VideoCore IV
Typed Assembly Language

Version 0.1

El Pin Al

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1. Syntax

To define the syntax, there are some primitive terms:

- i denotes an integer.
- $i4$ denotes a 4-bit integer in range $[0, 15]$.
- $i4^*$ denotes a 4-bit integer in range $[1, 16]$.
- $i7^*$ denotes a 7-bit integer in range $[1, 128]$.
- f denotes a floating-point number.

- l denotes a label.
- nat denotes an integer which is greater than or equal to 0.
- In general, ε denotes an empty construct.
- α denotes a variable.

The syntax is given below.

$\nu ::= n \mid rr$	operands
$n ::= i \mid f$	numbers
$r ::= rr \mid wr$	registers
$rr ::= rwr \mid \mathrm{uniform} \mid \mathrm{element\_number} \mid \mathrm{vpmr}$	readable registers
$urr ::= a \mid \mathrm{element\_number}$	unconstrained readable registers
$wr ::= rwr \mid \mathrm{uniforms\_address} \mid tmu \mid \mathrm{broadcast} \mid \mathrm{vpmw}$	writable registers
$rwr ::= ga \mid rf$	both readable and writable registers
$a ::= ga \mid sa$	accumulators
$ga ::= \mathrm{r0} \mid \mathrm{r1} \mid \mathrm{r2} \mid \mathrm{r3}$	general-purpose accumulators
$sa ::= \mathrm{r4} \mid \mathrm{r5}$	special-purpose accumulators
$rf ::= A \mid B$	register files
$A ::= \mathrm{ra0} \mid \ldots \mid \mathrm{ra31}$	locations in register file A
$B ::= \mathrm{rb0} \mid \ldots \mid \mathrm{rb31}$	locations in register file B
$tmu ::= \mathrm{tmu0} \mid \mathrm{tmu1}$	TMU
$vpmq ::= \varepsilon \mid (Y, i) :: vpmq$	VPM read queues
$Y ::= \mathrm{y}_0 \mid \ldots \mid \mathrm{y}_{63}$	VPM Y

$$\iota ::= \text{rotate}(wr, rr, i4) \mid \text{mov}(wr, \nu) \mid \text{setup\_vpm\_read}\left(Y, i4^{*}\right)$$
$$\mid \text{setup\_vpm\_write}(Y) \mid \text{setup\_dma\_load}\left(Y, i4^{*}\right) \mid \text{start\_dma\_load}(rr)$$
$$\mid \text{wait\_dma\_load} \mid \text{setup\_dma\_store}\left(Y, i7^{*}\right) \mid \text{start\_dma\_store}(rr)$$
$$\mid \text{wait\_dma\_store} \qquad \text{instructions}$$

$$cc ::= \mathrm{Z} \mid \mathrm{N} \mid \mathrm{C} \qquad \text{condition classes}$$

$$c ::= \text{set}(cc) \mid \text{clear}(cc) \qquad \text{conditions}$$

$$qc ::= \text{all}(c) \mid \text{any}(c) \qquad \text{quantified conditions}$$

$$s ::= \varepsilon \mid \text{load}\langle tmu \rangle \qquad \text{signals}$$

$$ci ::= (\iota, \varepsilon) \mid (\iota, c) \qquad \text{conditional instructions}$$

$$csi ::= (ci, s) \qquad \text{conditional instructions with signals}$$

$$I ::= \text{jmp}(l)\ ;\ csi\ ;\ csi\ ;\ csi \mid (ci, \text{thread\_end})\ ;\ csi\ ;\ csi \mid csi\ ;\ I$$
$$\mid \text{if } qc \text{ jmp}(l)\ ;\ csi\ ;\ csi\ ;\ csi\ ;\ I \qquad \text{instruction sequences}$$

$$R ::= \varepsilon \mid R[rr \mapsto arr] \qquad \text{register contexts}$$

$$arr ::= [n, n, n, n, n, n, n, n, n, n, n, n, n, n, n, n] \qquad \text{arrays}$$

$$U ::= \varepsilon \mid n \circ U \qquad \text{uniforms}$$

$$T ::= \varepsilon \mid arr \mid arr :: arr \mid arr :: arr :: arr \mid arr :: arr :: arr :: arr \qquad \text{TMU}$$

$$V ::= \varepsilon \mid V[Y \mapsto at] \qquad \text{VPM states}$$

$$P ::= \varepsilon \mid P[l \mapsto I] \qquad \text{programs}$$

$$M ::= (R, U, T, V, vpmq, P, I) \qquad \text{machines}$$

2. Type system

The type syntax is defined as follows:

