

A note on homology:

$H_n : \text{Ch}(R) \longrightarrow R\text{-mod}$ is an additive functor

But is not exact!

$A \rightarrow B \rightarrow C \Rightarrow H_n(A) \rightarrow H_n(B) \rightarrow H_n(C)$
exact is a Complex

but not necessarily exact!

Examples:

$$0 \rightarrow A \xrightarrow{f} B$$

$$\begin{array}{ccccccc}
2 & 0 & \xrightarrow{\circ} & 0 & & & \\
1 & \downarrow & & \downarrow & & & \\
1 & 0 & \xrightarrow{\circ} & \mathbb{Z} & \xrightarrow{H_0} & 0 \rightarrow H_0(A) \xrightarrow{H_0(f)} H_0(B) \\
0 & \downarrow & & \parallel & & & \\
0 & \mathbb{Z} & = & \mathbb{Z} & & & \\
-1 & \downarrow & & \downarrow & & & \\
-1 & 0 & \xrightarrow{\circ} & 0 & & & \text{not exact!}
\end{array}$$

$$B \xrightarrow{g} C \rightarrow 0$$

$$\begin{array}{ccccccc}
2 & 0 & \xrightarrow{o} & 0 & & & \\
\downarrow & & & \downarrow & H_1 & & \\
1 & \mathbb{Z} & = & \mathbb{Z} & \rightsquigarrow & H_1(B) & \xrightarrow{H_2(g)} H_1(C) \rightarrow 0 \\
& \parallel & & \downarrow & & & \\
0 & \mathbb{Z} & \xrightarrow{o} & 0 & & & \\
\downarrow & & & \downarrow & & & \\
-1 & 0 & \xrightarrow{o} & 0 & & &
\end{array}$$

not exact

What's really going on!

$$0 \rightarrow A \xrightarrow{f} B \xrightarrow{g} C \rightarrow 0$$

?

induces a LES in homology

the connecting homomorphism is not 0 (!)
(in some degrees)

$$H_{n+1}(C) \xrightarrow{o} H_n(A) \rightarrow H_n(B) \rightarrow H_n(C) \xrightarrow{\quad} \gamma^{(A)}$$

*usually
not 0*

*usually
not 0*

Precviously, on Homological Algebra:

$\rightarrow P$ projective $\Leftrightarrow \text{Hom}_R(P, -)$ is exact

\uparrow
 P free

$$\begin{array}{c} P \\ \downarrow \\ A \xrightarrow{\quad} B \xrightarrow{\quad} 0 \end{array}$$

\Leftrightarrow \exists \downarrow

- Every R -module M is a quotient of a (free \Rightarrow) projective module

$\rightarrow E$ injective $\Leftrightarrow \text{Hom}_R(-, E)$ is exact

$$\begin{array}{c} E \\ \uparrow \\ 0 \rightarrow I \xrightarrow{\quad} R \\ \text{ideal} \end{array}$$

- Every R -module M embeds into some injective module

Projective Resolutions

Slogan: Approximate M by projectives

A projective resolution of M is a complex

$$\dots \rightarrow \overset{2}{P_2} \rightarrow \overset{1}{P_1} \rightarrow \overset{0}{P_0} \rightarrow 0 \rightarrow \dots$$

with $H_i(P_0) = 0$ for $i > 0$ and $H_0(P_0) = M$.

$$\Leftrightarrow \text{an exact complex } \dots \rightarrow \overset{2}{P_2} \rightarrow \overset{1}{P_1} \rightarrow \overset{0}{P_0} \rightarrow M \rightarrow 0$$

A free resolution of M is a projective resolution where all the P_i are free.

