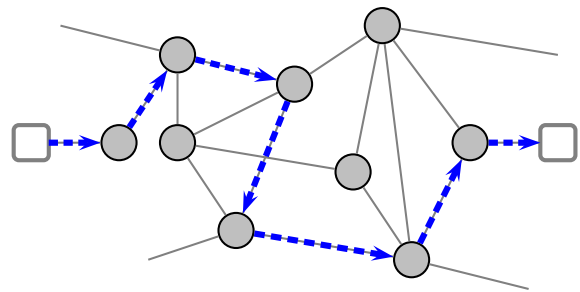
- ▶ information is transmitted in discrete units called **datagrams**

- **Connectionless service**

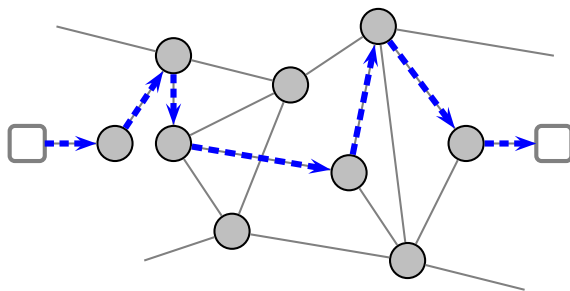
- ▶ a datagram is a **self-contained message**
- ▶ treated independently by the network
- ▶ no connection setup/tear-down phase

- **“Best-effort” service**

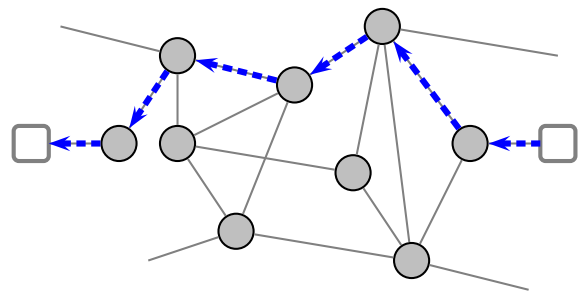
- ▶ delivery guarantee : none
- ▶ maximum latency guarantee : none
- ▶ bandwidth guarantee : none
- ▶ in-order delivery guarantee : none
- ▶ congestion indication : none



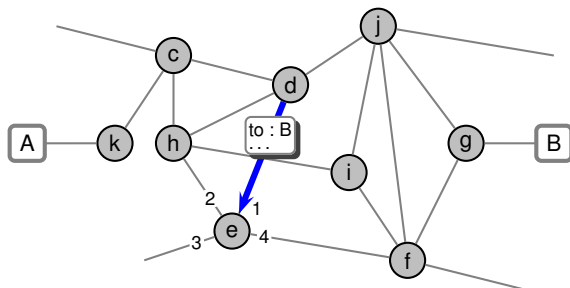
- Potentially **multiple paths** for the same source/destination



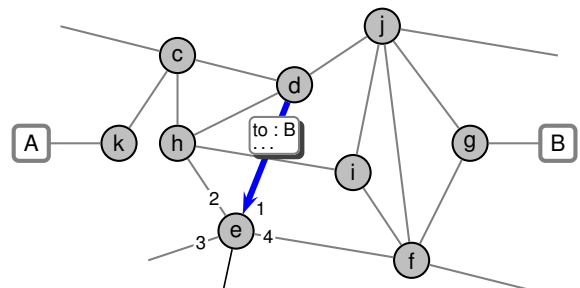
- Potentially **multiple paths** for the same source/destination



- Potentially **multiple paths** for the same source/destination
- Potentially **asymmetric paths**



- A sends a datagram to B
- The datagram is **forwarded** towards B



forwarding table	
dest.	output
...	...
B	port 4
...	...