$$\tau ::= b \mid at \mid vt \mid \Psi \qquad \text{types}$$

$p ::= nat \mid \alpha \mid p + p \mid nat \times p$	pointers
$b ::= \mathrm{int}(i) \mid \mathrm{int}(?) \mid \mathrm{float} \mid \mathrm{ptr}(p) \mid \mathrm{code}(\Theta)$	basic types
$at ::= [b, b, b, b, b, b, b, b, b, b, b, b, b, b, b, b]$	array types
$vt ::= \mathrm{vec}(b, \alpha)$	vector types
$\Gamma ::= \varepsilon \mid \Gamma[rr \mapsto at]$	register context types
$\Psi ::= \varepsilon \mid b \circ \Psi$	uniforms types
$\Sigma ::= \varepsilon \mid at \mid at :: at \mid at :: at :: at \mid at :: at :: at :: at$	TMU types
$\Pi ::= \varepsilon \mid \Pi[Y \mapsto at]$	VPM state types
$vpmw\_addr ::= \varepsilon \mid Y$	VPM write addresses
$C ::= (vpmq, \Pi, vpmw\_addr)$	VPM compound types
$u ::= \mathrm{u0} \mid \mathrm{u1} \mid \mathrm{u2}$	uniform-access countdowns
$\Theta ::= (\Gamma, \Psi, u, \Sigma, C, \Omega, wr)$	state types
$\Phi ::= \varepsilon \mid \Phi[l \mapsto \Theta]$	program types
$DLS ::= \varepsilon \mid \left(Y, i4^*\right)$	DMA load setups
$DSS ::= \varepsilon \mid \left(Y, i7^*\right)$	DMA store setups
$DL ::= \varepsilon \mid \left(Y, i4^*, \tau\right)$	DMA loads
$DS ::= \varepsilon \mid \left(Y, i7^*, \tau\right) \mid \mathrm{type\_preserving}$	DMA stores

2.1. Propositions

$p_1 .. p_2$ represents a range $[p_1, p_2)$.

$\varphi ::= p .. p \mid \varphi \vee \varphi$	propositions

2.2. Memory representaion

$\Omega ::= \{\Xi \mid \varphi\}$ memory subset types

$\Xi ::= \varepsilon \mid \Xi[p \mapsto \tau]$ memory types

The evaluation rules of memory subset types and memory types are given below. The equality rules for p are not defined here. There are the abuses of notations of a form $x[y \mapsto z]$.

$$\frac{p \text{ is in } \varphi}{\{\Xi \mid \varphi\}(p) \rightarrow_\Omega \Xi(p)}$$

$$\frac{p = p_1}{\Xi\left[p_1 \mapsto \tau\right](p) \rightarrow_\Xi \tau}$$

$$\frac{p \neq p_1}{\Xi\left[p_1 \mapsto \tau\right](p) \rightarrow_\Xi \Xi(p)}$$

2.3. Typing rules

Typing rules are defined as follows. Note that $dom(\Gamma)$ represents the domain of a context Γ, a map from read / write registers to array types.

numbers:

$\vdash i : \mathrm{int}(i) \qquad \vdash f : \mathrm{float}$

subtype relations:

$\mathrm{int}(i) <: \mathrm{int}(?) \qquad \mathrm{int}(?) <: \mathrm{int}(?) \qquad \dfrac{\forall i \;.\; at_1[i] <: at_2[i]}{at_1 <: at_2}\text{(S-Array)}$

$$\frac{m \leq n}{\mathrm{vec}(b,m) <: \mathrm{vec}(b,n)}\text{(S-Vector)}$$

$\Gamma <: \Gamma$

$\dfrac{\Gamma_1 <: \Gamma_2 \qquad rr \notin dom(\Gamma_1)}{\Gamma_1[rr \mapsto at] <: \Gamma_2}\text{(S-Ctx-Width)} \qquad \dfrac{\Gamma_1 <: \Gamma_2 \qquad at_1 <: at_2}{\Gamma_1[rr \mapsto at_1] <: \Gamma_2[rr \mapsto at_2]}\text{(S-Ctx-Depth)}$

pointers:

$$\frac{p \in dom(\Xi)}{\Xi \vdash p : \Xi(p)}$$

well-formed memory subsets:

$$\frac{\forall p \in dom(\Xi) \mathbin{.} p \mathbin{..} \big(p + size\_of\big(\Xi(p)\big)\big) \text{ is in } \varphi}{\vdash \{\Xi \mid \varphi\}}$$