We sometimes write $P_0 \rightarrow M$ or $P \rightarrow M$

Remark projective resolution

$$\begin{array}{ccccccc} \dots & \rightarrow & \overset{2}{P_2} & \rightarrow & \overset{1}{P_1} & \rightarrow & \overset{0}{P_0} & \rightarrow 0 \\ & & \circ \downarrow & & \circ \downarrow & & \downarrow & \\ & & 0 & \rightarrow & 0 & \rightarrow & M & \rightarrow 0 \end{array}$$

quasiiso

Every module has a free resolution:

Step 1 Find a surjection from a free module $\mathbb{P}_0 \xrightarrow{\pi_0} M$

Step 2 Look at the kernel of π_0

Find a free module surjecting onto $k_0 := \ker \pi_0$

$$\begin{array}{ccccc} \mathbb{P}_1 & \xrightarrow{i_0 \circ \pi_1} & \mathbb{P}_0 & \xrightarrow{\pi_0} & M \longrightarrow 0 \\ \pi_1 \downarrow & & \nearrow i_0 & & \\ 0 & \longrightarrow k_0 & \downarrow & & 0 \end{array}$$

Note $\ker \pi_1 = \ker (i_0 \circ \pi_1)$

\downarrow
is injective

$$\begin{array}{ccccccc} 0 & \longrightarrow & k_1 & \longrightarrow & 0 \\ & \nearrow & \downarrow & \searrow & & \\ \mathbb{P}_2 & \dashrightarrow & \mathbb{P}_1 & \longrightarrow & \mathbb{P}_0 & \longrightarrow & M \longrightarrow 0 \\ & & \searrow & \nearrow & & & \\ & & k_0 & & & & \\ & & \downarrow & & & & \\ & & 0 & & & & 0 \end{array}$$

Repeat. If $k_n = 0$ at some point, stop

$$\operatorname{pd}_{\mathcal{R}}(M) := \inf \left\{ d \mid \exists \mathbb{P} \rightarrow M \text{ with } \mathbb{P}_i = 0 \text{ for } i > d \right\}$$

Minimal Free resolution

Setup:

(R, \mathfrak{m}) Noetherian local ring
or

N -graded k -algebra, $R_0 = k$, $\mathfrak{m} = \bigoplus_{n \geq 1} R_n$

(so $R = \frac{k[x_1, \dots, x_d]}{\mathfrak{I}}$, \mathfrak{I} homogeneous, $\mathfrak{m} = (x_1, \dots, x_d)$)
 M fg (graded) R -module

Note: We can find a free resolution of M where all the \mathfrak{I}_i are fg

Recall $\mu(M) := \text{minimal } \# \text{ of generators of } M = \dim_k (M/\mathfrak{m}M)$

Can find a surjection $R^{\mu(M)} \rightarrow M = Rf_1 + \dots + Rf_n$
 $(r_1, \dots, r_n) \mapsto r_1 f_1 + \dots + r_n f_n$

In the graded case, we can take all the maps to be graded

$R(-f_1) \oplus \dots \oplus R(-f_n) \rightarrow M = Rf_1 + \dots + Rf_n$ $\deg(f_i) = d_i$
is a degree graded R -module map

$(R(-s))_t = R_{t-s}$ so R_0 lives in degree s

A minimal free resolution of M is one where each $\mathfrak{I}_i \cong R^{n_i}$
has n_i the smallest possible. In the graded case, we also
ask for the maps in the resolution to be degree preserving. So
 $\mu(\mathfrak{I}_0) = \mu(M)$, $\mu(\mathfrak{I}_i) = \mu(k_0)$, $\mu(\mathfrak{I}_{i+1}) = \mu(k_i)$

Will show: Minimal free resolutions are unique!

Betti numbers $\beta_i(M)$ = rank of F_i in a minimal free resolution

Graded betti numbers: $\beta_{ij}(M) := \# \text{copies of } R(-j) \text{ in homological degree } i$

