

0101  
0111  
0001 ✓

$$\frac{4}{2} = 2 > 1,7 \rightarrow \log_2 3 = 1,5$$

Split

(logo, aumenta  
a qtd de bits  
significativos)

00 0000  
01 0101  
10 1010

$R=3$   
 $N=4$   
 $i = \log_2 3 = 1,5$

0111  $\rightarrow N = 5/3 = 1,666$

00 0000  
01 0101  
10 1010

$$i = \log_2 3 = 1,5$$

0001  $\rightarrow 6/3 = 2 > 1,7$

Split

$\log_2 4 = 2$

00 0000  
01 0001  
10 1010  
11 1111  
0111

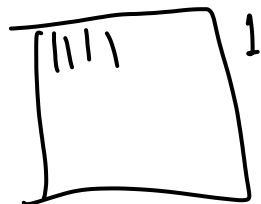
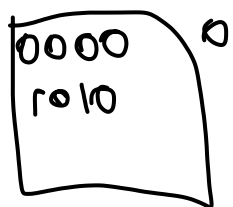
Fim

0000  
1010  
1111

$N = buckets = 2$   
 $R = chaves = 3$   
 $i = bits = 1$

$$\left. \begin{array}{l} 2^1 = 2 < 3 \\ 2^2 = 4 > 3 \end{array} \right\} 1,5 \rightarrow 2^{1,5} =$$