registers:

$$\frac{rr \in dom(\Gamma)}{\Gamma \vdash rr : \Gamma(rr)} \qquad \frac{\Psi_1 \equiv b \circ \Psi_2}{\vdash \text{uniform} : b \mid \Psi_1 \to \Psi_2}$$

$$\vdash \text{element\_number} : [\text{int}(0), \text{int}(1), ..., \text{int}(15)]$$

$$\frac{vpmq_1 \equiv \big(vpmq_2 :: (Y, i)\big) \qquad i \geq 2 \qquad \Pi \vdash Y : at}{\Pi \vdash \text{vpmr} : at \mid vpmq_1 \to vpmq_2 :: \big(inc(Y), i - 1\big)}$$

$$\frac{vpmq_1 \equiv \big(vpmq_2 :: (Y, 1)\big) \qquad \Pi \vdash Y : at}{\Pi \vdash \text{vpmr} : at \mid vpmq_1 \to vpmq_2}$$

VPM:

$$\frac{Y \in dom(\Pi)}{\Pi \vdash Y : \Pi(Y)}$$

$$\frac{\Pi \vdash Y : at \qquad at <: vt}{\Pi \vdash (Y, 1) : vt}$$

$$\frac{\Pi \vdash Y : at \qquad at <: vt_1 \qquad n \leq 62 \qquad i7^* \geq 2 \qquad \Pi \vdash \big(\mathrm{y}_{n+1}, i7^* - 1\big) : vt_2}{\Pi \vdash \big(\mathrm{y}_n, i7^*\big) : concat\big(vt_1, vt_2\big)}$$

instructions:

$$\frac{rr \not\equiv wr_{before} \qquad rr \in dom(\Gamma) \qquad at \equiv rotate\big(\Gamma(rr), i4\big)}{wr_{before} \vdash \text{rotate}\big(rwr, rr, i4\big) : \Gamma \to \Gamma[rwr \mapsto at] \ ; \ rwr}$$

$$\frac{\vdash \text{uniform} : b \mid \Psi_1 \to \Psi_2 \qquad at \equiv \mathit{array}(b)}{\text{u0} \vdash \text{rotate}(\mathit{rwr}, \text{uniform}, i4) : \Gamma \to \Gamma[\mathit{rwr} \mapsto at] \mid \Psi_1 \to \Psi_2 \; ; \mathit{rwr}}$$

$$\frac{\vdash \text{element\_number} : at}{\vdash \text{rotate}(\mathit{rwr}, \text{element\_number}, i4) : \Gamma \to \Gamma\big[\mathit{rwr} \mapsto \mathit{rotate}(at, i4)\big] \; ; \mathit{rwr}}$$

$$\frac{\Pi \vdash \text{vpmr} : at \mid \mathit{vpmq}_1 \to \mathit{vpmq}_2}{\Pi \vdash \text{rotate}(\mathit{rwr}, \text{vpmr}, i4) : \mathit{vpmq}_1 \to \mathit{vpmq}_2 \mid \Gamma \to \Gamma\big[\mathit{rwr} \mapsto \mathit{rotate}(at, i4)\big] \; ; \mathit{rwr}}$$

$$\frac{\mathit{rr} \not\equiv \mathit{wr}_{\mathit{before}} \qquad \Gamma \vdash \mathit{rr} : at \qquad \mathit{fst}\big(\mathit{rotate}(at, i4)\big) \equiv b}{\Gamma; \mathit{wr}_{\mathit{before}} \vdash \text{rotate}(\text{broadcast}, \mathit{rr}, i4) : \Gamma \to \Gamma\big[\text{r5} \mapsto \mathit{array}(b)\big]}$$

$$\frac{\vdash \text{uniform} : b \mid \Psi_1 \to \Psi_2}{\vdash \text{rotate}(\text{broadcast}, \text{uniform}, i4) : \Gamma \to \Gamma\big[\text{r5} \mapsto \mathit{array}(b)\big] \mid \Psi_1 \to \Psi_2}$$

$$\frac{\vdash \text{element\_number} : at \qquad b \equiv \mathit{fst}\big(\mathit{rotate}(at, i4)\big)}{\vdash \text{rotate}(\text{broadcast}, \text{element\_number}, i4) : \Gamma \to \Gamma\big[\text{r5} \mapsto \mathit{array}(b)\big]}$$