Betti table has $\beta_{ij}(M)$ in position $(i, i+j)$

Example $R = k[x, y, z]$ $M = R/(xy, xz, yz)$

$$0 \rightarrow R^2 \xrightarrow{\begin{pmatrix} z & 0 \\ -y & y \\ 0 & x \end{pmatrix}} R^3 \xrightarrow{(xy \ xz \ yz)} R \rightarrow M$$

$$\beta_1(M) = 3 \quad \beta_2(M) = 2 \quad \beta_0(M) = 1$$

Graded resolution: $0 \rightarrow R(-3)^2 \xrightarrow{\begin{pmatrix} z & 0 \\ -y & y \\ 0 & x \end{pmatrix}} R(-2)^3 \xrightarrow{\begin{matrix} (xy \ xz \ yz) \\ \downarrow \\ \text{degree 2} \end{matrix}} R \rightarrow M$

$$\beta(M) = \begin{array}{c|ccc} & 0 & 1 & 2 \\ \hline 0 & 1 & & \\ 1 & & 3 & 2 \\ 2 & & \brace{3} & 2 \end{array} \xrightarrow{\beta_{23}} \beta_{12}$$

$$\beta_{12}(M) = 3 \quad \beta_{23}(M) = 2$$

Example $R = k[x, y]$ $M = R/(x^2, xy, y^3)$

$$0 \rightarrow R(-3) \oplus R(-4) \xrightarrow{\begin{pmatrix} y & 0 \\ -x & y^2 \\ 0 & x \end{pmatrix}} R(-2)^2 \xrightarrow{\begin{matrix} (x^2 \ xy \ y^3) \\ \oplus \\ R(-3) \end{matrix}} R \rightarrow M$$

Note: $\begin{pmatrix} 0 \\ y^2 \\ x \end{pmatrix}$ lands in $\begin{array}{c} \text{deg 2} \\ \text{deg 2} \\ \text{deg 3} \end{array}$ so $\begin{array}{c} \text{deg 2+2=4} \\ \text{deg 1+3=4} \end{array} \checkmark$

Equivalently, F_{\cdot} is a minimal free resolution of M if it is a minimal complex, meaning

$$\dots \rightarrow F_2 \xrightarrow{\partial_2} F_1 \xrightarrow{\partial_1} F_0 \rightarrow 0$$

$$\text{im } \partial_i \subseteq \mathfrak{m} F_{i-1}.$$

\Leftrightarrow all the entries in the matrices are in \mathfrak{m} .

Proof: Suppose $\text{im } \partial_n \subseteq \mathfrak{m} F_{n-1}$ for all n , but F_n is the free module on non-minimal generators f_1, \dots, f_s generators for $k_{n-1} := \ker \partial_{n-1}$

$$\begin{array}{ccc} F_n & \longrightarrow & M \\ e_i & \longmapsto & f_i \end{array}$$

f_1, \dots, f_s linearly dependent in $M/\mathfrak{m}M$

$$\Rightarrow r_1 f_1 + \dots + r_s f_s = 0 \text{ for some } r_i \in R, \text{ not all in } \mathfrak{m}$$

Assume wlog that r_1 is invertible, and multiply by its inverse

$$\Rightarrow e_s - r_1 e_1 - \dots - r_{s-1} e_{s-1} \in \ker \partial_n = \text{im } \partial_{n-1} \subseteq \mathfrak{m} F_n$$

$\Rightarrow e_1, \dots, e_s$ are linearly dependent! \hookleftarrow

Suppose $\mu(F_n) = \mu(k_n)$ for all n , where $k_{n-1} = \ker \partial_{n-1}$ but $\text{im } \partial_{n+1} \not\subseteq \mathfrak{m} F_n$ for some n .

$\tilde{F}_n = \bigoplus_{i=1}^s \mathbb{R} e_i, \{ \partial_n(e_i) \}$ minimal generating set for k_{n-1}

$$\begin{array}{l} \exists r_1 e_1 + \dots + r_s e_s \in \ker \partial_n \\ \not\in \mathfrak{m} \tilde{F}_n, \quad \xrightarrow[\substack{r_i \notin \mathfrak{m}}]{\text{wlog}} \quad e_1 - c_2 e_2 - \dots - c_s e_s \in \ker \partial_n \end{array}$$

$$\Rightarrow \partial_n(e_1) = c_2 \partial_n(e_2) + \dots + c_s \partial_n(e_s) \quad \hookleftarrow$$