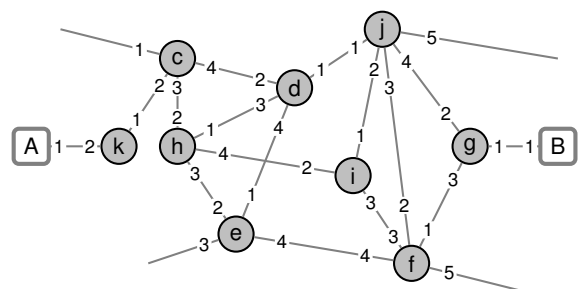
- **Input** : datagram destination

- **Output** : output port

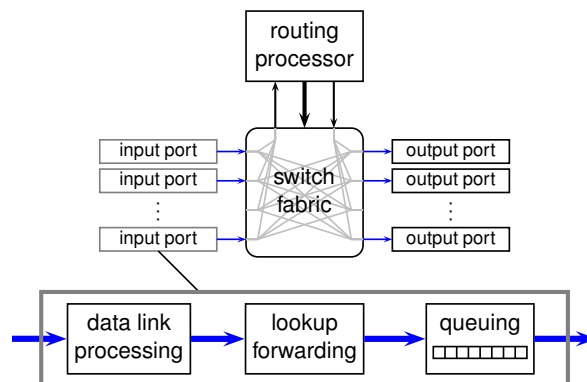
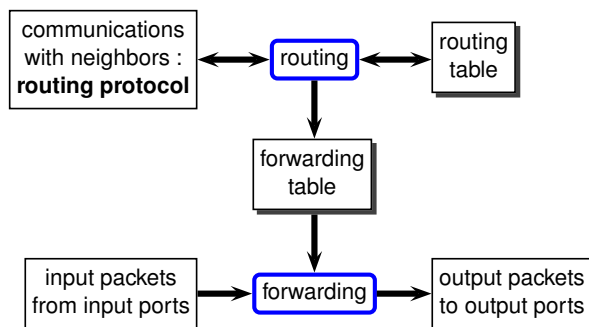
- Simple design : “forwarding table”

- **Issues**

- ▶ how big is the forwarding table ?
- ▶ how fast does the router have to forward datagrams ?
- ▶ how does the router build and maintain the forwarding table ?

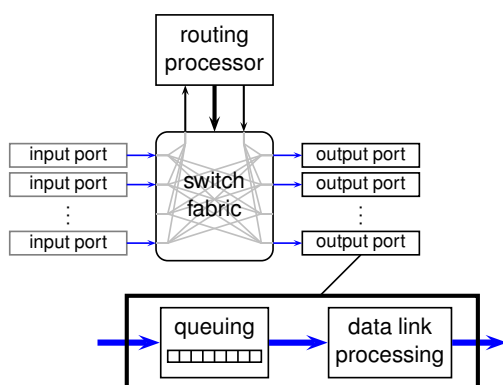


router k	
A	2
B	1



Anatomy of a Router

Queuing



- Where does queuing occur ?
- Input ports
 - queuing may occur here if the switching fabric is slower than the aggregate speed of all the input lines. i.e., $R_S < nR_{in}$
- Output ports
 - queuing may occur here because of the limited throughput of the output link. i.e., $R_{out} < \min(R_S, nR_{in})$

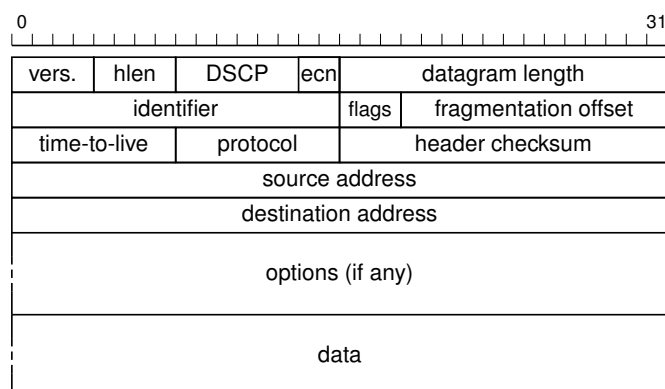
Queuing

Internet Network Layer

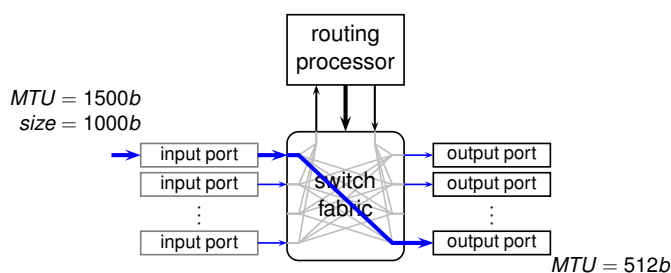
- What happens when packets queue up in a router ?
- **Scheduling** : deciding which packets to process
 - **first-come-first-served**
 - **weighted fair queuing** : the router tries to be balance traffic evenly among the different end-to-end connections. Essential to implement **quality-of-service guarantees**
- Deciding when to drop packets, and which packets to drop
 - **drop tail** : drop arriving packets when queues are full
 - **active queue management** : a set of policies and algorithms to decide when and how to drop or mark packets in the attempt to prevent congestion

- **Routing** : defining paths and compiling forwarding tables
 - RIP
 - OSPF
 - BGP
- IP
 - addressing
 - datagram format
 - fragmentation and packet handling
- ICMP
 - error reporting
 - signaling

IPv4 Datagram Format



Fragmentation



- How does the router handle cases where the size of an input datagram exceeds the maximum transmission unit (MTU) of the output link ?
- The datagram is **fragmented**
 - Today, fragmentation rarely occurs