$$\frac{\Pi \vdash \text{vpmr} : at \mid \mathit{vpmq}_1 \to \mathit{vpmq}_2 \qquad \mathit{fst}\big(\mathit{rotate}(at, i4)\big) \equiv b}{\Pi \vdash \text{rotate}(\text{broadcast}, \text{vpmr}, i4) : \Gamma \to \Gamma\big[\text{r5} \mapsto \mathit{array}(b)\big] \mid \mathit{vpmq}_1 \to \mathit{vpmq}_2}$$

$$\frac{\mathit{rr} \not\equiv \mathit{wr}_{\mathit{before}} \qquad \Gamma \vdash \mathit{rr} : at \qquad \mathit{fst}\big(\mathit{rotate}(at, i4)\big) \equiv \text{ptr}(p) \qquad \Xi \vdash p : \Psi_2}{\Gamma; \Xi \; ; \mathit{wr}_{\mathit{before}} \vdash \text{rotate}(\text{uniforms\_address}, \mathit{rr}, i4) : \Psi_1 \to \Psi_2 \;\mid u \to \text{u2}}$$

$$\frac{\vdash \text{uniform} : \text{ptr}(p) \mid \Psi_1 \to \Psi_2 \qquad \Xi \vdash p : \Psi_3}{\Xi \vdash \text{rotate}(\text{uniforms\_address}, \text{uniform}, i4) : \Psi_1 \to \Psi_3 \;\mid \text{u0} \to \text{u2}}$$

$$\frac{\Pi \vdash \text{vpmr} : at \mid \mathit{vpmq}_1 \to \mathit{vpmq}_2 \qquad \mathit{fst}\big(\mathit{rotate}(at, i4)\big) \equiv \text{ptr}(p) \qquad \Xi \vdash p : \Psi_2}{\Pi; \Xi \vdash \text{rotate}(\text{uniforms\_address}, \text{vpmr}, i4) : \Psi_1 \to \Psi_2 \mid \mathit{vpmq}_1 \to \mathit{vpmq}_2 \mid u \to \text{u2}}$$

$$\frac{\mathit{rr} \not\equiv \mathit{wr}_{\mathit{before}} \qquad \Gamma \vdash \mathit{rr} : at \qquad \mathit{notfull}\big(\Sigma^{tmu}\big)}{\Gamma; \; \Xi \; ; \mathit{wr}_{\mathit{before}} \vdash \text{rotate}(\mathit{tmu}, \mathit{rr}, i4) : \Sigma^{tmu} \to \mathit{rotate}\Big(\mathit{map}\Big(\mathit{unwrap}_{\Xi}, at\Big), i4\Big) :: \Sigma^{tmu}}$$

$$\frac{\vdash \text{uniform} : \text{ptr}(p) \;\mid \Psi_1 \to \Psi_2 \qquad \Xi \vdash p : b \qquad \mathit{notfull}\big(\Sigma^{tmu}\big)}{\text{u0}; \Xi \vdash \text{rotate}(\mathit{tmu}, \text{uniform}, i4) : \Sigma^{tmu} \to \mathit{array}(b) :: \Sigma^{tmu} \mid \Psi_1 \to \Psi_2}$$

$$\frac{\Pi \vdash \text{vpmr} : at \mid vpmq_1 \to vpmq_2 \qquad notfull\left(\Sigma^{tmu}\right)}{\Pi; \Xi \vdash \text{rotate}(tmu, \text{vpmr}, i4) : \Sigma^{tmu} \to rotate\left(map\left(unwrap_{\Xi}, at\right), i4\right) :: \Sigma^{tmu}}$$

$$\frac{rr \not\equiv wr_{before} \qquad rr \in dom(\Gamma) \qquad at \equiv rotate(\Gamma(rr), i4)}{wr_{before} \vdash \text{rotate}(\text{vpmw}, rr, i4) : \Pi \to \Pi[Y \mapsto at] \mid Y \to inc(Y)}$$

$$\frac{\vdash \text{uniform} : b \mid \Psi_1 \to \Psi_2 \qquad at \equiv array(b)}{\text{u0} \vdash \text{rotate}(\text{vpmw}, \text{uniform}, i4) : \Pi \to \Pi[Y \mapsto at] \mid \Psi_1 \to \Psi_2 \mid Y \to inc(Y)}$$

$$\frac{\vdash \text{element\_number} : at}{\vdash \text{rotate}(\text{vpmw}, \text{element\_number}, i4) : \Pi \to \Pi\left[Y \mapsto rotate(at, i4)\right] \mid Y \to inc(Y)}$$