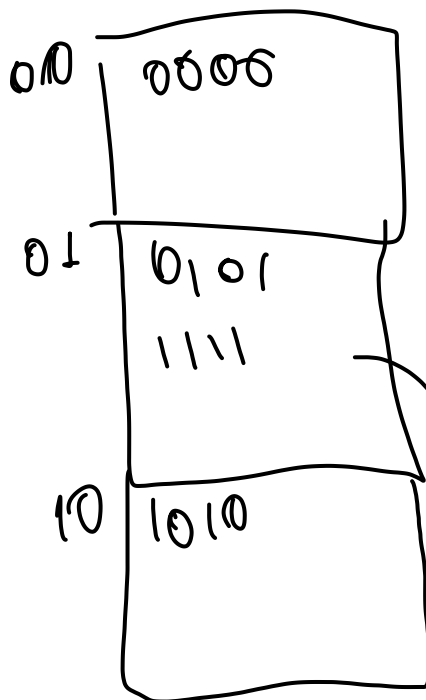




0101 →  $R = 2$   
 $N = 4$   
 $i = 1$

$R = 2$   
 $N = 3$   
 $i = 1$

$R = \text{Btd bucket}$   
 $N = \text{Char}_i$   
 $i = \text{bits}$



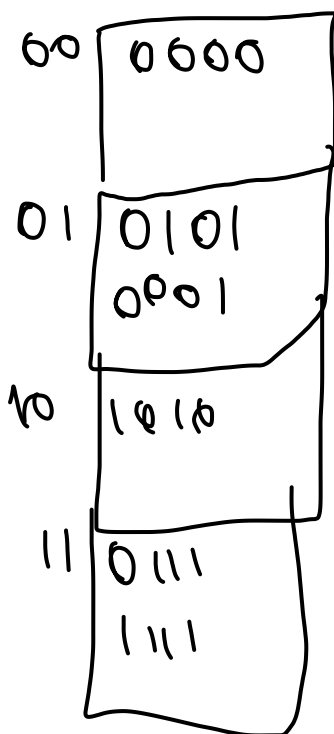
→ 2.6  
 → 0111

$R = 3$   
 $N = 5$   
 $i = 2$



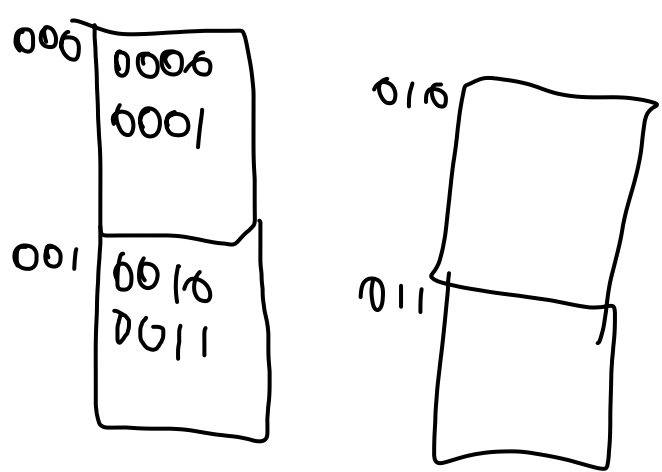
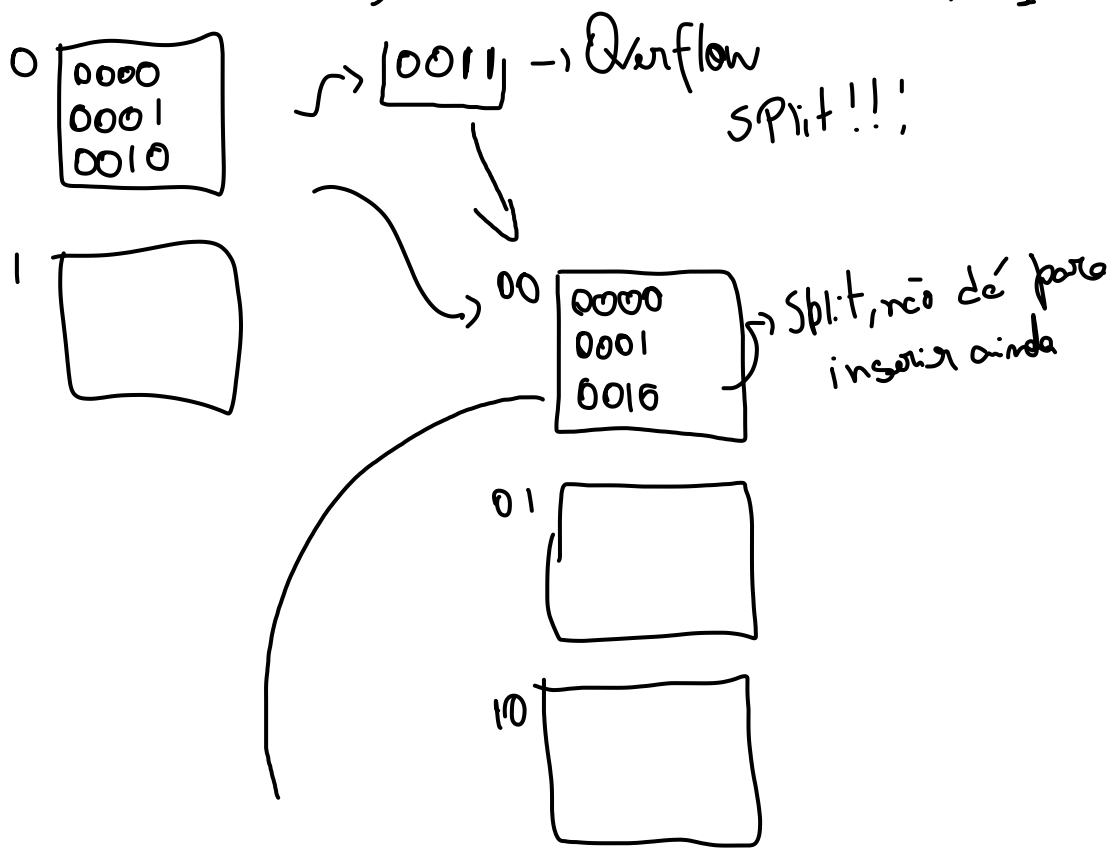
0001 → 2.7

$R = 6 > 2$   
 $N = 3$   
 $i = \textcircled{2}$



$\swarrow$  a) Linear  $\swarrow$  Direita  
 Extensível  $\rightarrow$  esquerda

buckets  $\begin{matrix} \text{chaves} \\ R = 4 \\ N = 2 \\ i = 1 \end{matrix}$



b) 0

|      |
|------|
| 0000 |
|------|

1

|      |
|------|
| 0001 |
|------|

split

0010  $\frac{3}{2} = 1,5$

00

|      |
|------|
| 0000 |
| 0100 |

01

|      |
|------|
| 0001 |
|------|

10

|      |
|------|
| 0010 |
|------|

11

|      |
|------|
| 0011 |
|------|

0101

$\frac{6}{4} = 1,5$   
split

R = Move  
N = Bucket  
i = 1  
taxa = 1,5

$\frac{4}{3} = 1,3$  → 0011

$\frac{5}{4} = 1,25$

→ 0100

|     |      |
|-----|------|
| 000 | 0000 |
| 001 | 0001 |
| 010 | 0010 |

→ 101 →  $\frac{6}{8} = 0,75$   
↑

|     |      |
|-----|------|
| 011 | 0011 |
| 100 | 0100 |
| 101 | 0101 |
| 110 |      |
| 111 |      |