$$\frac{\vdash n : b}{\vdash \text{mov}(rwr, n) : \Gamma \to \Gamma\left[rwr \mapsto array(b)\right] ; rwr}$$

$$\frac{\Gamma \vdash urr : at}{\vdash \text{mov}(rwr, urr) : \Gamma \to \Gamma[rwr \mapsto at] ; rwr}$$

$$\frac{\Gamma \vdash rf : at \qquad rf \not\equiv wr_{before}}{wr_{before} \vdash \text{mov}(rwr, rf) : \Gamma \to \Gamma[rwr \mapsto at] ; rwr}$$

$$\frac{\vdash \text{uniform} : b \mid \Psi_1 \to \Psi_2}{\text{u0} \vdash \text{mov}(rwr, \text{uniform}) : \Gamma \to \Gamma\left[rwr \mapsto array(b)\right] \mid \Psi_1 \to \Psi_2 ; rwr}$$

$$\frac{\Pi \vdash \text{vpmr} : at \mid vpmq_1 \to vpmq_2}{\Pi \vdash \text{mov}(rwr, \text{vpmr}) : \Gamma \to \Gamma[rwr \mapsto at] \mid vpmq_1 \to vpmq_2 ; rwr}$$

$$\frac{\vdash n : b}{\vdash \text{mov}(\text{broadcast}, n) : \Gamma \to \Gamma\left[\text{r5} \mapsto array(b)\right]}$$

$$\frac{\Gamma \vdash urr : at \qquad fst(at) \equiv b}{\Gamma \vdash \text{mov}(\text{broadcast}, urr) : \Gamma \to \Gamma\left[\text{r5} \mapsto array(b)\right]}$$

$$\frac{\Gamma \vdash rf : at \qquad fst(at) \equiv b \qquad rf \not\equiv wr_{before}}{\Gamma; wr_{before} \vdash \text{mov}(\text{broadcast}, rf) : \Gamma \to \Gamma\left[\text{r5} \mapsto array(b)\right]}$$

$$\frac{\vdash \text{uniform} : b \mid \Psi_1 \to \Psi_2}{\text{u0} \vdash \text{mov}(\text{broadcast}, \text{uniform}) : \Gamma \to \Gamma\left[\text{r5} \mapsto array(b)\right] \mid \Psi_1 \to \Psi_2}$$

$$\frac{\Pi \vdash \mathrm{vpmr} : at \mid vpmq_1 \to vpmq_2}{\Pi \vdash \mathrm{mov}(\mathrm{broadcast}, \mathrm{vpmr}) : \Gamma \to \Gamma[\mathrm{r5} \mapsto at] \mid vpmq_1 \to vpmq_2}$$

$$\frac{\Gamma \vdash a : at \qquad fst(at) \equiv \mathrm{ptr}(p) \qquad \Xi \vdash p : \Psi_2}{\Gamma; \Xi \vdash \mathrm{mov}(\mathrm{uniforms\_address}, a) : \Psi_1 \to \Psi_2 \mid u \to \mathrm{u2}}$$

$$\frac{\Gamma \vdash rf : at \qquad fst(at) \equiv \mathrm{ptr}(p) \qquad \Xi \vdash p : \Psi_2 \qquad rf \not\equiv wr_{before}}{\Gamma;\ \Xi\ ; wr_{before} \vdash \mathrm{mov}(\mathrm{uniforms\_address}, rf) : \Psi_1 \to \Psi_2 \mid u \to \mathrm{u2}}$$

$$\frac{\vdash \mathrm{uniform} : \mathrm{ptr}(p) \mid \Psi_1 \to \Psi_2 \qquad \Xi \vdash p : \Psi_3}{\Xi \vdash \mathrm{mov}(\mathrm{uniforms\_address}, \mathrm{uniform}) : \Psi_1 \to \Psi_3 \mid \mathrm{u0} \to \mathrm{u2}}$$

$$\frac{\Pi \vdash \mathrm{vpmr} : at \mid vpmq_1 \to vpmq_2 \qquad fst(at) \equiv \mathrm{ptr}(p) \qquad \Xi \vdash p : \Psi_2}{\Pi; \Xi \vdash \mathrm{mov}(\mathrm{uniforms\_address}, \mathrm{vpmr}) : vpmq_1 \to vpmq_2 \mid \Psi_1 \to \Psi_2 \mid u \to \mathrm{u2}}$$

$$\frac{\Gamma \vdash a : at \qquad notfull\left(\Sigma^{tmu}\right)}{\Gamma; \Xi \vdash \mathrm{mov}(tmu, a) : \Sigma^{tmu} \to map\left(unwrap_{\Xi}, at\right) :: \Sigma^{tmu}}$$

$$\frac{\Gamma \vdash rf : at \qquad notfull\left(\Sigma^{tmu}\right) \qquad rf \not\equiv wr_{before}}{\Gamma;\ \Xi\ ; wr_{before} \vdash \mathrm{mov}(tmu, rf) : \Sigma^{tmu} \to map\left(unwrap_{\Xi}, at\right) :: \Sigma^{tmu}}$$

$$\frac{\vdash \mathrm{uniform} : \mathrm{ptr}(p) \ \mid\ \Psi_1 \to \Psi_2 \qquad \Xi \vdash p : b \qquad notfull\left(\Sigma^{tmu}\right)}{\mathrm{u0}; \Xi \vdash \mathrm{mov}(tmu, \mathrm{uniform}) : \Sigma^{tmu} \to array(b) :: \Sigma^{tmu} \ \mid \Psi_1 \to \Psi_2}$$

$$\frac{\Pi \vdash \mathrm{vpmr} : at \mid vpmq_1 \to vpmq_2 \qquad notfull\left(\Sigma^{tmu}\right)}{\Pi; \Xi \vdash \mathrm{mov}(tmu, \mathrm{vpmr}) : vpmq_1 \to vpmq_2 \mid \Sigma^{tmu} \to map\left(unwrap_{\Xi}, at\right) :: \Sigma^{tmu}}$$

$$\frac{\vdash n : b}{\vdash \mathrm{mov}(\mathrm{vpmw}, n) : \Pi \to \Pi\left[Y \mapsto array(b)\right] \mid Y \to inc(Y)}$$

$$\frac{\Gamma \vdash urr : at}{\vdash \mathrm{mov}(\mathrm{vpmw}, urr) : \Pi \to \Pi[Y \mapsto at] \mid Y \to inc(Y)}$$

$$\frac{\Gamma \vdash rf : at \qquad rf \not\equiv wr_{before}}{wr_{before} \vdash \mathrm{mov}(\mathrm{vpmw}, rf) : \Pi \to \Pi[Y \mapsto at] \mid Y \to inc(Y)}$$

$$\frac{\vdash \text{uniform} : b \mid \Psi_1 \to \Psi_2}{\text{u0} \vdash \text{mov}(\text{vpmw}, \text{uniform}) : \Pi \to \Pi\big[Y \mapsto \mathit{array}(b)\big] \mid \Psi_1 \to \Psi_2 \mid Y \to \mathit{inc}(Y)}$$

$$\frac{\Pi \vdash Y : at}{\Pi \vdash \text{setup\_vpm\_read}\Big(Y, i4^*\Big) : \mathit{vpmq} \to \Big(Y, i4^*\Big) :: \mathit{vpmq}}$$

$$\vdash \text{setup\_vpm\_write}(Y) : \mathit{vpmw\_addr} \to Y$$

$$\frac{i + i4^* \le 64}{\vdash \text{setup\_dma\_load}\Big(\text{y}_i, i4^*\Big) : \mathit{DLS} \to \Big(\text{y}_i, i4^*\Big)}$$

$$\frac{\Gamma \vdash rr : at \qquad \mathit{fst}(at) \equiv \text{ptr}(p) \qquad \Xi \vdash p : \tau \qquad i4^* \times 16 \times 4 \le \mathit{size\_of}(\tau)}{\Big(Y, i4^*\Big); \Xi \vdash \text{start\_dma\_load}(rr) : \varepsilon \to \Big(Y, i4^*, \tau\Big)}$$

$$\vdash \text{wait\_dma\_load} : \Big(Y, i4^*, \tau\Big) \to \varepsilon \mid \Pi \to \Pi\ \mathit{dma\_load}\Big(Y, i4^*, \tau\Big)$$

$$\frac{i + i7^* \le 64}{\vdash \text{setup\_dma\_store}\Big(\text{y}_i, i7^*\Big) : \mathit{DSS} \to \Big(\text{y}_i, i7^*\Big)}$$

$$\frac{\Gamma \vdash rr : at_1 \qquad \mathit{fst}(at_1) \equiv \text{ptr}(p) \qquad \Omega \vdash p : \tau \qquad \Pi \vdash \Big(Y, i7^*\Big) : \tau_1 \qquad \tau_1 <: \tau}{\Big(Y, i7^*\Big); \Omega \vdash \text{start\_dma\_store}(rr) : \varepsilon \to \text{type\_preserving}}$$

$$\frac{\Gamma \vdash rr : at_1 \qquad \mathit{fst}(at_1) \equiv \text{ptr}(p) \qquad \Pi \vdash \Big(Y, i7^*\Big) : \tau \qquad p \notin \mathit{dom}(\Omega) \qquad \vdash \Omega[p \mapsto \tau]}{\Big(Y, i7^*\Big); \Omega \vdash \text{start\_dma\_store}(rr) : \varepsilon \to (p, \tau)}$$

$$\vdash \text{wait\_dma\_store} : \text{type\_preserving} \to \varepsilon$$

$$\vdash \text{wait\_dma\_store} : (p, \tau) \to \varepsilon \mid \Omega \to \Omega[p \mapsto \tau]$$

conditional instructions:

$$\frac{\begin{array}{c}\vdash \iota : \Gamma_1 \to \Gamma_2 \mid u_1 \to u_2 \mid \Psi_1 \to \Psi_2 \mid \Sigma_1^{tmu} \to \Sigma_2^{tmu} \mid C_1 \to C_2 \mid wr_1 \to wr_2 \\ \big[rr \mapsto at_1\big] \in \Gamma_1 \qquad \Gamma_2 \equiv \Gamma_{21}\big[rr \mapsto at_2\big] \\ at_1 <: at_2\end{array}}{\vdash (\iota, c) : \Gamma_1 \to \Gamma_2 \mid u_1 \to u_2 \mid \Psi_1 \to \Psi_2 \mid \Sigma_1^{tmu} \to \Sigma_2^{tmu} \mid C_1 \to C_2 \mid wr_1 \to wr_2}$$

$$\frac{\vdash \iota : \Theta_1 \to \Theta_2}{\vdash (\iota, \varepsilon) : \Theta_1 \to \Theta_2}$$

signals:

$$\vdash \text{load}\langle tmu \rangle : \Sigma^{tmu} :: at \to \Sigma^{tmu} \mid \Gamma \to \Gamma[\text{r4} \mapsto at]$$

conditional instructions with signals:

$$\frac{\vdash ci : \Theta_1 \to \Theta_2}{\vdash (ci, \varepsilon) : \Theta_1 \to \Theta_2}$$

$$\frac{\begin{array}{c}\vdash ci : \Gamma \to \Gamma[rr_1 \mapsto at_1] \mid u_1 \to u_2 \mid \Psi_1 \to \Psi_2 \mid C_1 \to C_2 \mid \Omega_1 \to \Omega_2 \mid wr_1 \to wr_2 \\ \vdash s : \Gamma \to \Gamma[\text{r4} \mapsto at_2] \mid \Sigma_1^{tmu} \to \Sigma_2^{tmu}\end{array}}{\begin{array}{c}\vdash (ci, s) : \Gamma \to \Gamma[rr_1 \mapsto at_1][\text{r4} \mapsto at_2] \mid u_1 \to u_2 \mid \Psi_1 \to \Psi_2 \\ \mid \Sigma_1^{tmu} \to \Sigma_2^{tmu} \mid C_1 \to C_2 \mid \Omega_1 \to \Omega_2 \mid wr_1 \to wr_2\end{array}}$$

labels:

$$\frac{l \in dom(\Phi)}{\Phi \vdash l : \Phi(l)}$$

instruction sequences:

$$\frac{\begin{array}{c}\vdash csi : \Theta_1 \to \Theta_2 \\ \Phi \vdash I : \Theta_2\end{array}}{\Phi \vdash csi \; ; I : \Theta_1}$$

$$\frac{\begin{array}{c}\vdash csi_1 : \Theta_1 \to \Theta_2 \\ \vdash csi_2 : \Theta_2 \to \Theta_3 \\ \vdash csi_3 : \Theta_3 \to \Theta_4 \\ \Phi \vdash l : \text{code}(\Theta_5) \\ \Theta_4 <: \Theta_5\end{array}}{\Phi \vdash \text{jmp}(l) \; ; csi_1 \; ; csi_2 \; ; csi_3 : \Theta_1}$$

$$\frac{\begin{array}{c}\vdash ci : \Theta_{11} \to \Theta_2 \qquad \vdash csi_1 : \Theta_2 \to \Theta_3 \\ \vdash csi_2 : \Theta_3 \to \Theta_4 \\ regctx(\Theta_2) \equiv regctx(\Theta_{11})[rwr \mapsto at] \\ rwr \not\equiv \mathrm{ra}n \\ rwr \not\equiv \mathrm{rb}n \qquad \Theta_1 \equiv \left(\Gamma_1, \Psi_1, u_1, \Sigma_1^{tmu}, C_1, \Omega_1, wr_1\right) \\ \Theta_{11} \equiv \left(\Gamma_1 \setminus \{\mathrm{ra14}, \mathrm{rb14}\}, \Psi_1, u_1, \Sigma_1^{tmu}, C_1, \Omega_1, wr_1\right) \\ \{\mathrm{ra14}, \mathrm{rb14}\} \notin regctx(\Theta_2) \\ \{\mathrm{ra14}, \mathrm{rb14}\} \notin regctx(\Theta_3) \\ \{\mathrm{ra14}, \mathrm{rb14}\} \notin regctx(\Theta_4)\end{array}}{\vdash (ci, \mathrm{thread\_end}) \ ; \ csi_1 \ ; \ csi_2 : \Theta_1}$$

$$\frac{\begin{array}{c}\vdash csi_1 : \Theta_1 \to \Theta_2 \qquad \vdash csi_2 : \Theta_2 \to \Theta_3 \\ \vdash csi_3 : \Theta_3 \to \Theta_4 \\ \Phi \vdash l : \mathrm{code}(\Theta_4) \qquad \Phi \vdash I : \Theta_5 \\ \Theta_4 <: \Theta_5\end{array}}{\Phi \vdash \mathrm{if}\ qc\ \mathrm{jmp}(l) \ ; \ csi_1 \ ; \ csi_2 \ ; \ csi_3 \ ; \ I : \Theta_1}$$

programs:

$$\frac{\forall l \in dom(P) \,.\, \Phi \vdash P(l) : \Theta_l}{\Phi \vdash P}$$

It is defined that when the elements of at all have the same basic type b, it is convertible with $\mathrm{vec}(b, 16)$.

2.4. Auxiliary functions

Note that all free meta-variables are assumed to be fresh.

$$notfull\left(\Sigma^{tmu}\right) \stackrel{\mathrm{def}}{=} \left(\Sigma^{tmu} \not\equiv at_1 :: at_2 :: at_3 :: at_4\right)$$

$$unwrap_{\Xi}\left(\mathrm{ptr}(p)\right) \stackrel{\mathrm{def}}{=} \Xi(p) \ \text{ if } p \in dom(\Xi)$$

$$fst\left([b_0, b_1, ..., b_{15}]\right) \stackrel{\mathrm{def}}{=} b_0$$

$$map\left(f, [b_0, b_1, ..., b_{15}]\right) \stackrel{\mathrm{def}}{=} [f(b_0), f(b_1), ..., f(b_{15})]$$

When an array type has the same 16 basic type, written $array(b)$:

$$array(b) \stackrel{\text{def}}{=} [b, b, ..., b]$$

$$inc\left(\mathrm{y}_{63}\right) \stackrel{\text{def}}{=} \mathrm{y}_0$$

$$inc\left(\mathrm{y}_n\right) \stackrel{\text{def}}{=} \mathrm{y}_{n+1} \text{ if } 0 \leq n \leq 62$$

$$regctx\big((\Gamma, \Psi, u, \Sigma, C, \Omega, wr)\big) \stackrel{\text{def}}{=} \Gamma$$

$size\_of(\tau)$ represents the size of a value of τ in bytes.

$$size\_of(b) \stackrel{\text{def}}{=} 4$$

$$size\_of(at) \stackrel{\text{def}}{=} 16 \times 4$$

$$size\_of(b \circ \Psi) \stackrel{\text{def}}{=} 4 + size\_of(\Psi)$$

$$size\_of(\varepsilon) \stackrel{\text{def}}{=} 0$$

$$size\_of\big(\mathrm{vec}(b, n)\big) \stackrel{\text{def}}{=} size\_of(b) \times n$$

$$dma\_load(Y, 1, \tau) \stackrel{\text{def}}{=} \left[Y \mapsto truncate(\tau)\right]$$

$$dma\_load\left(Y, i4^*, at\right) \stackrel{\text{def}}{=} [Y \mapsto at] \text{ if } i4^* \geq 2$$

$$truncate(at) \stackrel{\text{def}}{=} at$$

$$concat\big(\mathrm{vec}(b, m), \mathrm{vec}(b, n)\big) \stackrel{\text{def}}{=} \mathrm{vec}(b, m + n)$$

3. Future

- Any properties are not proved.
- There are many implicitness.
- The current definition is so conservative that it cannot serve practical use.
- The current definition may be incorrect or inconsistent.